

- ✓ LIS in  $n \log n$
- ✓ Interleaving Strings
- ✓ Regular Expression - II
- ✓ Poisonous graph.
- ✓ matrix and Absolute Difference.
  - Auto-complete (Hint)

Contest → Revise assignments.

5 Questions.

- 2G
- 2m
- 1 hand.

# ① Longest Increasing Subsequence

Eg → 

✓	✓	✓	✓	✓	✓	✓	✓	✓
10	22	9	33	21	50	41	45	48
0	1	2	3	4	5	6	7	8

list = 

<del>10</del>	22	33	<del>50</del>	45	48
1	2	3	4	5	6

# code →

```
list <=> lis;
```

```
lis.add(arr[0]);
```

```
for (i = 1; i < N; i++) {
```

```
    if (arr[i] > lis.get(lis.size() - 1)) {
```

```
        lis.add(arr[i]);
```

```
    } else {
```

// find idx of the element which is just greater than arr[i]; (Binary Search)

```
        lis[idx] = arr[i];
```

```
    }
```

```
return lis.size();
```

T.C →  $O(N \log N)$   
S.C →  $O(N)$

## ② Interleaving Strings →

Given A, B, C find whether C is formed by the interleaving of A and B.

A = "aabcc"

B = "dbbca"

C = "aadbcbcbcac"

ans → true.

a a b c c

d b b c a

a d a b c b b c c a

A → ~~a~~<sub>0</sub> ~~a~~<sub>1</sub> b<sub>2</sub> c<sub>3</sub> c<sub>4</sub> ← i

B → ~~d~~<sub>0</sub> b<sub>1</sub> b<sub>2</sub> c<sub>3</sub> a<sub>4</sub> ← j

C → ~~d~~<sub>0</sub> ~~a~~<sub>1</sub> ~~d~~<sub>2</sub> b<sub>3</sub> b<sub>4</sub> c<sub>5</sub> b<sub>6</sub> c<sub>7</sub> a<sub>8</sub> c<sub>9</sub>

↑  
i+j

s1, s2, s3, i, j

~~$s1[i] == s3[i+j]$~~

↙

s1, s2, s3, i+1, j

$s2[j] == s3[i+j]$ 

↘

s1, s2, s3, i, j+1

# code →

$dp[m][n], \forall i, j \quad dp[i][j] = -1$

```
int interleaving ( s1, s2, s3, 0i, 0j, int dp ) {  
    if ( i == s1.length() && j == s2.length() ) return 1;  
    if ( dp[i][j] != -1 ) return dp[i][j];  
    f1 = 0, f2 = 0;  
    if ( s1[i] == s3[i+j] ) {  
        f1 = interleaving ( s1, s2, s3, i+1, j, dp );  
    }  
    if ( s2[j] == s3[i+j] ) {  
        f2 = interleaving ( s1, s2, s3, i, j+1, dp );  
    }  
    dp[i][j] = Max ( f1, f2 );  
    return Max ( f1, f2 );  
}
```

T.C →  $O(N \times m)$   
S.C →  $O(N \times m)$

[ Bottom - up → # todo ]

Bottom-up idea.

A = "aabcc"

B = "dbbca"

C = "aadbcbcac"

ans  $\rightarrow$  true.

		-	d	b	b	c	a
		0	1	2	3	4	5
-	0	t	f	f	f	f	f
a	1	t	f	f	f	f	f
a	2	t	t				
b	3	f					
c	4	f					
c	5	f					

$\left[ \begin{array}{l} T.C \rightarrow O(N \times m) \\ S.C \rightarrow O(N \times m) \end{array} \right] \checkmark$

③

## Regular Expression - II

Implement wildcard pattern matching with support for '?' and '\*' for strings A and B.

