

Frankie



A Frankenstein's Monster of Architectures

The Accumulators

- Started out with one accumulator and a stack
- Realized that it was hard to keep track of multiple variables
- Fixed by adding second accumulator
- Named Mary and Shelley after author of Frankenstein

Why Shelley?

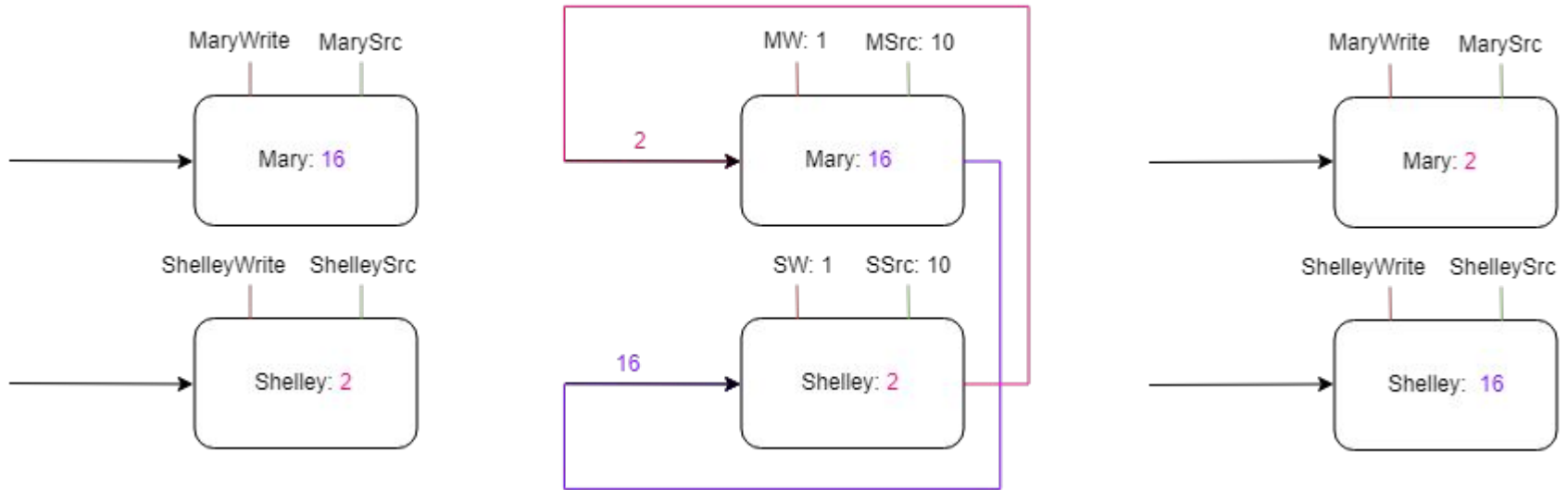
- Try this:

