

Computational Intelligence



LUDWIG-
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LMU Munich
winter term 2024/2025

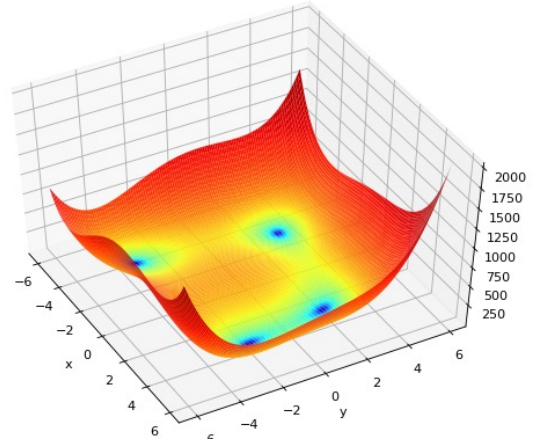
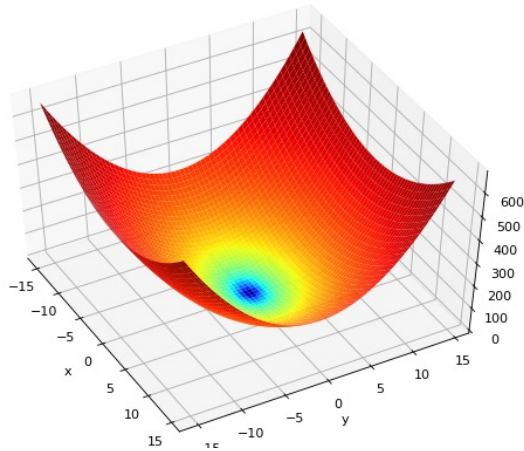
Thomas Gabor
Claudia Linnhoff-Popien

Definition 2 (optimization). Let \mathcal{X} be an arbitrary state space. Let \mathcal{T} be an arbitrary set called target space and let \leq be a total order on \mathcal{T} . A total function $\tau : \mathcal{X} \rightarrow \mathcal{T}$ is called target function. Optimization (minimization/maximization) is the procedure of searching for an $x \in \mathcal{X}$ so that $\tau(x)$ is optimal (minimal/maximal). Unless stated otherwise, we assume minimization. An optimization run of length $g + 1$ is a sequence of states $\langle x_t \rangle_{0 \leq t \leq g}$ with $x_t \in \mathcal{X}$ for all t .

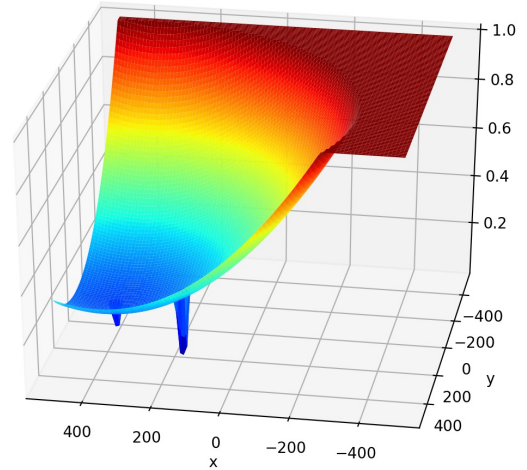
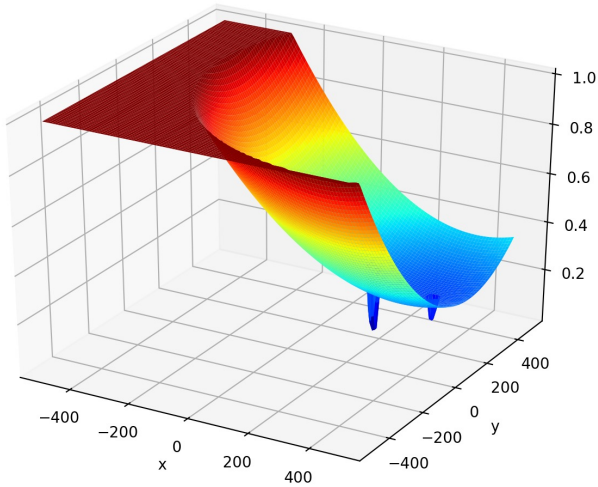
Let $e : \langle \mathcal{X} \rangle \times (\mathcal{X} \rightarrow \mathcal{T}) \rightarrow \mathcal{X}$ be a possibly randomized or non-deterministic function so that the optimization run $\langle x_t \rangle_{0 \leq t \leq g}$ is produced by calling e repeatedly, i.e., $x_{t+1} = e(\langle x_u \rangle_{0 \leq u \leq t}, \tau)$ for all t , $1 \leq t \leq g$, where x_0 is given externally (e.g., $x_0 =_{def} 42$) or chosen randomly (e.g., $x_0 \sim \mathcal{X}$). An optimization process is a tuple $(\mathcal{X}, \mathcal{T}, \tau, e, \langle x_t \rangle_{0 \leq t \leq g})$.

