Complete notes on

Operating System

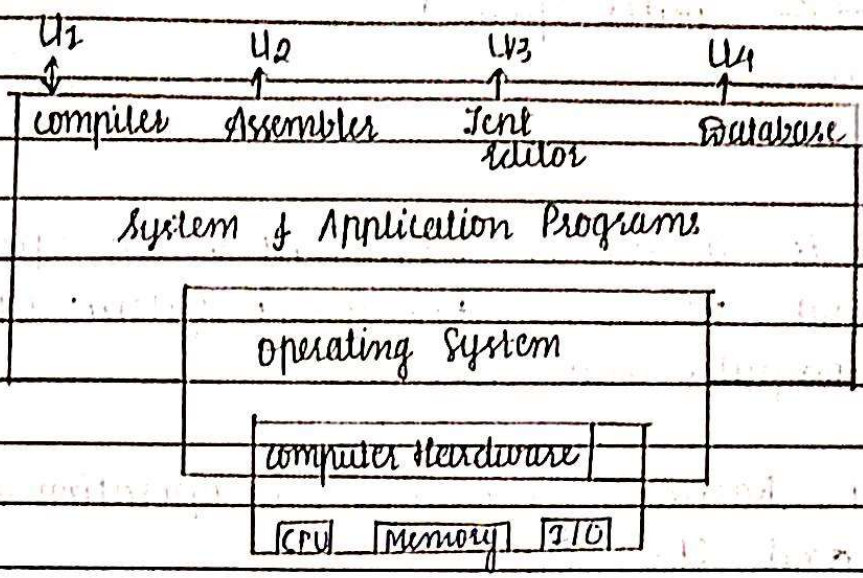
Written By:-

Sanskriti Satpute

(Software engineer)

DATE : / /

OPERATING SYSTEMS



Structure of computer System.

- Computer systems can be divided into types.

1] Hardware - It provides basic computing resources. (CPU | memory | IO) devices.

2] OS operating system : O.S controls & co-ordinates the use of hardware among various app. & users.

3] Application programs : defines the way in which system resources are used to solve computing of users (compilers ..) (assemblers, editor, database)

4] Users : people, machine (printer, scanner or other computers)

DATE : / /

PAGE No :

Os prioritizes the work of application (eg- phone in letter game)

Kernel -

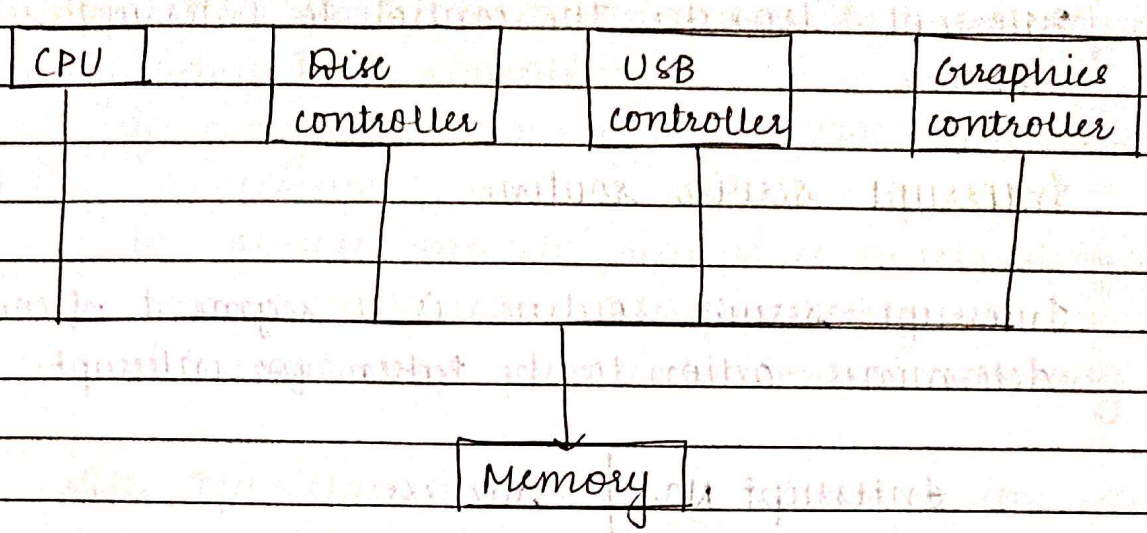
- It is a control component of an operating system that manages ~~operating system~~ operations of computer and hardware.
- It basically manages the operation of memory and CPU type.
- It is ~~more~~ core component of OS.
- It acts as bridge between applied & data processing performed at hardware level using interprocess communication & system call.

Bootstrap Program.

- It is a program which is loaded at power up or reboot.
- It is typically stored in ROM & EPROM
- It is generally known as firmware
- It initialises all aspects of system.
- It loads the operating system kernel & start its execution.

DATE : __/__/____ PAGE No : _____

Computer system organisation



1] one or more CPUs, device controllers connects through common bus providing access to shared memory concurrent execution of CPUs & devices competes for memory cycles.

2] I/O devices & CPU can execute concurrently

3] Each device controller is incharge of particular device type

4] Each device controller has local buffer

5] CPU moves the data from/ I/O to main memory to/from local buffer

6] I/O is from device to local buffer of controller

7] Device controller informs CPU that it has finish its operation by causing an interrupt.

DATE : __/__/__ PAGE No :

Interrupts

Interrupt transfer the control to interrupt service routine.

Interrupt service routine

Interrupt service routine is a segment of code that determines action to be taken for interrupt.

Interrupt No.	Address
0	003h
1	002h
2	hexadecimal
3	

Determining Type of Interrupt

- Polling -

same interrupt handler called for all interrupts; which then polls all devices to figure out the reason for interrupt.

2] Interrupt vector Table

Different interrupt handler will be executed for diff interrupts

DATE : __/__/__ PAGE No : _____

I/O Structure

1] After I/O starts, control returns to user programs only upon I/O completion.
   a] wait instruction - idols the CPU until next interrupt.
   b] Almost one I/O request is outstanding at a time no simultaneous I/O processing.
   c] wait loop (contentsation for memory access)

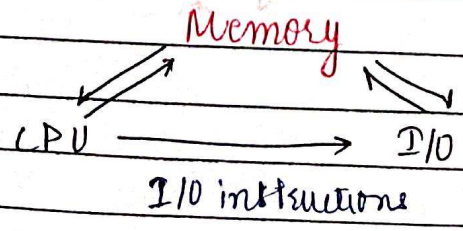
2] After I/O starts, the control returns to the user program w/o waiting for I/O completion.

System call - request to O.S to allow user to wait for I/O completion.

Device status Table
a] contains entry of each I/O device indicating its type address & state.

b] OS indexes into I/O devices table to determine device status & to modify table entry to include the interrupt.

Direct Memory Access.



DATE : __/__/____ PAGE No :

1] Typically used for I/O devices with lot of data to transfer in order to reduce load on CPU

2] Device controller transforms the block of data from buffer storage directly to main memory w/o CPU interruption.

3] Device controller interrupts the cpu on completion of I/O.

Storage Hierarchy.

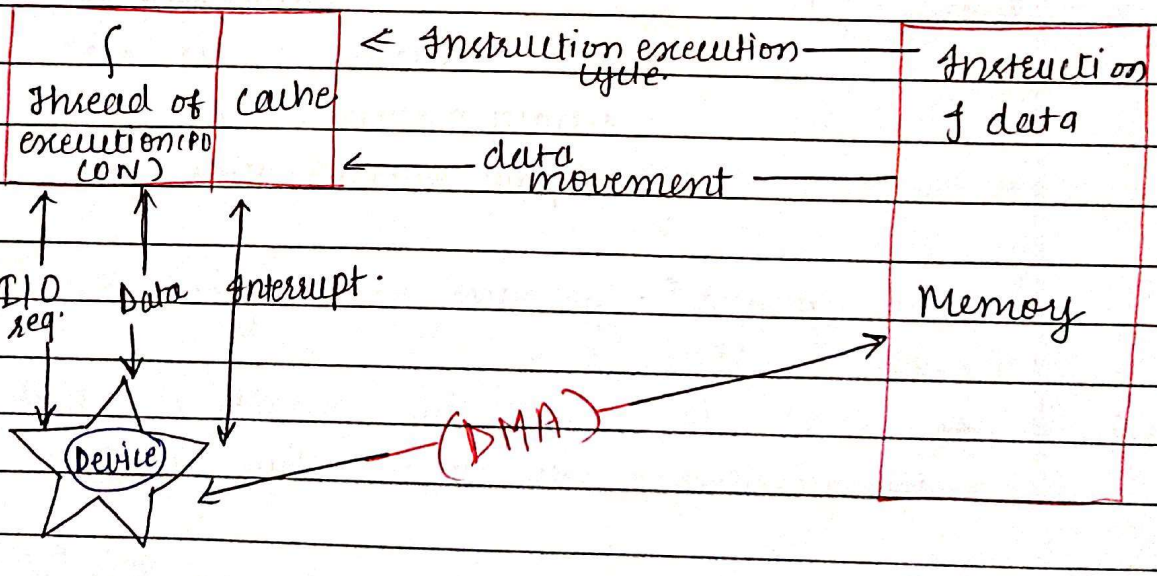
Registers
↑
cache memory
↑
main memory
↑
(SSDs) ← electronic disc
↑
magnetic disc
↑
optical disc
↑
magnetic tapes
↑

DATE : __/__/__ PAGE No : _____

Caching

1] It is an imp principle performed at many levels in comp. (hardware, software, O.S)
2] Information in use is copied from slower to faster storage temporarily.
3] Faster storage (cache) check first to determine if information is there.
   a] If it is, informationalised directly from cache.
   b] If not, data copied to cache & used there.
4] Cache are smaller than the storage being cached.
5] O.S manages the cache size & replacement policy.

★ Computer System Architecture.



DATE : 20 / 1 / 22.

PAGE No :

Kernal entry points

→ what are diff event in OS.

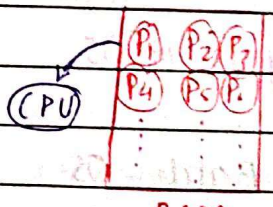
In computing environment event in an action or occurence detected by the program that may be handled by the program.

System call

20/01/22.

Types of operating system.

1] Multiprogrammed OS



→ Non premptive
→ Idleness is decreased.
→ If CPU will execute a process, it will execute completely.

2] Multitasking / Time sharing → mostly used. (Round Robin)

→ Time is shared, will execute the process for specific question, will execute the remaining after.

→ not ideal

→ But response time is decreased.

DATE : __/__/__

3] Real time OS < Hard / Soft

—xox—

21|01|22

Introduction to OS.

- An OS is a program that manages computer hardware.
- It also provides a basis for application programs and acts as an intermediary betn computer users & hardware.

→ windows, linux, ubuntu, Mac OS x IOS, Andio

Types of OS.

i] Batch OS
ii] Time sharing OS
iii] Distributed OS
iv] Network OS
v] Real time OS
vi] Multiprogramming
vii] Multiprocessing / tasking

Goals of OS

i] convenience
ii] Efficiency
iii] Both.

→ Functions.

- It is an interference betn User & Hardware
- Alloration of Resources
→ Management of memory, security etc.

