

What is Recursion?

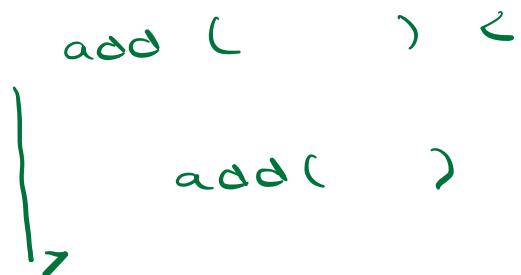
How to write recursive code?

Function call Tracing

3 problems

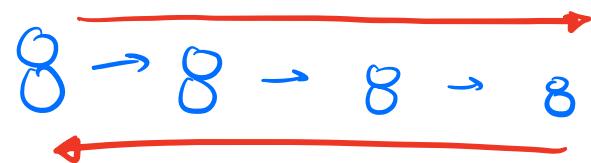
TC and SC

Recursion \rightarrow function calling itself





1. bigger doll \rightarrow smaller doll
2. similar
3. end doll



Recursion \rightarrow solving larger problems
using smaller similar subproblem

$$\text{sum}(N) = 1 + 2 + 3 + 4 + \dots + N-1 + N$$

$$\text{sum}(N) = \text{sum}(N-1) + N \quad \text{→ Recursion solution}$$

- Q. Write recursive fn for sum of
N natural no.

1. Assumption : Decide what your fn does

2. Main Logic :

Recursive relation to code /

breaking big problem into smaller
problem

3. Base condition

terminating condition /

smallest prob for which you already
know answer

// given N, return sum of first N natural nos.

int sum (n) <

Base
condn

if ($n == 1$) return 1

return sum(n-1) + n

sum(4)



sum(3) → 4



sum(2) + 3



sum(1) + 2

'C'

Fn call Tracing

seq of fn calls that are made
when a program is executed



fn calls, arguments, return values

```
int add(int x, int y) {  
    return x + y;  
}  
  
int mul(int x, int y, int z) {  
    return x * y z;  
}  
  
int sub(int x, int y) {  
    return x - y;  
}  
  
void print(int x) {  
    cout << x << endl;  
}  
  
int main() {  
    int x = 10;  
    int y = 20;  
    print(sub(mul(add(x, y), 30, 75)));  
    return 0;  
}
```

main() <
x = 10, y = 20 ✓
→ print (sub(mul(add(x, y), 30, 75)))

print (sub (mul (add (x, y), 30, 75)))

↓↑ 825

sub (mul (add (x, y, 30), 75))

↓↑ 900

mul (add (x, y, 30))

↓↑ 30

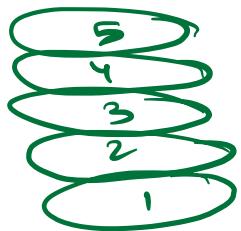
add(x, y) <

|
return x + y

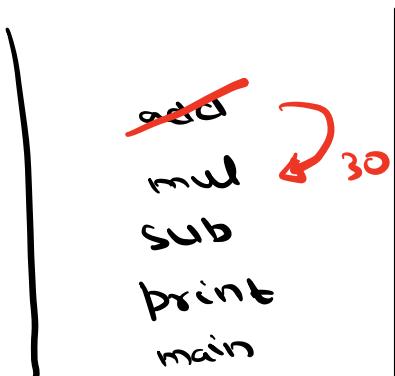
O/P
825

Fn call stack

↑
LIFO
Last in First
out



stack
fn calls
arg
return values



Top of the stack
fn is executed

- Given a +ve integer N , find factorial of N

$$5! = 120$$

$$5! = 5 \times 4 \times 3 \times 2 \times 1$$

$$5! = 5 \times 4!$$

①

$$\text{fact}(N) = N \times \text{fact}(N-1)$$

recusive relation

$$N! = N \times (N-1)!$$

②

$$\text{if } n == 1 \text{ return 1}$$

$$0! = 1$$

Memory limit exceeded /
stack overflow

```
// Given input N, return N!      fact(0)
int factorial (int N) {
    if (N == 0) {
        return 1;
    } else {
        return N * factorial(N-1);
    }
}
```

