

INSTRUCTION DIVISION SECOND SEMESTER 2015-2016 Course Handout - Part II

In addition to part-I (General Handout for all courses appended to the time table) this portion gives further specific details regarding the course

Course No.: CS F364

Course Title: Design & Analysis of Algorithms

Instructor-in-Charge: Shan SundarBalasubramaniam (email: sundarb)

Instructor(s): JagatSeshChella (email: jagatsesh)

Course Website: access viahttp://nalanda.bits-pilani.ac.in

1. a. Course Objective:

The objective of the course is to impart students

- a working knowledge of how to choose and apply one of several Algorithm Design techniques to solve a given (computational)problem and
- the ability to study, analyze, and understand a given problemso as to be able to formulate it computationally, relate it to other known problems, and characterize the difficulty in solving it.

1.b. Course Scope

This course will cover the broad subject of Algorithms with emphasis on Design Techniques under the assumption that the student has already learnt the basics of data structures and algorithms. The scope of the course would broadly include the following:

- Problems (Types of problems, Formulation of problems, and Reductions between problems; Specific
 Problem Domains such as Graphs, Text Processing, and Number Theory)
- Algorithm Design (Basic techniques Divide & Conquer, Dynamic Programming, Greedy, and Randomization; Techniques for handling hard problems – Backtracking, Branch & Bound, and Approximation Algorithms).
- **Complexity Theory** (Lower Bounding of complexity of problems, Complexity Classes including Approximation Classes, and NP-completeness)

2. Text and References:

a.Text Book:

T1. Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest and Clifford Stein. *Introduction to Algorithms*. MIT Press. 3rd Edition.

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b.References:



R1. Ellis Horowitz, SartajSahni and SanguthevarRajasekaran. *Computer Algorithms.* Universities Press. 2007.

AR. Additional references to be posted on the course website where applicable.

3. Course Plan:

3. a. Pre-requisite Topics:

During the course it will be assumed that the student has systematically studied and understood the following topics:

S.No.	Topics	Nature and level of understanding required and assumed
1	Discrete Structures	[Basic] Properties and theoretical results on Trees and Graphs
2	Recurrence Relations	[Rigorous] Formulation and Solution Techniques
3	Data Structures	[In-depth and breadth -typical end-to-end coverage of a Data
		Structures & Algorithms course]: linearly ordered and
		partially ordered data and data structures, graphs –
		representation and basic algorithms; Sorting algorithms;
4	Divide-and-Conquer	[Basic]
5	Order Complexity	[Basic] Notation (big-O and big-Omega and theta)
		[Rigorous] Growth rate, and Asymptotic complexity;
6	Logic	[Rigorous]Boolean algebra / Propositional Logic: Satisfiability
		Problem
7	Theory of Computation	[Basic] Turing Machines and Computability; Context Free
		Grammars

3.b. Lecture Schedule:

Lec. #	Topic	Learning Outcome(s)	Reading
		[The student should be able to:]	
L1.0	Course Introduction	-	
L1.1	Machine Model: RAM, Uniform Cost	• statesalient features of the RAM	-
	Model vs. Logarithmic Cost Model;	model and its implications for	
		designing / analyzing algorithms	
L2.1	Design:Review of Divide-and-Conquer:	apply divide-and-conquer to design	T1
	Basic Pattern and Example	algorithms for simple problems	Sec.4.1
L2.2	Design:Divide-and-Conquer: (Structural)	• use structural induction on different	-
	Induction:Examples and Typical Cases	forms of data (or data structures) for	
		designing algorithms	
L2.3	Analysis: Locality of Reference	• explain how locality (or the lack of it)	-
		may impact the performance of an	
		algorithm	