DATE : __/__/____ PAGE No : _____

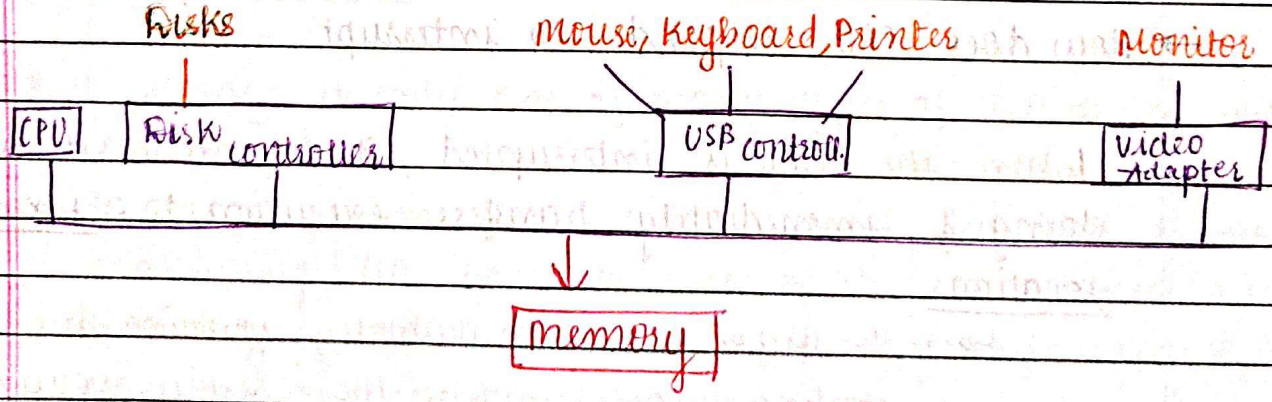
Basic of operating system (computer system operation)

→ a modern general-purpose computer system consists of one or more CPUs and a no. of device controllers connected through a common bus that provided access to shared memory.

→ Each device controller is in charge of a specific type of device.

→ The CPU and the device controllers can execute concurrently, competing for memory cycles.

→ To ensure orderly access to all the shared memory, memory controller is provided whose fun' is to synchronize access to the memory.



Some important terms

1) Bootstrap Program- The initial program that runs when a computer is powered up or rebooted.
   - It is stored in the ROM (secondary mem
   - It must know how to load the OS

• and start executing that system.
• It must locate and load into memory the OS kernel.
(heart of OS)

2] Interrupt

• The occurence of an event is usually signalled by an Interrupt from Hardware or Software
• Hardware may trigger an interrupt at any time by sending a signal to the CPU, usually the way of system bus.

3] System call (monitor call)

• Software may trigger by executing a special operation called system call.

→ How does CPU responds to interrupt

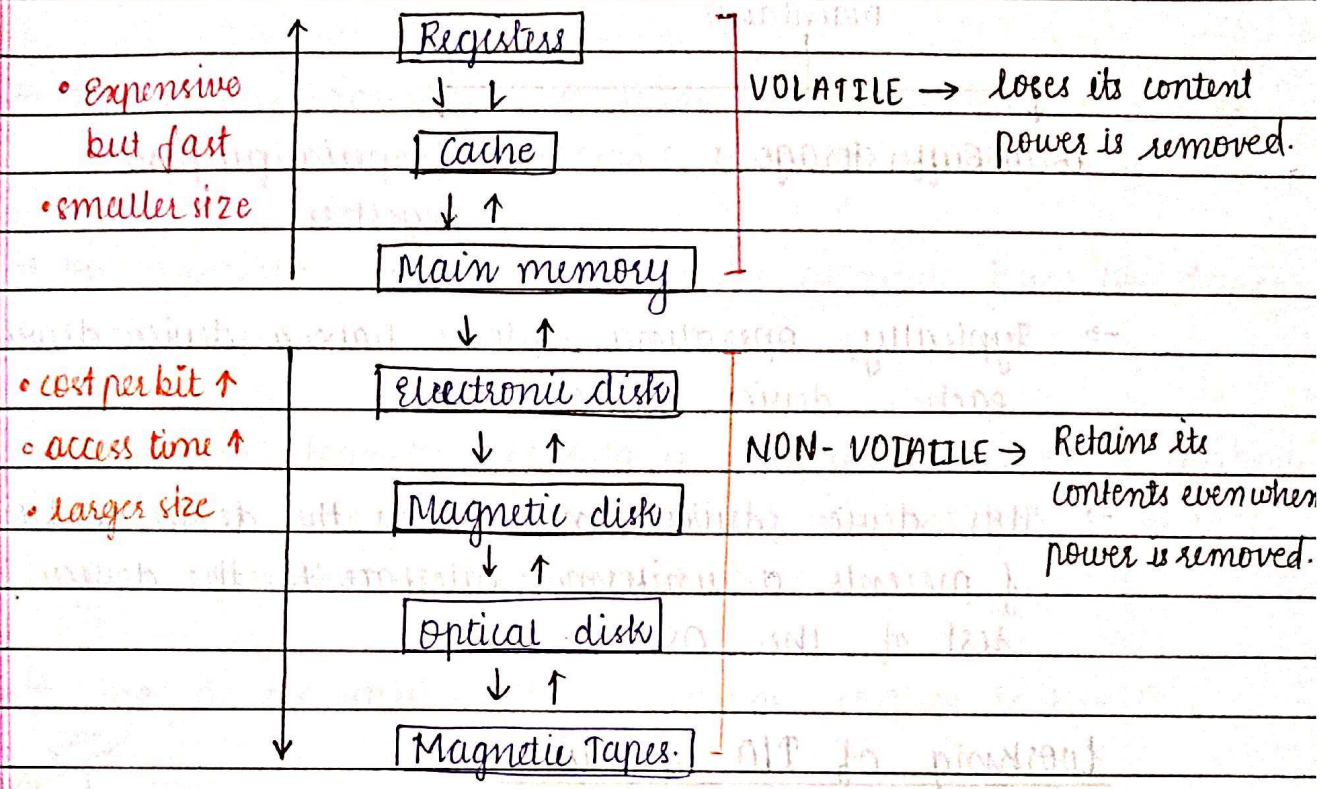
When the CPU is interrupted, it stops what it is doing & immediately transfers execution to a fixed location.
↳ The fixed location usually contains the starting address where the service routine of the interrupt is located. (place where what interrupt want to do is written)
Interrupt service routine executes

on compl^{n}, the CPU resumes the interrupted computation

DATE : __/__/____ PAGE No : _____

Basic of OS [storage structure]

☆ special case (NVRAM → non volatile)



(I/O Structure)

→ A storage is only one of many types of I/O devices within a computer.

→ A large portion of OS. code is ~~deed~~ dedicated to ~~mang~~ managing I/O both because of its importance to the realibility and performance of a system & because of the varying nature of the devices.

→ A general purpose comp. system of CPUs & multiple devices controllers that are connected through a common bus.

→ Each

DATE : __/__/____

Each device controller is in charge of a specific type of device

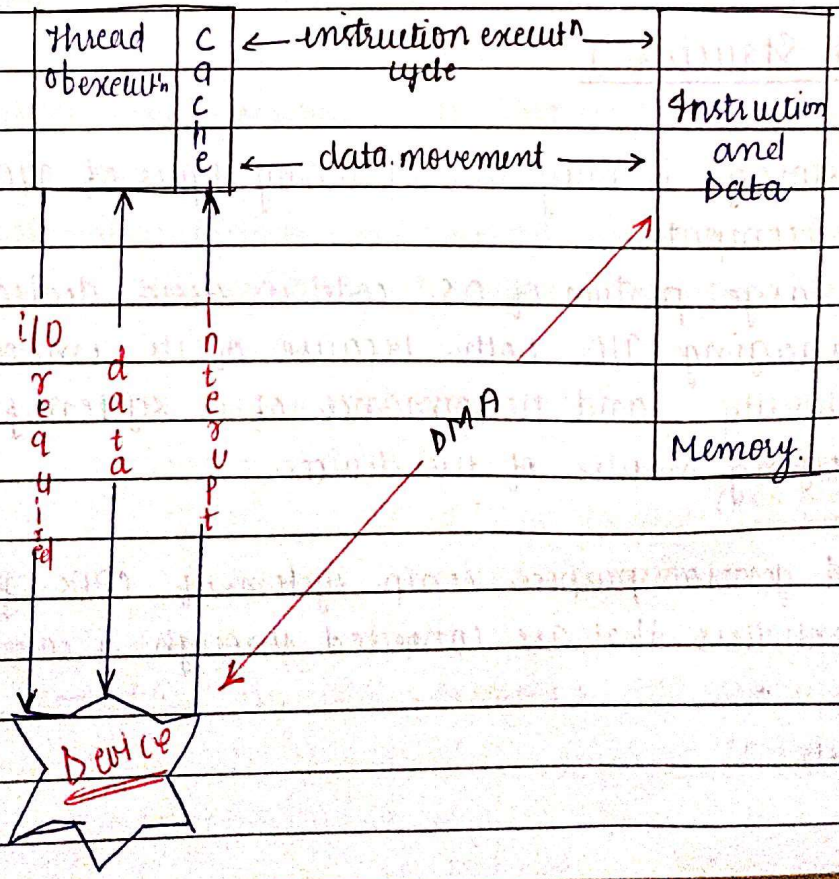
maintains

- local Buffer storage
- set of special purpose Registers

→ Typically, operating systems have a device driver for each device controller.

→ This device driver understands the device controller & presents a uniform interface to the device to the rest of the OS.

Working of I/O operation



DATE : __/__/____ PAGE No : _____

→ To start an I/O operation, the device driver loads the appropriate registers within the device controllers.

→ The device controller, in turn, examines the contents of these registers to determine what action to take.

→ The controller starts the transfer of data from the device to its local buffer.

→ Once the transfer of data is complete, the device controller informs the device driver via an interrupt that it has finished its operation.

→ The device driver then returns control to the OS.

Dis Adv.

→ This forms of interrupt-driven I/O is fine for moving small amounts of data but can produce high overhead when used for bulk data movement.

- To solve this prob. Direct Memory Access is used.

→ after setting up buffers, pointers and counters for the I/O device, the device controller transfers entire block of data directly to or from its own buffer storage to memory, with no intervention by the CPU.

→ Only one interrupt per block, to tell the device driver the op. is completed.

→ while the device controller is performing these op. the CPU is available accomplish other work.

DATE : __/__/__ PAGE No. :

Computer System Architecture

Types of computer system on no. of general purpose processors :-

1] Single Processor System.

→ one main CPU capable of executing a general purpose instruction set including instructions from user processes.

→ Other special purpose processors are also present which perform device specific tasks
(micro processor chip in keyword)

2] Multi Processor System / Parallel or tightly coupled sy

→ Has two or more processors in close communication sharing the computer bus and sometimes the clock, memory and pheripheral devices.

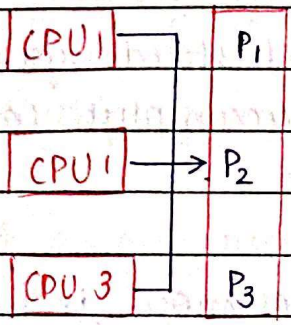
Adv.

