

# Natural Computing

LMU Munich  
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**Definition 3** (Conway's game of life (standard)). Let  $G = (V, E)$  be a graph with vertices  $V$  and (undirected) edges  $E \subseteq V \times V$  with a fixed degree of 8 for all nodes. We define surroundings:  $V \rightarrow \wp(V)$  via

$$\text{surroundings}(v) = \{w \mid (v, w) \in E\}$$

so that  $|\text{surroundings}(v)| = 8$  and  $v \notin \text{surroundings}(v)$  for all  $v \in V$ . A state  $x \in \mathcal{X}$  is a mapping of vertices to the labels  $\{\text{dead}, \text{alive}\}$ , i.e., the state space  $\mathcal{X}$  is given via  $\mathcal{X} = (V \rightarrow \{\text{dead}, \text{alive}\})$ . Let  $\underline{x_t}$  be a state that exists at time step  $t \in \mathbb{N}$ . We define

$$|v|_{x_t} = |\{w \mid w \in \text{surroundings}(v) \wedge x_t(w) = \text{alive}\}|.$$

In the game of life, the evolution of a state  $x_t$  to its subsequent state  $x_{t+1}$  is given deterministically via

$$x_{t+1}(v) = \begin{cases} \text{dead} & \text{if } |v|_{x_t} \leq 1, \\ x_t(v) & \text{if } |v|_{x_t} = 2, \\ \text{alive} & \text{if } |v|_{x_t} = 3, \\ \text{dead} & \text{if } |v|_{x_t} \geq 4, \end{cases}$$

for all  $v \in V$ . A tuple  $(G, x_0)$  is called an instance of the game of life for initial state  $x_0 \in \mathcal{X}$ .

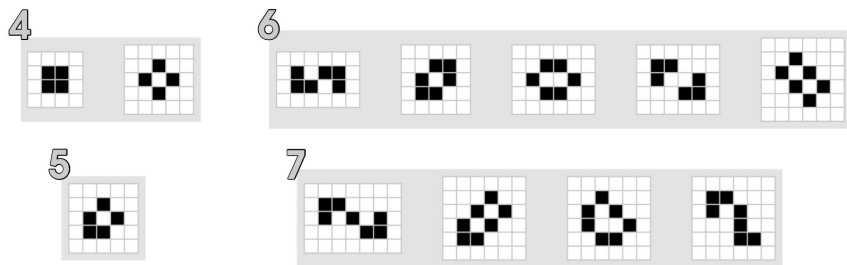


$|v|_{x_t}$

Let's try

[conwaylife.com](http://conwaylife.com)

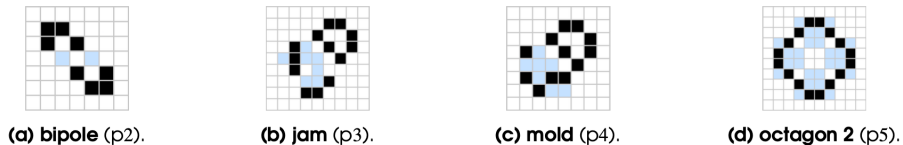
# Still Lifes



**Figure 2.1:** All still lifes with 7 or fewer live cells, arranged by their cell count. From left to right: (4 cells) block, tub, (5 cells) boat, (6 cells) snake, ship, beehive, aircraft carrier, barge, (7 cells) long snake, long boat, loaf, and eater 1 (which is sometimes called **fishhook**).

source: [conwaylife.com/book/](http://conwaylife.com/book/)

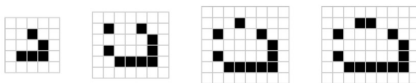
# Oscillators



**Figure 3.1:** Some more small naturally occurring (but rare) oscillators that can be found via computer-assisted soup searches. These oscillators were found by (a) early Lifenthusiasts at M.I.T. in 1970, (b,c) Achim Flammenkamp in 1988,<sup>2</sup> and (d) Sol Goodman and Arthur Taber in 1971.

source: [conwaylife.com/book/](http://conwaylife.com/book/)