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श्रीनं प्रमं बर्ल			
L3.1	Design: Divide-and-Conquer – Example: Matrix Multiplication – Classic algorithm and its complexity	 analyze the cost of multiplying two matrices using the classic algorithm or its variants 	T1 Sec. 4.2
L3.2	Design: Divide-and-Conquer – Example: Matrix Multiplication – Blocking and its implication, Strassen's algorithm and its complexity	 design divide-and-conquer algorithms on matrices using blocking explain the efficiency achieved by Strassen's algorithm 	T1 Sec. 4.2
L4.1	Design:Issues in Top-Down Design: Overlapping Sub-problems	 identify overlapping sub-problems, if they exist, in any top down design calculate the impact of overlapping sub-problems on the cost of the solution 	T1 Sec.15.1
L4.2	Design:Top-Down Design: Memoization – Example.	 memoize the results in a given top-down-design solution and analyze the impact of memoization on the cost of the solution 	T1 Sec.15.1 and Sec. 15.3
L5.1	Design:Top-down Design vs. Bottom-up Design	 state the pros and cons of top-down design and bottom-up-design choose between top-down-design and bottom-up-design as a design approach to a given problem 	-
L5.2 L6.1	Design: Dynamic Programming – Function Problems – Examples Design: Top-Down Design: Optimal Sub-	 apply dynamic programming to solve function problems discover and argue that an 	- T1
	structure Property – Example.	 optimization problem exhibits optimal substructure formulate a recurrence for an optimization problem 	Sec.15.2 - 15.3
L6.2	Design: Dynamic Programming – Optimization Problems – Motivating Example 1	 formulate a recurrence relation based on optimal sub-structure; write a dynamic programming (DP) algorithm based on the recurrence; optimize the space requirement for the DP algorithm 	T1 Sec.15.3- 15.5



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L7.1	Design: Dynamic Programming –		- ditto -	T1
	Optimization Problems – Motivating			Sec.15.3-
	Example 2			15.5
L7.2	Design: Dynamic Programming -	•	formulate and solve typical	AR
	Problem Domain: Text/String		text/string processing problems	
	Processing Problems: Application:		using DP	
	DNA Sequence Alignment			
L7.3	Design: Dynamic Programming -	•	recognize variants of string matching	AR, T1
	Problem Domain: Text/String		/ text matching problems	Sec. 15.4
	Processing Problems: Approximate	•	adapt DP solution for a given variant	
	String Matching and Text Matching –			
	Application(s)			
L8.1	Design: Dynamic Programming –	•	formulate and solve typical graph	T1 Sec.
	Problem Domain: Graph Problems: All		problems using DP	25.2
	Pairs Shortest Paths			
L8.2	Design: Dynamic Programming –	•	recognize variants of a graph	-
	Problem Domain: Graph Problems:		problem	
	Transitive Closure	•	adapt DP solution for a given variant	
L8.3	Design: Dynamic Programming –		- ditto -	-
	Problem Domain: Graph Problems:			
	Application: Translating Finite			
	Automata to Regular Expressions			
L9.1	Design: Greedy Algorithms:	•	recognize problems where Greedy	T1 Sec.
	Motivating Examples		technique would be applicable	16.1
		•	apply Greedy Technique to solve	
			problems	
L9.2	Design: Greedy Choice Property and	•	verify whether Greedy Choice applies	T1 Sec.
	Optimal Substructure Property;		to a given problem	16.2-
	Limitation of Greedy Technique.	•	distinguish between Optimal	16.3
			Substructure and Greedy Choice for	
			a given problem	
		•	recognize problems where Greedy	
			technique would not be applicable	
L10.1	Design:Greedy Technique: Graph	•	apply Greedy Technique to solve	T1 Ch 23
	Problems - Example: Algorithms for		Graph problems	
	Minimum Spanning Tree			
L10.2	Design:Greedy Technique: Counter-	•	recognizeproblems where Greedy	-
	Example: Steiner Tree problem		technique would not be applicable	