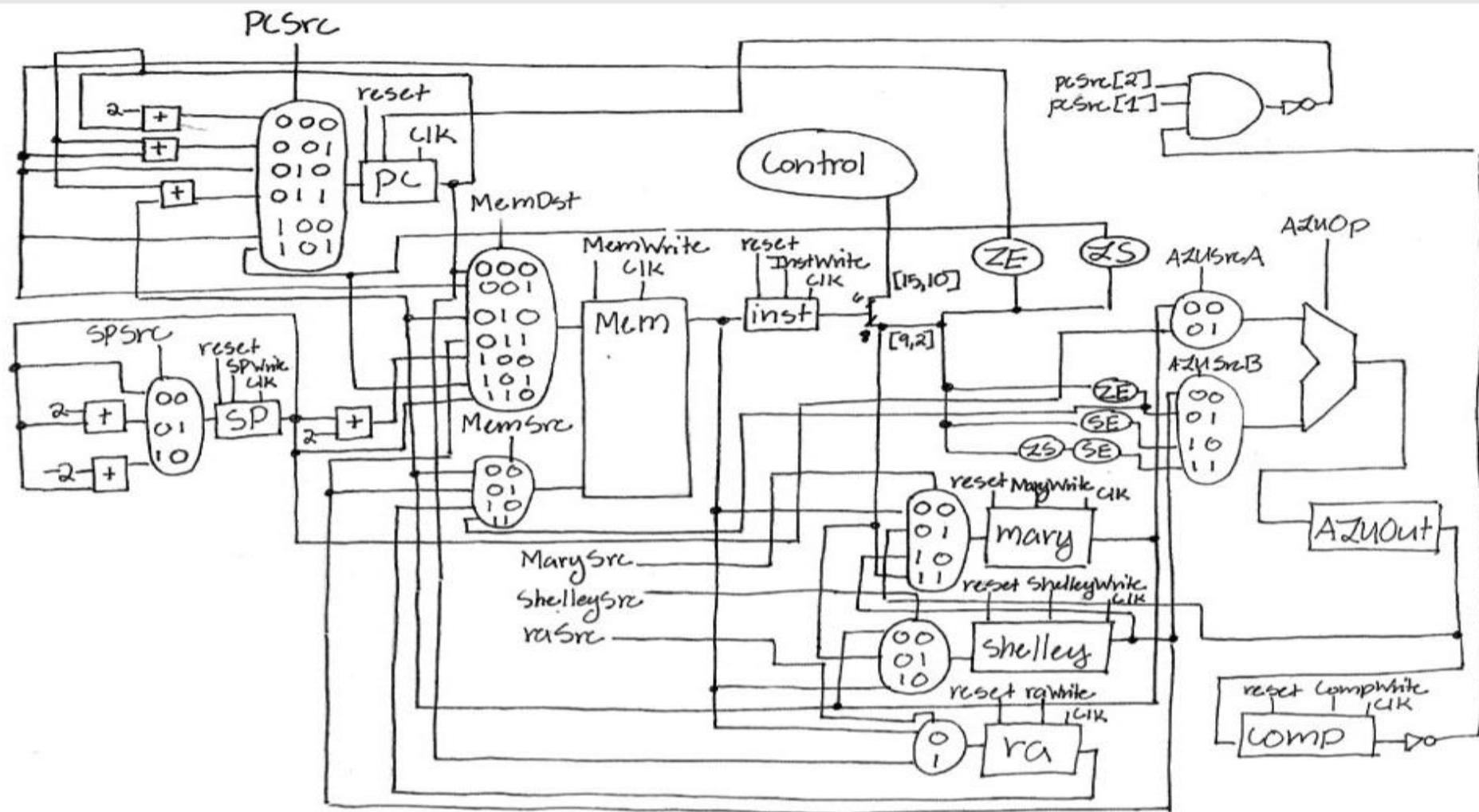
```
for(int i = 0; i < 10; i++)
```

```
    sum += i;
```

- sum goes in Mary
- Where do you keep i? On the stack? Should we add right off the stack, then?
 - This only fixes the problem if the thing you want from the stack doesn't get buried
 - you can't pop without losing Mary's contents

Two Accumulators: (Swap)





Instruction Set

- Instructions are grouped together into common operation groups
 - Arithmetic/Logical
 - Stack
 - Jump
 - Compare
 - Memory-based
 - Backup



flag	op code	immediate	unused
1	5	8	2

Flagbit

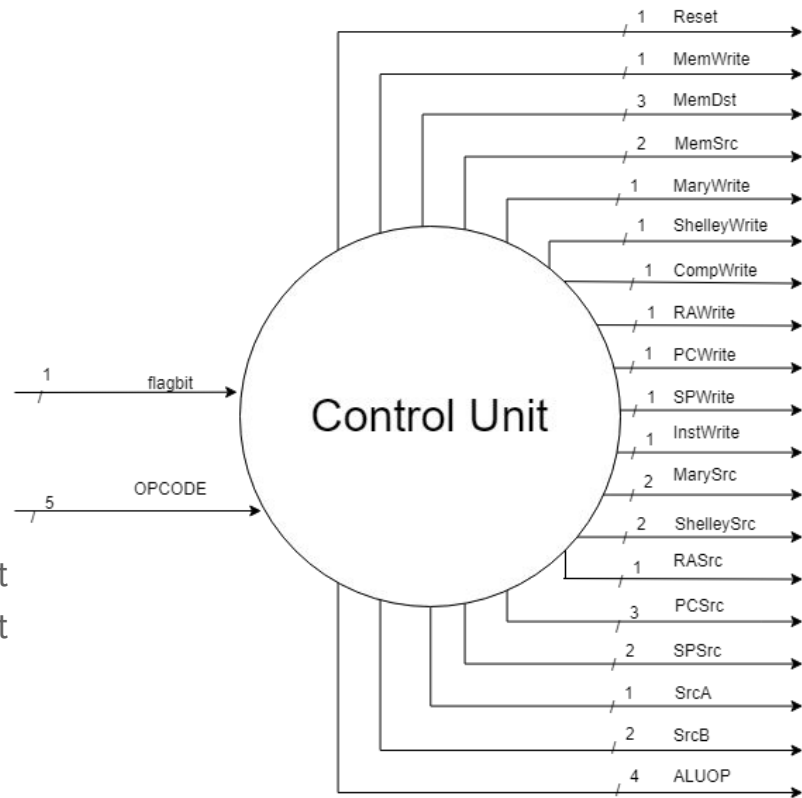
- The flagbit provides an alternate form of most instructions
 - Solution to duplication of instructions w/ two accumulators
- Indicated in assembly code using the '@' symbol
- Flagbit 0: Operation is performed on Mary and an immediate
- Flagbit 1: Operation is performed on Mary and Shelley
- Example: aadd instruction
 - `aadd 5` `//adds 5 to the value in Mary`
 - `aadd@` `//adds the value in Shelley to the value in Mary`