$$\textcircled{1} \quad | \quad g = 6 \text{ cm/sec}$$

② return factorial(N-1) * N

fact(0)
↓
fact(-1)
↓
fact(-2)
↓
-3
↓

$$\begin{array}{c} f(1) \\ \downarrow \\ f(0) \times 1 \\ \downarrow \\ f(-1) \end{array}$$

$$f(2) \rightarrow 2$$

\downarrow

$$\cancel{f(1)} + 2$$

```
main() <  
print(factorial(5))  
|  
|
```

```
print(factorial(5))
```

```

int factorial (int n) {
    if (n == 1)
        return 1;
    return factorial (n-1) * n;
}

```

~~int factorial (int n) {~~

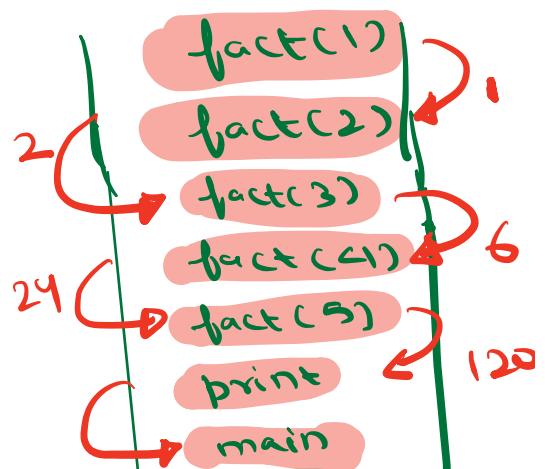
if (n == 1)
return 1
return ~~factorial (N-1)~~ × N

②

4

. 9 (B)

```
int factorial (int N) {
    if (n == 1)
        return 1
    return factorial (N-1)
```



```

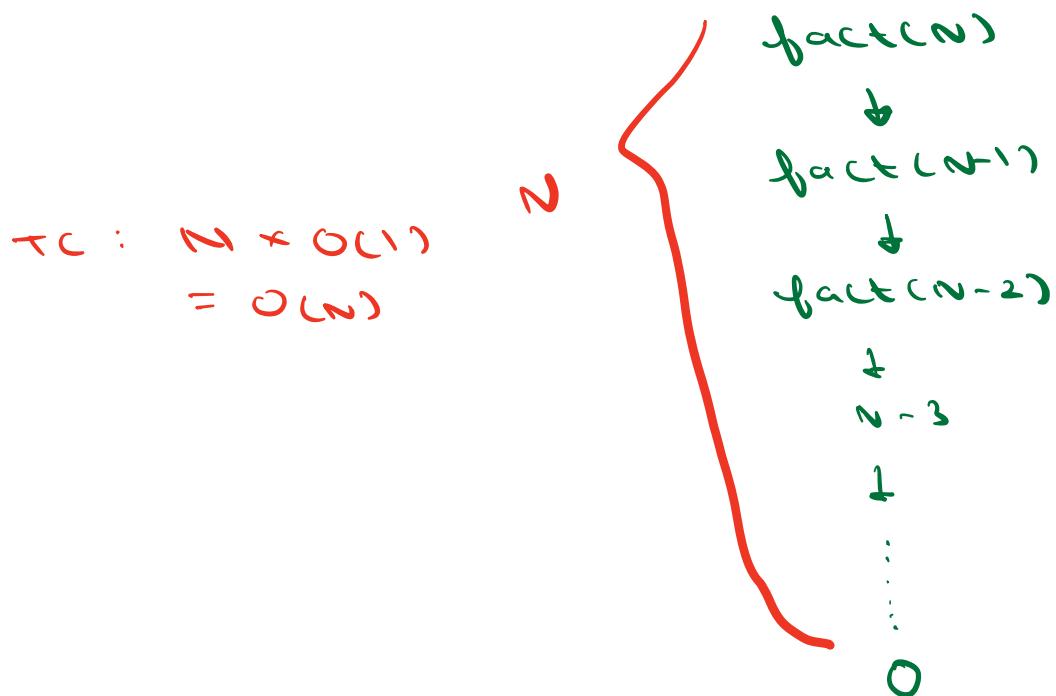
int factorial (int n) {
    if (n == 1)
        return 1
    else
        return factorial (n-1) * n
}

int factorial (int n) {
    if (n == 1)
        return 1
    else
        return factorial (n-1) * n
}

```

The diagram shows two identical recursive function definitions for 'factorial'. The first definition has red annotations: a red circle labeled '①' is placed over the 'if' condition, and a red circle labeled '②' is placed over the 'return' statement. A red arrow points from '①' to the 'if' condition in the first definition. Another red arrow points from '②' to the 'return' statement in the first definition. The second definition is identical but lacks these annotations.