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L11.1	Design: Theory of Greedy	•	formulate a given problem using	T1 Sec.
	Algorithms: Matroids: Introduction and		matroids	16.4
	Examples.			
L11.2	Design: Theory of Greedy	•	formulate a given problem using	T1 Sec.
	Algorithms:Weighted Matroids –		weighted matroids	16.4
	Definition and Examples.			
L11.3	Design:Theory of Greedy Algorithms:	•	state the relation between	T1 Sec.
	Weighted Matroids and Greedy		weigtedmatroids and greey	16.4
	Algorithms		algorithms	
L11.4	Design:Theory of Greedy Algorithms:	•	formulate a given problem using	T1 Sec.
	Template Greedy Algorithm		matroids and instantiate a greedy	16.4 –
			algorithm for the specific problem	16.5
L12.1	Problems: Online Problems and	•	distinguish between online and	AR
	Algorithms: Example Problem and		offline problems	
	Typical Algorithms	•	distinguish between online	
			algorithms and offline algorithms	
L12.2	Analysis: Analysis of Online Algorithms:	•	distinguish between worst-case	AR
	Worst-case analysis vs. Competitiveness		analysis of algorithms and	
	Analysis		competitiveness analysis of	
	,		algorithms	
		•	perform competitiveness analysis of	
			algorithms for online problems	
L13	Analysis: Amortized Analysis:	•	perform amortized analysis of	AR
	Introduction and Examples		algorithms for online problems	,
1141	Design: Randomization: Motivation for	•	identify scenarios where	T1 Sec.
	Randomized Algorithms and Example		randomization would be applicable	7.3
	(QuickSort)		(in designing an algorithm)	7.3
1142	Analysis: Analysis of Randomized	•	differentiate analysis of randomized	T1 Sec.
L14.2	Algorithms – Example.		algorithms from that of deterministic	
	Augoritimis Example.		algorithms	7.5
11/12	Design: Randomized Algorithms:	•	distinguish between Las-Vegas and	
L14.5	Models : Las Vegas vs. Monte Carlo	•	Monte-Carlo algorithms	-
115 1	<u> </u>		Worte-Carlo algorithms	A D
L15.1	Design: Game Tree Evaluation: Context,	•	-	AR
	Definition, and a Deterministic			
1153	Algorithm.		design a Lan Margara de Collega Co	AD
L15.2	Design: Randomized Algorithms: Las	•	design a Las Vegas algorithm for a	AR
	Vegas Algorithms: Example (Game Tree		given problem	
	Evaluation)	•	analyze a Las Vegas algorithm	

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L16.1	Problem Domain: Number Theory:	• st	ate the correctness and complexity	T1 Sec.
	Computing GCD and Euclid's Algorithm	of	f Euclid's algorithm	31.2
L16.2	Problem Domain: Number Theory:	• st	cate the correctness and complexity	T1 Sec.
	Aryabhatia's algorithm for computing	of	f Aryabhatia's algorithm	31.2
	GCD and Linear coefficients			
L16.3	Problem Domain: Number Theory:	• fc	ormulate modular arithmetic	T1 Sec.
	Groups and Modular Arithmetic;	рі	roblems as groups	31.3
	Groups Z_n and Z_n^* , Size of Z_n^* and		esign an algorithm to compute the	
	computing the size of Z _n *		ze of Z _n *	
			elate the complexity of factoring an	
			nteger and that of computing the	
			ze of Z _n *	
L17.1	Problem Domain: Number Theory:		elate properties of Z _n * to Euler's	T1 Sec.
	Properties of Z _n *; Euler's Theorem and		heorem.	31.3
	Fermat's Theorem		elate Euler's Theorem to Fermat's	
			heorem	
			tate the issue in using Fermat's	
			heorem in distinguishing between	
			rime numbers and composite	
117.2	Problem Demains Counted graphs		umbers	T1 Sec.
LI/.Z	Problem Domain: Cryptography: Secrecy and Encryption Systems: Public		rate the secrecy or confidentiality	31.7
	Key Cryptography		equirement as a computational roblem	31.7
	incy dryptography		istinguish between Shared Key	
			ncryption and Public Key Encryption	
			tate the computational	
			equirements and the pros and cons	
			f using Public Key Key Encryption	
			3 3 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	
L17.3	Problem Domain: Cryptography: RSA	• ar	rgue the communication	T1 Sec.
	Encryption Protocol: Design and		orrectness of RSA	31.7
	Correctness			
L18.1	Problem Domain: Cryptography: RSA	• ar	rgue the computational	T1 Sec.
	Encryption Protocol: Efficiency and	re	equirements for sender and receiver	31.7
	Security	in	using RSA	
		• pı	rovide a rigorous – even if informal	
		_	argument of the security of RSA	
		-1		

