# **Chapter-3 Programming with 8086 microprocessor**

**Internal Architecture and Features of 8086 Microprocessor** 

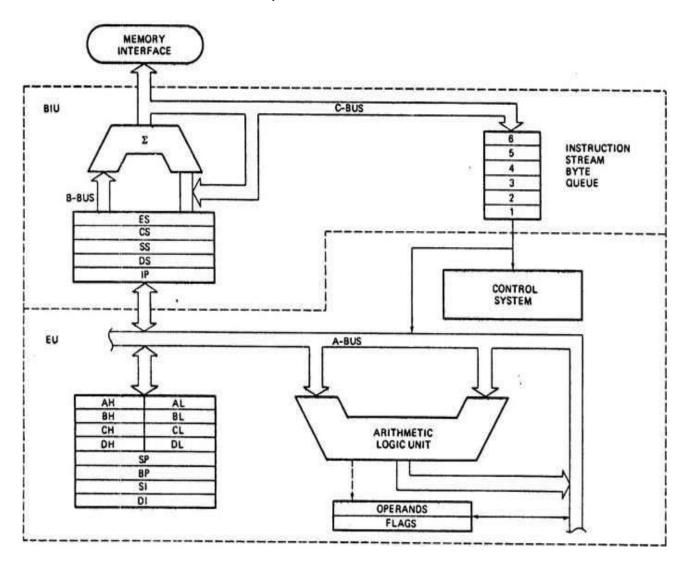


Fig: Internal Block Diagram of 8086 Microprocessor

## Features of 8086 microprocessor

- Intel 8086 is a widely used 16 bit microprocessor.
- The 8086 can directly address 1MB of memory.
- The internal architecture of the 8086 microprocessor is an example of register based microprocessor and it uses segmented memory.
- It pre-fetches up to 6 instruction bytes from the memory and queues them in order to speed up the instruction execution.

- It has data bus of width 16 bits and address bus of width 20 bits. So it always accesses a 16 bit word to or from memory.
- The 8086 microprocessor is divided internally into two separate units which are Bus interface unit (BIU) and the execution unit (EU).
- The BIU fetches instructions, reads operands and write results.
- The EU executes instructions that have already been fetched by BIU so that instructions fetch overlaps with execution.
- A 16 bit ALU in the EU maintains the MP status and control flags, manipulates general register and instruction operands.

## Bus Interface Unit(BIU) and its Components

The BIU sends out addresses, fetches instructions from memory reads data from memory or ports and writes data to memory or ports. So it handles all transfers of data and address on the buses for EU. It has main 2 parts instruction queue and segment registers.

- The BIU can store up to 6 bytes of instructions with FIFO (First in First Out) manner in a register set called a queue. When EU is ready for next instruction, it simply reads the instruction from the queue in the BIU. This is done in order to speed up program execution by overlapping instruction fetch with execution. This mechanism is known as pipelining.
- The BIU contains a dedicated address, which is used to produce 20 bit address. Four segment registers in the BIU are used to hold the upper 16 bits of the starting address of four memory segments that the 8086 is working at a particular time. These are code segment, data segment, stack segment and extra segment. The 8086's 1 MB memory is divided into segments up to 64KB each.
- Code segment register and instruction pointer (IP): The CS contains the base or start of the current code segment. The IP contains the distance or offset from this address to the next instruction byte to be fetched. Code segment address plus an offset value in the IP indicates the address of an instruction to be fetched for execution.

## - Data Segment

Data segment Contains the starting address of a program's data segment. Instructions use this address to locate data. This address plus an offset value in an instruction, causes a reference to a specific byte location in the data segment.

## - Stack segment (SS) and Stack Pointer (SP)

Stack segment Contains the starting address of a program's stack segment. This segment address plus an offset value in the stack pointer indicates the current word in the stack being addressed.

## Extra Segment(ES)

It is used by some string (character data) to handle memory addressing. The string instructions always use the ES and destination index (DI) to determine 20 bit physical address.

## **Execution Unit (EU)**

The EU decodes and executes the instructions. The EU contains arithmetic and logic (ALU), a control unit, and a number of registers. These features provide for execution of instructions and arithmetic and logical operations. It has nine 16 bit registers which are AX, BX, CX, DX, SP, BP, SI, DI and flag. First four can be used as 8 bit register (AH, AI, BH, BL, CH, DH, DL)

## AX Register

AX register is called 16 bit accumulator and AL is called 8 bit accumulator. The I/O (IN or OUT) instructions always use the AX or AL for inputting/Outputting 16 or 8 bit data from or to I/O port.

## BX Register

BX is known as the base register since it is the only general purpose register that can be used as an index to extend addressing. The BX register is similar to the 8085's H, L register. BX can also be combined with DI or SI as C base register for special addressing.

## - CX register:

The CX register is known as the counter register because some instructions such as SHIFT, ROTATE and LOOP use the contents of CX as a Counter.

#### - DX register:

The DX register is known as data register. Some I/O operations require its use and multiply and divide operations that involve large values assume the use of DX and AX together as a pair. DX comprises the rightmost 16 bits of the 32-bit EDX.

## - Stack Pointer (SP) and Base Pointer (BP):

Both are used to access data in the stack segment. The SP is used as an offset from the current stack segment during execution of instructions. The SP's contents are automatically updated (increment/decrement) during execution of a POP and PUSH instructions.

The BP contains the offset address in the current stack segment. This offset is used by instructions utilizing the based addressing mode.

## - Index register:

The two index registers SI (Source index) and DI (Destination Index) are used in indexed addressing. The instructions that process data strings use the SI and DI index register together with DS and ES respectively, in order to distinguish between the source and destination address.

## Flag register:

The 8086 has nine 1 bit flags. Out of 9 six are status and three are control flags. The control bits in the flag register can be set or reset by the programmer.