Assembler

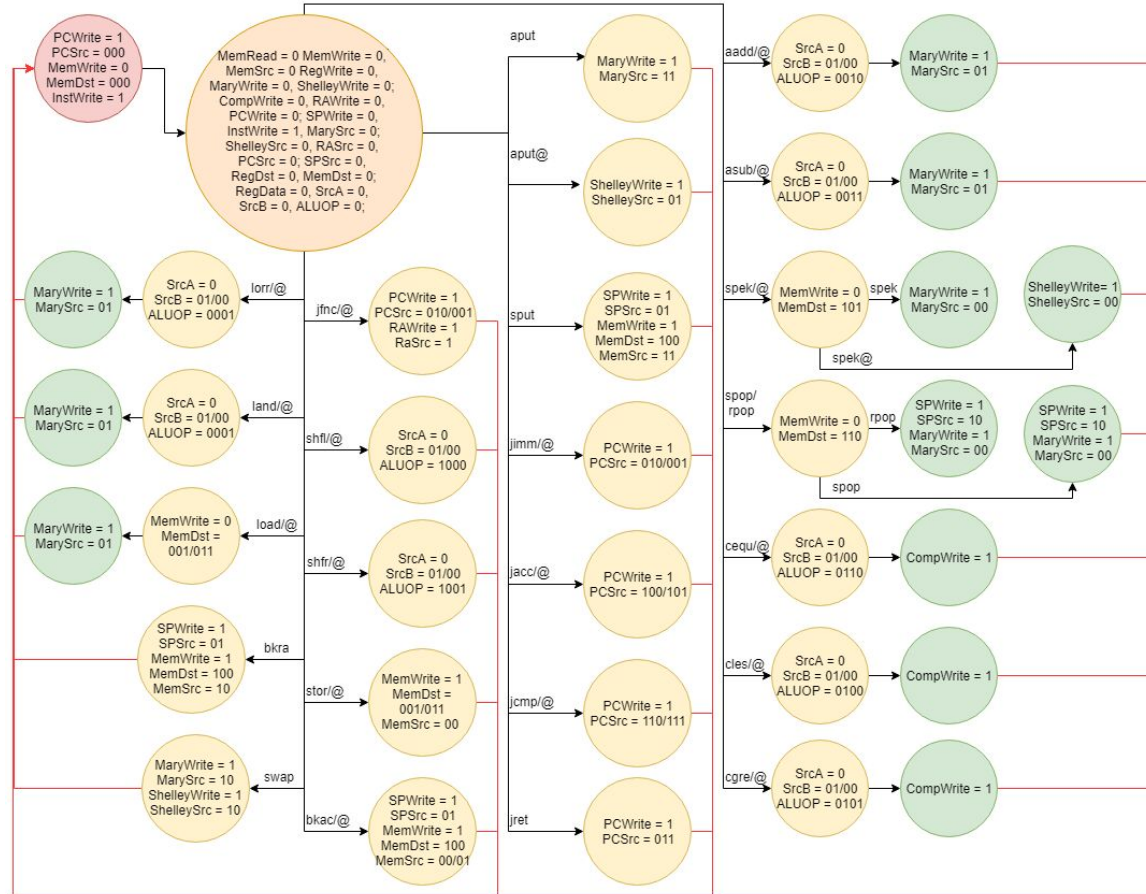
- Written in python
- Supports the entirety of our instruction set
- Auto-injects our kernel into the memory space for interrupts
- Can output debug data during assembling step to ease debugging

Control Unit

- For regular functionality...
 - Controls everything except I/O
 - 2 inputs (OPCODE, flagbit)
 - 19 outputs
- States:
 - 4 general states:
 - Fetch-- fetches instruction
 - Decode-- what OPCODE, flagbit?
 - 3-- determined by OPCODE and flagbit
 - 4-- determined by OPCODE and flagbit



Finite State Machine:



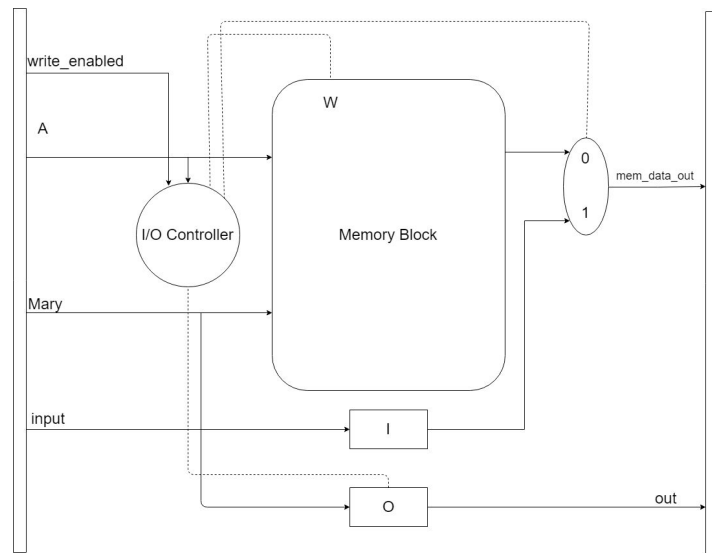
Memory

- We didn't use the Xilinx block memory generator
- Didn't have to use .coe files; used .txt files instead
- Verilog provides a command to read binary data (instructions) from files
 - `$readmemb("file.txt", memory);`



Memory-Mapped I/O:

- Controlled via a “sub-control unit” inside our memory block
- Uses memory register 255
 - `stor 255` //store value from Mary into output register
 - `load 255` //load value from input register into Mary



Interrupts

- Frankie provides support for input via interrupts
- Interrupt is triggered when the value in the input register changes
- Control unit handles the interrupt by jumping to the kernel
- Program returns from kernel upon reaching address 255 in memory

OUR KERNEL

load 255

Euclid's

```
prep:
  aput@ 2 //int m; m = 2;
  bkac@ //back up shelley's initial value
relprime_loop:
  jfnc 6 //gcd
  cequ 1 //if (gcd(n, m) == 1
  jcmp@ 23 //relprime_end
  spop //restore shelley's initial value
  aadd 1 //m=m+1
  bkac //back up shelley
  swap //put a in front (prep args)
gcd:
  cequ 0 //if (a == 0)
  jimm@ 2
  swap //prep b for return
  jret //return b;
  jcmp@ -3
  bkac //back up mary
gcd_loop:
  jcmp@ 10 //break
  cgre@ //if a > b
  jcmp@ 5 //jump past a=a-b to b=b-a (gcd_2)
  swap
  asub@ //else b=b-a
  cequ 0 //if b == 0
  swap
  jimm@ -8 //gcd_loop
gcd_2
  asub@ //a=a-b
  cequ 0 //if b == 0
  jimm@ -11 //gcd_loop
gcd_end:
  spop@ //pop backup of mary back into mary
  jret //return a
relprime_end:
  spop //put m into mary to prep for return

kernel:
  load 255 //load input into mary
```

```
int
relPrime(int n)
{
```

```
    int m;
```

```
    m = 2;
```

```
    while (gcd(n, m) != 1) {
        m = m + 1;
    }
```

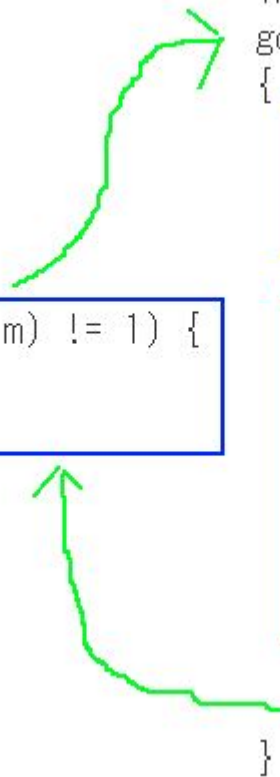
```
    return m;
}
```

```
int
gcd(int a, int b)
{
```

```
    if (a == 0) {
        return b;
    }
```

```
    while (b != 0) {
        if (a > b) {
            a = a - b;
        } else {
            b = b - a;
        }
    }
```

```
    return a;
}
```



Performance Data

- Memory space: 62 bytes
- Instruction count: 61286
- Cycle count: 214466
- CPI: 3.4994
- Cycle time: 9.247 ns = 108.143 MHz
- Runtime for Euclid's, $n=5040$:

2.2634 ms

Future Modifications

- Exception handling with interrupts
- Able to handle larger programs
- More memory
- FPGA board
- Even faster!