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L18.2	Problem Domain: Number Theory:	define the problem of testing for	
	Primality Testing: Motivation, Cost of typical deterministic algorithms	 primes derive the complexity of typical deterministic algorithms for testing primality 	
L18.3	Problem Domain: Number Theory: Primality Testing: Cost of Finding a prime number	 compute the expected cost of finding a prime number assuming the cost of testing for primes 	
L18.4	Problem Domain: Number Theory: Primality Testing: A pseudo-primality testing algorithm	 explain the design and limitations of a primality testing algorithm based on Fermat's Theorem 	T1 Sec. 31.8
L19	Problem Domain: Number Theory: Primality Testing: Miller-Rabin Test: Algorithm, Time Complexity, and Error Bounds.	 explain the Miller-Rabin Test argue the time complexity and error bounds of the Miller-Rabin Test 	T1 Sec. 31.8
L20.1	Complexity: Lower Bound Analysis of Problems – Examples (Searching and Sorting).	 provide intuitive explanations of lower bounds for searching and sorting problems 	R1 Sec. 10.1
L20.2	Complexity: Lower Bound Analysis of Problems: Decision Tree method	 apply the Decision Tree method for analyzing problems for lower bounds 	R1 Sec. 10.1
L20.3	Complexity: Reduction: Models: Karp Reduction and 1-1 Reduction — Examples	 explain the notion of reduction and its relation to complexity of problems apply reductions for specific problems distinguish between the use of Karp reduction and 1-1 reduction 	-
L21.1	Complexity: Lower Bound Analysis of Problems: Reduction method	 apply reductions for analyzing problems for lower bounds 	R1 Sec. 10.3
L21.2	Complexity: Reduction: Models: Turing Reduction and the relation between different reductions	 explain Turing reduction and its use apply Turing reduction on problems relate other reductions to Turing reduction 	-

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	Complexity: Computability Review: Turing Machine Model, Church-Turing Hypothesis, Equivalence of TM and RAM models	 explain Church-Turing hypothesis and its significance argue equivalence between Turing Machines and Random Access Machines 	-
L22.2	Complexity: Computability Review: Non-computable Functions – Examples and Cardinality	 provide examples of non-computable functions argue that a given function is not computable by using reduction argue that the set of non-computable functions form an uncountable set 	
L23.1	Complexity: Complexity Classes: Time Complexity Classes and Space Complexity Classes	 prove membership of a given problem in a specific complexity class argue the relation between time complexity classes and space complexity classes 	T1 Sec. 34.1
L23.2	Complexity: Tractability: Polynomial Time vs. Exponential Time, Classes P and EXP.	 provide examples of problems in P and EXP prove membership of problems in P 	T1 Sec. 34.1
L24.1	Complexity: Non-deterministic Computation: Ideas and Examples	write non-deterministic algorithms for given problems	T1 Sec. 34.2
L24.2	Complexity: Non-deterministic Computation: Certificate Verification model	 explain non-deterministic algorithms using the Certificate Verification model 	T1 Sec. 34.2
L24.3	Complexity: Non-deterministic Computation: Non-deterministic Time Complexity Classes, Class NP, Example Problems in NP.	 relate deterministic time complexity classes to non-deterministic time complexity classes prove membership of problems in NP 	T1 Sec. 34.2
L25.1	Complexity: P vs. NP	 explain why P is a subset of NP explain the implication of P being a proper subset of NP explain the implication of NP being a subset of P 	-

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L25.2	Complexity: NP-Completeness	 explain the notion of NP- completeness explain the implication of a problem being NP-complete 	T1 Sec. 34.3
L25.3	Complexity: NP-Completeness of Circuit-SAT	explain intuitively why Circuit-SAT is NP-complete	T1 Sec. 34.3
L26.1	Complexity: Transitivity of Reductions; Proving NP-completeness via reduction	 explain why reductions are transitive and how transitivity helps in proving NP-completeness of a problem 	T1 Sec. 34.3
L26.2	Complexity: NP-completeness via reduction: Examples (SAT and its variants)	 prove that SAT and some of its variants and special cases are NP- complete 	T1 Sec. 34.3- 34.5
L27	Complexity: NP-completeness via reduction: Techniques for Reduction - examples	apply different techniques to reduce a NP-hard problem to another	T1 Sec. 34.4
L28.1	Complexity: The class NPC	state the nature and properties of the class NPC	T1 Sec. 34.4- 34.5
L28.2	Complexity: NP-completeness: Structure of Problems – What makes problems hard?	 explain intuitively why apparently simple variants of problems in P are NP-complete 	-
L29.1	Complexity: Complexity Classes: NP and coNP: Complements of Problems, Closures of Complexity Classes	explain the relation and difference between classes NP and coNP	T1 Sec. 34.2, AR
L29.2	Complexity: Complexity Classes: coNP, Problems in coNP, Intersection of NP and coNP	 prove that a given problem is in coNP explain the implication of membership in the intersection of NP and coNP explain the relation between P vs. NP and NP vs. co-NP 	T1 Sec. 34.2, AR
L30.1	Design: Dealing with Hard problems	characterize outcomes / trade-offs of a given approach to deal with a hard problem	

