

Agenda :

Target Sum

Min Jumps

N digit number

2.5 hrs → 5 ques

Contest → 8<sup>th</sup> Feb Thursday

Discussion → 9<sup>th</sup> Feb

Pass →  $\geq 60\%$ . (3 ques)

| R1     | R2     | R3      |
|--------|--------|---------|
| 2 Days | 9 Days | 20 Days |

5 Ques → 1 DP, 1 Graphs

Placement

- ① Mock Interview
- ② Contest

} 1 month

Revision ← PSP  $\geq 70\%$ .

5 classes → Feb

1. Given an integer array  $A$  of size  $N$  and  $B$ .  
Find whether there exist a subset in  $A$  whose sum equal  $B$ .

$A = [3, 34, 4, 12, 5, 2]$

$B = 9$

(T)

$\langle 4, 5 \rangle$

$\langle 3, 4, 2 \rangle$

$B = 30$

(F)

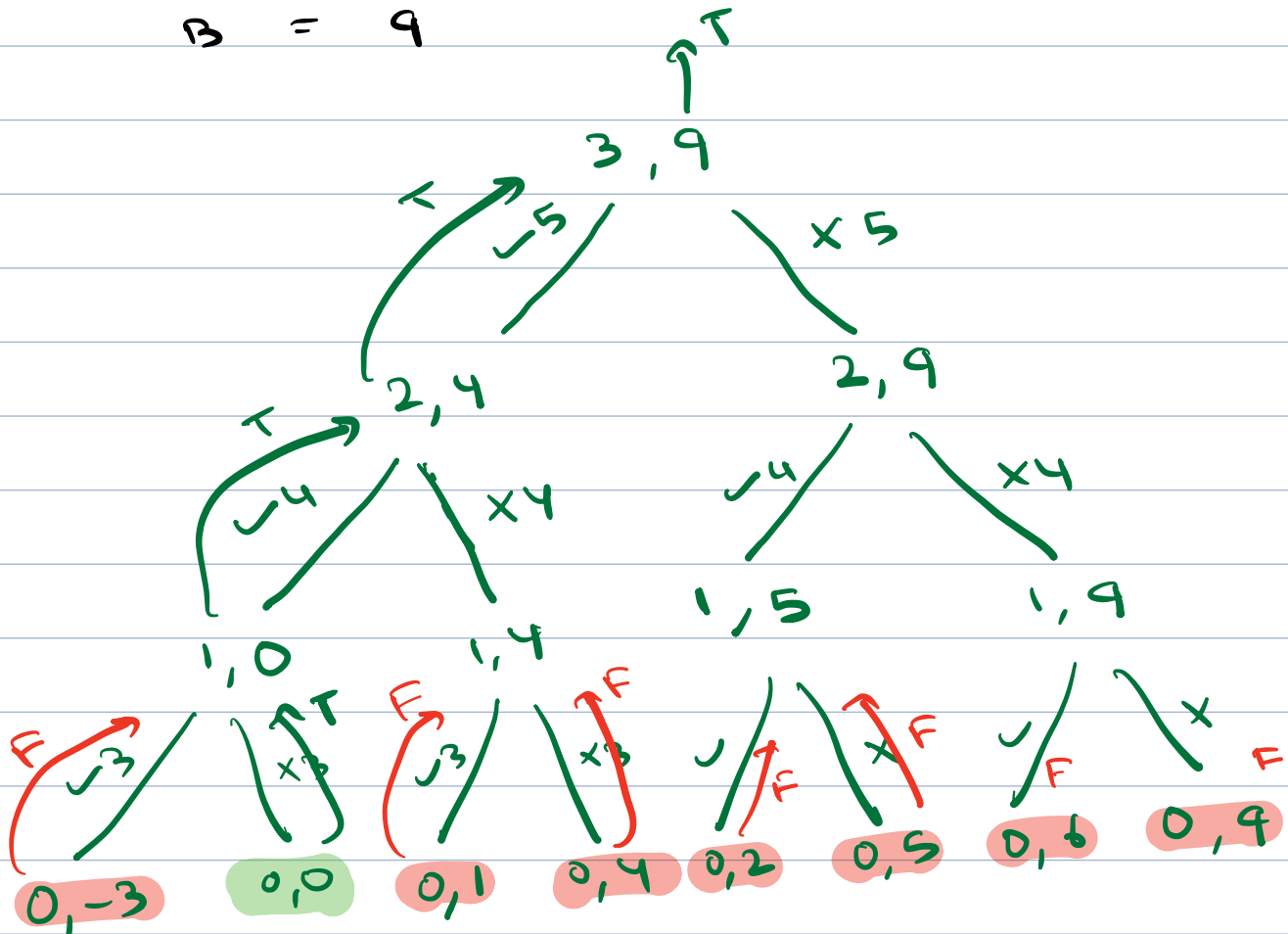