	 	0 -0						1 0	-		
		$\cap$	D	T	Т	ς	7		Δ	P	C
			$\mathcal{D}$	1	1	5			11	1	
											_
U <sub>15</sub>											$D_0$

- O- Overflow flag This flag is set if an arithmetic overflow occurs, i.e. if the result of a signed operation is large enough to be accommodated in a destination register.
- **D-Direction Flag** This is used by string manipulation instructions. If this flag bit is '0', the string is processed beginning from the lowest address to the higher address, i.e. auto incrementing mode otherwise the string is processed from the highest address towards the lowest address, i.e. autodecrementing mode.
- **I-Interrupt flag** If this flag is set the maskable interrupts are recognized by the CPU, otherwise they are ignored.
- T- Trap flag If this flag is set the processor enters the single step execution mode. In other words, a trap interrupt is generated after execution of each instruction. The processor executes the current instruction and the control is transferred to the Trap interrupt service routine.
- **S Sign flag:** This flag is set when the result of any computation is negative. For signed computations, the sign flag equals the MSB of the result.
- **Z- Zero** This flag is set when the result of the computation is or comparison performed by the previous instruction is zero. 1 for zero result, 0 fir nonzero result
- **A- Auxiliary Carry** This is set if there is a carry from the lowest nibble, i.e. bit three during the addition or borrow for the lowest nibble i.e. bit three, during subtraction.
- **P- Parity flag** This flag is set to 1 if the lower byte of the result contains even number of 1s otherwise reset.
- **C-Carry flag** This flag is set when there is a carry out of MSB in case of addition or a borrow in case of subtraction.

#### **SEGMENT AND OFFSET ADDRESS:**

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- Segments are special areas defined in a program for containing the code, data and stack. A segment begins on a paragraph boundary. A segment register is of 16 bits in size and contains the starting address of a segment.
- A segment begins on a paragraph boundary, which is an address divisible by decimal 16 or hex 10. Consider a DS that begins at location 038EOH. In all cases, the rightmost hex digit is zero, the computer designers decided that it would be unnecessary to store the zero the zero digit in the segment register. Thus 038E0H is stores in register as 038EH.
- The distance in bytes from the segment address to another location within the segment is expressed as an offset or displacement. Suppose the offset of 0032H for above example of data segment. Processor combines the address of the data segment with the offset as:

**SA: OA** (segment address: offset address)

038E: 0032 H = 038E \* 10 +0032

= 038EO + 0032

Physical address = 03912H

## **Instructions in 8086**

- 1) Arithmetic Instructions
- a) ADD reg<sub>8</sub>/mem<sub>8</sub>, reg<sub>8</sub>/mem<sub>8</sub>/ Immediate<sub>8</sub>

ADD reg<sub>16</sub>/mem<sub>16</sub>, reg<sub>16</sub>/ mem<sub>16</sub>/ Immediate<sub>16</sub>

E.g. ADD AH, 15; It adds binary number

ADD AH, NUM1

ADD AI, [BX]

ADD [BX], CH/CX ADD AX,[BX]

- b) ADC: Addition with Carry ADC reg/mem, reg/mem/Immediate data
- c) SUB: Subtract 8 bit or 16 bit binary numbers SUB reg/mem, reg/mem/Immediate
- d) **SBB:** Subtract with borrow SBB reg/mem, reg/mem/Immediate
- e) **MUL**: unsigned multiplication

MUL reg<sub>8</sub>/mem<sub>8</sub> (8 bit accumulator – AL)

 $MUL reg_{16} / mem_{16}$  (16 bit accumulator-Ax)

E.g. MUL 
$$R_8$$
 (multiplier)  $\equiv R_8 \times AL \rightarrow AX$  (16 bit result)

MUL 
$$R_{16}(multiplier) \equiv R_{16} \times AL \rightarrow DX:AX$$
 (32 bit result)

IMUL - signed multiplication

Same operation as MUL but takes sign into account

f) **DIV** reg/mem

E.g. DIV 
$$R_8 \equiv AX/R_8 \rightarrow (Remainder \rightarrow AH) \& (Quotient \rightarrow AL)$$
  
DIV  $R_{16} \equiv DX:AX/R_{16} \rightarrow (Remainder \rightarrow DX) & (Quotient \rightarrow AX)$ 

**IDIV- Signed division** 

Same operation as DIV but takes sign into account.

g) INC/DEC (Increment/Decrement by 1)

INC/DEC reg./mem. (8 bit or 16bit)

E.g. INC AL DEC BX

INC NUM1

## h) **NEG**- Negate (2's complement)

## i) ASCII-BCD Conversion

AAA: ASCII Adjust after addition AAS: ASCII Adjust after subtraction AAM: Adjust after multiplication AAD: Adjust after division

DAA: Decimal adjust after addition
DAS: Decimal adjust after subtraction

## 2) Logical/shifting/comparison instructions

## a) Logical

AND/OR/XOR reg/mem, reg/mem/immediate
NOT reg/mem
E. g. AND AL, AH
XOR [BX], CL

## b) Rotation

## **ROL- rotate left, ROR-rotate right**

E.g. ROL AX, 1 ; rotated by 1 ROL AX, CL ; if we need to rotate more than one bit

## RCL-rotate left through carry RCR-rotate right through carry

E.g. RCL AX, 1

RCL AX, CL ; Only CL can be used

## c) Shifting

**SHL** -logical shift left

SHR - logical shift right

Shifts bit in true direction and fills zero in vacant place

E.g. SHL reg/mem, 1/CL

## arithmetic shift left

**SAR**- arithmetic shift right

Shifts bit/word in true direction, in former case place zero in vacant place and in later case place previous sign in vacant place.