$TC = \text{No. of recursive calls} \times \text{time taken of each recursive call}$



2. Fibonacci series

0	1	2	3	4	5	6	7
→ 0	1	1	2	3	5	8	13

$$\text{Fib}(6) = \text{Fib}(5) + \text{Fib}(4)$$

↗
 5 + 3
 = 8

$$\text{Fib}(n) = \text{Fib}(n-1) + \text{Fib}(n-2)$$

```

if (n==0) return 0
if (n==1) return 1
  
```

// Given N, return Nth fibonacci no.

```

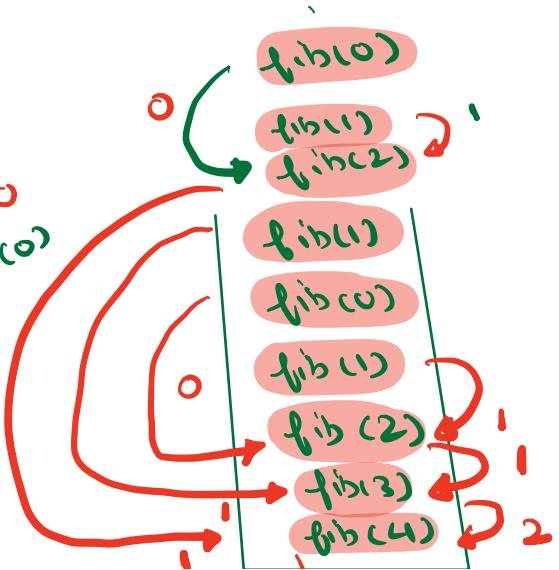
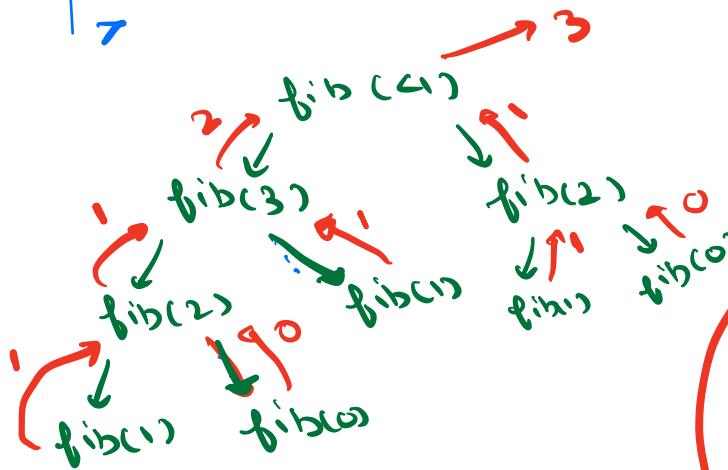
int fibonacci (int n) {
  