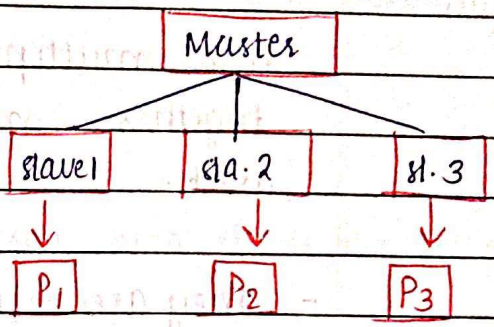
→ increased throughput (Measure of perform. of system
→ economy of scale. (as here sharing of resources)
→ increased reliability.

DATE : __/__/____ PAGE No : _____

→ Types of multiprocessor system.

- Symmetric
- Asymtric Multiprocessing





→ all the processors are similar to each other

→ commonly used

→ all processors share physical memory.

Master → slave approach

The master allocates processor schedules & allocate the work to the worker processors

→ many processes can run simultaneously (N CPUs → N Processes)	→ Multiprocessing adds CPUs to increase computing power.
→ one CPU sits idle while other overload	→ M.P can cause a system to change its memory access model from uniform memory access to (NUMA)

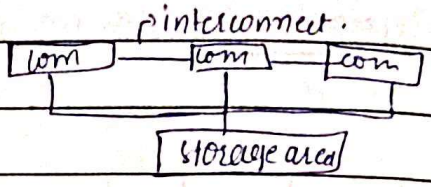
→ UMA is defined as the situation in which access to RAM from any CPU takes the same amount of time.

★ All multicore sys → multiprocessors (not vice versa).

DATE : __/__/____ PAGE No :

- Blade servers → multi processor board, I/O boards & network boards are placed in the same chassis.

3| Clustered System.



- like multiprocessor systems, clustered systems gather together multiple CPUs to accomplish computational work.

- They are composed of two or more individual systems coupled together.

→ Provides high availability.

→ can be structured :

asymmetrically
↓
- One machine in Hot-standby mode.
- Other runs application.

✔(more efficient)
symmetrically
↓
- Two or more host run application
- Monitor each other.

- Other forms are parallel clusters → multiple. allows host to access the same data on shared storage.

Clustering over WAN

DATE : __/__/____ PAGE No : _____

Operating System Structures

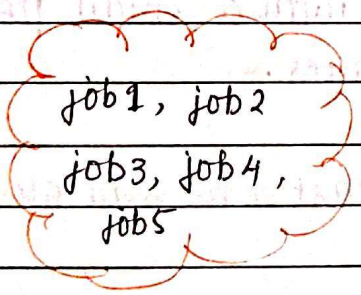
→ (Multiprogramming & Multitasking) → commonalities

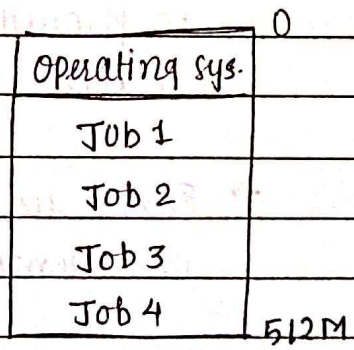
→ Operating system vary greatly in their makeup internally

i) Multiprogramming.

→ A single user cannot, in general, keep either the CPU or the I/O devices busy at all times.

→ Multiprogramming increases CPU utilization by organizing jobs (code and data) so that the CPU always has one to execute.





Job pool

→ CPU doesn't remain idle

→ Efficient.

→ Multiprogramming systems provide an environment in which the various system resources (for ex- CPU, memory, and peripheral devices) are utilized effectively but not provide interaction with computer system.

DATE : __/__/__

iii] Time Sharing (Multitasking)

→ CPU executes multiple jobs by switching among them (virtual memory)

→ Switches occur so frequently that the users can interact with each program while it is running

→ Time sharing requires an interactive (or hand on) computer system, which provides direct communication bet^n the user & the system.

→ A time shared system allows many users to share the computer simultaneously.

→ Uses CPU scheduling & multiprogramming to provide each user with a small portion of a time-shared computer.

→ Each user has at least one seperate program in memory.

→ A program loaded into memory & executing is called a "PROCESS"

Operating System Services

- An OS provides an environment for the execution of programs
- It provides certain services to program and to users of those programs.

DATE : __/__/____ PAGE No :

iii) Time Sharing (Multitasking)

→ CPU executes multiple jobs by switching among them (virtual memory)

→ Switches occur so frequently that the users can interact with each program while it is running.

→ Time sharing requires an interactive (or hand on) computer system, which provides direct communication bet^n the user & the system.

→ A time shared system allows many users to share the computer simultaneously.

→ Uses CPU scheduling & multiprogramming to provide each user with a small portion of a time-shared computer.

→ Each user has at least one seperate program in memory.

→ A program loaded into memory & executing is called a "PROCESS"

Operating System Services

- An OS provides an environment for the execution of programs
- It provides certain services to program and to users of those programs.

DATE : __/__/____ PAGE No : _____

i] User Interface.

→ interaction betⁿ users & computer
ex- command line , graphic

ii] Program Execution

source code → compiler → objectcode → executor → output

iii] I/O operations

Ex- Keyboards → CPU → monitor

iv] File system manipulation.

create, delete, modify, search, access restriction of file.

v] Communications

vi] Error detection

vii] Resource Allocation

viii] Accounting

ix] Protection & Security.

DATE : __/__/____ PAGE No :

User Operating System Interface

i) Provide a common-line interface or command Interpretes that allows users to directly enter commands that are to be performed by the operating system.

ii) Allows the user to interface with the O.S via a graphical user Interference or GUI

Command Interpreter

- Some OS include the command interpreter in the kernel.
- Others, such as Windows XP & UNIX, treat the command interpreter as a special program.
  ex → in windows → code prompt
- On systems, multi command interpreters to choose ~~form~~ from, the interpreters are called as shells
  Eg:-
  Bourne shell, C shell, Bourne-Again shell, Korn shell.

Example:
-We want to perform a task → create a file
first approach → code included in command interpreter itself
→ that code will be executed.

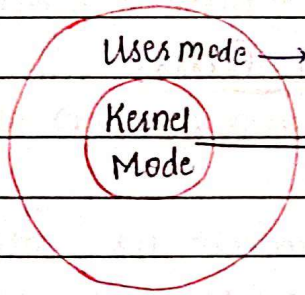
* Command interpreter itself doesn't contain any code But the codes are written in certain program

Second app:
command interpreter calls the program where code is written to create file.

DATE : __/__/__ PAGE No : _____

System Calls

System calls provides an interface to the service made available by an operating system.



(safer mode)

Uses mode → (doesn't have access to memory, hardware resources)

Kernel Mode → priveleged mode

- Strict access
- program is executed in kernel mode crashes down, whole system would crash.

- System call is a programming way in which a computer program requests a service from the kernel of the OS.
- These calls are general available as routines written in C and C++.

Ex:-

Source file → Destination file

> Acquire input filename
> write prompt to screen
> Access input
> Acquire output filename
> write prompt to screen
> Accept input
> Open Input file.
> If file does'nt (ABORT)

→ create output file
If file exists, ABORT
~~Read from input file~~

loop { Read the input file.
write to output file
until Read fails } Close output file
write completion to screen
Terminate normally.

> Types Of System Calls.

(5 types)

- Process control
- File manipulation
- Device management
- Information Maintenance
- Communications

1] Process Control

- end, abort
- load, execute
- create process, terminate process
- get process attributes, set process attributes
- wait for time
- wait event, signal event
- allocate and free memory.

2] File Manipulation

- create file, delete file.
- open, close
- read, write, reposition
- get file attributes, set file attributes

DATE : __/__/__ PAGE No : _____

3] Device Manipulation

- request device, release device
- read, write, reposition
- get device attributes, set device attributes
- logically attach or detach devices.

4] Information Maintenance

- get time or data, set time or data
- get system code data, set system data
- get process, file or device attributes
- set process, file or device attributes.

5] Communication System calls

- create, delete communication connection
- send, receive messages
- transfer status information
- attach or detach remote devices

System Programs

> an important aspect of a modern system is the collection of system programs.

- System programs provide a convenient environment for ~~to~~ program development and execution.

DATE : __/__/____ PAGE No :

• Some of them are simply user interfaces to system calls others are considerably complex.

i) File management

> create > copy > print
> delete > rename > dump
> list & generally manipulate files & directories

ii) Status Information

Ask the system for-

• Data, time
• amount of available memory or disk space.
• Number of user
• detailed performance
• logging & debugging information etc.

iii) File modification.

> several text editors may be available to create & modify the content of files stored on disk or other storage devices.

> There may be also be special commands to search contents of file or perform transformation of the text.

DATE : __/__/__ PAGE No : _____

iv] Programming - language Support

- compilers
- assemblers
- Debuggers &
- Interpreters. (such as C, C++, Java, Visual Basic and PERL)

v] Program loading and Execution

Once a program is assembled or compiled, it must be loaded into memory to be executed.
The system may provide:
- Absolute loaders
- Relocatable loaders
- Linkage editors and
- overlay loaders.

Debugging system for either high-level language or machine languages are needed as well

vi] Communications

These program provide mechanism for:-
i] Creating virtual conn^c among processes, users & com. sys.
ii] Allowing users to send messages to one → another comp.
iii] To browse webpages
iv] To send electronic-mail messages
v] To log in remotely or to transfer files from one → another.

DATE : __/__/____ PAGE No :

Operating System Design & Implementation

- Design goals.

1st problem : Defining goals & specifications

- choice of hardware
- type of system.

Beyond the high design level, the requirement may be much harder to achieve.

Requirements

- Uses goals :
  - convenient to use
  - easy to learn & use
  - reliable, safe & fast

- System goals
(Designer, eng. req.)
  - easy to design, implement
  - maintain, operate.
  - It should be flexible, reliable, error free & efficient

Mechanisms & Policies.

↓ (How to do things)

↓ (determine what will be done)

DATE : __/__/__ PAGE No : _____

* one imp. principle that ; the seperation of policy from mechanism.

Implementations

- once an OS is designed, it must be implemented
- Traditionally, Os have written in assembly language
- Now, however the are most commonly written in higher level languages such as C or C++.

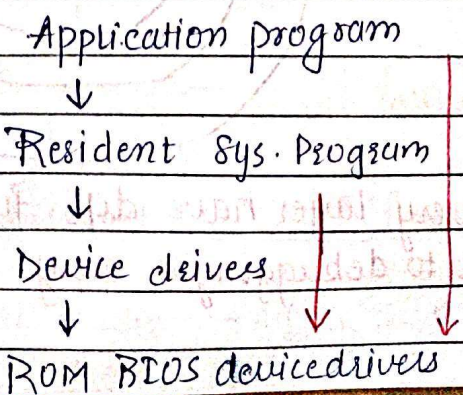
- Advantages.
  - The code can be written fastes
  - It is more compact
  - It is easies to understand
  - It is easier to port

- Disadvantages.
  - written in assembly lang. can't operated on many devices.

  Ex: MS-DOS in Intel 8088 A.L. available only on Intel family of CPUs.

Structure of OS.

