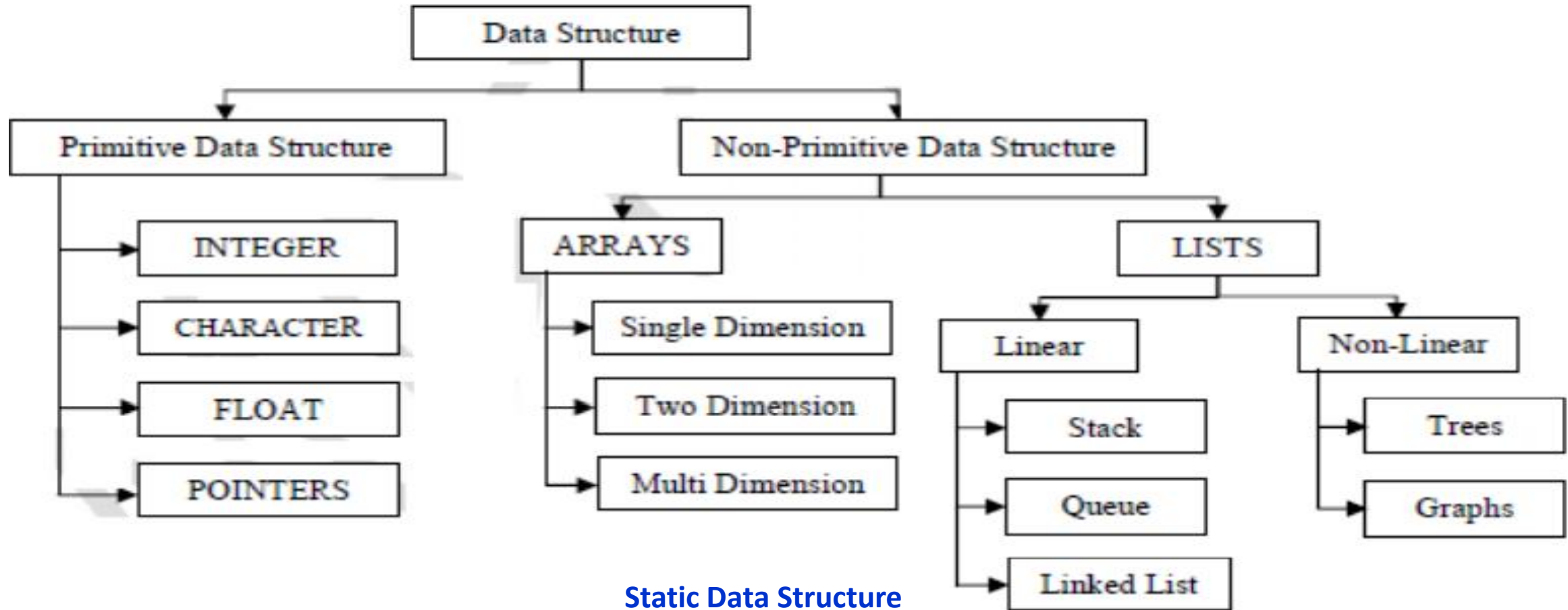


Algorithms & Data Structure

Kiran Waghmare

Classification of Data Structure



Static Data Structure

Dynamic data structure

Non-Primitive Data Structure

- The most commonly used operation on data structure are broadly categorized into following types:
 - Traversal
 - Insertion
 - Selection
 - Searching
 - Sorting
 - Merging
 - Destroy or Delete

Home work

HighArray
public HighArray()//Constructor
public boolean find (int key) public void insert(int value) public boolean delete(int long) public void display()

HighArrayApp
main() create object
insert()// all elements
display() find() delete()

Algorithm Complexity:-

Two main factors, which decide the efficiency of X.

- **Time Factor** – Time is measured by counting the number of key operations such as comparisons in the sorting algorithm.
- **Space Factor** – Space is measured by counting the maximum memory space required by the algorithm.

Asymptotic Notations

Asymptotic analysis of an algorithm refers to defining the mathematical boundation of its run-time performance. Using asymptotic analysis, we can very well conclude the best case, average case, and worst case scenario of an algorithm.

Usually, the time required by an algorithm falls under three types –

- **Best Case** – Minimum time required for program execution.
- **Average Case** – Average time required for program execution.
- **Worst Case** – Maximum time required for program execution.

Asymptotic Notations

Following are the commonly used asymptotic notations to calculate the running time complexity of an algorithm.

- **O Notation**
- **Ω Notation**
- **θ Notation**

Linear Search

- Find 37?