Definition 3 (optimization (policy)). Let $\mathcal{X} = \Pi$ be a policy space. Let $\mathcal{D} = (\Pi, \mathcal{T}, \tau, e, \langle x_t \rangle_{0 \leq t \leq g})$ be an optimization process according to Definition 2. \mathcal{D} is called a policy optimization process.

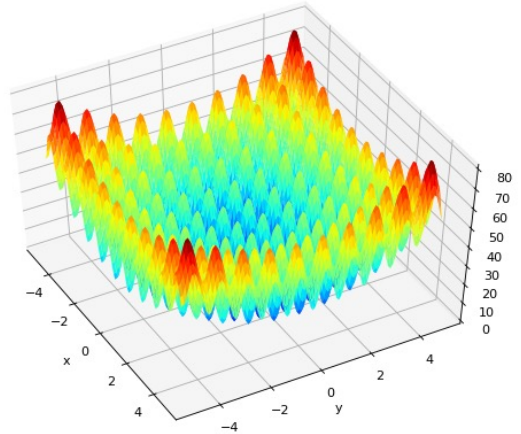
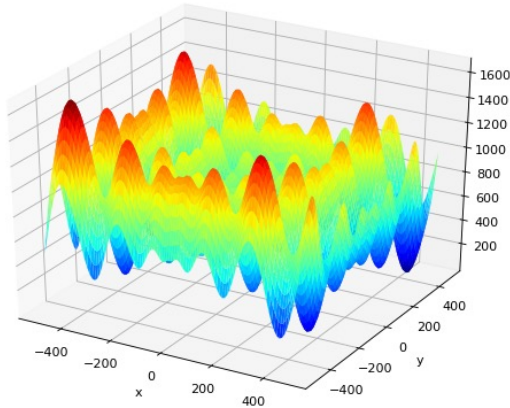
The Solution Landscape Metaphor



The Solution Landscape Metaphor



The Solution Landscape Metaphor



Algorithm 3 (simulated annealing). Let $\mathcal{D} = (\mathcal{X}, \mathcal{T}, \tau, e, \langle x_u \rangle_{0 \leq u \leq t})$ be an optimization process. Let $neighbors : \mathcal{X} \rightarrow \wp(\mathcal{X})$ be a function that returns a set of neighbors of a given state $x \in \mathcal{X}$. Let $\kappa : \mathbb{N} \rightarrow \mathbb{R}$ be a temperature schedule, i.e., a function that returns a temperature value for each time step. Let $A : \mathcal{T} \times \mathcal{T} \times \mathbb{R} \rightarrow \mathbb{P}$ with $\mathbb{P} = [0; 1] \subset \mathbb{R}$ be a function that returns an acceptance probability given two target values and a temperature. Commonly, we use

$$A(Q, Q', K) = e^{\frac{-(Q' - Q)}{K}}$$

for $\mathcal{T} \subseteq \mathbb{R}$. The process \mathcal{D} continues via simulated annealing if e is of the form

$$e(\langle x_u \rangle_{0 \leq u \leq t}, \tau) = x_{t+1} = \begin{cases} x'_t & \text{if } \tau(x'_t) \leq \tau(x_t) \text{ or } r \leq A(\tau(x_t), \tau(x'_t), \kappa(t)), \\ x_t & \text{otherwise,} \end{cases}$$

where $x'_t \sim neighbors(x_t)$ and $r \sim \mathbb{P}$ are drawn at random for each call to e .