## d) Comparison

CMP -compare

CMP reg/mem, reg/mem/immediate

E.g. CMP BH, AL

Operand1		Operand 2	CF	SF	ZF
	>		0	0	0
	=		0	0	1
	<		1	1	0

TEST: test bits (AND operation)

TEST reg/mem, reg/mem/immediate

## 3) Data Transfer Instructions:

MOV reg./mem , reg./mem./immediate

LDS: Load data segment register

LEA: load effective address

LES: Load extra segment register LSS: Load stack segment register

E.g. LEA BX, ARR = MOV BX, OFFSET ARR

LDS BX, NUM1

Segment address → DS

Offset address → BX

XCHG reg/mem, reg/mem

E.g. XCHG AX, BX

XCHG AL, BL

XCHG CL,[BX]

IN AL, DX ; DX: Port address, AH also in AL

OUT DX, AL/AH

## 4) Flag Operation

CLC: Clear carry flag

CLD: Clear direction flag

CLI: Clear interrupt flag

STC: Set Carry flag

STD: Set direction flag

STI: Set Interrupt flag

CMC: Complement Carry flag

LAHF: Load AH from flags (lower byte)

SAHF: Store AH to flags

PUSHF: Push flags into stack

POPF: Pop flags off stack

## 5) **STACK Operations**

PUSH reg<sub>16</sub> POP reg<sub>16</sub>

## 6) Looping instruction (CX is automatically used as a counter)

LOOP: loop until complete LOOPE: Loop while equal LOOPZ: loop while zero

LOOPNE: loop while not equal LOOPNZ: loop while not zero

## 7) Branching instruction

## a) Conditional

JA: Jump if Above

JAE: Jump if above/equal

JB: Jump if below

JBE: Jump if below/equal

JC: Jump if carry JNC: Jump if no carry JE: Jump if equal JNE: Jump if no equal

JZ: Jump if zero
JNZ: Jump if no zero
JG: Jump if greater
JNG: Jump if no greater

JL: Jump if less
JNL: Jump if no less
JO: jump if overflow
JS: Jump if sign
JNS: Jump if no sign

JP: jump if plus

JPE: Jump if parity even JNP: Jump if no parity JPO: Jump if parity odd

## b) **Unconditional**

CALL: call a procedure RET: Return

INT: Interrupt IRET: interrupt return

JMP: Unconditional Jump RETN/RETF: Return near/Far

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## 8) Type conversion

CBW: Convert byte to word

CWD: Convert word to double word

 $AX \rightarrow DX: AX$ 

## 9) String instructions

a) MOVS/ MOVSB/MOVSW; Move string

DS: SI source DS: DI destination CX: String length

b) CMPS/ CMPSB/CMPW ; Compare stringc) LODS /LODSB/LODW ; Load stringd) REP ; Repeat string

#### Operators in 8086

- Operator can be applied in the operand which uses the immediate data/address.
- Being active during assembling and no machine language code is generated.
- Different types of operators are:
  - 1) Arithmetic: + , , \* , /
  - 2) Logical: AND, OR, XOR, NOT
  - 3) SHL and SHR: Shift during assembly
  - 4) []: index
  - 5) HIGH: returns higher byte of an expression
  - 6) **LOW**: returns lower byte of an expression.

E.g. NUM EQU 1374 H

MOV AL HIGH Num ; ( [AL]  $\leftarrow$  13 )

- 7) **OFFSET**: returns offset address of a variable
- 8) **SEG**: returns segment address of a variable
- 9) **PTR**: used with type specifications

BYTE, WORD, RWORD, DWORD, QWORD

E.g. INC BYTE PTR [BX]

## 10) Segment override

MOV AH, ES: [BX]

- 11) **LENGTH**: returns the size of the referred variable
- 12) **SIZE:** returns length times type

E.g.: BYTE VAR DB?

WTABLE DW 10 DUP (?)

 $\begin{array}{lll} \text{MOV AX, TYPE BYTEVAR} & ; \text{ AX = } 0001\text{H} \\ \text{MOV AX, TYPE WTABLE} & ; \text{ AX = } 0002\text{H} \\ \text{MOV CX, LENGTH WTABLE} & ; \text{ CX = } 000\text{AH} \\ \text{MOV CX, SIZE WTABLE} & ; \text{ CX = } 0014\text{H} \\ \end{array}$ 

## **Coding in Assembly language:**

Assembly language programming language has taken its place in between the machine language (low level) and the high level language.

- High level language's one statement may generate many machine instructions.
- Low level language consists of either binary or hexadecimal operation. One symbolic statement generates one machine level instructions.

## **Advantage of ALP**

- They generate small and compact execution module.
- They have more control over hardware.
- They generate executable module and run faster.

## **Disadvantages of ALP:**

- Machine dependent.
- Lengthy code
- Error prone (likely to generate errors).

## **Assembly language features:**

The main features of ALP are program comments, reserved words, identifies, statements and directives which provide the basic rules and framework for the language.

## **Program comments:**

- The use of comments throughout a program can improve its clarity.
- It starts with semicolon (;) and terminates with a new line.
- E.g. ADD AX, BX ; Adds AX & BX

## Reserved words:

- Certain names in assembly language are reserved for their own purpose to be used only under special conditions and includes
- Instructions : Such as MOV and ADD (operations to execute)
- Directives: Such as END, SEGMENT (information to assembler)
- Operators: Such as FAR, SIZE
- Predefined symbols: such as @DATA, @ MODEL

## **Identifiers:**

- An identifier (or symbol) is a name that applies to an item in the program that expects to reference.
- Two types of identifiers are Name and Label.
- Name refers to the address of a data item such as NUM1 DB 5, COUNT DB 0
- Label refers to the address of an instruction.
- E. g: MAIN PROC FAR
- L1: ADD BL, 73

#### **Statements:**

- ALP consists of a set of statements with two types
- Instructions, e. g. MOV, ADD
- Directives, e. g. define a data item

Identifiers operation operand comment
Directive: COUNT DB 1 ; initialize count
Instruction: L30: MOV AX, 0 ; assign AX with 0

## **Directives:**

The directives are the number of statements that enables us to control the way in which the source program assembles and lists. These statements called directives act only during the assembly of program and generate no machine-executable code. The different types of directives are:

## 1) The page and title listing directives:

The page and title directives help to control the format of a listing of an assembled program. This is their only purpose and they have no effect on subsequent execution of the program.

The page directive defines the maximum number of lines to list as a page and the maximum number of characters as a line.

PAGE [Length] [Width] Default: Page [50][80]

TITLE gives title and place the title on second line of each page of the program.