BF: Go to all subsets

$n \rightarrow 2^n$

$3 \rightarrow 2^3$

$A = \begin{matrix} 0 & 1 & 2 \\ [3 & 4 & 5] \end{matrix}$

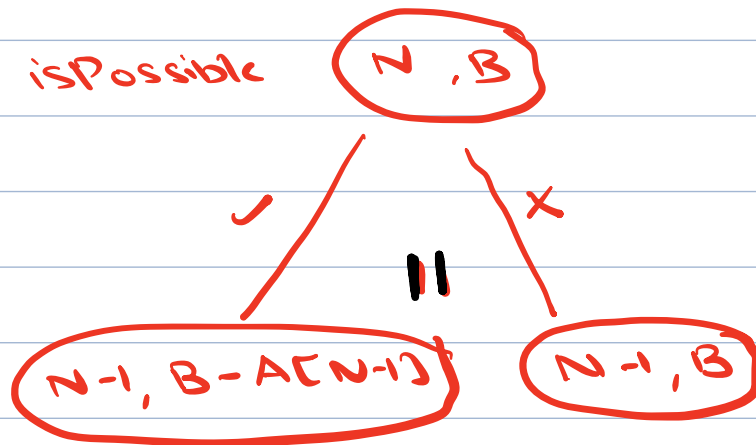
$B = 9$



TC:  $O(2^N)$

$A : [ 0 \ 1 \ \dots \ N-1 ]$

$B :$



$\text{int dp}[N+1][B+1] = \{-1\}$

$\text{dp}[N][B]$

$\text{int arr}[]$

$\text{bool isPossible}(\text{int } N, \text{int } B) \{$

    if  $(B == 0)$  return true

    if  $(B < 0)$  return false

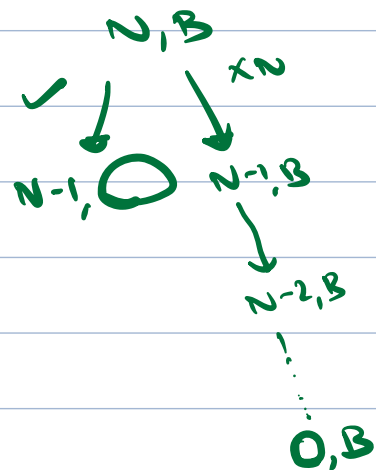
    if  $(N == 0)$  return false

    if  $(\text{dp}[N][B] \neq -1)$  return  $\text{dp}[N][B]$

$\text{dp}[N][B] = (\text{isPossible}(N-1, B - \text{arr}[N-1]) \ || \ \text{isPossible}(N-1, B))$