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		 decide and/orchoose an approach for dealing with a hard problem 	
L30.2	Design: Dealing with Hard problems: Exact Solutions: Backtracking: Motivating Example and Solution	 explain how search works in general and the role of backtracking in search 	-
L31.1	Design: Dealing with Hard problems: Exact Solutions: Backtracking: Template solution and example solution	 explain how a generic backtracking template works design algorithms for specific problems using the backtracking template 	R1 Ch. 7
L31.2	Design: Dealing with Hard problems: Exact Solutions: Backtracking: Application Domain	solve problems using search with backtracking as a generic strategy	R1 Ch. 7
L32	Design: Dealing with Hard problems: Exact Solutions: Branch-and-Bound	 relate Branch-and-Bound to Backtracking design algorithms for specific problems using branch-and-bound 	R1 Ch. 8
L33	Design: Dealing with Hard problems: Exact Solutions: Branch-and-Bound: Use of Heuristics	 explain and illustrate the role of heuristics in solving specific problems using branch-and-bound design heuristics to improve the performance of branch-and-bound algorithms in solving specific problems 	R1 Ch. 8
L34.1	Problems: Optimization Problems: Examples	formulate optimization problems crisply	-
L34.2	Complexity: Complexity Classes: PO and NPO, Example problems in PO and NPO.	decide membership of specific problems in PO and NPO	AR

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L34.3	Complexity: Complexity Classes: NP-hard Optimization Problems – Examples.	•	prove NP-hardness of optimization problems using reductions	AR
L35.1	Design: Dealing with Hard problems: Approximation: Introduction to Approximation Algorithms – Vertex Cover as an Example.	•	explain the notion of approximating the solution and that of approximation algorithms	T1 Sec. 35.1
L35.2	Design: Dealing with Hard problems: Approximation: Lower-Bounding as a technique for Designing Approximation Algorithms - Example	•	explain the significance of lower- bounding in designing approximation algorithms	AR T1 Sec. 35.1
L35.3	Design: Dealing with Hard problems: Approximation Algorithms: Absolute Error and Absolute Approximation; Example and Counter-example.	•	assess – informally – whether absolute approximation is feasible for a given problem design simple absolute error approximation algorithms for specific problems explain how absolute approximation is an inadequate approach prove that certain problems are not approximable with an absolute error	AR
L36.1	Design: Dealing with Hard problems: Approximation Algorithms: Relative Error and Approximation within Constant Factor - Example	•	explain how to bound the error of an approximation algorithm by a constant factor	AR
L36.2	Design: Dealing with Hard problems: Approximation Algorithms – Design Techniques: Greedy Technique	•	design approximation algorithms using the greedy technique	AR T1 Sec. 35.3
L36.3	Design: Dealing with Hard problems: Approximation Algorithms – Design Techniques: Sequencing	•	design approximation algorithms using sequencing	AR
L37.1	Design: Dealing with Hard problems: Approximation Algorithms – Design Techniques: Randomization	•	design approximation algorithms using randomization	T1 Sec. 35.4

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L37.2	Design: Dealing with Hard problems Approximation Algorithms – - Problem Domain: Graph Problems: Steiner Tree Problem	 design approximation algorithms for specific graph problems 	AR
L38.1	Design: Dealing with Hard problems Approximation Algorithms – - Problem Domain: Graph Problems: Metric TSP	design approximation algorithms for specific graph problems	T1 Sec. 35.2
L38.2	Complexity: Complexity Classes: Class APX – Examples	prove that specific problems are in APX	AR
L38.3	Complexity: Non-approximability: TSP is not r-approximable.	prove that specific problems are not r-approximable	T1 Sec. 35.2.2
L39	Design: Dealing with Hard problems: Approximation Algorithms – Polynomial Time Approximation Schemes	design PTAS for NP-hard optimization problems	AR
L40	Summary and Conclusion	 state what you have learnt in this course and where you can apply what you have learnt state what you can learn further on your own based on what you have learnt 	-

4. Evaluation

4. a. Evaluation Model

Component	Nature of Assessment	Intent	Form / Format
Quizzes (3)	Design / Analysis of	Review of immediate	In class,
	algorithm(s) for a specific problem	material; Focus on one or	Written,
	 Application of basic 	a few topics	Timed, Open
	concepts		Book.