TITLE text [comment]

## 2) **SEGMENT** directive

It gives the start of a segment for stack, data and code.

Seg-name Segment [align][combine]['class']

Seg-name ENDS

- Segment name must be present, must be unique and must follow assembly language naming conventions.
- An ENDS statement indicates the end of the segment.
- Align option indicates the boundary on which the segment is to begin; PARA is used to align the segment on paragraph boundary.
- Combine option indicates whether to combine the segment with other segments when they are linked after assembly. STACK, COMMON, PUBLIC, etc are combine types.
- Class option is used to group related segments when linking. The class code for code segment, stack for stack segment and data for data segment.

#### 3) PROC Directives

The code segment contains the executable code for a program, which consists of one or more procedures, defined initially with the PROC directives and ended with the ENDP directive.

PROC - name PROC [FAR/NEAR]

.....

PROC - name ENDP

- FAR is used for the first executing procedure and rest procedures call will be NEAR.
- Procedure should be within segment.

## 4) END Directive

- An END directive ends the entire program and appears as the last statement.
- ENDS directive ends a segment and ENDP directive ends a procedure. END PROC-Name

## 5) ASSUME Directive

- An .EXE program uses the SS register to address the stack, DS to address the data segment and CS to address the code segment.
- Used in conventional full segment directives only.
- Assume directive is used to tell the assembler the purpose of each segment in the program.
- Assume SS: Stack name, DS: Data Segname CS: codesegname

## 6) Processor directive

- Most assemblers assume that the source program is to run on a basic 8086 level computer.
- Processor directive is used to notify the assembler that the instructions or features introduced by the other processors are used in the program.
  - E.g. .386 program for 386 protected mode.

## 7) Dn Directive (Defining data types)

Assembly language has directives to define data syntax [name] Dn expression The Dn directive can be any one of the following:

DB Define byte 1 byte DW Define word 2 bytes Define double DD 4 bytes DF defined farword 6 bytes Define quadword 8 bytes DQ DT Define 10 bytes 10 bytes

#### **VAL1** DB **25**

**ARR** DB 21, 23, 27, 53 MOV AL, ARR [2] or

MOV AL, ARR + 2; Moves 27 to AL register

## 8) The EQU directive

- It can be used to assign a name to constants.
- E.g. **FACTOR** EQU **12**

- MOV BX, FACTOR ; MOV BX, 12
- It is short form of equivalent.
- Do not generate any data storage; instead the assembler uses the defined value to substitute in.

## 9) DUP Directive

- It can be used to initialize several locations to zero.
  - e. g. SUM DW 4 DUP(0)
- Reserves four words starting at the offset sum in DS and initializes them to Zero.
- Also used to reserve several locations that need not be initialized. In this case (?) is used with DUP directives.
  - E. g. **PRICE** DB 100 DUP(?)
- Reserves 100 bytes of uninitialized data space to an offset PRICE.

## Program written in Conventional full segment directive

```
Page 60,132
TITLE SUM program to add two numbers
<u>-----</u>
STACK SEGMENT PARA STACK 'Stack'
DW 32 DUP(0)
STACK ENDS
;-----
DATA SEG SEGMENT PARA 'Data'
NUM1 DW 3291
NUM 2 DW 582
SUM DW?
DATA SEG ENDS
;-----
CODE SEG SEGMENT PARA 'Code'
MAIN PROC FAR
  ASSUME SS: STACK, DS:DATASEG, CS:CODESEG
  MOV AX, @DATA
  MOV DS, AX
  MOV AX, NUM1
  ADD AX, NUM2
  MOV AX, 4C00H
  INT 21H
MAIN ENDP
CODESEG ENDS
END MAIN
```

#### **Description for conventional program:**

- STACK contains one entry, DW (define word), that defines 32 words initialized to zero, an adequate size for small programs.
- DATASEG defines 3 words NUM1, NUM2 initialized with 3291 and 582 and sum uninitialized.
- CODESEG contains the executable instructions for the program, PROC and ASSUME generate no executable code.
- The ASSUME directive tells the assembler to perform these tasks.
- Assign STACK to SS register so that the processor uses the address in SS for addressing STACK.
- Assign DATASEG to DS register so that the processor uses the address in DS for addressing DATASEG.
- Assign CODESEG to the CS register so that the processor uses the address in CS for addressing CODESEG.

When the loading a program for disk into memory for execution, the program loader sets the correct segment addresses in SS and CS.

## Program written using simplified segment directives:

.Model memory model

Memory model can be

TINY, SMALL, MEDIUM, COMPACT, LARGE, HUGE or FLAT

TINY for .com program

FLAT for program up to 4 GB

- **Assume** is automatically generated

.STACK [size in bytes]

Creates stack segment

.DATA: start of data segment

.CODE: start of code segment

- DS register can be initialized as

MOV AX, @DATA

MOV DS, AX

## **ALP written in simplified segment directives:**

Page 60, 132

TITLE Sum program to add two numbers.

.MODEL SMALL

.STACK 64

.DATA

NUM1 DW 3241

NUM 2 DW 572

SUM DW ? .CODE MAIN PROC FAR MOV AX, @ DATA ; set address of data segment in DS MOV DS, AX MOV AX, NUM1 ADD AX, NUM2 MOV SUM, AX MOV AX, 4C00H ; End processing INT 21H MAIN ENDP ; End of procedure ; End of program END MAIN

## DOS Debug(TASM)

- 1) Save the code text in **.ASM** format and save it to the same folder where masm and link files are stored.
- 2) Open dos mode and reach within that folder.

3)	\> tasm filename.asm	> makes.obj
4)	\> tlink filename	>makes .exe
5)	\> filename.exe	run the code
6)	\> td filename.exe	debug the code [use F7 and F8]

## Assembling, Linking and Executing

## 1) Assembling:

- Assembling converts source program into object program if syntactically correct and generates an intermediate **.obj** file or module.
- It calculates the offset address for every data item in data segment and every instruction in code segment.
- A header is created which contains the incomplete address in front of the generated **obj** module during the assembling.
- Assembler complains about the syntax error if any and does not generate the object module.
- Assembler creates **.obj .lst** and **.crf** files and last two are optional files that can be created at run time.
- For short programs, assembling can be done manually where the programmer translates each mnemonic into the machine language using lookup table.
- Assembler reads each assembly instruction of a program as ASCII character and translates them into respective machine code.