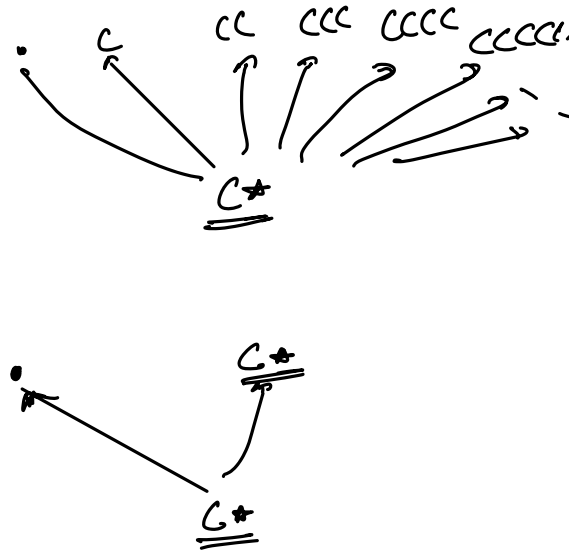
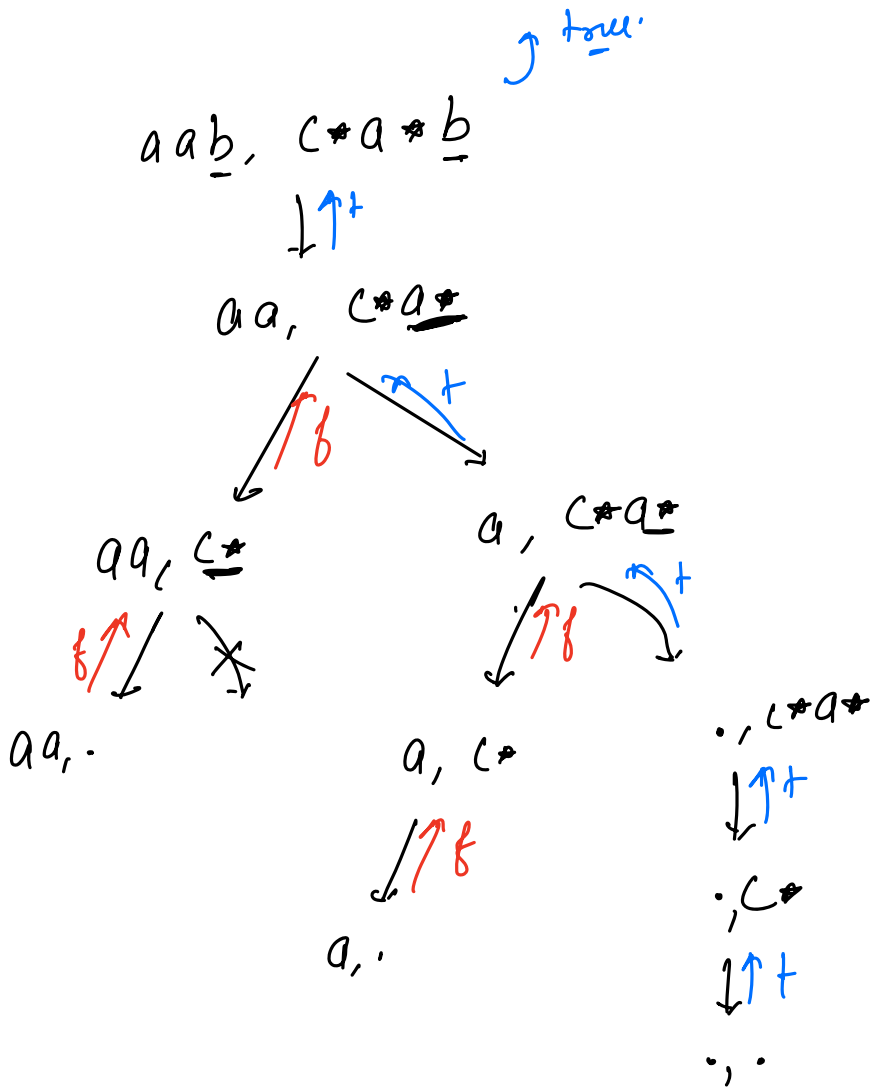
' . ' : Matches any single character.

' \* ': Matches zero or more of the preceding element.

The matching should cover the entire input string (not partial).

A = "aab"

B = "c\*a\*b"



s1  $\rightarrow$  a b c b b x y t

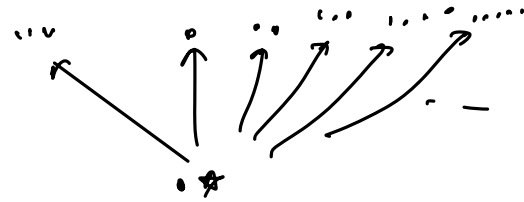
s2  $\rightarrow$  a b \* c . \* t

		- a b * c . * t $\rightarrow$ j							
		0	1	2	3	4	5	6	7
i - a b c b b x y t	0	t	f	f	b	f	f	f	f
	1	f	t	f	t	f	f	f	f
	2	f	f	t	t	f	f	f	f
	3	f	f	f	f	t	f	t	f
	4	f	f	f	f	f	t	t	f
	5	f	f	f	f	f	f	t	f
	6	f	f	f	f	f	f	t	f
	7	f	f	f	f	f	f	t	f
	8	f	f	f	f	f	f	t	t