Simple Structure



DATE : / /

PAGE No :

Monolithic Structure. → (UNIX)

The users

Shells and commands

compilers and interpreters

System libraries

System-call interface to the Kernel

Kernel

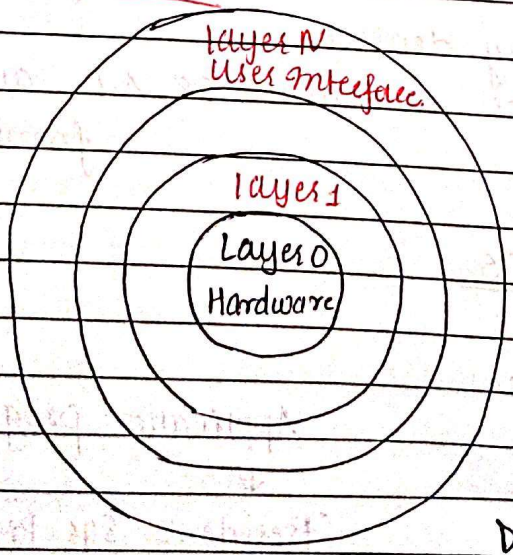
signal, terminal handling	file system	CPU scheduling
character, I/O sys.	swapping block I/O	page replacement
terminal devices	du system	demand paging
	disk & tape drivers	virtual memory

Kernel interface to the hardware

terminal controllers	device controllers	memory controllers
terminals	disk & tapes	physical memory

diff. maintain & implemented.

Layered structure.



Adv

→ Every layer have diff. func. so to debugging is easy.

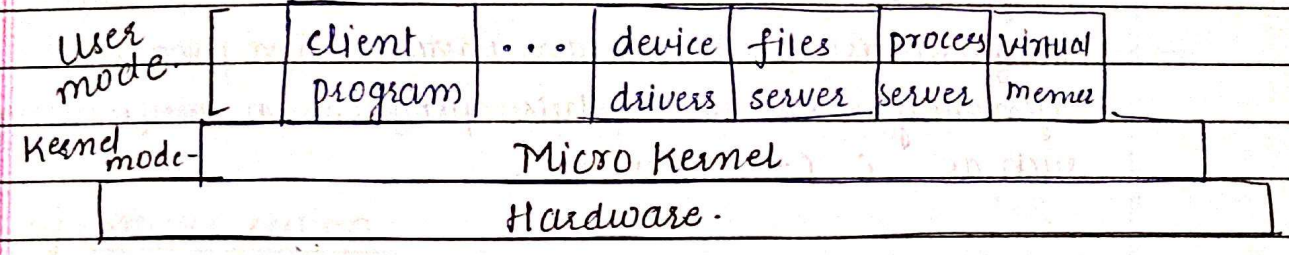
Disadv

Deciding the position of layer as layers can use layer below it

DATE : __/__/____ PAGE No : _____

↦ not every efficient
↦ as it has go layer by layer one by one.

Microkernels



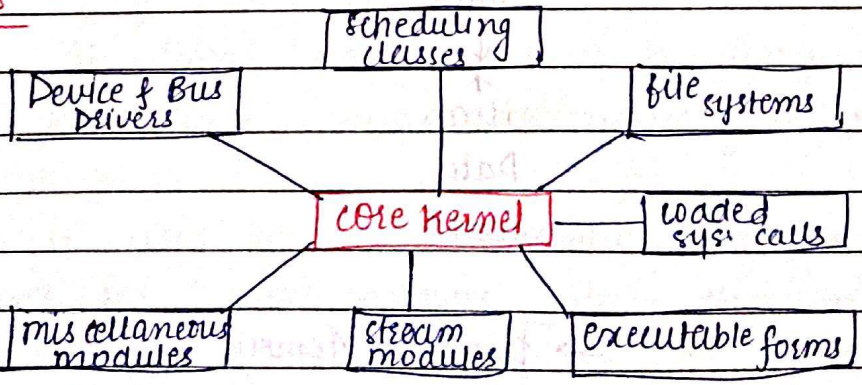
Adv

- To Remove non-essential component from system. (only core imp others implemented as system programs.
- most of the programs are in user mode.
- Crashing of entire program won't happen.

Disadv

- Since the communication needs to be done always, system overhead will take place

Modules



→ resembles micro-kernel
→ each module can communicate directly → By far The Best.
→ loaded directly. no system overhead.

threads, multireading model, goals of rescheduling & diff

DATE: 27/01/22

PAGE No.:

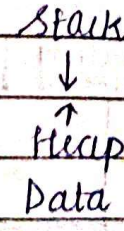
Unit-2 Process & It's Scheduling

2.1] Process

- Program ~~is~~ a under execution.
- Program is consider as an instance of a program
- Generally users writes their program in specific lan such as C, C++, Java etc.

Program → compile → execute → Process

- Program consists of multiple threads, where thread nothing but light weight process.
- Program is a passing entity but process is an activi active entity when it loads into memory.
- Several process may be associated with same progra eg. opening up several instances of



→ Process in Memory

DATE : __/__/__ PAGE

Four sections are -

i) Text

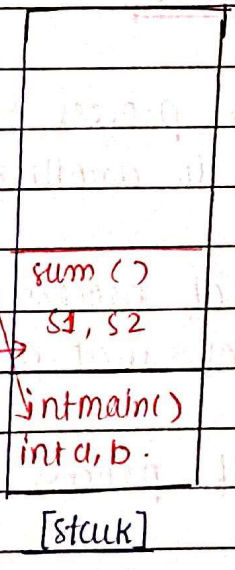
- consists of source code.
- current activity are represented by value of program counter and content of processor's registes.

ii) Stack section.

- stack memory hold temporary data. such as fun', parameters, return address, local variable

Ex -
```
int main()
{ int a, b;
  sum();
}

void sum()
{
  int s1, s2;
}
```



when sum(); is executed, the memory will be cleared from stack memory & would go to main memory.

- In Iteration - no new stack memory will form
- In Recursion - new stack memory is formed

→ Stack is used for local variables, space on the stack is reserved for local variables when they are declared.

Space is freed up when variables go out of scope.

DATE : __/__/____

iii) Data section.

- Global variables, static variables which needs to be allocated & initialized prior to executing main program.

iv) Heap section

It is used for dynamic memory allocation, and is managed via calls to new, delete malloc free etc.

2.2) Process States

- During the process execution, it changes it's state from one state to another state
- The state of process is defined in part by the current activity of that process.
- State of process represents it's current status
- Each process may be in one of the foll" state.

DATE : __/__/____ PAGE No : _____

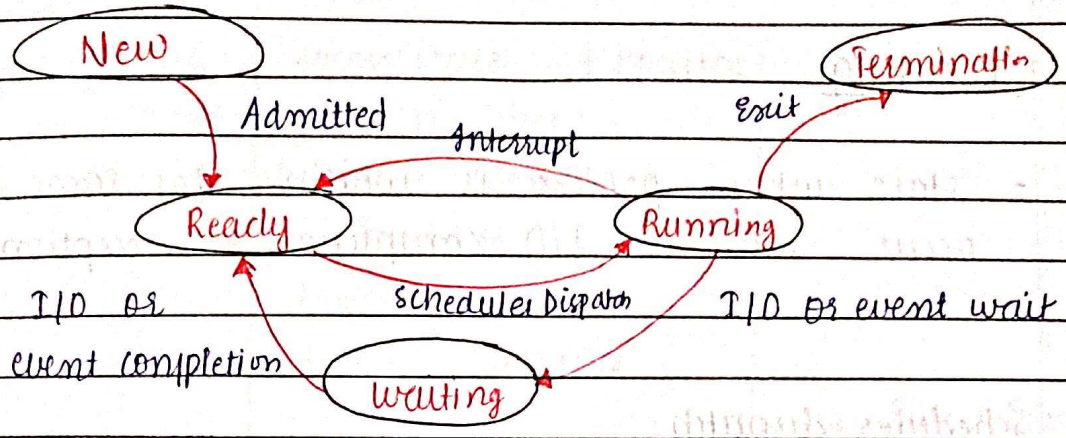


Fig: Process state

i) New

This is the first state of process in which process belongs when it is created

Ready

This is the second state, in which process belongs when it is ready to execute. Scheduler ~~may~~ pick up the processes from ready queue and submit it to CPU for execution

Running

This is the state where actual ~~exec~~ execution of process is carried out. From running state process may go to one of following state.

i) Terminated state: If process successfully executed
ii) Ready state: If interrupt occurs
iii) Waiting state: When any I/O event occurs.

DATE : __/__/____ PAGE No :

Waiting

- State where process is waiting for some event to occur such as I/O completion or reception of a signal.

Scheduler dispatch:

- Scheduler dispatches from ready state to running state.

Terminated.

Process enters into this state when it finishes execution.

2.3] Process Control Block

- The OS needs some information related to process, which is stored in data structure called process control block.

- Each process in OS represented by process control block.

- PCB is also called Task control block
- PCB contains many pieces of information associated with specific process.

DATE : __/__/__ PAGE No : _____

Process state	Pointer
Process number	
Program counter	
Registers	
Memory limits	
list of open files	
- - - - - - - - -	

Fig: Process Control Block

1] Process State

The state may be new, ready, running, waiting and halted.

ii] Program counter

The counter indicates the address of next instⁿ to be executed for this process.

iii] CPU registers.

The registers vary in number and type of depending on computer architecture.
It includes accumulators, index registers, stack pointers and general purpose registers.

iv] Process Number

It represents identifes of the process

v] Memory management information.

It include information as value of base and limit registers, the page tables or segement table based on memory system.

vi] Pointer

It include pointer to next PCB that is pointer to PCB of the next process to run.

vii] I/O Status Information

It includes the list of I/O devices allocated to process a list of open files

2.4] 3. Process Scheduling

- Process scheduling is the technique adopted by OS which governs the order in which process need to be executed.
- An order to achieve multiprogramming environment in OS ~~which governs~~ we require that some process is always running, which helps in maximizing CPU utilization.
- To achieve this CPU need to switch from one process to another process
- To keep CPU busy we need some entity that assigns Job's to CPU and helps in decreasing idle times of CPU & increasing CPU utilization.

DATE : __/__/__ PAGE No : _____

Scheduling Queue.

- This is the data structure which keep set of processes
- scheduling queue refers to queue of processor or devices.
- when the process enters into system then this process is put into queue.
- O.S maintains diff. types of queues which are as follows:
  i]

i] Job Queues.

Initially when processes enters into system they are always put into queue called as Job queue. It consists of all processes in the system.

2] Ready Queue.

The processes that are residing in main memory and are ready and waiting to execute are kept on list called ready queue.