    return  $\text{dp}[N][B]$

}



TC :  $O(N * B)$

SC :  $O(N * B)$

RS  $\rightarrow \min(N, B)$

$A = \langle 3, 1, 4, 9, 7, 2 \rangle$

$B = 6$

$N=6$

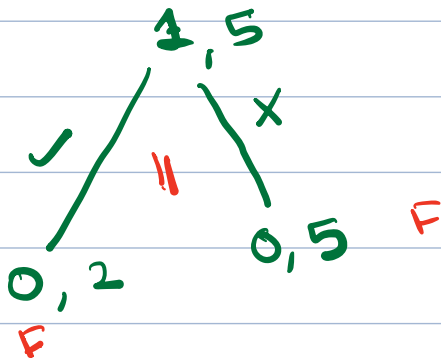
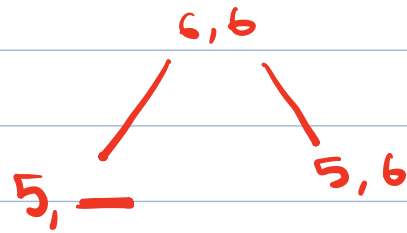
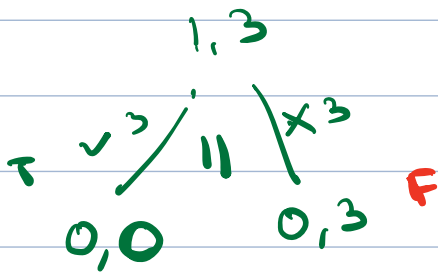
$B=6$

$dp[7][7]$

sum  $\rightarrow$

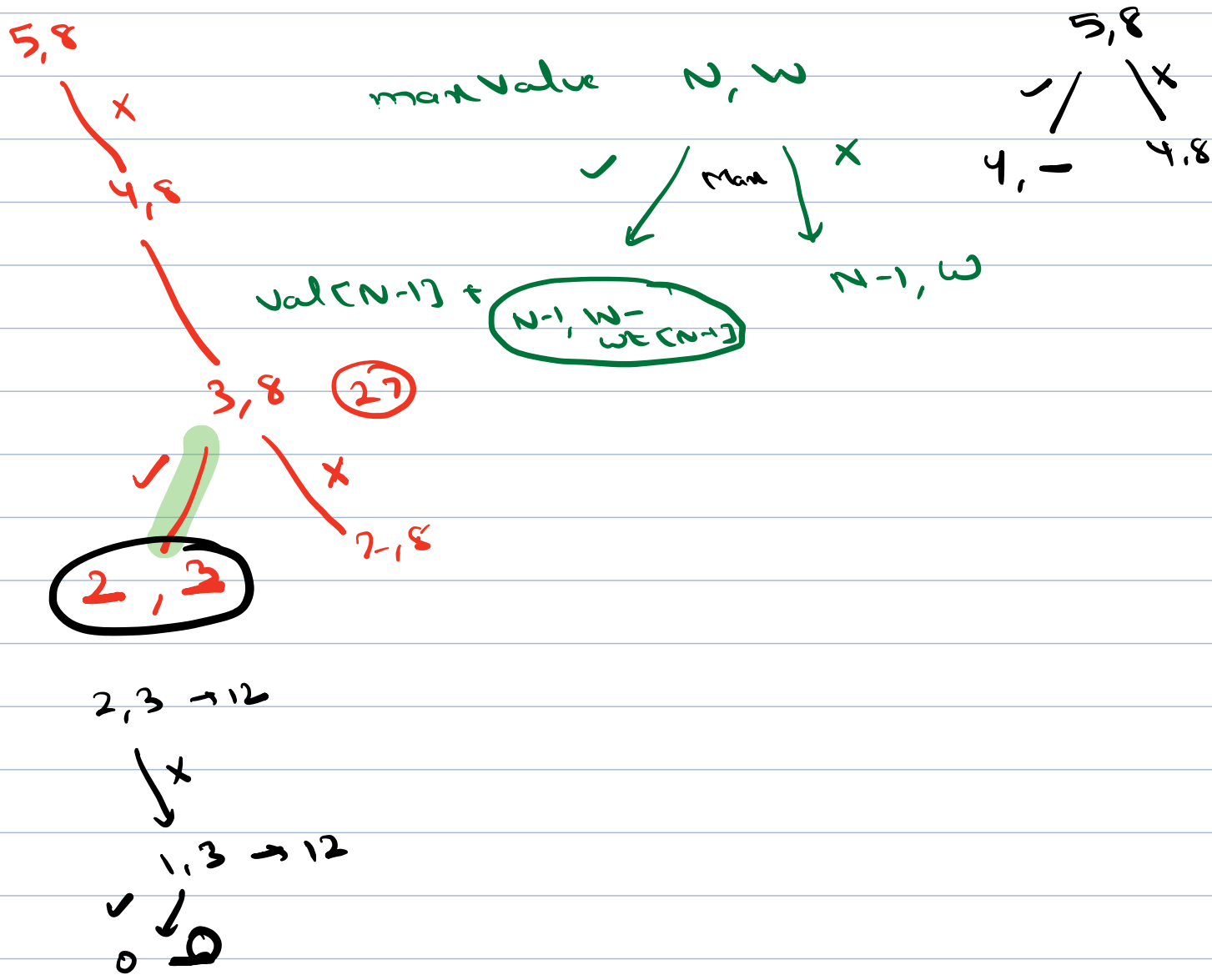
|   |   | 0 | 1 | 2 | 3 | 4 | 5   | 6 |
|---|---|---|---|---|---|---|-----|---|
|   | 0 | T | F | F | F | F | F   | F |
| 3 | 1 | T | F | F | T | F | * F | F |
| 1 | 2 | T | T | F | T | T | F   | F |
| 4 | 3 | T |   |   | * |   | *   |   |
| 9 | 4 | T |   |   |   |   |     |   |
| 7 | 5 | T |   |   |   |   |     |   |
| 2 | 6 | T |   |   |   |   |     | * |

