

Hash Tables

CS 1.3 - Core Data Structures



Linear Data Structures



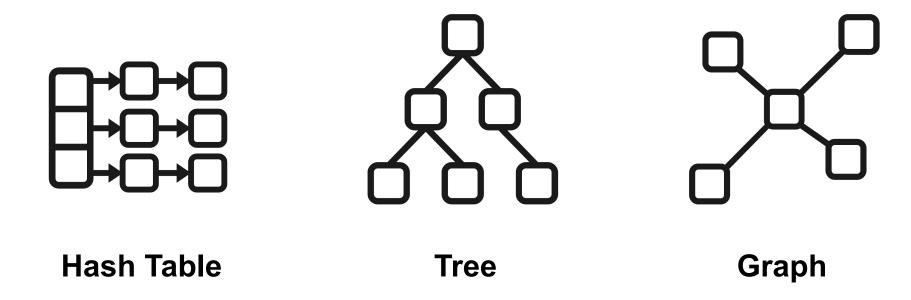
Array Linked List Queue Stack

A linear data structure has a specific order or sequence of its elements. There is a first element and a last element (or a top and a bottom).



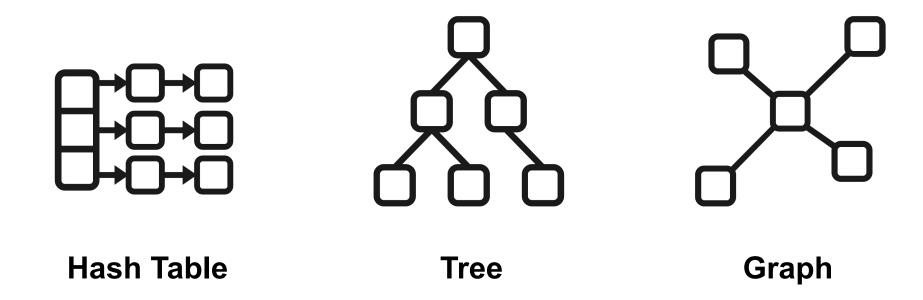
Non-Linear Data Structures





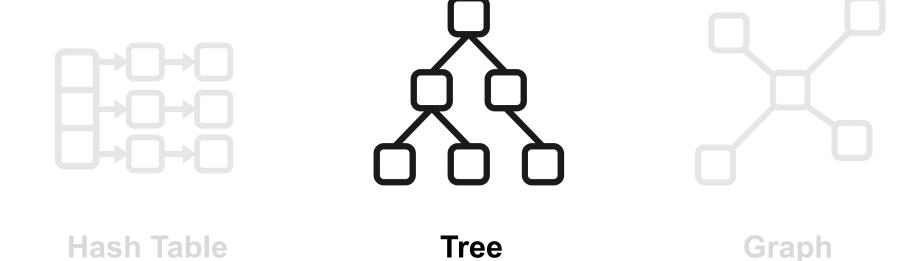
A non-linear data structure has no specific order or sequence of elements. Some structures can have multiple paths to connect to other elements.





Non-linear data structures can support multi-level storage and often cannot be traversed in single run.



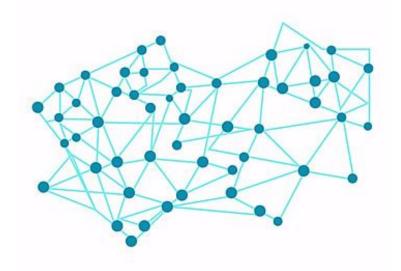


For example, try tracing the 3 different paths we can follow from the top level to the bottom level of the tree above.





What are examples of apps or technologies you have used that might use a non-linear data structure to organize their data?







Remember Dictionaries?





A dictionary data type is a collection which is unordered and mutable (changeable).



It is a set of key and value pairs.



| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | 281054 |



| key | value |
|-----------------|---------------------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | <mark>281054</mark> |

"What is the population (value) of city (key) Newark?" **Lookup Operation**



| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | 281054 |



"Irvington" 54233

"Add the city of Irvington and its 54,233 population." **Insert Operation**



| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| | |
| | |
| "Atlantic City" | 37999 |



"Camden" 74002

"Delete the city of Camden and its 74,002 population." **Delete Operation**



In Python, an easy way to create dictionaries using curly braces {}.



Hash Tables



What is a hash table?

A hash table is a data structure that maps keys to values.

A value can be **retrieved** from the hash table using its key.



Book metaphor

Book pages = array of buckets

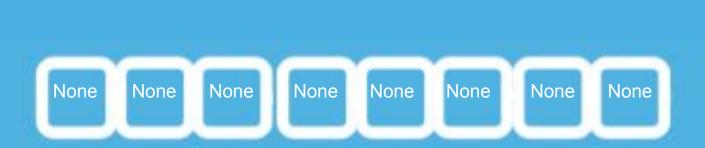
Book index = hash function



How to Build a Hash Table:

The create an array. Because we do not use arrays in Python, we create a list with "None" as the placeholder for each element.

For example, when creating a Hash Table of size 8:



[None] * 8



| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | 281054 |

"Paterson" 145710 "Atlantic City" 37999

"Camden" 74002 "Jersey City" 261940

"Newark" 281054



How to Build a Hash Table:

Create a hash function that turns each key into a number (called a hash code) that we can use to decide where in our array of buckets each key → value pair should be stored.

For example, we could make a hash function that looks at the **first letter** of each key and calculates how many characters away it is from the letter 'A'. Here are some examples: hash("Apple") \rightarrow 0, hash("Fig") \rightarrow 5, hash("Kiwi") \rightarrow 10



For example:

Let's find the hash code for "Paterson"

"Paterson" starts with P, which is 15 characters away from the letter 'A'.

| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Jersey City" | 261940 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | 281054 |



But wait... our Hash Table only has 8 buckets. Meaning it only has buckets with indexes ranging from index 0 to 7.

