Extra Credit 2 Answer

Due: May 9, 2022 **Points:** 20

Remember, you must justify all your answers.

1. (20 points) Prove that Belady's anomaly cannot occur in a stack algorithm.

Answer: Let n be the number of frames and M(n) the set of their contents.

The key here is to characterize Belady's anomaly generally; it is simply that as the number of page frames increases, page faulting increases. This means that the page being brought in was not in memory before it is brought in as a result of the memory size allocation changing to n + 1. If all pages in the memory M(n) were present in the memory of n + 1 frames M(n + 1), this could not happen; the number of page faults would either be constant or decrease.

To see this, consider what happens when a page fault occurs in a stack algorithm. One of 3 cases arise:

- (a) The page fault occurs in the M(n) memory but not in the M(n+1) memory. This means that the page coming into the M(n) memory is in the M(n+1) memory, as before the page fault, every page in M(n) is in M(n+1) by definition of stack algorithm. So there is one more page fault in the M(n) memory than in the M(n+1) memory.
- (b) The page fault occurred in the M(n+1) memory but not in the M(n) memory. This means the process needed a page not in the M(n+1) memory but that was present in the M(n) memory. This means that $M(n) \nsubseteq M(n+1)$ because if the page was in M(n), by definition of stack algorithm, it must be present in M(n+1). This contradicts the assumption that we are dealing with a stack algorithm, and hence cannot happen.
- (c) The page fault occurs in both M(n) and M(n+1). Then one page fault occurs in both.

So, for a stack algorithm, the number of page faults in a memory M(n) is equal to or greater than the number of page faults in M(n+1). But Belady's anomaly is exactly the opposite, that is, the number of page faults in M(n) is greater than the number of page faults in M(n+1),

Hence Belady's anomaly cannot occur in a stack algorithm.