

# CSC 348 – Homework #5

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## 1 Questions

1.  $f(n) = n^2 + 1$  is not 1-1.

*Proof.* Consider  $-1, 1 \in \mathbb{R}$ . Note that  $1 \neq -1$ , but that  $f(-1) = f(1) = 1^2 + 1 = 2$ . Thus,  $f$  is not 1-1 by definition of 1-1. d(^\_^)>

2.  $f(x) = n^5$  is 1-1.

*Proof.* Suppose for some  $x, y \in \mathbb{R}$ ,  $f(x) = f(y)$ .

$$\begin{aligned}x^5 &= y^5 \\ \sqrt[5]{x^5} &= \sqrt[5]{y^5} \\ x &= y\end{aligned}\tag{1}$$

Thus, by definition of 1-1,  $f$  is 1-1. d(^\_^)>

3.  $f(n) = n - 1$  is onto.

*Proof.* Let  $y \in \mathbb{R}$ . Note that  $y = (y + 1) - 1 = f(y + 1)$ . Since  $(y + 1) - 1 \in \mathbb{R}$ ,  $f(x)$  is onto by definition of onto. d(^\_^)>

4.  $f(n) = n^2 + 1$  is not onto.

*Proof.* Consider  $0 \in \mathbb{R}$ . Note that  $x = \sqrt{0 - 1} \notin \mathbb{R}$ . Because of this, no  $x \in \mathbb{R}$  such that  $f(x) = x^2 + 1 = 0$  is possible. Therefore, by definition of onto,  $f$  is not onto. d(^\_^)>

5.  $f(n) = \sqrt[5]{n}$  is onto.

*Proof.* Let  $y \in \mathbb{R}$ . Note that  $y = (\sqrt[5]{y})^5 = f(\sqrt[5]{y})$ . Since  $(\sqrt[5]{y})^5 \in \mathbb{R}$ ,  $f(x)$  is onto by definition of onto. d(∩\_∩)>

6.  $f : \mathbb{N} \rightarrow \mathbb{N}; f(n) = 2n$  is injective, but not surjective.

*Proof of injectivity.* Suppose  $n = 3$  d(∩\_∩)>

*Proof of non-surjectivity.* Suppose d(∩\_∩)>

7.  $f : \mathbb{N} \rightarrow \mathbb{N}; f(n) = \sqrt{n}$  is surjective, but not injective.

*Proof of non-injectivity.* Suppose d(∩\_∩)>

*Proof of surjectivity.* Suppose d(∩\_∩)>

8.  $f : \mathbb{N} \rightarrow \mathbb{N}; f(n) = (-1)^n + n$  is one-to-one and onto.

*Proof of injectivity.* Suppose  $a, b \in \mathbb{N}$  and  $f(a) = f(b)$ .

By the definition of those functions,

$$(-1)^a + a = (-1)^b + b \quad (2)$$

Note that the terms  $(-1)^a$  and  $(-1)^b$  will either be 1 or  $-1$ , which are both odd. Thus,  $(-1)^a$  and  $(-1)^b$  are always odd.

**Case 1.**  $f(a) = f(b)$  is even.

By Lemma 1, since this value is even, its addends have the same sign. Thus,  $a$  and  $b$  must be odd.

By definition of odd, there exists some  $i, j \in \mathbb{Z}$  such that  $a = 2i + 1$  and  $b = 2j + 1$ .

$$\begin{aligned} (-1)^{2i+1} + a &= (-1)^{2j+1} + b \\ ((-1)^2)^i \cdot (-1)^1 + a &= ((-1)^2)^j \cdot (-1)^1 + b \\ -1 + a &= -1 + b \\ a &= b \end{aligned} \quad (3)$$

Thus,  $f(n)$  is 1-1 for even  $n$ .

**Case 2.**  $f(a) = f(b)$  is odd.

By Lemma 1, since this value is even, its addends have the same sign. Thus,  $a$  and  $b$  must be odd.

By definition of odd, there exists some  $i, j \in \mathbb{Z}$  such that  $a = 2i$  and  $b = 2j$ .

$$\begin{aligned} (-1)^{2i} + a &= (-1)^{2j} + b \\ ((-1)^2)^i + a &= ((-1)^2)^j + b \\ 1 + a &= 1 + b \\ a &= b \end{aligned} \tag{4}$$

Thus,  $f(n)$  is 1-1 for even  $n$ .

Since both cases are true,  $f(n)$  is a 1-1 function.

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*Proof of surjectivity.* Suppose  $y \in \mathbb{N}$ . We will prove that there always exists some  $n \in \mathbb{N}$  such that  $f(n) = y$ .

**Case 1.**  $y$  is even.

By definition of even, for some  $i \in \mathbb{N}$ ,  $y = 2i$ .

Consider  $f(y + 1)$ .

$$\begin{aligned} f(y + 1) &= (-1)^{y+1} + (y + 1) \\ &= (-1)^{2i+1} + y + 1 \\ &= ((-1)^2)^i \cdot (-1)^1 + y + 1 \\ &= 1 \cdot (-1) + y + 1 \\ &= y \end{aligned} \tag{5}$$

Thus, if  $y$  is even, then  $f(y + 1) = y$ .

**Case 2.**  $y$  is odd.

By definition of odd, for some  $i \in \mathbb{N}$ ,  $y = 2i + 1$ .

Consider  $f(y - 1)$ .

$$\begin{aligned} f(y - 1) &= (-1)^{y-1} + (y - 1) \\ &= (-1)^{(2i+1)-1} + y - 1 \\ &= (-1)^{2i} + y - 1 \\ &= ((-1)^2)^i + y - 1 \\ &= 1 + y - 1 \\ &= y \end{aligned} \tag{6}$$

Thus, if  $y$  is odd, then  $f(y - 1) = y$ .