\*  $\rightarrow$  dp[i][j-2]

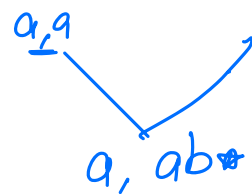
s1  $\rightarrow$  a b c d d m

s2  $\rightarrow$  a b . \* m true



s1  $\rightarrow$  a b c c d e m m t y

s2  $\rightarrow$  . \* true



# code  $\rightarrow$   $dp[n+1][m+1]; \quad \forall i, j \quad dp[i][j] = false;$

for ( $i=0; i \leq n; i++$ ) {

for ( $j=0; j \leq m; j++$ ) {

if ( $i==0$  &  $j==0$ ) {  $dp[i][j] = true;$  }

else if ( $j==0$ ) {  $dp[i][j] = false;$  }

else if ( $i==0$ ) {

if ( $s2[j-1] == '*'$ ) {

{  $dp[i][j] = dp[i][j-2];$

else {  $dp[i][j] = false;$

}

else {

if ( $s1[i-1] == s2[j-1]$  ||  $s2[j-1] == '.'$ ) {

{  $dp[i][j] = dp[i-1][j-1];$

else if ( $s2[j-1] == '*'$ ) {

$f1 = dp[i][j-2];$

$f2 = false;$

if ( $s2[j-2] == '.'$  ||  $s2[j-2] == s1[i-1]$ ) {

{  $f2 = dp[i-1][j];$

$dp[i][j] = (f1 || f2);$

}

}

}

}

return  $dp[n][m];$

$T.C \rightarrow O(n \times m)$   
 $S.C \rightarrow O(n \times m)$



Break → 10:41 → 10:48

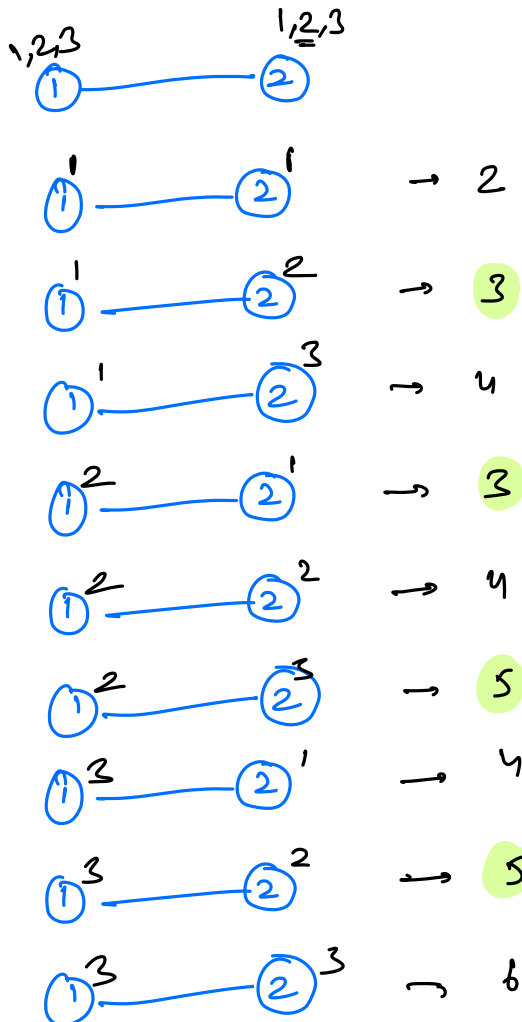
## ④ Poisonous Graph

You are given an undirected unweighted graph consisting of  $A$  vertices and  $M$  edges given in a form of 2D Matrix  $B$  of size  $M \times 2$  where  $(B[i][0], B[i][1])$  denotes two nodes connected by an edge.

You have to write a number on each vertex of the graph. Each number should be 1, 2 or 3. The graph becomes Poisonous if for each edge the sum of numbers on vertices connected by this edge is odd.

Calculate the number of possible ways to write numbers 1, 2 or 3 on vertices so the graph becomes poisonous. Since this number may be large, return it modulo 998244353.