What can we do to change the value 15 to be within the range of 0 to 7?







15 mod 8 = 7

We could use the mod operator.



Check for Understanding



Compute the index for the key:value pairs in the table use the hash function we discussed on the previous slide.

| key | value | bucket |
|-----------------|--------|--------|
| "Paterson" | 145710 | 7 |
| "Jersey City" | 261940 | |
| "Camden" | 74002 | |
| "Atlantic City" | 37999 | |
| "Newark" | 281054 | |

So far we know...



To create a Hash Table we need:

- 1. Data in key \rightarrow value pairs.
- 2. An array of buckets with a specific fixed size.
- 3. A hash function to determine where each key value pair will be stored in the array of buckets.



Hash Collision



For example:

Let's find an index for "Paterson."

"Paterson" starts with P, which is 15 characters away from the letter A.

But what if we want to add "Princeton" to our hash table...

| key | value |
|-----------------|--------|
| "Paterson" | 145710 |
| "Princeton" | 31000 |
| "Camden" | 74002 |
| "Atlantic City" | 37999 |
| "Newark" | 281054 |

"Princeton" starts with P, which is 15 characters away from the letter A.

Check for Understanding



Compute the bucket for the key:value pairs in the table use the hash function we discussed on the previous slide.

| key | value | bucket |
|--------------|--------|--------|
| "Paterson" | 145710 | 7 |
| "Princeton" | 31000 | 7 |
| "Moorestown" | 20355 | 4 |
| "Montclair" | 38676 | 4 |



What is a hash collision?





What is a hash collision?

A hash collision is when multiple keys land in the same bucket in the array. When determining which hash function to use for your data, you should avoid collisions to keep your hash table efficient.



Example: DJB2



Handling Collision

Methods of Handling Collision



- 1. Linear Probing
- 2. Chaining
- 3. Double Hashing (not covered in CS 1.3)

Linear Probing

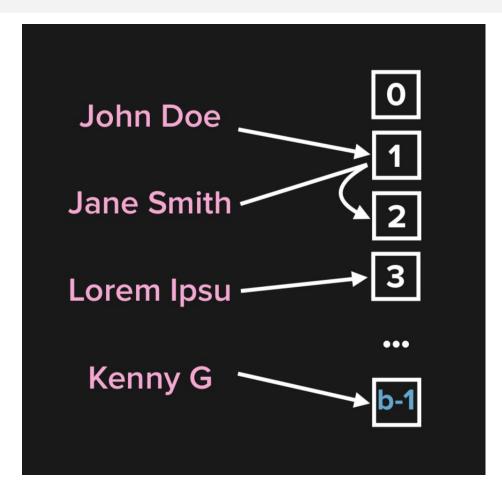


Each bucket contains at most one entry

On collision - find next open bucket, add new entry there

To retrieve - find bucket, if that's not entry, try next bucket until you find entry or empty bucket

Python's dict uses probing



Chaining

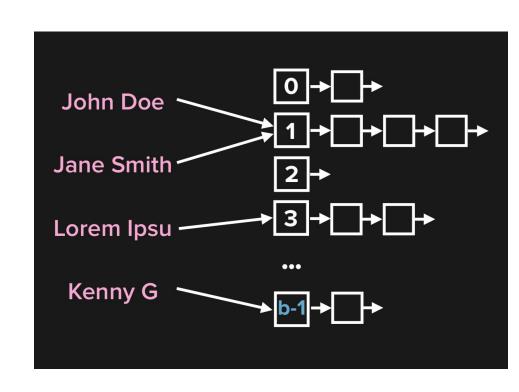


Each bucket contains a **linked list** of entries

On collision - add new entry to the bucket's linked list

To retrieve - find bucket, find entry in linked list

We will use chaining with linked lists to implement our hash table





Let's Draw a Hash Table

Load Factor

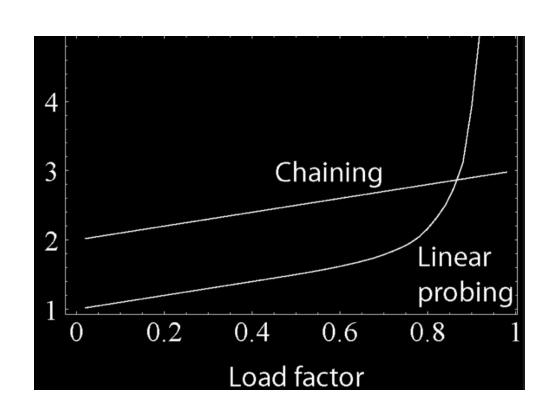


Load Factor = entries / buckets

Load Factor affects performance

Collision Resolution method also

affects performance



Resizing



Once the load factor reaches a certain threshold and begins to impact performance we need to resize the hash table

- 1. Store what's already in our hash table
- 2. Make a new hash table with twice the number of buckets
- 3. Rehash our stored items

Amortized Time Complexity



"Amortized time is the way to express the time complexity when an algorithm has the very bad time complexity only once in a while besides the time complexity that happens most of time"

 $O(1)^*$

https://medium.com/@satorusasozaki/amortized-time-in-the-time-complexity-of-an-algorithm-6dd9a5d38045

Time Complexity Analysis Review



- items
- get
- set
- delete

Hash table with chaining animation

Time Complexity Analysis Review



| Operation | Amortized | Worst |
|-----------|-----------|-------|
| items() | O(n) | O(n) |
| get() | O(1) | O(n) |
| set() | O(1) | O(n) |
| delete() | O(1) | O(n) |



Module 5: Hash Tables