## **Assembler Types:**

There are two types of assemblers:

## a) One pass assembler:

- This assembler scans the assembly language program once and converts to object code at the same time.

- This assembler has the program of defining forward references only.
- The jump instruction uses an address that appears later in the program during scan, for that case the programmer defines such addresses after the program is assembled.

## b) Two pass assembler

- This type of assembler scans the assembly language twice.
- First pass generates symbol table of names and labels used in the program and calculates their relative address.
- This table can be seen at the end of the list file and here user need not define anything.
- Second pass uses the table constructed in first pass and completes the object code creation.
- This assembler is more efficient and easier than earlier.

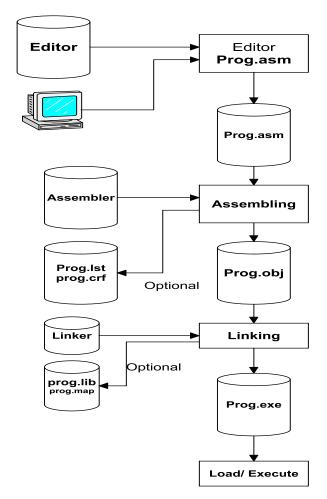


Fig: Steps in assembling, linking & Executing

## 2) Linking:

- This involves the converting of **.OBJ** module into **.EXE**(executable) module i.e. executable machine code.
- It completes the address left by the assembler.

- It combines separately assembled object files.
- Linking creates .EXE, .LIB, .MAP files among which last two are optional files.

## 3) Loading and Executing:

- It Loads the program in memory for execution.
- It resolves remaining address.
- This process creates the program segment prefix (PSP) before loading.
- It executes to generate the result.

Sample program <u>assembling</u> object Program <u>linking</u> executable program

## **Writing .COM programs:**

- It fits for memory resident programs.
- Code size limited to 64K.
- .com combines PSP, CS, DS in the same segment
- SP is kept at the end of the segment (FFFF), if 64k is not enough, DOS Places stack at the end of the memory.
- The advantage of **.com** program is that they are smaller than **.exe** program.
- A program written as .com requires ORG 100H immediately following the code segment's SEGMENT statement. The statement sets the offset address to the beginning of execution following the PSP.

```
.MODEL TINY
.CODE
ORG 100H
                        ; start at end of PSP
                         ;Jump Past data
BEGIN:JMP MAIN
      VAL1 DW 5491
      VAL2 DW 372
      SUM DW?
MAIN: PROC NEAR
      MOV Ax, VALL
      ADD AX, VAL2
      MOV SUM, AX
      MOV AX, 4C00H
      INT 21H
MAIN ENDP
END BEGIN
```

#### **Macro Assembler:**

- A macro is an instruction sequence that appears repeatedly in a program assigned with a specific name.
- The macro assembler replaces a macro name with the appropriate instruction sequence each time it encounters a macro name.

- When same instruction sequence is to be executed repeatedly, macro assemblers allow the macro name to be typed instead of all instructions provided the macro is defined.
- Macro are useful for the following purposes:
  - o To simplify and reduce the amount of repetitive coding.
  - o To reduce errors caused by repetitive coding.
  - o To make an assembly language program more readable.
  - o Macro executes faster because there is no need to call and return.
  - Basic format of macro definition:

	Macro name	MACRO [Parameter list]	; Define macro
		[Instructions]	; Macro body
		ENDM	; End of macro
E.g.	Addition	MACRO IN AX, PORT ADD AX, BX OUT PORT, AX ENDM	

## Passing argument to MACRO:

- To make a macro more flexible, we can define parameters as dummy argument

```
Addition MACRO VALL1, VAL2
      MOV AX, VAL1
      ADD AX, VAL2
      MOV SUM, AX
      ENDM
.MODEL SMALL
.STACK 64
.DATA
     VAL1 DW
                 3241
                 571
     VAL2 DW
      SUM DW
                 ?
.CODE
MAIN PROC FAR
      MOV AX, @ DATA
      MOV DS, AX
```

Addition VAL1, VAL 2 MOV AX, 4C00H INT 21H MAIN ENDP END MAIN

## Addressing modes in 8086:

Addressing modes describe types of operands and the way in which they are accessed for executing an instruction. An operand address provides source of data for an instruction to process an instruction to process. An instruction may have from zero to two operands. For two operands first is destination and second is source operand. The basic modes of addressing are register, immediate and memory which are described below.

## 1) Register Addressing:

For this mode, a register may contain source operand, destination operand or both.

E.g. MOV AH, BL MOV DX, CX

#### 2) Immediate Addressing

In this type of addressing, immediate data is a part of instruction, and appears in the form of successive byte or bytes. This mode contains a constant value or an expression.

E.g. MOV AH, 35H MOV BX, 7A25H

## 3) <u>Direct memory addressing:</u>

In this type of addressing mode, a 16-bit memory address (offset) is directly specified in the instruction as a part of it. One of the operand is the direct memory and other operand is the register.

E.g. ADD AX, [5000H]

**Note:** Here data resides in a memory location in the data segment, whose effective address may be computed using 5000H as the Offset address and content of DS as segment address. The effective address, here, is **10H\*DS + 5000H**.

## 4) Direct offset addressing

In this addressing, a variation of direct addressing uses arithmetic operators to modify an address.

E.g. ARR DB 15, 17, 18, 21

MOV AL, ARR [2] ; MOV AL, 18 ADD BH, ARR+3 ; ADD BH, 21

## 5) Indirect memory addressing:

Indirect addressing takes advantage of computer's capability for segment: offset addressing. The registers used for this purpose are base register (BX and BP) and index register (DI and SI)

E.g. MOV [BX], AL

ADD CX, [SI]

## 6) Base displacement addressing:

This addressing mode also uses base registers (BX and BP) and index register (SI and DI), but combined with a displacement (a number or offset value) to form an effective address.