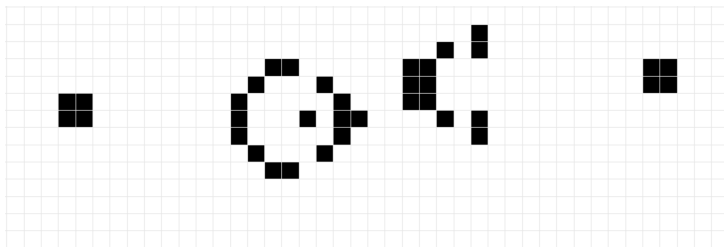
# Space Ships



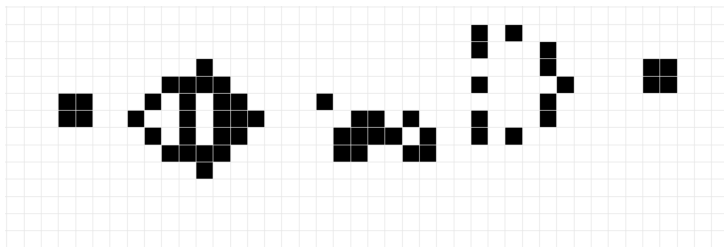
**Figure 4.1:** The four basic spaceships in Conway's Game of Life. From left to right, these are the glider (which moves diagonally at a speed of  $c/4$ ) and the light/middle/heavyweight spaceships (which each move orthogonally<sup>1</sup> at a speed of  $c/2$ ).

source: [conwaylife.com/book/](http://conwaylife.com/book/)

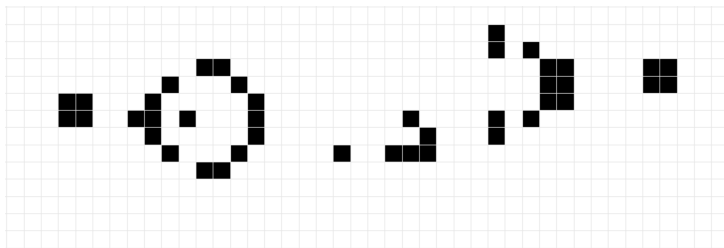
# Guns



t = 1



t = 14

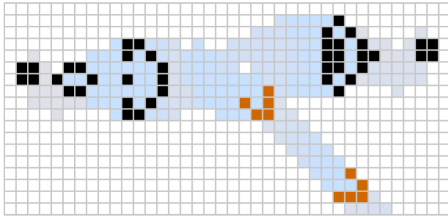


t = 15

[try online](#)

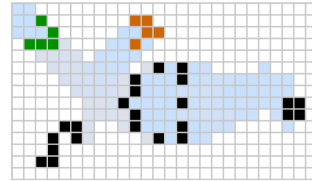
source: conwaylife.com

# Space Ship Maneuvering



**Figure 6.2:** A Gosper glider gun producing gliders at a spacing of 30 generations.

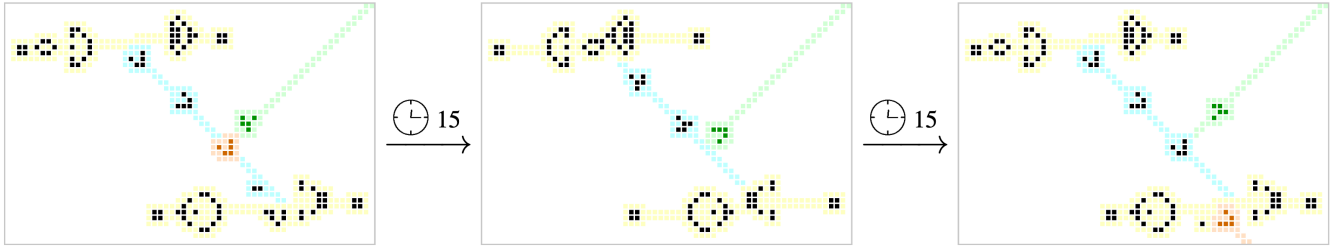
source: [conwaylife.com/book/](http://conwaylife.com/book/)



**Figure 6.3:** A buckaroo can reflect a glider by 90 degrees, from the position marked in green to the one in orange 30 generations later.



# More Space Ship Maneuvering



**Figure 6.7:** Two Gosper glider guns can be placed near each other so as to create a finite stream of gliders (highlighted here in aqua) between them. Here, we bounce a single glider (in green) off of one of those gliders (in orange) so that it is destroyed and thus released by the inline inverter to the southeast.

source: [conwaylife.com/book/](http://conwaylife.com/book/)

[conwaylife.com/book/periodic\\_circuitry](http://conwaylife.com/book/periodic_circuitry)

- No initial pattern with simply proven unlimited growth
- Initial patterns with apparently unlimited growth
- Simple initial patterns with behavior for a long period of time

Gardner, 1970.