Cnt  $\downarrow$



capacity / W  $\rightarrow$

| val | wt | cnt | 0 | 1 | 2 | 3  | 4  | 5  | 6  | 7  | 8  |
|-----|----|-----|---|---|---|----|----|----|----|----|----|
|     |    | 0   | 0 | 0 | 0 | 0  | 0  | 0  | 0  | 0  | 0  |
| 12  | 3  | 1   | 0 | 0 | 0 | 12 | 12 | 12 | 12 | 12 | 12 |
| 20  | 6  | 2   | 0 | 0 | 0 | 12 | 12 | 12 | 20 | 20 | 20 |
| 15  | 5  | 3   | 0 | 0 | 0 | 12 | 12 | 15 | 20 | 20 | 27 |
| 6   | 2  | 4   | 0 | 0 | 6 | 12 | 12 | 18 | 20 | 21 | 27 |
| 10  | 4  | 5   | 0 | 0 | 6 | 12 | 12 | 18 | 20 | 22 | 27 |



// dp array

$dp[N+1][W+1]$

$i = N, j = W$

while ( $i > 0$  &  $j > 0$ ) <

$dp(i, j)$

/x

$dp(i-1, j)$

if ( $dp[i][j] == dp[i-1][j]$ ) <

i--

>

else < // i<sup>th</sup> elem selected  $\rightarrow$  idk

print (i-1)

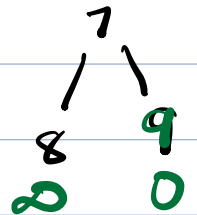
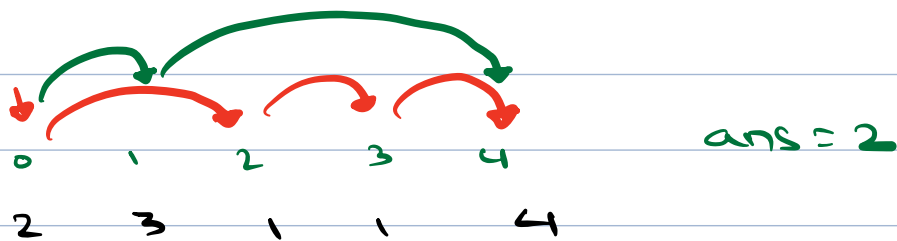
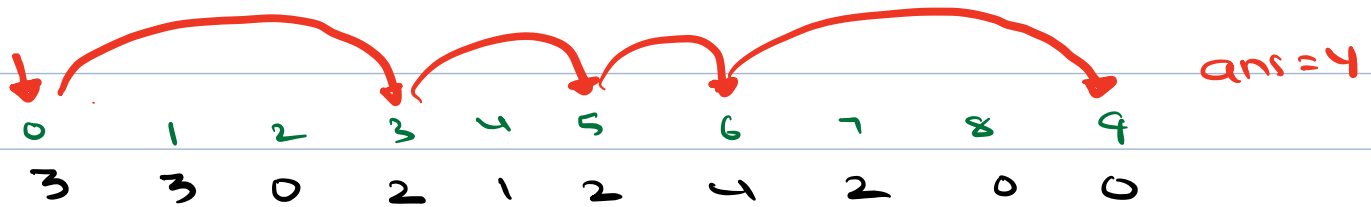
j -= wt[i-1]