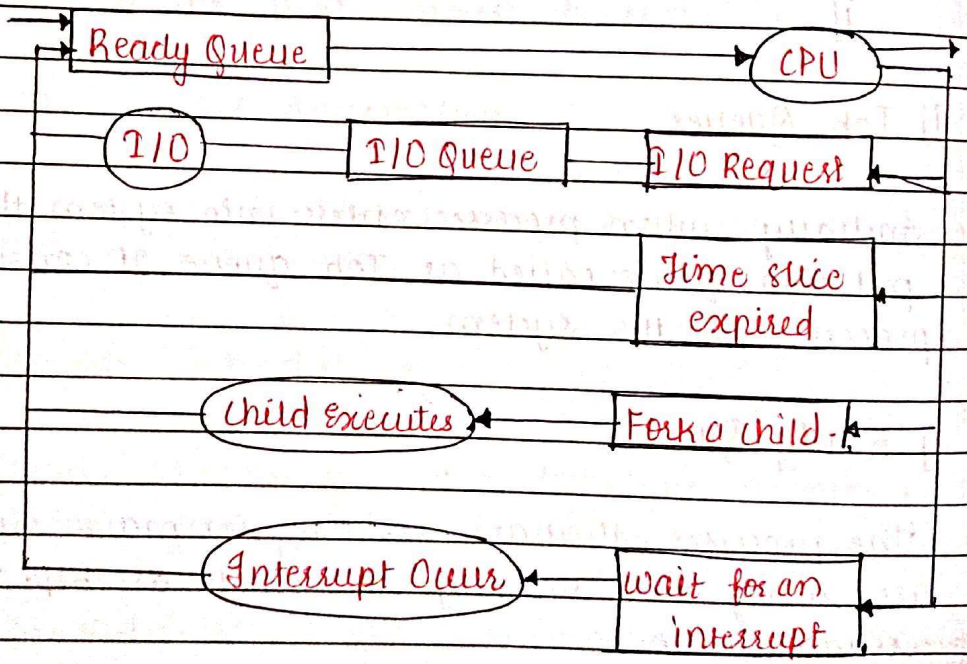
- This queue is generally stored as linked list.
- A ready queue header contains pointers to first and final PCB in the list.

DATE : __/__/____ PAGE No :

(iii) Device queue.

The list of processes waiting for a particular I/O devices is called device Queue. Each device has its own device-queue.

- A common representation of process scheduling is a queuing diagram.



In above fig. rectangular box: represents a box queue as ready queue and device queue.

The circle represents the resources that serve the queue and arrow indicate flow of processes in the system.

DATE : 1 / 02 / 22 PAGE No :

Schedulers.

- Schedular is an entity responsible for arranging the processes in proper order and then submitting it to CPU for execution.

- Schedular arranges the processes into an approximate seq.

- As process migrates among various scheduling queues throughout it is life time, this task of migsation is carried out by scheduler.

- Scheduler in an intermediate level of scheduling.

- There are diff types of scheduler

i] Long Term Schedulers

- a.k a Job schedules.
- responsible for selecting the process from Job queue to ready queue
- Simply the scheduler pick the process from secondary memory and load it into primary memory.
- works as a gate keeper
- Also called as admission schedular.
- It is very important in the large system like super comp.
- It also controls the degree of multiprogramming
- Unix system doesn't make use of long term scheduler.
- Long term schedular executes less frequently.

ii) Short Term Scheduler or CPU scheduler

- It select the process from ready queue that are ready to execute and submit the process to CPU for execution.

- This scheduler temporarily remove process from main memory and submit it to CPU.

- short term scheduler make scheduling decision more frequently.

- This scheduler can be a pre-emptive that means it is capable of forcibly removing process from CPU when it decide to allocate CPU to another process

iii) Medium Term Scheduler.

It select the process from ready queue that are ready to execute and submit the process to CPU for execution

This scheduler temporarily remove process from main memory and submit it to CPU.

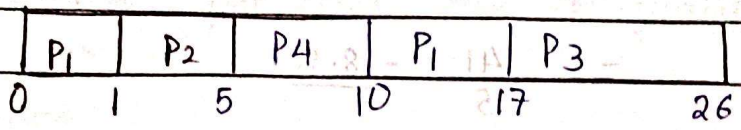
short term schedulers make scheduling decision

ii) Shortest Job First Scheduling.

→ pre-emptive approach.
→ shortest job response approach.

Ex. 3)

Process	Arrival time	Burst time
P_1	0	8
P_2	1	4
P_3	2	9
P_4	3	5



W.T = P_1 = final exe time − Previous time − arrival time.

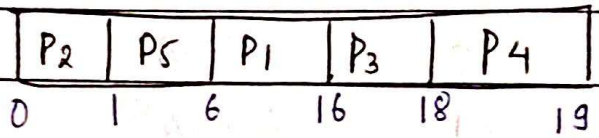
P_1 = 10 − 1 − 0 = 9
P_2 = 1 − 1 = 0
P_3 = 17 − 2 = 15
P_4 = 5 − 3 = 2.

Ex. 4) Priority Scheduling.

Process	Arrival time

↳ P.T.O

Process	Burst time	Priority
P_1	10	3
P_2	1	1
P_3	2	4
P_4	1	5
P_5	5	2



$$A.W.T = \frac{0+1+6+16+18}{5}$$

$$= \frac{41}{5} = 8.2$$

→ Priority can't be same for any processes But Burst and arrival time can be.

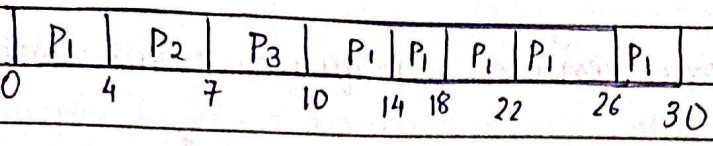
→ Round - Robin Algorithm.

→ Pre-emptive

Process	Burst time
P_1	24
P_2	3
P_3	3.

Time quantum : 4ms

DATE : __/__/____ PAGE No : _____



$P_1 = 26 - 20 = 6$
$P_2 = 4$
$P_3 = 7$
} W.T.

Avg. wt → $\frac{17}{3}$ = 5.66 ms.

Multi-level Queue Scheduling.

- Another type of scheduling algo. created for situation in which processes are ~~d~~ qualified into diff. groups.

~~Groups~~ →

- Two types of processes are ~~foregro~~.
  1) foreground (interactive processes)
  2) Background (Batch processes)

→ System processes → highest Priority processes

→ Interaction processes

→ Interactive editing processes

→ Batch processes

→ ~~Student~~ Student Process. → lowest Priority.

DATE : __/__/__

→ Process once entered an queue, can't switch the queue

Multilevel Feedback Queue Scheduling

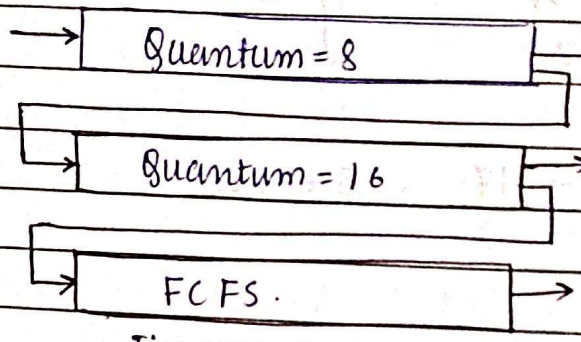


Fig: MFQS

This algo is very similar to ordinary MQSA, except job may move from one → another for variety of reasons. such as.

i] If characteristics of job change betn CPU intensive and I/O intensive.

ii] Job has waited for long time can get bumped into highest priority queue.

iii] Multilevel feed Qs is most flexible as it can be tuned for any situation.

iv] It is most complex to implement because of all adjustable parameters.

v] Some of the parameters which define one of these systems include
a] no. of queues
b] scheduling algo. for each queue

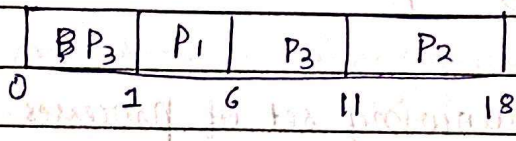
DATE : __/__/____ PAGE No : _____

c) methode to upgrade / ~~degrade~~ demote process from oneque → another
d) method used to determine which a process enters initially

Ex-1

Process	Burst time	Priority	Arrival time	
P_1	5	1	1	Timequant - 2ms.
P_2	7	2	5	
P_3	6	3	0	

① SJF (pre-empitive)



W.T = P_1 = 1 - 1 = 0

P_2 = 11 - 5 = 6

~~P_3 = 11 - 0 = 11~~ P_3 = 6 - 1 - 0 = 5.

Avg. wt = 11/3 = 3.66 ms

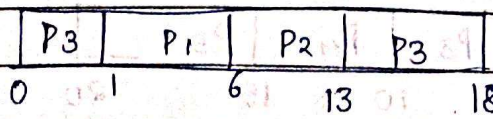
TAT = P_1 = 6 - 1 = 5

P_2 = 18 - 5 = 13

P_3 = 11 - 0 = 11

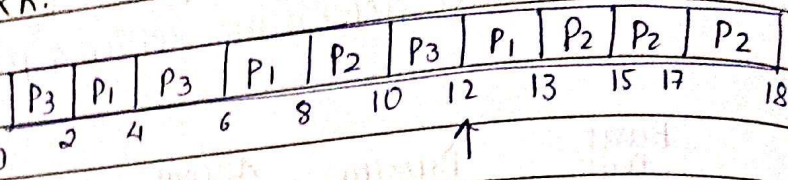
Avg. TAT = 29/3 = 9.6 ms.

② Pre-emptive Priority based schedule.



W.T → P_1 = 0 P_2 = 1 P_3 = 13 - 0 = ~~13~~ 12.

③ RR.



W.T : $P_1 = 12-4-1 = 7$
$P_2 = 17-6-5 = 6$
$P_3 = 10-4-0 = 6.$

Avg. wt = 6.33 ms.

TAT : $P_1 = 13-1 = 12$
$P_2 : 18-5 = 13$
$P_3 = 12-0 = 12$

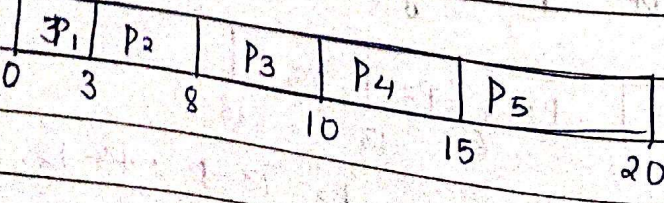
Avg. wt : 12.33 ms.

2] Consider the following set of processes.

Process	Burst time	Arrival time.	Priority
P_1	3	0	5
P_2	5	1	3
P_3	2	2	1
P_4	5	3.	2
P_5	5	4	4

Calculate avg. waiting time & turn around time

|| FCFS.



DATE : __/__/__ PAGE No : _____

waiting time.

$P_1 = 0$

$P_2 = (3-1) = 2$

$P_3 = (8-2) = 6$

$P_4 = (10-3) = 7$

$P_5 = (15-4) = 11$

Average waiting time $= \frac{2+6+7+11}{5}$

$= 5.2$ ms

Turn around time = completion time - submission time.

$P_1 = 3-0 = 3$

$P_3 = 10-2 = 8$

$P_2 = 8-1 = 7$

$P_3 = 10-2 = 8$

$P_4 = 15-3 = 12$

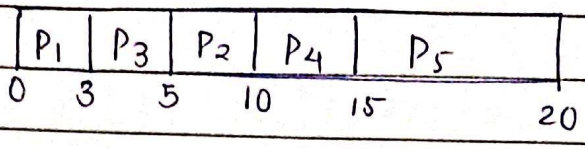
Avg. T.A.T $= \frac{(3+7+8+12+16)}{5}$

$= 9.2$ ms

DATE : __/__/____

PAGE No :

2] SJF.