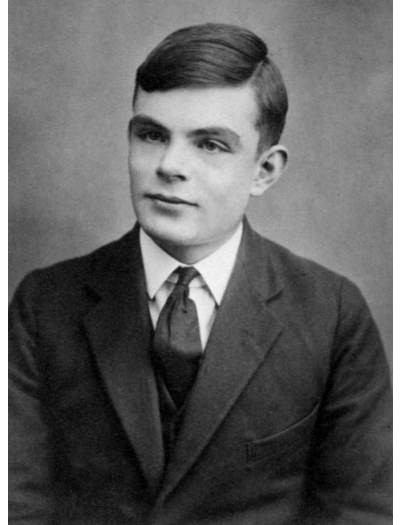
## Why is this interesting?

# Turing-Completeness

# Alan Turing

1912 – 1954

- invented the field of (theoretical) computer science and artificial intelligence
- influential in breaking Enigma and thus shortening WWII



source: Wikipedia

# The Turing Machine

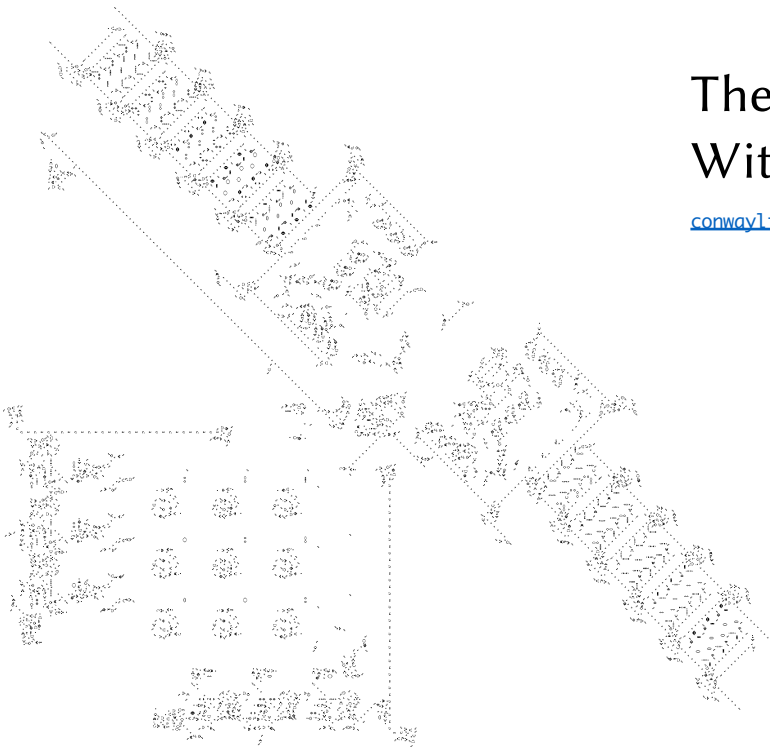
# The Turing Machine

**Definition 4** (Turing machine [1]). A Turing machine is a tuple  $(Q, \Sigma, \Gamma, \sqcup, q_S, q_F, \delta)$  so that

- $Q$  is a finite set of states,
- $\Sigma$  is a finite input/output alphabet (e.g.,  $\Sigma = \{0, 1\}$ ),
- $\Gamma$  is a finite tape alphabet with  $\Sigma \subset \Gamma$ ,
- $\sqcup \in \Gamma$  is a blank symbol with  $\sqcup \in \Sigma$ ,
- $q_S \in Q$  is a start state,
- $q_F \in Q$  is a stop state,
- and  $\delta : Q \setminus \{q_F\} \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, N\}$  is a partially defined transition function.

# The Turing Machine Within the Game of Life

[conwaylife.com/wiki/Turing\\_machine](http://conwaylife.com/wiki/Turing_machine)



But... why?



**Theorem 1** (Church-Turing thesis). Any computable function can be computed by a Turing machine.

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**Theorem 2** (extended Church-Turing thesis). Any computable function can be computed by a Turing machine *with at most polynomial overhead*.