i--

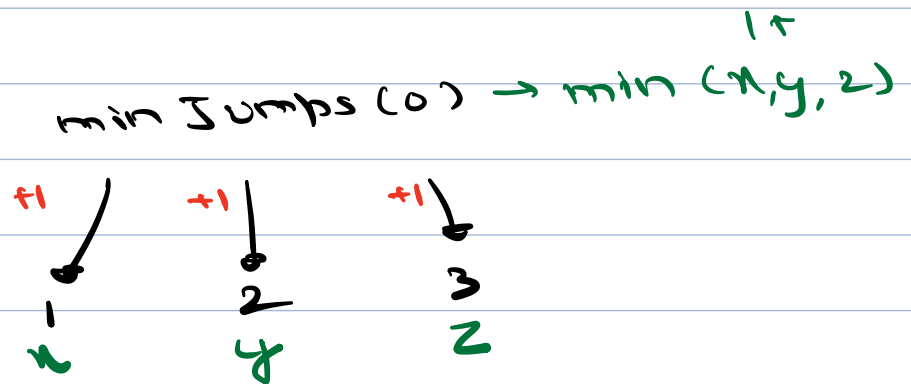
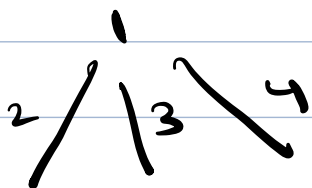
>

10:26

2. Given an integer array nums of length N. Initially, we're at 0, return min. no. of jumps to reach n-1.



|       | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | dp[n] |
|-------|---|---|---|---|---|---|---|---|---|---|-------|
|       | 3 | 3 | 0 | 2 | 1 | 2 | 4 | 2 | 0 | 0 |       |
| dp[i] | 4 | 4 | ∞ | 3 | 3 | 2 | 1 | 1 | ∞ | 0 |       |



dp[i] → Min jumps from i<sup>th</sup> idx to reach the end

```
int dp[n]
dp[n-1] = 0
```

i    i+1    i+2    i+3  
3

```
for (i = n-2; i ≥ 0; i--) <
```

```
    int minsteps = INT_MAX
```

```
    for (j = i+1; j ≤ min(N-1, i+A[i]); j++) <
```

```
        minsteps = min(minsteps, dp[j])
```

```
    if (minsteps != INT_MAX)
```

```
        dp[i] = minsteps + 1
```

```
    else
```

```
        dp[i] = minsteps
```

```
>
```

```
return dp[0]
```