## 7) Base index addressing:

This addressing mode combines a base registers (BX or BP) with an index register (SI or DI) to form an effective address.

```
E.g. MOV AX, [BX +SI]
ADD [BX+DI], CL
```

## 8) Base index with displacement addressing

This addressing mode, a variation on base- index combines a base register, an index register, and a displacement to form an effective address.

```
E.g. MOV AL, [Bx+SI+2]
ADD TBL [BX +SI], CH
```

## 9) String addressing:

This mode uses index registers, where SI is used to point to the first byte or word of the source string and DI is used to point to the first byte or word of the destination string, when string instruction is executed. The SI or DI is automatically incremented or decremented to point to the next byte or word depending on the direction flag (DF). E.g. MOVS, MOVSB, MOVSW

## **Examples:**

## TITLE Program to add ten numbers

```
.MODEL SMALL
.STACK 64
.DATA
ARR DB 73, 91, 12, 15, 79, 94, 55, 89
SUM DW ?
.CODE
MAIN PROC FAR
MOV AX, @DATA
MOV DS, AX
```

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```
MOV CX, 10
MOV AX, 0
LEA BX, ARR
L2: ADD AI, [BX]
JNC L1
INC AH
L1: INC BX
LOOP L2
MOV SUM, AX
MOV AX, 4C00H
INT 21H
MAIN ENDP
END MAIN
```

## DOS FUNCTIONS AND INTERRUPTS (KEYBOARD AND VIDEO PROCESSING)

The Intel CPU recognizes two types of interrupts namely hardware interrupt when a peripheral devices needs attention from the CPU and software interrupt that is call to a subroutine located in the operating system. The common software interrupts used here are INT 10H for video services and INT 21H for DOS services.

#### <u>INT 21H:</u>

It is called the DOS function call for keyboard operations follow the function number. The service functions are listed below:

## # 00H- It terminates the current program.

- Generally not used, function 4CH is used instead.

## # 01H- Read a character with echo

- Wait for a character if buffer is empty
- Character read is returned in AL in ASCII value

## # 02H- Display single character

- Sends the characters in DL to display
- MOV AH, 02H
- MOV DL, 'A'; move Dl, 65
- INT 21H

## # 03H and 04H - Auxiliary input/output

- INT 14H is preferred.

#### # 05H - Printer service

- Sends the character in DL to printer

## # 06H- Direct keyboard and display

Displays the character in DL.

## # 07H- waits for a character from standard input

does not echo

## # 08H- keyboard input without echo

- Same as function 01H but not echoed.

## # 09H- string display

- Displays string until '\$' is reached.
- DX should have the address of the string to be displayed.

## # 0AH - Read string

## # OBH- Check keyboard status

- Returns FF in AL if input character is available in keyboard buffer.
- Returns 00 if not.

# OCH- Clear keyboard buffer and invoke input functions such as 01, 06, 07, 08 or 0A.

- AL will contain the input function.

## **INT 21H Detailed for Useful Functions**

#### # 01H

MOV, AH 01H; request keyboard input INT 21H

- Returns character in AL. IF AL= nonzero value, operation echoes on the screen. If Al= zero means that user has pressed an extended function key such as F1 OR home.

#### # 02H

```
MOV AH, 02H; request display character
MOV DL, CHAR; character to display
INT 21H
```

- Display character in D2 at current cursor position. The tab, carriage return and line feed characters act normally and the operation automatically advances the cursor.

## # 09H

```
MOV Ah, 09H; request display
LEA DX, CUST_MSG; local address of prompt
INNT 21H
CUST_MSG DB "Hello world", '$'
```

- Displays string in the data area, immediately followed by a dollar sign (\$ or 24H), which uses to end the display.

#### # OAH

```
MOV AH, 0AH ; request keyboard input LEA DX, PARA_ LIST ; load address of parameter list INT 21H
```

## Parameter list for keyboard input area:

PARA\_LIST LABEL BYTE; start of parameter list MAX LEN DB 20; max. no. of input character

ACT \_ LEN DB ?; actual no of input characters
KB-DATA DB 20 DUP ('); characters entered from keyboard

- LABEL directive tells the assembler to align on a byte boundary and gives location the name PARA \_LIST.
- PARA\_LIST & MAX\_LEN refer same memory location, MAX\_LEN defines the maximum no of defined characters.
- ACT\_LEN provides a space for the operation to insert the actual no of characters entered.
- KB\_DATA reserves spaces (here 20) for the characters.

## **Example:**

```
TITLE to display a string
.MODEL SMALL
.STACK 64
.DATA
      STR DB 'programming is fun', '$'
.CODE
MAIN PROC FAR
      MOV AX, @DATA
      MOV DS, AX
      MOV AH, 09H
                          ;display string
      LEA DX, STR
      INT 21H
      MOV AX, 4C00H
      INT 21H
MAIN ENDP
END MAIN
```

#### INT 10H

It is called video display control. It controls the screen format, color, text style, making windows, scrolling etc. The control functions are:

```
# 00H – set video mode

MOV AH, 00H ; set mode

MOV AL, 03H ; standard color text
INT 10H ; call interrupt service

# 01H- set cursor size

MOV AH, 01H

MOV CH, 00H ; Start scan line

MOV CL, 14H ; End scan line
INT 10H ; (Default size 13:14)
```

```
# 02H – Set cursor position:
       MOV AH, 02H
       MOV BH, 00H
                           ; page no
                           ; row/y (12)
       MOV DH, 12H
       MOV DL, 30H
                           ; column/x (30)
       INT 10H
# 03H - return cursor status
       MOV AH, 03H
       MOV BH, 00H;
       INT 10H
       Returns: CH- starting scan line, CL-end scan line, DH- row, DL-column
# 04H- light pen function
# 05H- select active page
       MOV AH, 05H
       MOV AL,page-no.
                           ; page number
       INT 10H
# 06H- scroll up screen
MOV AX, 060FH
                    ; request scroll up one line (text)
MOV BH, 61H
                    ; brown background, blue foreground
MOV CX, 0000H
                    ; from 00:00 through
MOV DX, 184F H
                    ; to 24:79 (full screen)
INT 10H
AL= number of rows (00 for full screen)
BH= Attribute or pixel value
CX= starting row: column
DX= ending row: column
# 07H-Scroll down screen
Same as 06H except for down scroll
# 08H (Read character and Attribute at cursor)
MOV AH, 08H
MOV BH, 00H
                    ; page number 0(normal)
INT 10H
AL= character
BH= Attribute
# 09H -display character and attribute at cursor
```

MOV AH, 09H

MOV AL, 01H

; ASCII for happy face display

MOV BH, 00H ; page number

MOV BL, 16H ; Blue background, brown foreground

MOV CX, 60 ; No of repeated character

INT 10H

## # 0AH-display character at cursor

MOV AH, 0AH

MOV Al, Char

MOV BH, page no

MOV BL, value

MOV CX, repetition

INT 10H

## # 0BH- Set color palette

✓ Sets the color palette in graphics mode

✓ Value in BH (00 or 01) determines purpose of BL

✓ BH= 00H, select background color, BL contains 00 to 0FH (16 colors)

✓ BH = 01H, select palette, Bl, contains palette

MOV AH, 0BH MOV AH, 0BH

MOV BH, 00H; background MOV BH, 01H; select palette

MOV BL, 04H; red MOV BL, 00H; black

INT 21H INT 21H

## #0CH- write pixel Dot

- Display a selected color

AL=color of the pixel CX= column BH=page number DX= row

MOV AH, 0CH

MOV AI, 03

MOV BH,0

**MOV CX, 200** 

MOV DX, 50

**INT 10H** 

It sets pixel at column 200, row 50

## #0DH- Read pixel dot

Reads a dot to determine its color value which returns in AL

MOV AH, 0DH

MOV BH, 0 ; page no MOV CX, 80 ; column MOV DX, 110 ; row

INT 10H

## **#OEH- Display in teletype mode**

- Use the monitor as a terminal for simple display

MOV AH, 0EH

MOV AL, char

MOV BL, color; foreground color

INT 10H

#### **#OF H- Get current video mode**

Returns values from the BIOS video .

AL= current video mode MOV AH, 0FH

AH= no of screen columns INT 10H

BH = active video page

#### **TITLE To Convert letters into lower case**

.MODEL SMALL

.STACK 99H

.CODE

MAIN PROC

MOV AX, @ DATA

MOV DS, AX

MOV SI, OFFSER STR

M: MOV DL, [SI]

MOV CL, DL

CMP DL, '\$'

JE N

CMP DL, 60H

JL L

K: MOV DL, CL

MOV AH, 02H

INT 21H

**INC SI** 

JMP M

L: MOV DL, CL

ADD DL, 20H

MOV AH, 02H

**INT 21H** 

**INC SI** 

JMP M

N: MOV AX, 4C00H

INT 21H

MAIN ENDP

.DATA

STR DB 'I am MR Rahul ", '\$'

#### **END MAIN**

```
TITLE to reverse the string
       .MODEL SMALL
       .STACK 100H
       .DATA
          STR1 DB " My name is Rahul", '$'
          STR2 db 50 dup ('$')
        .CODE
   MAIN PROC FAR
          MOV BL,00H
          MOV AX, @ DATA
          MOV DS, AX
          MOV SI, OFFSER STR1
          MOV DI, OFFSET STR2
   L2:
          MOV DL, [SI]
          CMP DI, '$'
          JE L1
          INC SI
          INC BL
          JMP L2
   L1:
          MOV CL, BL
          MOV CH, 00H
          DEC SI
   L3:
          MOV AL, [SI]
          MOV [DI], AL
          DEC SI
          INC DI
          LOOP L3
          MOV AH,09H
          MOV DX, OFFSET STR2
          INT 21H
          MOV AX, 4C00H
          INT 21H
   MAIN ENDP
   END MAIN
TITLE to input characters until 'q' and display
   .MODEL SMALL
   .STACK 100H
   .DATA
          STR db 50 DUP ('$')
    .CODE
   MAIN PROC FAR
```

MOV AX, @ DATA MOV DS, AX

MOV SI, OFFSET STR

L2: MOV AH, 01H

INT 21H

CMP AL, 'q'

JE L1

MOV [SI], AL

**INC SI** 

JMP L2

L1: MOV AH, 09H

MOV DX, OFFSET STR

INT 21H

MOV AX, 4C00H

INT 21H

MAIN ENDP END MAIN

## **Calling procedure/subroutine**

Compiled by: Er. Hari Aryal

Procname PROC FAI
Procname FNDP

Here the code segment consists only one procedure. The FAR operand in this case informs the assembler and linker that the defined procedure name is the entry point for program execution, whereas the ENDP directive defines the end of the procedure. A code segment however, may contain any number of procedures, each distinguished by its own PROC and ENDP directives.

A called procedure is a section of code that performs a clearly defined task known as subroutine which provides following benefits.

- Reduces the amount of code because a common procedure can be called from any number of places in the code segment.
- Encourage better program organization.
- Facilitates debugging of a program because defects can be more clearly isolated.
- Helps in the ongoing maintenance of programs because procedures are readily identified for modification.

A CALL to a procedure within the same code segment is NEAR CALL<. A FAR CALL calls a procedure labeled FAR, possibly in another code segment.