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नि प्रमं बर्ल			
Assignment - Phase I Assignment - Phase II	Reading, Understanding Writing, and Explaining Research Adapting and Applying Research	Self Learning Self Learning	Take Home; Teams of 2 or 3; Individual contribution will be assessed as well ditto -
Mid-Term Test	 Design / Analysis of algorithm(s) for a specific problem Application of techniques in formulating problems, designing and analyzing problems. Analyzing problems, Performing reductions. 	Sound and extensive understanding of topics taught in class; Ability to apply and relate what is taught in class to problems not known apriori.	Scheduled centrally; Written, Timed, Open Book.
Comprehen sive Exam	 Design / Analysis of algorithm(s) for a specific problem Application of techniques in formulating problems, designing and analyzing problems. Analyzing problems, Performing reductions. Analyzing and proving complexity of problems and relations between complexity classes Approaching hard problems 	Sound and comprehensive understanding of topics taught in class; Ability to apply and relate what is taught in class to problems in the external world. Ability to apply algorithmic thinkingfor understanding and solving problems.	Scheduled centrally; Written, Timed, Open Book.

4. b. Evaluation Scheme:

The Evaluation selection				
Component	Weight	Date& Time	Duration	Remarks
Quizzes (3)	3x15M=	In Jan., Feb. and	30 to 45 min.	Actual date will be announced one
	45M	Apr.	each	lecture prior to the quiz (i.e. one or
				two days notice).
				Read make-up policy below.



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न पर्म वर्ग				
Assignment	20M	In Feb.	10 to 15 days	Actual date will be announced one day
– Phase I				before the start.
				Read late submission policy below.
Assignment	35M	In Mar. – Apr.	15 to 25 days	Actual date will be announced one day
– Phase II				before the start.
				Read late submission policy below.
Mid-Term	33M	16/3 9:00 -	90 minutes	Read Make-up Policy below.
Test		10:30 AM		
Comprehens	67M	7/5 FN	180 minutes	Read Make-up Policy below.
ive Exam				
TOTAL	200M			

4. c. Make-up Policy:

- **Quizzes:**
 - 1. A student may avail at most one make-up for all three quizzes put together.
 - 2. The make-up quiz will be scheduled after all the regular quizzes.
 - 3. Coverage for the make-up quiz will be that of all three quizzes put together.

Quizzes and Mid-Term:

- 1. Make-up for quizzes or for the mid-term test will be granted only for genuine reasons when the student is physically unable to appear for the quiz/test.
- 2. It is the responsibility of the student to communicate a make-up request to one of the instructor(s) before or during the test/quiz.
- 3. Decision of the instructor-in-charge with respect to 1. and 2. above is final.

Assignment:

- 1. Late submission of assignment will incur a penalty of 25% up to 24 hours and a penalty of 50% up to 48 hours past the deadline for that phase.
- 2. No further submissions will be entertained 48 hours past the deadline.

Comprehensive Exam:

- 1. Permission for a Make-up for the comprehensive exam will have to be obtained from Dean Instruction and
- 2. Make-up for the comprehensive exam will usually be scheduled centrally.

4.d. Fairness Policy:

- Fairness policies of BITS Pilani, as stated in Academic Regulations and Guidelines to Students, are all applicable for this course.
- In addition, the following are applicable in the context of take-home assignment for this course:
 - Teams are expected to work on their own on assignments.



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- Students are allowed to consult/discuss with other students/teams for the take-home assignment but such consultation/discussion should be explicitly acknowledged and reported to the instructor prior to evaluation.
- All students are expected to contribute equally within a team:
 - Individual contributions should be identified and documented in qualitative and quantitative terms by the team members.
 - The instructor's assessment regarding the contributions of team members would be final.
- o All material referred to should be explicitly included in Bibliography / References of the reports submitted as part of the assignment.
- o Contents of reports should not be included as is from other sources.
- Unfair means would include:
 - o copying from other students or enabling copying by other students;
 - o copying / borrowing material from the Web or from other sources of information including all electronic sources.
- Any use of unfair means in quizzes, assignment, or test/exam will be handled strictly:
 - o The minimum penalty would be loss of full weight of the component.
 - In addition, students involved in such activity are liable for further sanctions including being formally reported to the Unfair Means committee of the Department as well as to the Examination Committee in which case they will be subject to penalties enabled by Unfair Means Rules of the Institute.
- 5. Consultation Hours: TuTh 2.30pm to 3.30pm (Room 6120-L, CSIS Dept., NAB).
- 6. Notices: All notices concerning this course will be displayed on the course website (on Nalanda) only. If there is a need email would be used on short notice (12 hours) - only BITS Pilani mail id of students would be used.

Instructor -In- Charge **CS F364**



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