other means of
biased sampling

Definition 4 (population-based optimization). Let \mathcal{X} be a state space. Let \mathcal{T} be a target space with total order \leq . Let $\tau : \mathcal{X} \rightarrow \mathcal{T}$ be a target function. A tuple $\mathcal{E} = (\mathcal{X}, \mathcal{T}, \tau, E, \langle X_t \rangle_{0 \leq t \leq g})$ is a population-based optimization process iff $X_t \in \wp^*(\mathcal{X})$ for all t and $E : \langle \wp^*(\mathcal{X}) \rangle \times (\mathcal{X} \rightarrow \mathcal{T}) \rightarrow \wp^*(\mathcal{X})$ is a possibly randomized, non-deterministic, or further parametrized function so that the population-based optimization run is produced by calling E repeatedly, i.e., $X_{t+1} = E(\langle X_u \rangle_{0 \leq u \leq t}, \tau)$ where X_0 is given externally or chosen randomly.

Definition 5 (population-based optimization (alternate)). An optimization process $\mathcal{E} = (\mathcal{X}, \mathcal{T}, \tau, E, \langle X_t \rangle_{0 \leq t \leq g})$ is called population-based iff \mathcal{X} has the form $\mathcal{X} = \wp^*(\mathcal{Y})$ for some other state space \mathcal{Y} .

Multi-Sets

Let X be a set with members from a space \mathcal{X} .

We can identify a set X with its membership function $\mu_X : \mathcal{X} \rightarrow \mathbb{B}$ such that

$$\mu_X(x) = 1 \iff x \in X.$$

For a multiset X' we generalize the membership function so that

$$\mu_{X'}(x) \geq 1 \iff x \in X'.$$

In both cases, the subset relation between (multi-)sets X, Y can be defined as

$$X \subseteq Y \iff \forall x \in X : \mu_X(x) \leq \mu_Y(x).$$

We then define $\wp(X) = 2^X = \mathbb{B}^X \cong X \rightarrow \mathbb{B}$,

and $\wp^*(X) = \mathbb{N}^X \cong X \rightarrow \mathbb{N}$,

to be the powerset functions for sets and multisets, respectively.

Multi-Sets

Evolutionary Algorithms

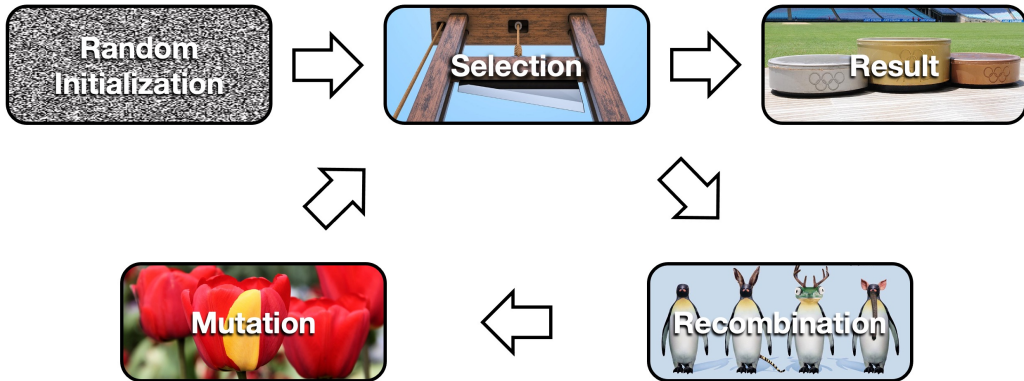


image sources:

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<https://phys.org/news/2019-10-guillotine-cruel-poisoning.html>

Computational Intelligence, winter term 2024/2025, LMU Munich

Algorithm 4 (basic evolutionary algorithm). Let $\mathcal{E} = (\mathcal{X}, \mathcal{T}, \tau, E, \langle X_u \rangle_{0 \leq u \leq t})$ be a population-based optimization process. The process \mathcal{E} continues via an evolutionary algorithm if E has the form

$$E(\langle X_u \rangle_{0 \leq u \leq t}, \tau) = X_{t+1} = \textit{selection}(X_t \uplus \textit{variation}(X_t))$$

where *selection* and *variation* are possibly randomized or non-deterministic functions so that for any $X \in \wp^*(\mathcal{X})$ it holds that $|\textit{selection}(X)| \leq |X|$ and $|\textit{selection}(X \uplus \textit{variation}(X))| = |X|$.

Which optimization algorithm
is the best?