Tc:  $O(N^2)$

Sc:  $O(N)$

↓

Tc:  $O(N)$

Sc:  $O(1)$



3. Given  $N$  and  $S$ , find  $N$  digit numbers with sum  $S$  (sum of  $N$  digits)

$$N = 2 \quad S = 4$$

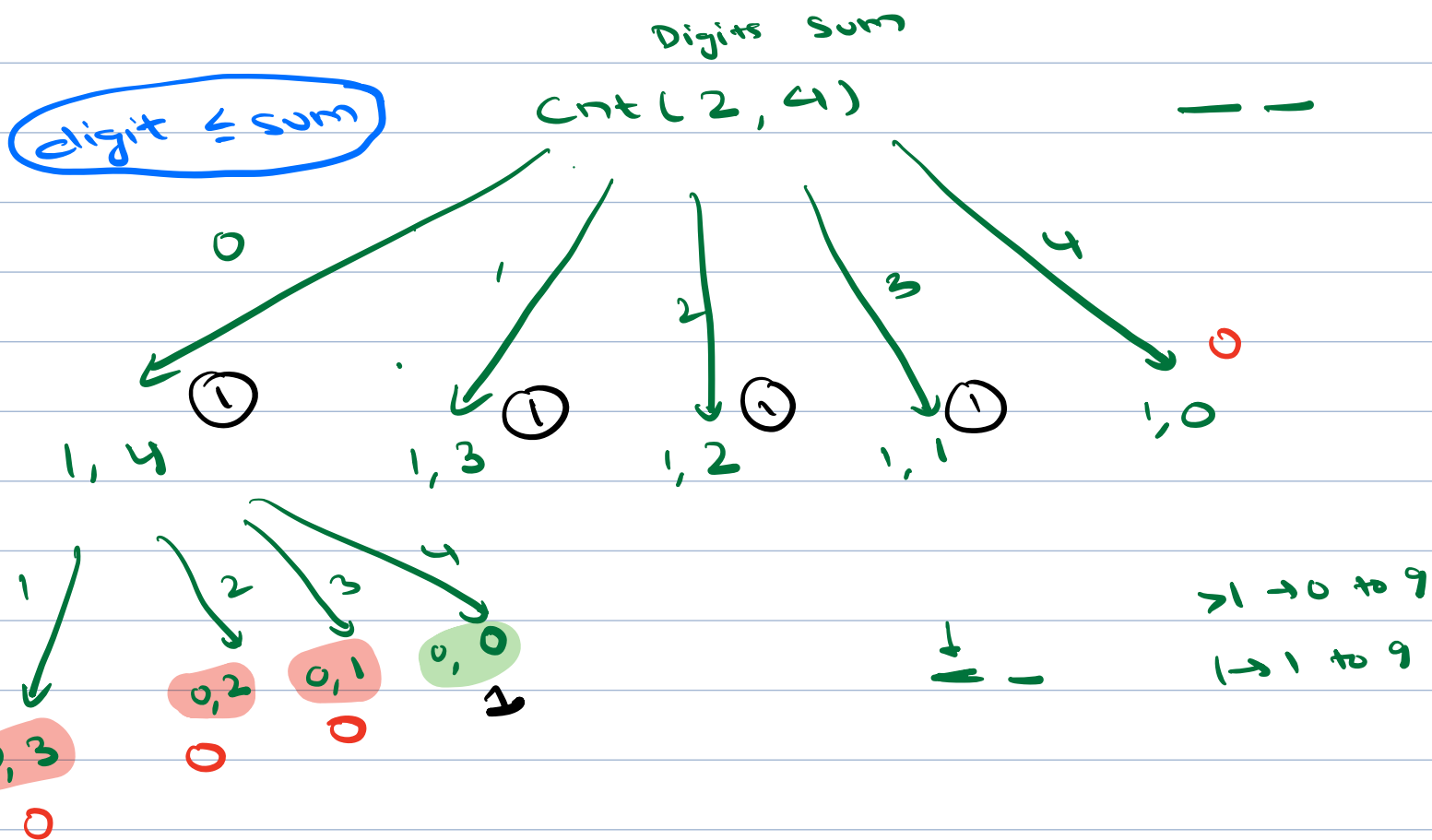
13, 31, 22, 40

ans = 4

$$N = 1 \quad S = 3$$

3

ans = 1

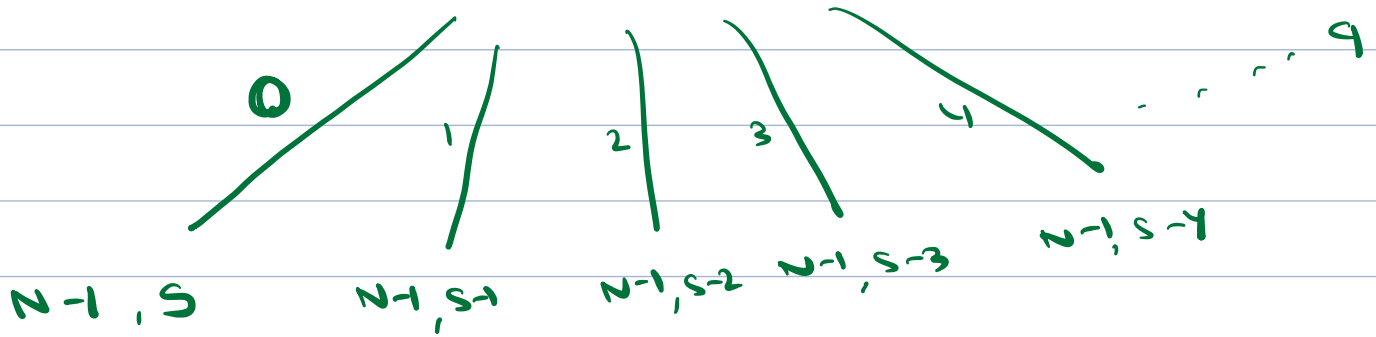


Sum = 0  $\rightarrow$  digit = 0      True

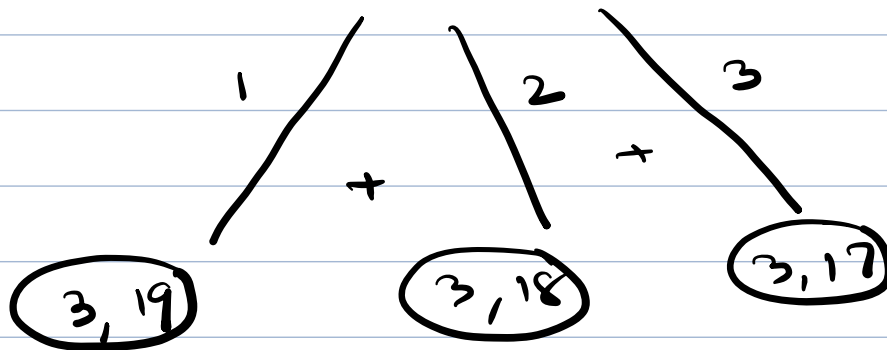
$\rightarrow$  digit = 0      False

digit = 0  $\rightarrow$  sum > 0      False

cnt(N, S)



→ cnt(4, 20)



19 ← 5

res → dp[N][S]

int dp[N+1][S+1] = {-1}

int cnt(int N, int S) {

if (N == 0 && S == 0) return 1