Therefore, for all  $y \in \mathbb{N}$ , there exists a value  $n \in \mathbb{N}$  such that  $f(n) = y$ . Thus,  $f$  is onto.

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9. (a)  $(f \circ g)(x) = (e^x)^2 + 1 = e^{2x} + 1$
- (b)  $(g \circ f)(x) = e^{x^2+1}$

10.

**Theorem 1.** Let  $f : B \rightarrow C$  and  $g : A \rightarrow B$ . If  $f$  and  $g$  are onto, then  $f \circ g$  is onto.

*Proof.* Suppose  $c \in C$ . By definition of  $f$  as an onto function, there exists a  $b \in B$  such that  $f(b) = c$ . It follows that there exists an  $a \in A$  such that  $g(a) = b$  by definition of  $g$  as an onto function. Therefore, there always exists some  $a \in A$  such that  $(f \circ g)(a) = c$ .

Thus, by definition of onto, if  $f$  and  $g$  are onto, then  $f \circ g$  is onto. d(^\_^)>

11. (a)

**Theorem 2.** Let  $f : B \rightarrow C$  and  $g : A \rightarrow B$ .  $f \circ g$  being onto does not imply that  $g$  is onto.

*Proof.* Consider the case when  $A = B = C = \mathbb{R}$ ,  $f(x) = \ln x$ , and  $g(x) = x^2$ . It follows that  $(f \circ g)(x) = \ln(x^2)$ .

**Claim 1.**  $f$  is onto.

*Proof.* Let  $y \in \mathbb{R}$ . Note that  $y = \ln(e^y) = f(e^y)$ . Since  $e^y \in \mathbb{R}$ ,  $f(x)$  is onto by definition of onto. d(^\_^)>

**Claim 2.**  $g$  is not onto.

*Proof.* Consider  $-1 \in \mathbb{R}$ . Note that  $x = \sqrt{-1} \notin \mathbb{R}$ . Therefore, there is no such  $x \in \mathbb{R}$  such that  $g(x) = -1$ . Thus,  $g$  is not onto. d(^\_^)>

**Claim 3.**  $f \circ g$  is onto.

*Proof.* Let  $y \in \mathbb{R}$ . Note that  $y = \ln((e^{\frac{y}{2}})^2) = f(e^{\frac{y}{2}})$ . Since  $e^{\frac{y}{2}} \in \mathbb{R}$ ,  $f(x)$  is onto by definition of onto. d(^\_^)>

Notice that  $f \circ g$  is onto, but  $g$  is not onto. Therefore,  $f \circ g$  being onto does not imply that  $g$  is onto. d(^\_^)>

(b)

**Theorem 3.** Let  $f : B \rightarrow C$  and  $g : A \rightarrow B$ .  $f \circ g$  being onto implies that  $f$  is onto.

*Proof.* Suppose  $c \in C$ . By definition of  $f \circ g$  as an onto function, there exists an  $a \in A$  such that  $(f \circ g)(a) = f(g(a)) = c$ . Note that the existence of  $f(g(a))$  requires  $g(a) \in B$ , so  $g(a) \in B$ .

Let  $b = g(a)$ . Thus,  $f(b) = c$ . Therefore, if  $f \circ g$  is onto then  $f$  is onto. d(^\_^)>

## 2 Additional Lemmas with Proofs

**Lemma 1.** Let  $a, b \in \mathbb{Z}$ .  $a + b$  is even if and only if  $a$  and  $b$  have the same parity, and odd if and only if  $a$  and  $b$  have different parities. (Reproduced from homework 4).

*Proof.* **Case 1.**  $a$  and  $b$  are both even.

By definition of even:

$$\begin{aligned} a &= 2k \\ b &= 2l \end{aligned} \tag{7}$$

Therefore,

$$\begin{aligned} a + b &= 2k + 2l \\ &= 2(k + l) \end{aligned} \tag{8}$$

Let  $m = k + l$ . Then

$$a + b = 2m \tag{9}$$

By definition of even,  $a + b$  is even.

**Case 2.**  $a$  and  $b$  are both odd.

By definition of odd:

$$\begin{aligned} a &= 2k + 1 \\ b &= 2l + 1 \end{aligned} \tag{10}$$

Therefore,

$$\begin{aligned} a + b &= 2k + 1 + 2l + 1 \\ &= 2(k + l + 1) \end{aligned} \tag{11}$$

Let  $m = k + l + 1$ . Then,

$$a + b = 2m \tag{12}$$

By definition of even,  $a + b$  is even.

**Case 3.**  $a$  and  $b$  have different parity.

WLOG let  $a$  be even and  $b$  be odd. By definition of even:

$$a = 2k \tag{13}$$

By definition of odd:

$$b = 2l + 1 \tag{14}$$

Therefore,

$$\begin{aligned} a + b &= 2k + 2l + 1 \\ &= 2(k + l) + 1 \end{aligned} \tag{15}$$

Let  $m = k + l$ . Then,

$$a + b = 2m + 1 \tag{16}$$

By definition of odd,  $a + b$  is odd.

**In Summary:**

$a$ parity	$b$ parity	$a + b$ parity
even	even	even
even	odd	odd
odd	even	odd
odd	odd	even

Therefore,  $a + b$  is only even when  $a$ 's parity =  $b$ 's parity, and only odd when  $a$ 's parity  $\neq$   $b$ 's parity, concluding the lemma. d(^\_^)>

### 3 Scratch Work