①

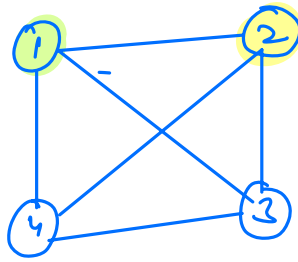


Ans = 4

②

4

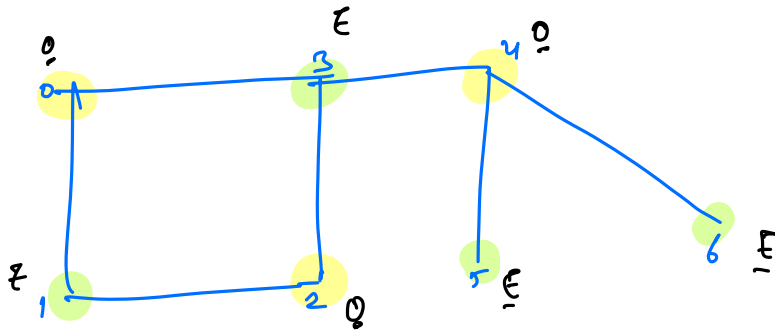
1	2
1	3
1	4
2	3
2	4
3	4



ans = 0

observation

- if graph is non bi-partite, then ans → 0
- if graph is bi-partite.



0	1	0	1	0	1	1
0	1	2	3	4	5	6

yellow → 3, green → 4

[1,2,3]

yellow colored nodes  $\rightarrow E$  and green colored nodes  $\rightarrow 0$

$\text{or}$   
yellow colored nodes  $\rightarrow 0$  and green colored nodes  $\rightarrow E$

$$\Rightarrow 1^3 \times 2^4 + 2^3 \times 1^4 = \underline{\underline{2^4 + 2^3}}$$

$$\left[ \therefore \text{ans} = 2^{\text{yellow-colored nodes}} + 2^{\text{green colored nodes}} \right]$$

$$\left[ \begin{array}{l} \text{T.C} \rightarrow O(N+E) \\ \text{S.C} \rightarrow O(N) \end{array} \right]$$

## ④ Matrix and Absolute Difference

Given a matrix  $C$  of integers, of dimension  $A \times B$ .

For any given  $K$ , you are not allowed to travel between cells that have an absolute difference greater than  $K$ .

Return the minimum value of  $K$  such that it is possible to travel between any pair of cells in the grid through a path of adjacent cells.

NOTE:

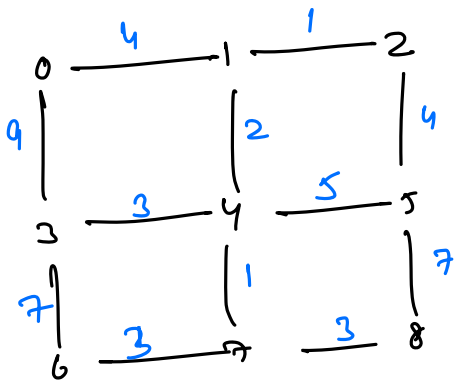
Adjacent cells are those cells that share a side with the current cell.

$C_{ij} = 1$

	0	1	2
0	1	4	5
1	9	2	4
2	10	3	2

⇒ Minimum Spanning Tree.

↙ ↘  
Prim's Algo      Kruskal's Algo



1, 2, 1 ✓

1, 4, 2 ✓

4, 7, 2 ✓

6, 7, 3 ✓

3, 4, 3 ✓

7, 8, 3 ✓

2, 5, 4 ✓

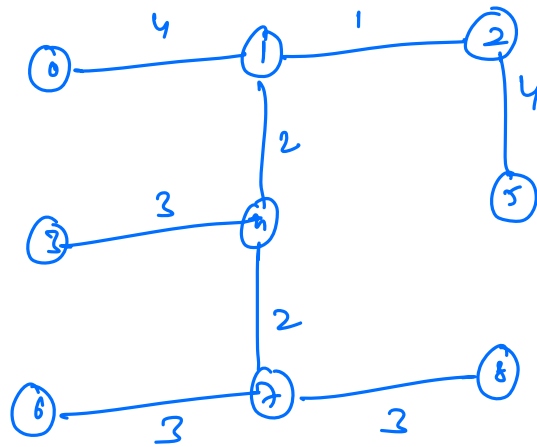
0, 1, 4 ✓

4, 5, 5 ✗

3, 6, 7 ✗

5, 8, 7 ✗

0, 3, 9 ✗



ans = 4.