Gantt Diagram

waiting time for $P_1 = 0$

$P_2 = (5-1) = 4$

$P_3 = (3-2) = 1$

$P_4 = (10-3) = 7$

$P_5 = (15-4) = 11$

Avg. $= \frac{0+4+1+7+11}{5} = \underline{4.6 \text{ ms.}}$

Turn around,

$P_1 = 3-0 = 3$

$P_2 = 10-1 = 9$

$P_3 = 5-2 = 3$

$P_4 = 15-3 = 12$

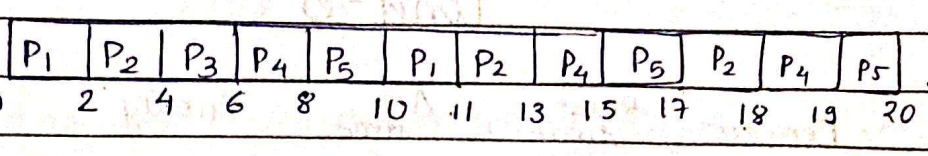
$P_5 = 20-4 = 16.$

Avg. Turnaround time $= \frac{3+3+12+9+16}{5}$

$= \frac{43}{5} = \underline{8.6 \text{ ms}}$

DATE : __/__/__ PAGE No : _____

3] RR (slide) (Pre-emp)



waiting time. $P_1 = (10-2) = 8$

~~$P_2 = (18-4) =$~~ $P_2 = (17-4-1) = 12$ ms

$P_3 = (4-3) = 1$ ms

$P_4 = (18-4-3) = 11$ ms

$P_5 = (19-4-4) = 11$ ms.

Avg. wt time $= \frac{12+1+11+11+8}{5} = \frac{43}{5} = 8.8$ ms

TAT = $P_1 = 11-0 = 11$ ms

$P_2 = 18-1 = 17$ ms

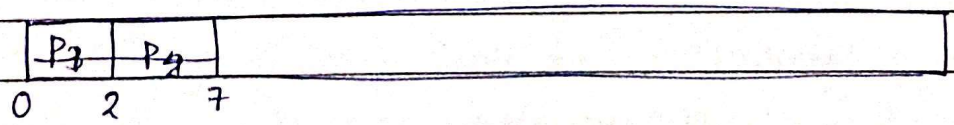
$P_3 = 6-2 = 4$ ms

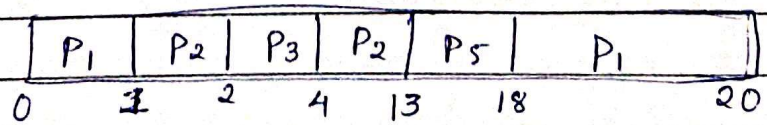
$P_4 = 19-3 = 16$ ms

$P_5 = 20-4 = 16$ ms

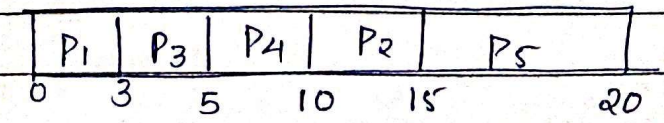
Avg. TAT time $= \frac{11+17+4+16+16}{5} = \frac{64}{5} = 12.8$ ms ~~$= 1.08$ ms~~.

4] Priority Based. (pre-emptive)





Non-preemptive



DATE : 22 / 2 / 22 PAGE No : _____

Unit-3

Process Management & Synchronization

There are 2 two types of process

- Co-operative processes → 1 printer shared by all comp.
- independent processes → bank acc. in SBI & HDFC. (exams)

int shared = 5

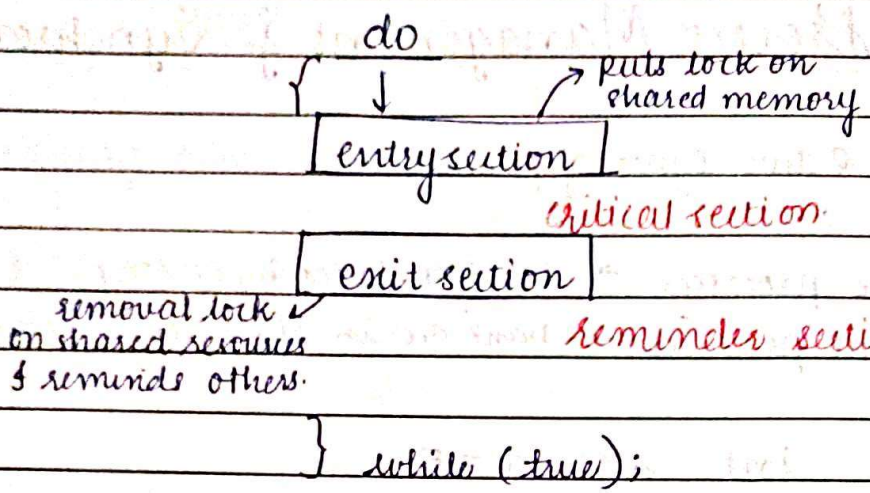
P1	P2
int x = shared	int y = shared
x++	y--
sleep (1)	sleep (1)
shared = x	shared = y
Terminate.	terminate

Race condition

- In OS we have no. of processes & is this processes require no. of ~~processes~~ resources.
- Now think of situation, we have 2 processes & this proses reading same variable, they are reading & updating the value of variable & finally writing the data in memory

DATE : / /

PAGE No :



i] Mutual exclusion

If a process in a critical section, then other process shouldn't be allowed to enter. i.e there must be m.e betn processes.

ii] Progress

It means that if one process doesn't needs to execute into critical section then it shouldn't stop other processes to get into critical section.

iii] Bounded waiting

- There must be some limit to no. of times a process can go into critical section i.e there must be some upper bound.
- If no upper bound is there, same process will be allowed to go critical section again & again

DATE : __/__/__ PAGE No : _____

& other processes will never get a chance to go into critical section.

24/02/24

~~Simphoi~~

Semaphore.

wait (S) - decreement
signal (S) - increement.

```
wait(S)
  {
     while (S==0); // There is a ';' here
        S--;
  }

signal(S)
  {
     S++;
  }
```

Types of semaphore

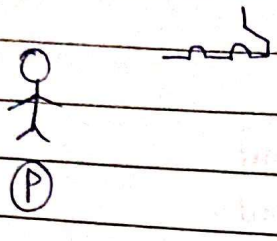
i] Binary semaphore
ii] Counting semaphore

DATE : __/__/____

PAGE No.

Produces - Consumes Problem.

Produces consumes



- multiple types of semaphore variables.

- semaphore (S) → mutual exclusion,
  $S \rightarrow 0/1$ ← initially
- semaphore (E) → empty space in buffer
  $E = n$
- semaphore (F) = space filled by producer
  $F \Rightarrow 0$

```
void producer ()
{
    while (T)
    { produce (x)
      wait (E)
      wait (S)
      append ()
      signal (S)
      signal (F)
    }
}
```

```
void consumer ()
{
    while (T)
    { wait (F)
      wait (S)
      take ()
      signal (S)
      signal (E)
      use ()
    }
}
```

DATE : / /

PAGE No :

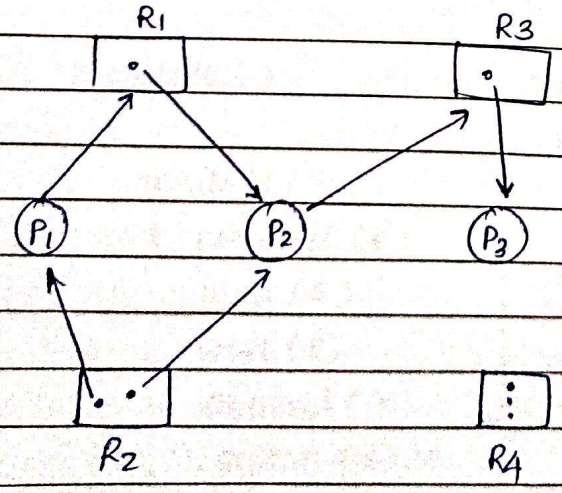
Unit-4 Deadlocks

- Deadlock is a situation where a set of processes are blocked because each process is allocated holding a resource & waiting for another ~~processes~~ resource acquired by some other process.

- The execution of two or more processes is blocked because each processes holds some resource & waiting for another resource holds some other process.

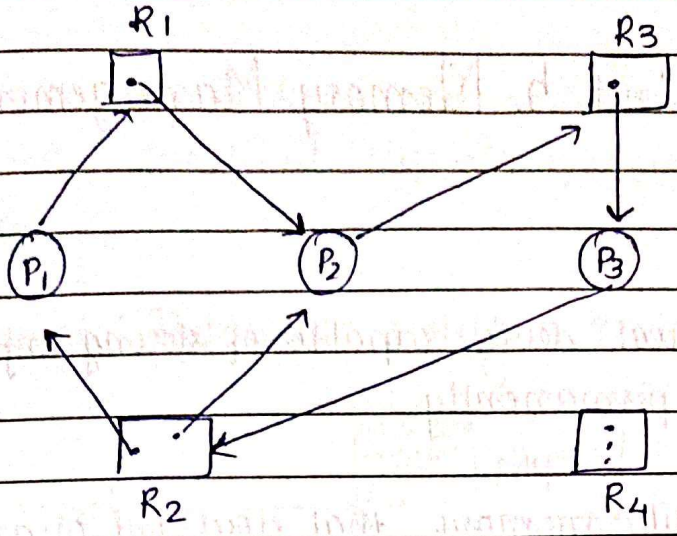
Necessary conditions

→ Mutual Exclusion
→ No preemption
→ Hold & wait
→ Circular wait.



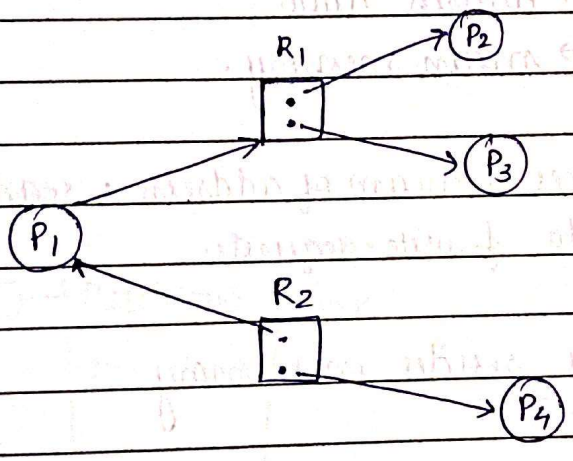
no deadlock.

DATE : __/__/____ PAGE No : _____



deadlock.

Resource Allocation Graph with cycle but no deadlock.



~~Allocation~~ Unit-5. Memory Management

Memory:

- is any physical device capable of storing information temporarily or permanently.

- RAM is a volatile memory, that stores inf on an integrated circuit used by the OS.

- cache memory → address, data
  → main memory.

- memory unit sees a stream of address + read requests or address + data & write requests.

Multiple processes resides in memory

- ensure correct operation.

1] Protect OS 2] protect users processes.