```

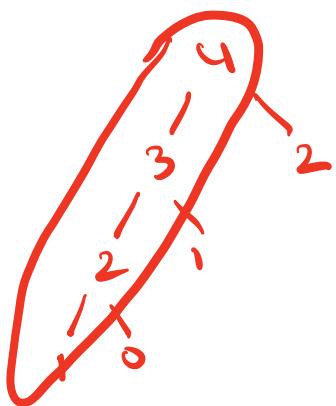
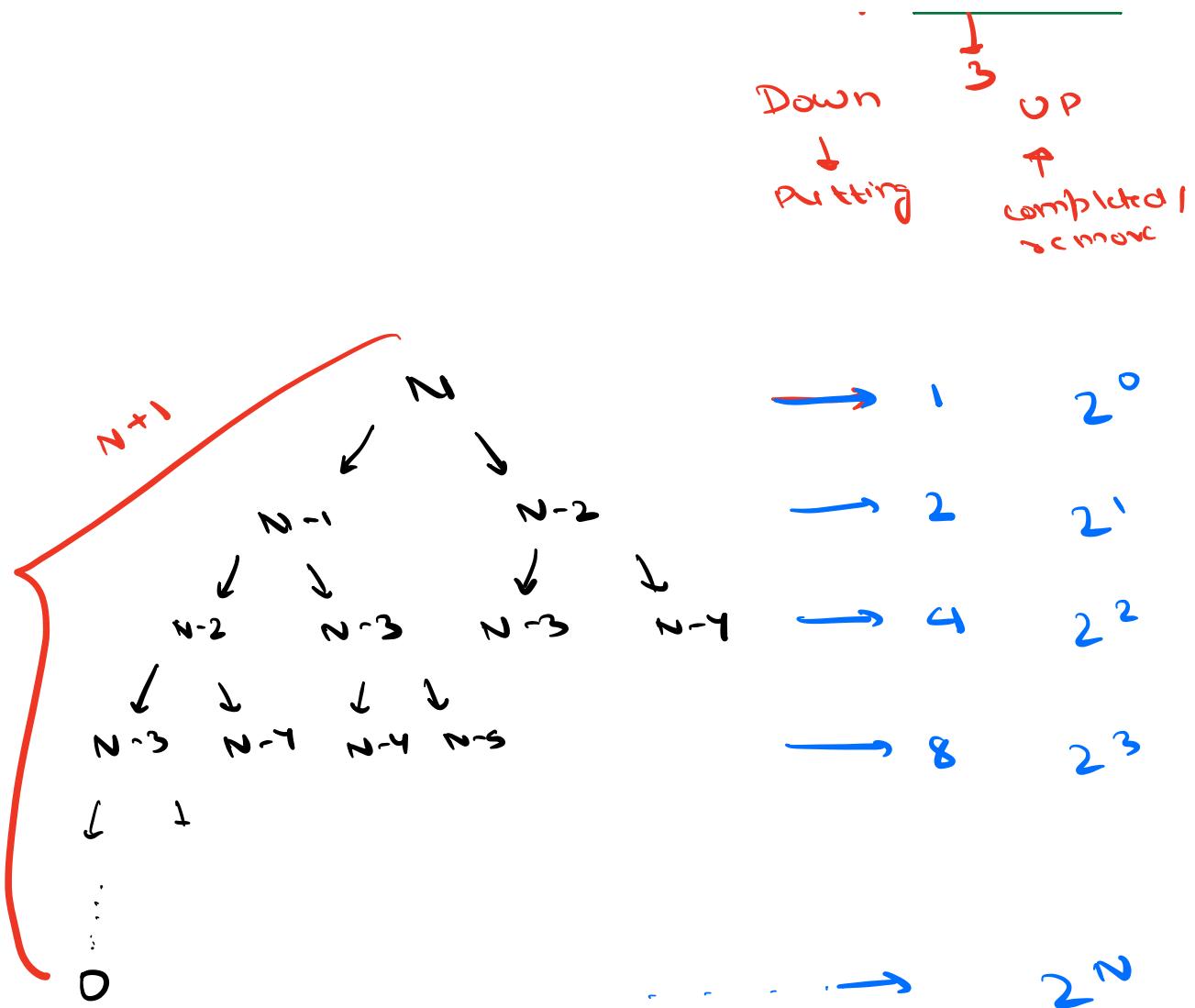
```

    if (n <= 1) return n
  
```

```

    return fibonacci (n-1) + fibonacci (n-2)
  