```
DISPLAY PROC NEAR

MOV AH, 09H

MOV DX, OFFSET STR

INT 21H

RET

DISPLAY ENDP

✓ To display number contained in [BX]

DISPLAY PROC NEAR

MOV DI, [BX]

ADD DI, 30

MOV AH, 02H

INT 21H
```

## **INT 10H Video service:**

RET

**DISPLAY ENDP** 

## <Video -modes>

Text mode	Row×column	Color	No.of	Resolution	colors
			Pages		
00	25×40	Color	8	360×400	16 colors
01	25×40	Color	8	360×400	16 colors
02	25×80	Color	4	720×400	16 colors
03(by	25×80	color	4	720×400	16 colors
default)					
07	25×80	Mono-hrome	0	720×400	16 colors

Graphic mode	Color	Pages	Resolution	No of colors	
04	Color	8	320×200	4	
05	Color	8	320×200	4	
06	Color	8	640×200	2	
0D	Color	8	320×200	16	
0E	Color	4	640×200	16	
0F	Mono chrome	2	640×350	1	
10	Color	2	640×350	16	
11	Color	1	640×480	2	
12	Color	1	640×480	16	
13	Color	1	320×200	256	

## <u>Attribute</u>

Background Foreground

 Attribute:
 BLRGB
 IRGB

 Bit number:
 7 6 5 4
 3 2 1 0

I – Intensity, BL - Blink

Color	Hex
	Value
Black	0
Blue	1
Green	2
Cyan	3
Red	4
Magnet	5
Brown	6
White	7
Gray	8
Light Blue	9
Light Green	A
Light cyan	В
Light red	С
Light magenta	D
Yellow	Е
Bright white	F

## TITLE sorting the numbers – descending order

DOSSEG

.MODEL SMALL

.STACK 100H

.CODE

MAIN PROC FAR

MOV AX, @ DATA

MOV DS, AX

MOV DX, 4H

DOPASS: MOV CX, 4H

MOV SI, 00H

CHECK: MOV AL, ARR [SI]

CMP ARP [SI+1], AL

**JC NOSWAP** 

MOV BL, ARR [SI + 1] MOV ARR[SI +1], AL

MOV ARR [SI], BL

NOSWAP: INC SI

LOOP CHECK

DEC DX

JNZ DOPASS

MOV AX, 4C00H

INT 21H

MAIN ENDP

.DATA

ARR DB 8,2,9,4,7

END MAIN

## Note: Display if numbers are with 1 digit

MOV CX, 05H

MOV SI, 00H

L: MOV DL, ARR[SI]

ADD DL, 30H

MOV AH, 02H

INT 21H

MOV DL,''

MOV AH, 02H

INT 21H

INC SI

LOOP L

MOV AX, 4C00H

INT 21H

MAIN ENDP

```
TITLE addition of 100 natural even numbers
   .MODEL SMALL
   .STACK 100H
   .DATA
          TEN DW 10
   .CODE
MAIN
          PROC FAR
          MOV AX, @ DATA
          MOV DS, AX
          MOV CX, 63H
          MOV AX, 02H
          MOV DX, 04H
   L1:
         ADD AX, DX
         ADD DX, 02H
          LOOP L1
   L2:
         MOV DX, 0000H
                       ; DX: AX /10
          DIV TEN
          INC CX
          ADD DX, 30H ; remainder
          PUSH DX
          CMP AX, 00H; quotient
          JE L3
          JMP L2
   L3:
          POP DX
          MOV AH, 02H
          INT 21H
          LOOP L3
          MOV AX, 4C00H
          INT 21H
   MAIN ENDP
   END MAIN
      TITLE to display string at (10,40) with green background and red foreground
          dosseg
          .Model small
          .Stack 100H
          .Code
      MAIN PROC FAR
          MOV AX, @ DATA
          MOV DX, AX
          MOV SI, OFFSET VAR1
   L2:
          MOV AH, 02H
                             ; Set cursor position
          MOV DH, ROW
          MOV DL, COL
```

```
INT 10H
      MOV AL, [SI]
      CMP AL, '$'
      JE L1
      MOV AH, 09H
      MOV DH, ROW
      MOV DL, COL
                         ;background & foreground
      MOV BL, 24H
      MOV BH, 00h
                         ; page
      MOV CX, 01H
                         ; no. of repeated characters
      INT 10H
      INC SI
      INC COL
      JMP L2
L1:
      MOV AX, 4C00H
      INT 21H
MAIN ENDP
      .DATA
      ROW DB 10
      COL DB 40
      VAR1 DB "video model", '$'
END MAIN
  TITLE TO GENERATE MULTIPLICATION TABLE
      .MODEL SMALL
      .STACK 32
      .DATA
             NUM1 DB 5
            NUM2 DB 1
            TAB DB 10 DUP (?)
      .CODE
   MAIN PROC FAR
      MOV AX, @ DATA
      MOV DS, AX
      MOV BX, 0
      MOV CX, 10
 L1: MOV AL, NUM1
      MUL NUM2
      MOV TAB [BX], AL
      INC BX
      INC NUM2
      LOOP L1
      MOV AX, 4C00H
```

INT 21H MAIN ENDP **END MAIN** 

## TITLE to add 10 sixteen bit Numbers in memory table

.MODEL SMALL .STACK 32

.DATA

NUM DW DUP (2) NUM DW DUP (3)

SUMH DW 0

SUML DW 0

.CODE

MAIN PROC FAR MOV AX, @ DATA

MOV DS, AX MOV CX, 10 MOV AX, 0 MOV BX, 0

ADD AX, NUM [BX] L1:

MOV SUML, AX

JNC L2

**INC SUMH** 

L2: ADD BX, 2 LOOP L1

MOV AX, 4C00H

INT 21H

MAIN ENDP **END MAIN** 

## SUBROUTINE TO CLEAR THE SCREEN

SCR CLEAR PROC NEAR

MOV AX, 0600H ; Request scroll

MOV BH, 61H; blue on brown for attribute on pixel(generally (07H) white on black

MOV CX, 0000 ; Full screen

MOV DX, 184FH

INT 10H

**RET** 

SCR CLEAR ENDP

→ AH-06h: Scroll upward of lines in a specified area of the screen.

AL- 00H caused entire screen to scroll up, effectively clearing it. Setting a nonzero

value in AL causes the number of lines to scroll up.

Note: To access the data of the memory i.e. table.

We use e.g. NUM[BX]

Increasing the BX register by 2