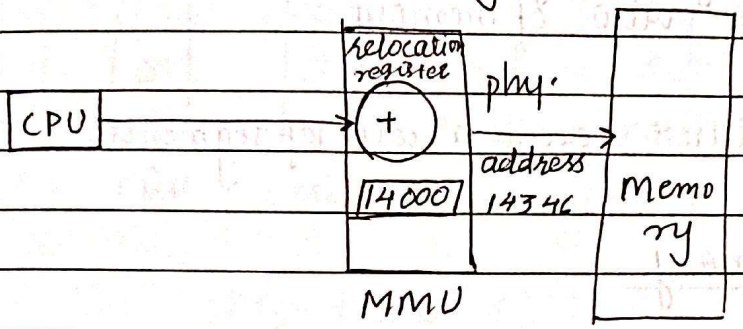
Base and limit Registers

2560000	operating systems	
300040	process	300040 (base)
420940	process	120900 (limit)
880000	process	

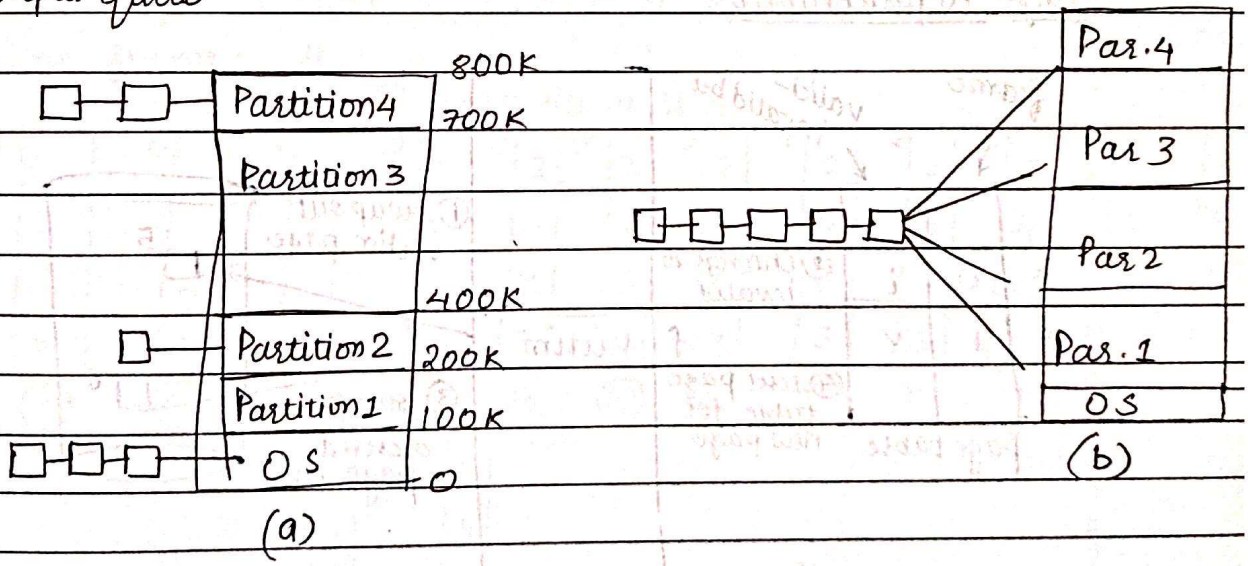
DATE : __/__/____ PAGE No : _____

The base register holds the smallest legal physical memory address; the limit register specifies the ~~for~~ size of the range.

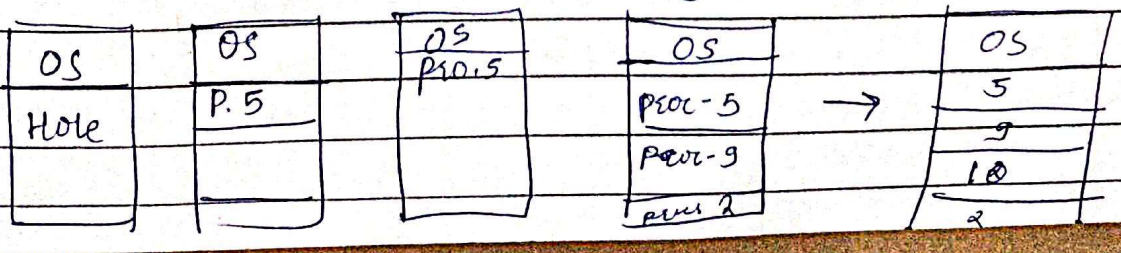
Dynamic Relocation using a relocation register



Multiple input queue



Hole: block of available memory; holes of various size are scattered throughout memory.



DATE : __/__/____

→ Fragmentation.

→ Segmentation.

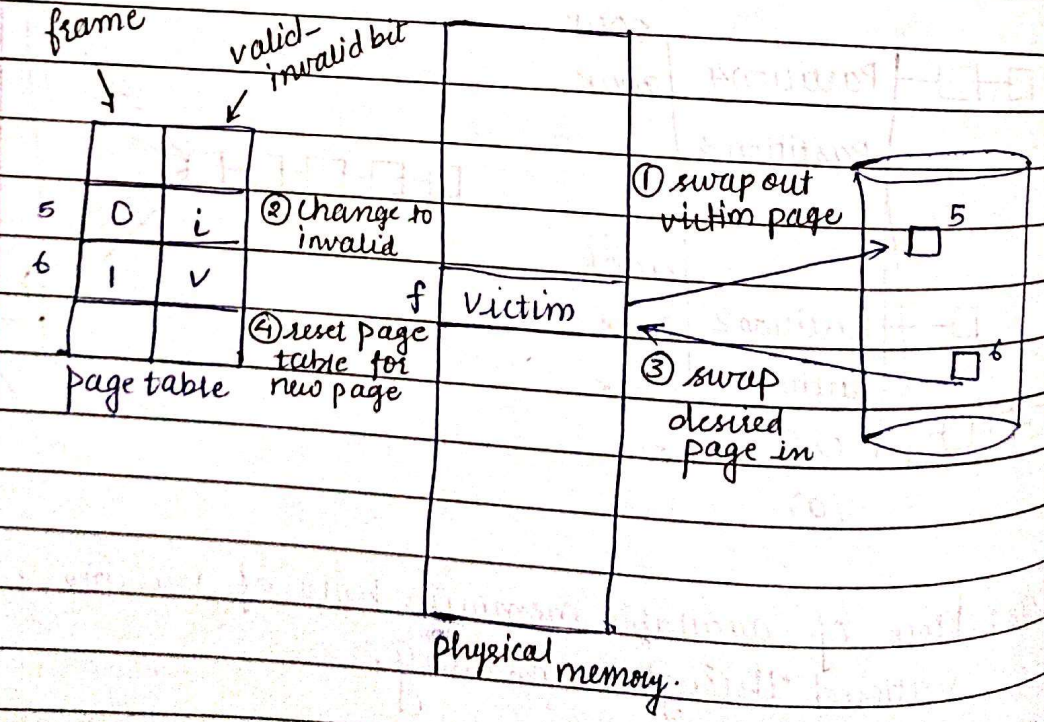
is a memory-management scheme that supports programmar view of memory.

A logical address space is a coll^n of segments:

Vistual Memory

→ TLB - Translation Lookaside Buffer.

Page Replacement



DATE : __/__/____ PAGE No : _____

Page Replacement Algorithms.

(1) FIFO

Page Ref : 1 3 0 3 5 6 3

Page frame : 3.

1	3	0	3	5	6	3
1	3	0	0	0	0	~~3~~ 3
	1	3	3	3	~~6~~ 6	6
		1	1	~~5~~ 5	5	5
miss	miss	miss	Hit	miss	miss	miss

Page faults : 6

Hits 1.

Page ref : 7, 0, 1, 2, 0, 3, 0, 4, 2, 3, 0, 3, 2, 3

Page frame : 4.

(2) Least Recent Used.

7	0	1	2	0	3	0	4	2	3	0	3	2	3
7	7	7	7	7	~~3~~ 3	3	3	3	3	3	3	3	3
	0	0	0	0	0	0	0	~~4~~ 0	0	0	0	0	0
		1	1	1	1	1	1	1	4	4	4	4	4
M	M	M	M	H	M	H	M	H	H	H	H	H	H

2	3
2	2
4	4
0	0
3	3
H	H

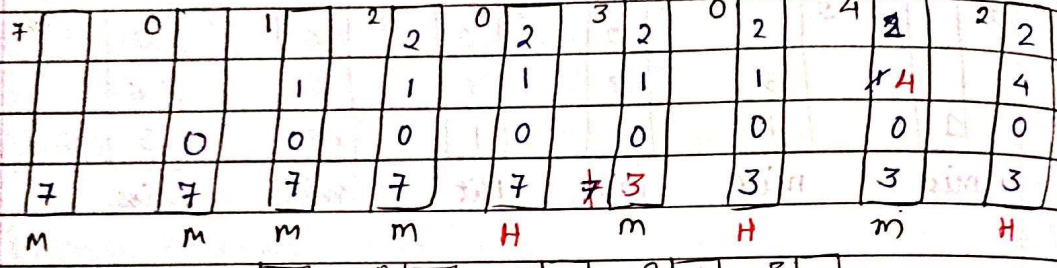
F → 6

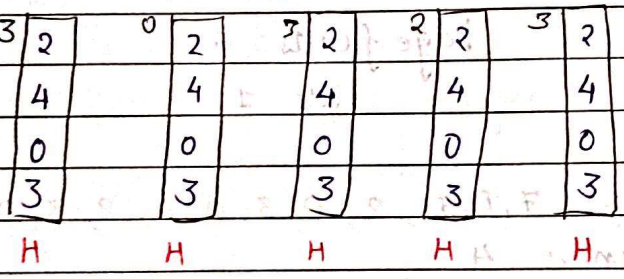
H → 8

DATE : __/__/____ PAGE No :

③ Optimal Page Replacement Algorithm

Page seq: 7, 0, 1, 2, 0, 3, 0, 4, 2, 3, 0, 3, 2, 3
Page frames : 4



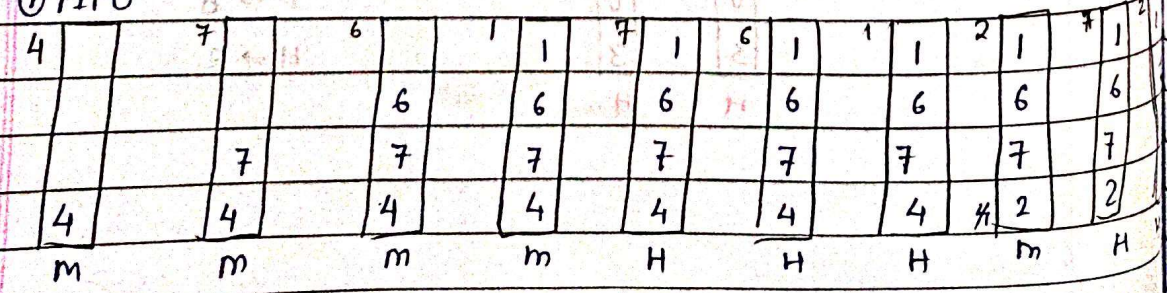


Page fault = 6
Page hit = 8

* This is efficient but not implementable.