```





$$\text{sum} = \frac{\alpha(2^N - 1)}{2 - 1}$$

Total fn calls =

$$2^0 + 2^1 + 2^2 + \dots + 2^N$$

$$= \frac{1}{2 - 1} (2^{N+1} - 1)$$

$$= 2^{N+1}$$

T.C.: $O(2^N)$

10:57

3. Find a^n using recursion.

$$a=2$$

$$n=3$$

ans $\rightarrow 8$

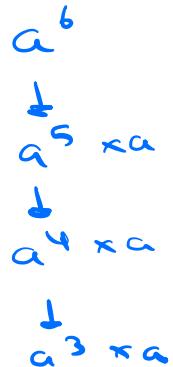
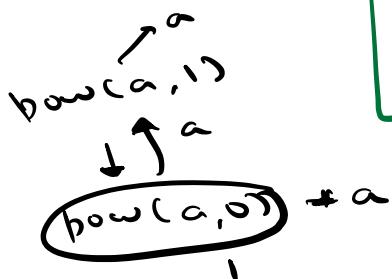
$$a^6 = a \times a \times a \times a \times a$$

$$a^5 = a^4 \times a$$

$$a^{10} = a^9 \times a$$

$$a^n = \underline{a^{n-1}} \times a$$

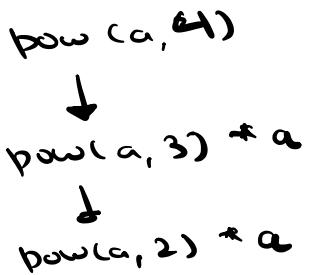
$$\boxed{\text{pow}(a, n) = \text{pow}(a, n-1) * a}$$

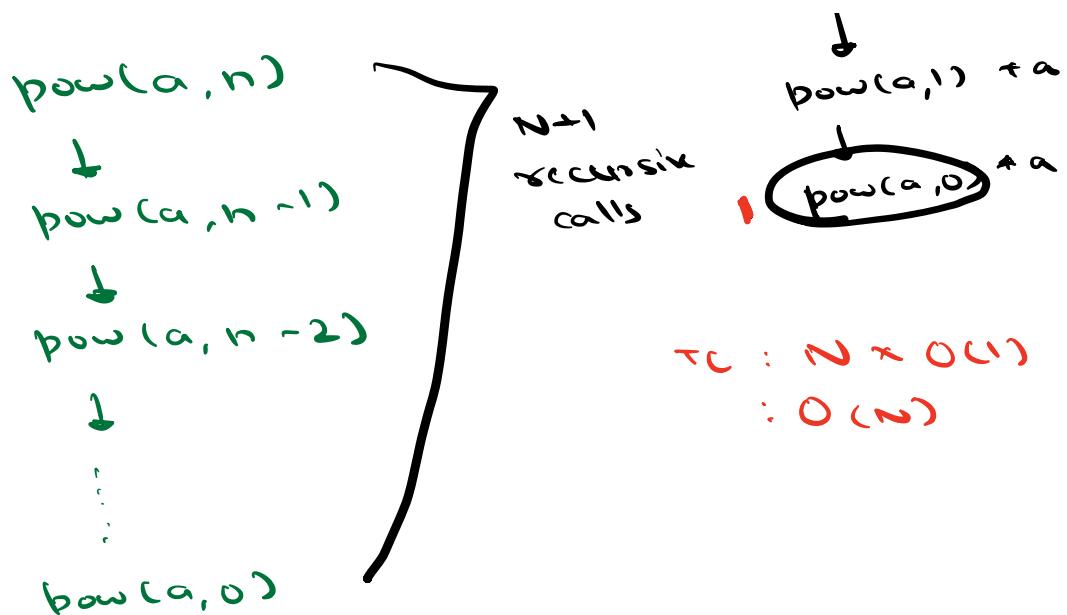


// Given a, n ; return a^n

```
int pow (int a,int n) {
```

```
    if (n == 0) return 1;
    return pow (a,n-1) * a;
```





$$a^{10} = a^5 \times a^5$$

$$a^{11} = a^5 \times a^5 \times a$$

$$a^6 = a^3 \times a^3$$

$$a^N = \begin{cases} a^{N/2} \times a^{N/2} & \text{if } N \text{ is even} \\ a^{N/2} \times a^{N/2} \times a & \text{if } N \text{ is odd} \end{cases}$$