if (N == 0 || S == 0) return 0

if (dp[N][S] != -1) return dp[N][S]

int ans = 0

if (N == 1) {

for (d = 1; d ≤ 9; d++) {  
if (d ≤ S)  
ans += cnt(N-1, S-d)  
}

else {

for (d = 0; d ≤ 9; d++) {  
if (d ≤ S)  
ans += cnt(N-1, S-d)  
}

dp[N][S] = ans

return ans

TC:  $O(N * S)$

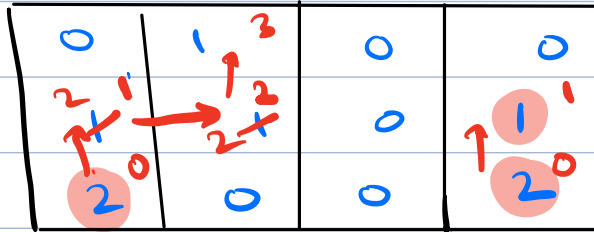
SC:  $O(N * S)$

#### 4. Rotten Oranges

mat[N][M]  $\rightarrow$  0 empty  
 $\rightarrow$  1 fresh  
 $\rightarrow$  2 rotten

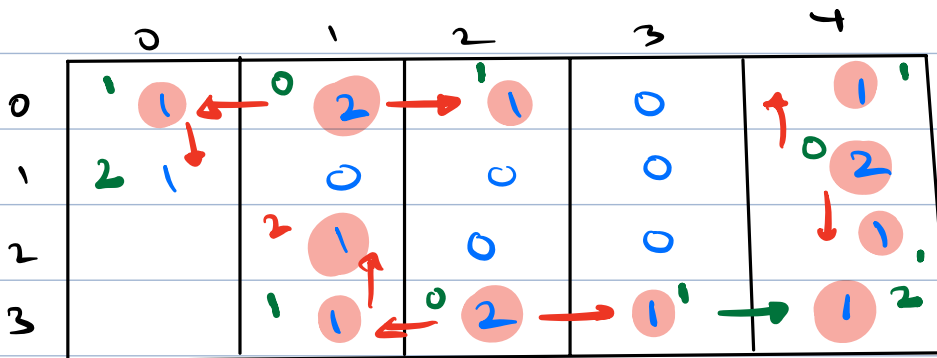
Every minute any fresh orange adjacent to a rotten orange becomes rotten, find min time when all oranges become rotten.