Q. P.R - 4, 7, 6, 1, 7, 6, 1, 2, 7, 2
P.f : 4.

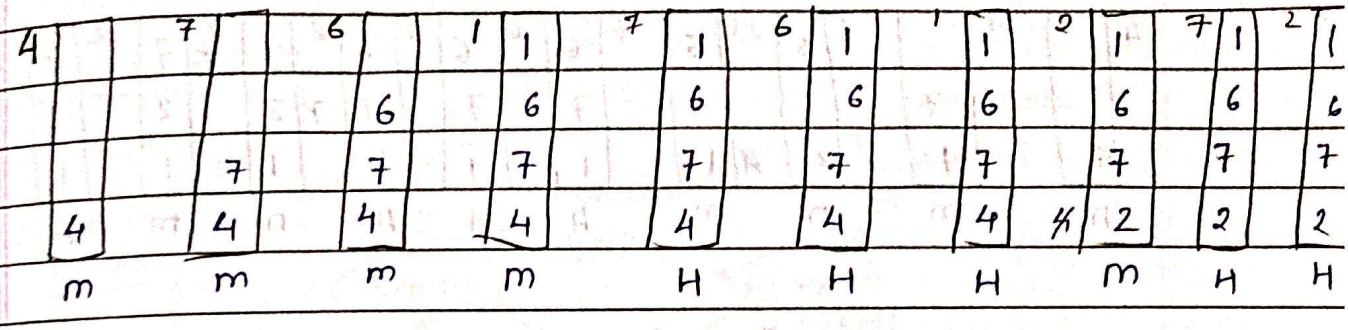
① FIFO



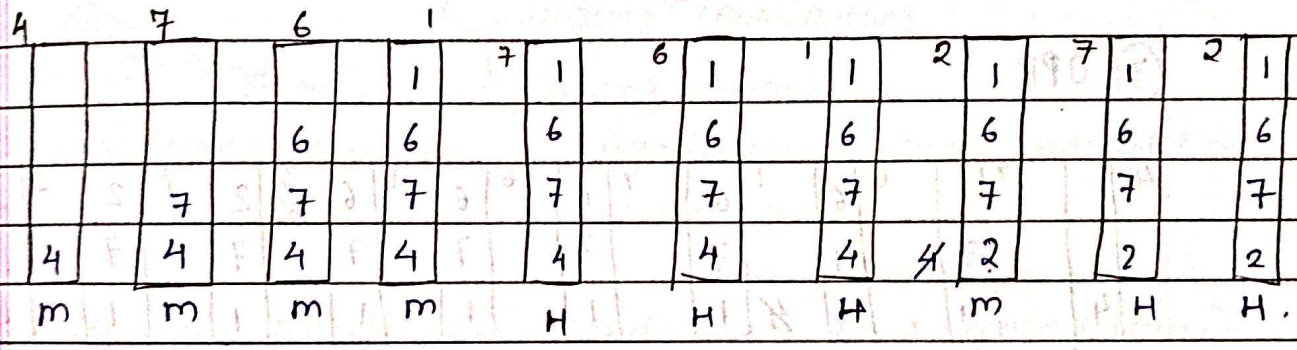
Hit : 5 Miss - 5.

DATE : __/__/____ PAGE No : _____

LRU



OPR



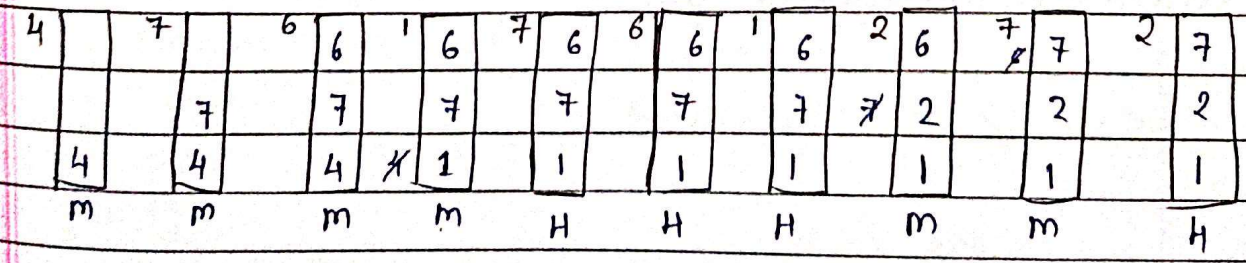
P-f : 3 frame

FIFO - Pf = 6

LRU - Pf = 6

OPR - Pf = 5

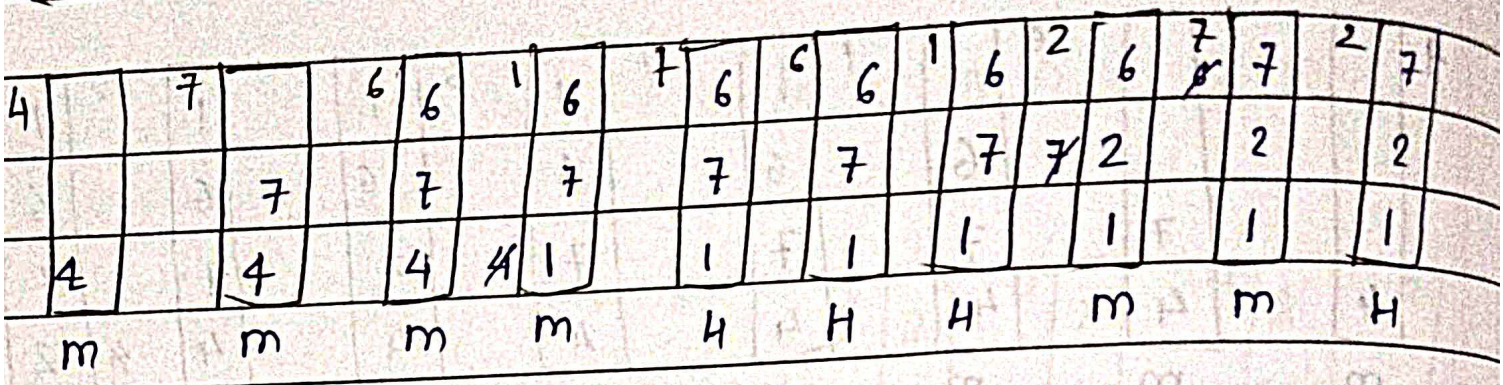
① FIFO.



Miss = 6

Hit = 4.

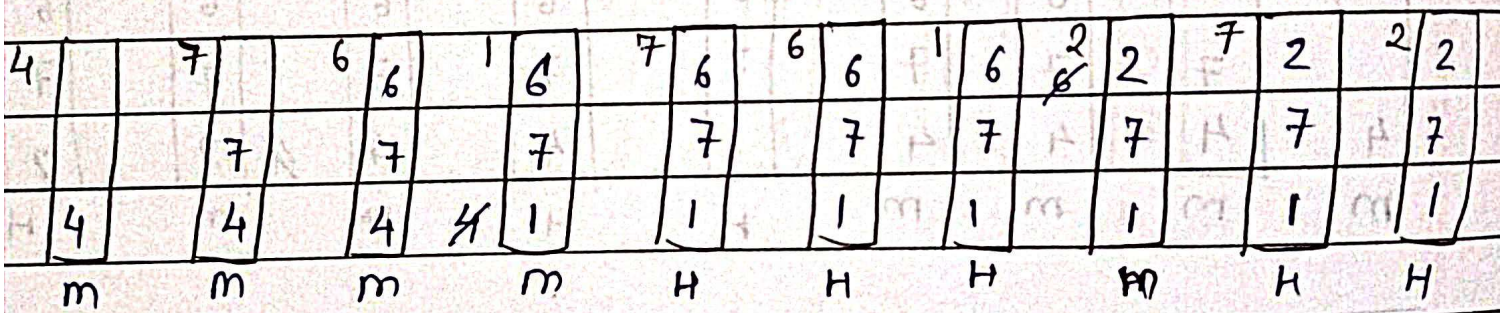
② LRU



Page fault = 6

Hit = 4

③ OPR.



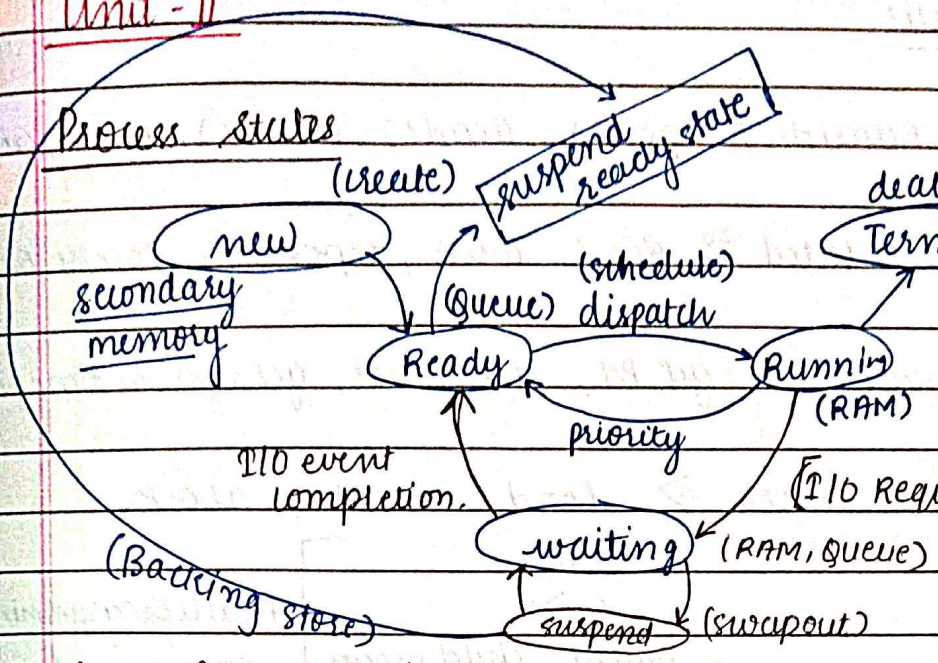
Page fault : 5

Hit = 5.

DATE : __/__/____ PAGE No : _____

Unit - II

Process States



long term scheduler - maximum process bought in ready queue.

Short term scheduler - (ready state to running state)

Medium term scheduler - ## When the waiting state is full, with I/O request of processes, then, the coming process is sent to (suspend state) to swapout. this is done by medi term scheduler.

- same with ready state.

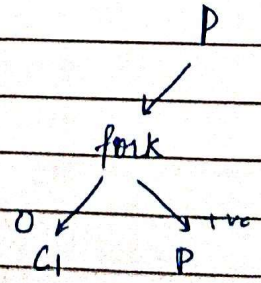
System Calls

- → file Related ⇒ open(), Read(), write(), close(), create(),
- → device related ⇒ Read, write, reposition, request, release
- → Information ⇒ get Pd, attributes, get system time/data
- → Process control ⇒ load, execute, abort,
  - Fork
    - ↙ parent process
    - ↘ child process.
  - multiprogramming environment.
  - wait, signal, allocate etc.
- → Inter process communication - pipe(), create()/delete connection, shmget().

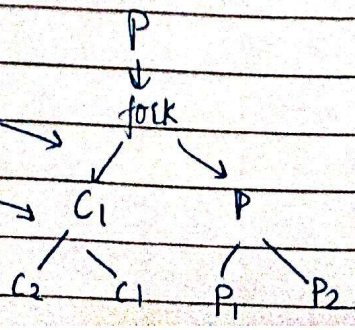
Fork()

- creates a child process.

Fork() → 0 child
→ +1 +ve Parent
→ -1 child




```
main() {
    fork();
    fork();
    printf("hello");
}
```



DATE : __/__/__ PAGE No : _____

no. of ~~parents~~ execution 2^n ← (no. of time fork written).
child process $2^n - 1$