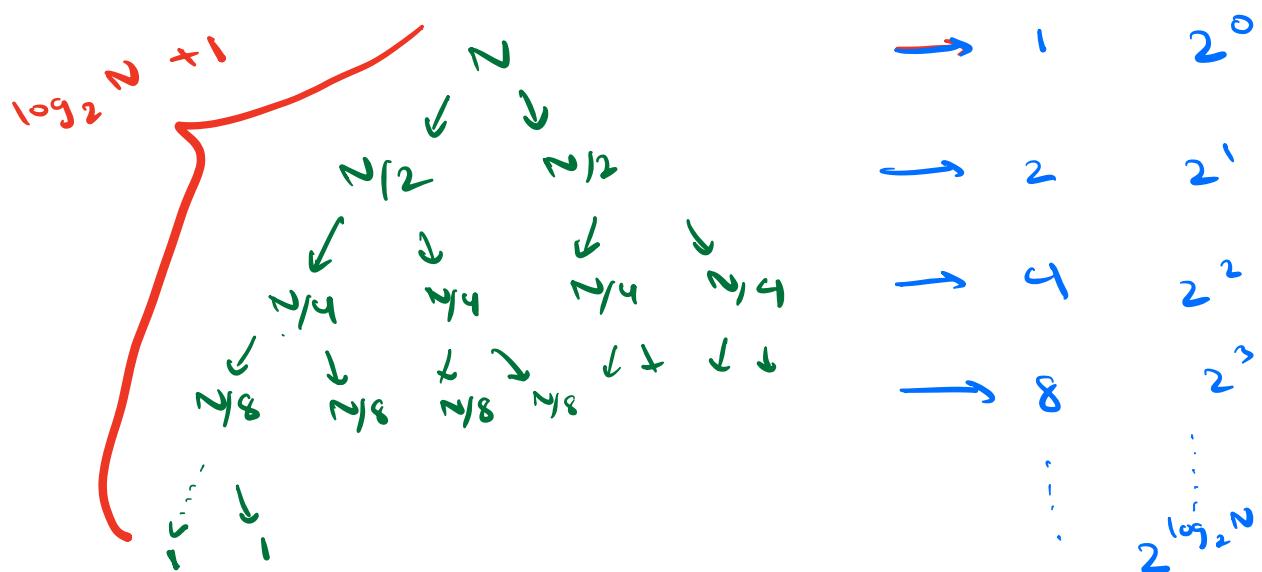
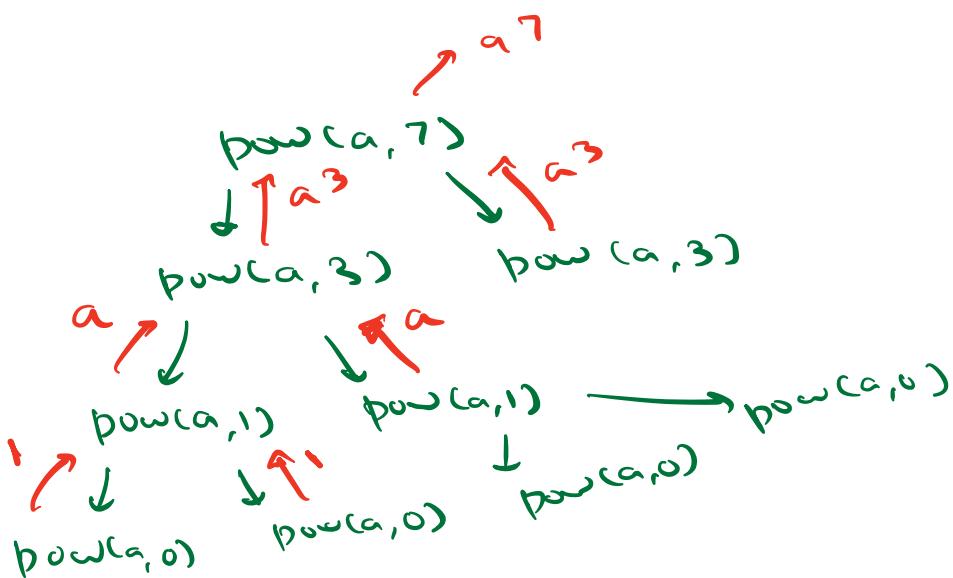
if ($N==0$) return 1

$$\boxed{a^0 = 1}$$

```

int pow (int a,int N) {
    if (n == 0) return 1
    if (n%2 == 0)
        return pow(a,n/2) * pow(a,n/2)
    else
        return pow(a,n/2) * pow(a,n/2) + a
}