If not possible, return -1.

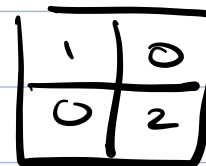


$T = 3$

3x4



$T = 2$



$T = -1$

|   | 0            | 1            | 2            | 3            | 4            |
|---|--------------|--------------|--------------|--------------|--------------|
| 0 | <del>2</del> | 2            | <del>2</del> | 0            | <del>2</del> |
| 1 | <del>2</del> | 0            | 0            | 0            | 2            |
| 2 | 0            | <del>2</del> | 0            | 0            | <del>2</del> |
| 3 | 0            | <del>2</del> | 2            | <del>2</del> | <del>2</del> |

Multi  
source  
BFS

~~(0,1,0)~~ ~~(1,4,0)~~ ~~(3,2,0)~~ ~~(0,0,1)~~ ~~(0,2,1)~~

~~(0,4,1)~~ ~~(2,4,1)~~ ~~(3,1,1)~~ ~~(3,3,1)~~

~~(1,0,2)~~ ~~(3,4,2)~~ ~~(2,1,2)~~

```
class triplet <
  int i, j, time
  >
```

queue < triplet > q

```
for (i=0 ; i < N ; i++) <
  for (j=0 ; j < M ; j++) <
    if (max [i] [j] == 2)
      q.enqueue (i, j, 0)
  >
>
```

```
int row[4] = {-1, 0, 1, 0}
```

```
int col[4] = {0, 1, 0, -1}
```

```
int maxtime = 0
```

```
while (!q.empty()) {
```

```
    Triplet t = q.front
```

```
    q.dequeue()
```

```
    maxtime = max(maxtime, t.time)
```

```
    for (int nbr = 0; nbr < 4; nbr++)
```

```
    {
```

```
        int nbri = t.i + row[nbr]
```

```
        int nbrj = t.j + col[nbr]
```

```
    }
```

```
    if (nbri >= 0 && nbri < n
```

```
        && nbrj >= 0 && nbrj < m
```

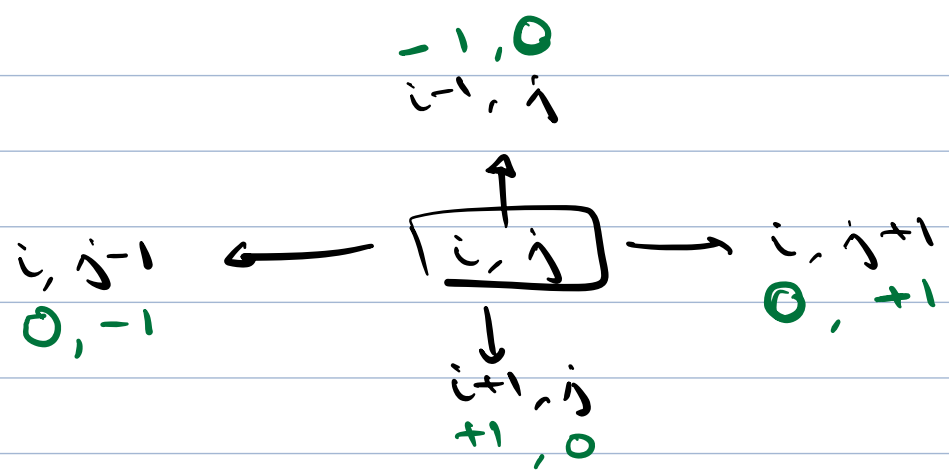
```
        && mat[nbri][nbrj] == 1) {
```

```
        q.enqueue(<nbri, nbrj, t.time + 1>)
```

```
        mat[nbri][nbrj] = 2
```

```
    }
```

```
}
```



int row [4] =  $\langle -1, 0, 1, 0 \rangle$

int col [4] =  $\langle 0, 1, 0, -1 \rangle$   
 U D L R