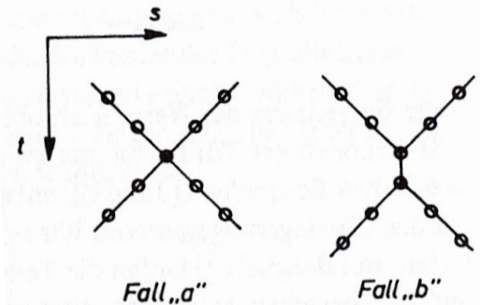
Can cellular automata  
explain the universe?

state space  $\mathcal{X} = (V \rightarrow \mathbb{Z} \times \mathbb{Z})$

$t_0$	$v$								
	$p$					+1	+1		
$t_1$	$v$				-1		+1		
	$p$				+1			+1	
$t_2$	$v$			-1				+1	
	$p$		+1						+1

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	$p$				+1			+1	
$t_2$	$v$				-1				+1
	$p$		+1						+1



# Konrad Zuse

1910 – 1995

- built first programmable computer (Zuse Z3)
- designed first high-level programming language (Plankalkül)



source: Wikipedia

**Example 1** (Cellular Automaton for 1D Particle Collision, inspired by [1]). Consider the following variation of a 1D cellular automaton (cf. Definitions 1 and 2 on the definition sheet): Let the state space  $\mathcal{X} = (V \rightarrow \mathbb{Z} \times \mathbb{Z})$ . Let

$$f((v_l, p_l), (v_c, p_c), (v_r, p_r)) = (v_l^+ + v_r^- + \lfloor \frac{(p_l - p_c)^+}{2} \rfloor - \lfloor \frac{(p_r - p_c)^+}{2} \rfloor, \\ p_c - |v_c| + v_l^+ - v_r^- + \lfloor \frac{(p_l - p_c)^+}{2} \rfloor + \lfloor \frac{(p_r - p_c)^+}{2} \rfloor - \lfloor \frac{(p_c - p_l)^+}{2} \rfloor - \lfloor \frac{(p_c - p_r)^+}{2} \rfloor)$$

where  $x^+ = \begin{cases} x & \text{if } x > 0, \\ 0 & \text{otherwise,} \end{cases}$  and  $x^- = \begin{cases} x & \text{if } x < 0, \\ 0 & \text{otherwise,} \end{cases}$  for any  $x \in \mathbb{Z}$ .

Note:  $\lfloor x \rfloor$  for  $x \in \mathbb{R}$  is  $x$  rounded down to an integer as implemented in a typical `floor` function (see Python, e.g.). From the start configuration of this 1d cellular automaton with 5 cells below, compute the next 3 time steps.

(0, 0)	(+1, +1)	(0, 0)	(−1, +1)	(0, 0)

```

1 from math import floor
2
3 def eplus(x):
4     return x if x > 0 else 0
5
6 def eminus(x):
7     return x if x < 0 else 0
8
9 def f(left,center,right):
10     vl,pl = left
11     vc,pc = center
12     vr,pr = right
13     pressure_from_left = floor(eplus(pl-pc)/2)
14     pressure_from_right = floor(eplus(pr-pc)/2)
15     pressure_towards_left = floor(eplus(pc-pl)/2)
16     pressure_towards_right = floor(eplus(pc-pr)/2)
17     vnew = eplus(vl) + eminus(vr) + pressure_from_left - pressure_from_right
18     pnew = pc - abs(vc) + eplus(vl) - eminus(vr) + pressure_from_left + pressure_from_right - pressure_towards_left - pressure_towards_right
19     return vnew, pnew
20
21 def apply(vector, f):
22     new_vector = []
23     for e,element in enumerate(vector):
24         new_element = f(vector[e-1], element, vector[(e+1)%len(vector)])
25         new_vector.append(new_element)
26     return new_vector
27
28 rounds = range(4)
29 vector = [(0,0), (+1,+1), (0,0), (-1,+1), (0,0)]
30 for _ in rounds:
31     print(vector)
32     vector = apply(vector,f)

```



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$$\text{where } x^+ = \begin{cases} x & \text{if } x > 0, \\ 0 & \text{otherwise,} \end{cases} \text{ and } x^- = \begin{cases} x & \text{if } x < 0, \\ 0 & \text{otherwise,} \end{cases} \text{ for any } x \in \mathbb{Z}.$$

Note:  $\lfloor x \rfloor$  for  $x \in \mathbb{R}$  is  $x$  rounded down to an integer as implemented in a typical `floor` function (see Python, e.g.).  
 From the start configuration of this 1d cellular automaton with 5 cells below, compute the next 3 time steps.

(0, 0)	(+1, +1)	(0, 0)	(−1, +1)	(0, 0)
(0, 0)	(0, 0)	(0, 2)	(0, 0)	(0, 0)
(0, 0)	(−1, +1)	(0, 0)	(+1, +1)	(0, 0)
(−1, +1)	(0, 0)	(0, 0)	(0, 0)	(+1, +1)

What about  
non-local  
information?

# Quantum Computing

# What is quantum computing?

What is quantum computing  
within natural computing?

two strange  
quantum effects