```





$$N \rightarrow N/2 \rightarrow N/4 \rightarrow \dots \cdot 1 \rightarrow 0$$

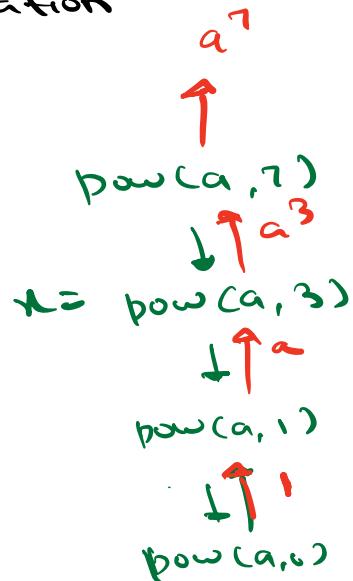
$\log_2 N + 1$

$$\begin{aligned} \text{No. of recursive calls} &= 2^0 + 2^1 + 2^2 + \dots + 2^{\log_2 N} \\ &= \frac{2^0 (2^{\log_2 N + 1} - 1)}{2 - 1} \\ &= 2^{\log_2 N} \cdot 2 \end{aligned}$$

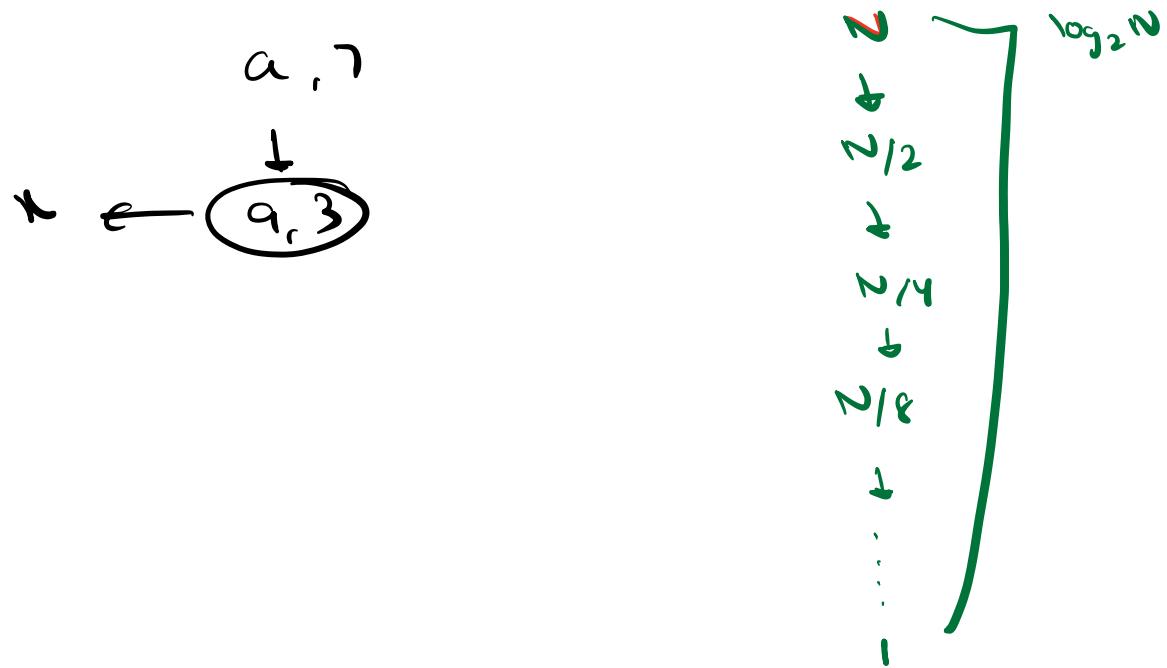
$\mathcal{T.C}: O(N)$

Fast Power / Fast Exponentiation

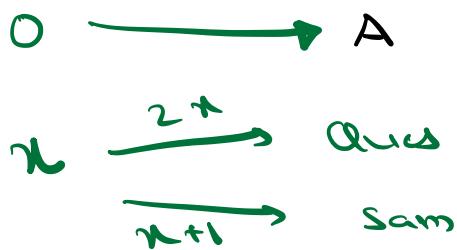
```
int pow (int a,int N) {
    if (n == 0) return 1
    int x = pow (a,N/2)
    if (n%2 == 0)
        return x*x
    else
        return x*x+a
}
```



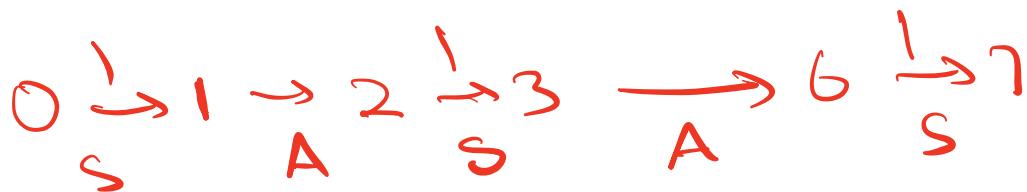
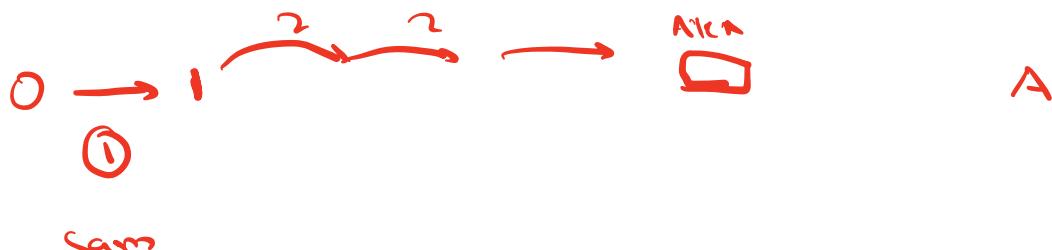
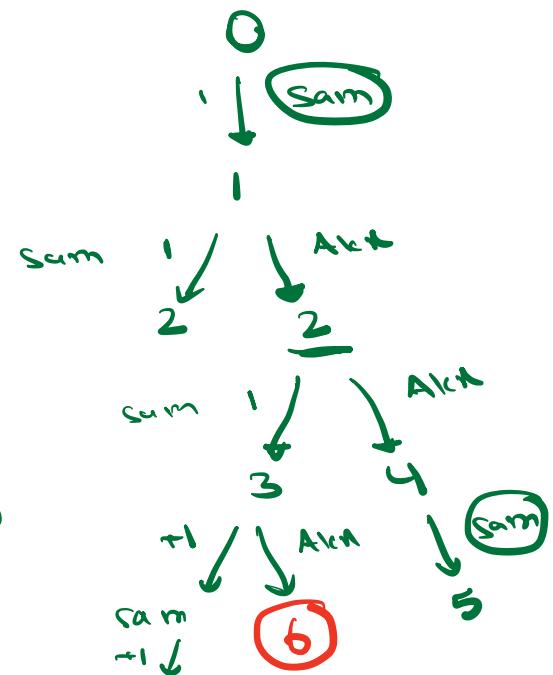
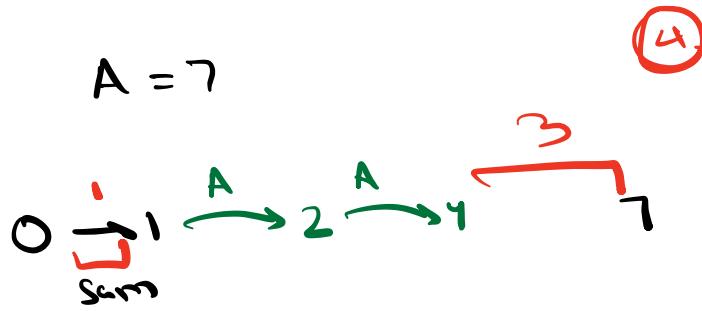
$\mathcal{T.C}: O(\log_2 N)$



Doubts



$$A = 5$$



arr of 1 and 0

→ 1 1 0 0 1 1 1

Find cnt of substrings where OR = 1

TC : O(N)