0	1	2	3	4	5	6	7	8
20	35	37	40	45	50	51	55	67



≠



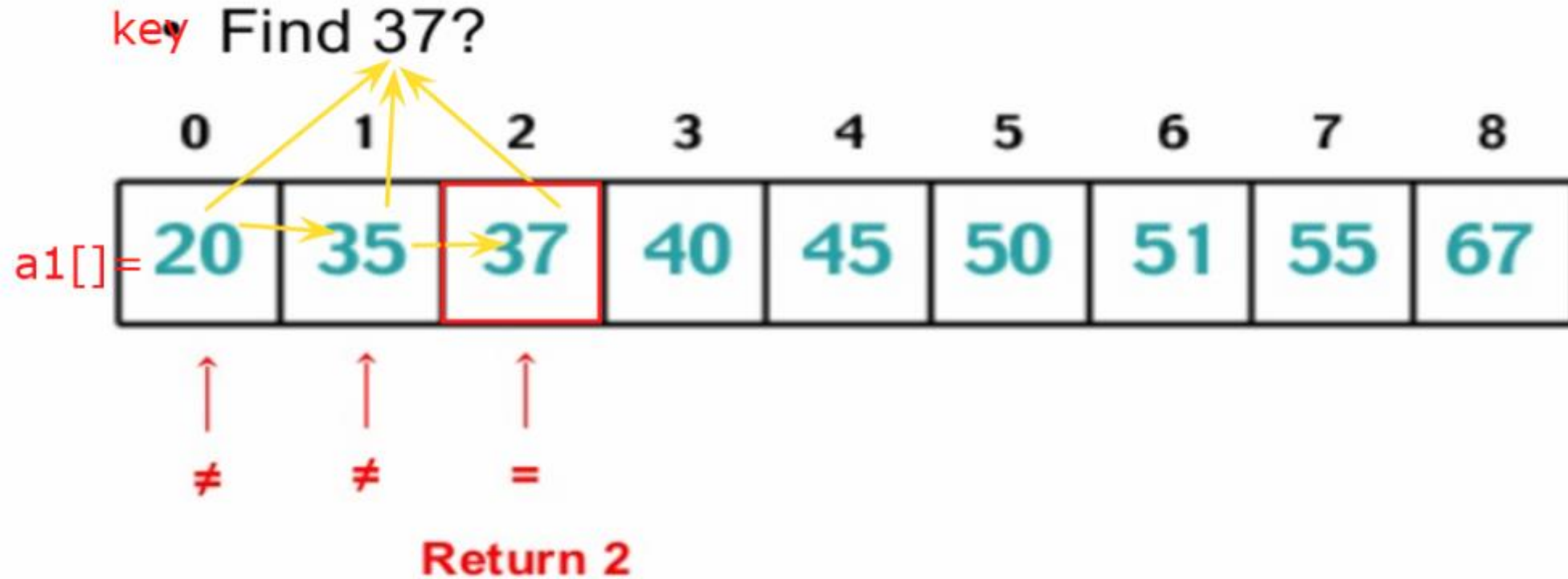
≠



=

Return 2

Linear Search



Linear Search

Algorithm

Consider **LA** is a linear array with **N** elements and **K** is a positive integer such that **$K \leq N$** . Following is the algorithm to find an element with a value of **ITEM** using sequential search.

1. Start
2. Set $J = 0$
3. Repeat steps 4 and 5 while $J < N$
4. IF $LA[J]$ is equal ITEM THEN GOTO STEP 6
5. Set $J = J + 1$
6. PRINT J , ITEM
7. Stop

Home work

Problem: Given an array `arr[]` of `n` elements, write a function to search a given element `x` in `arr[]`.

Examples :

Input : `arr[] = {10, 20, 80, 30, 60, 50, 110, 100, 130, 170}`

`x = 110;`

Output : 6

Element x is present at index 6

Input : `arr[] = {10, 20, 80, 30, 60, 50, 110, 100, 130, 170}`

`x = 175;`

Output : -1

Element x is not present in `arr[]`.

Binary Search

- Find 37?
 - Sort Array.

0	1	2	3	4	5	6	7	8
20	35	37	40	45	50	51	55	67



$$\text{mid} = (\text{low} + \text{high}) / 2$$
$$= (0 + 8) / 2 = 4$$

mid = 35

mid = 37

key = 37

Iteration 1:
if(a1[mid]==key)
return mid

else

if key > mid --> high

else

if key < mid --> low

- Array Applications
- Stack
- Stack Applications

Binary Search

```
Procedure binary_search
  A ← sorted array
  n ← size of array
  x ← value to be searched

  Set lowerBound = 1
  Set upperBound = n

  while x not found
    if upperBound < lowerBound
      EXIT: x does not exists.

    set midPoint = lowerBound + ( upperBound - lowerBound ) / 2

    if A[midPoint] < x
      set lowerBound = midPoint + 1

    if A[midPoint] > x
      set upperBound = midPoint - 1

    if A[midPoint] = x
      EXIT: x found at location midPoint
  end while
end procedure
```

Problem Statement 1: Find the Missing Number

You are given a list of $n-1$ integers and these integers are in the range of 1 to n . There are no duplicates in the list. One of the integers is missing in the list. Write an efficient code to find the missing integer.

Example:

Input: `arr[] = {1, 2, 4, 6, 3, 7, 8}`

Output: 5

Explanation: The missing number from 1 to 8 is 5

Input: `arr[] = {1, 2, 3, 5}`

Output: 4

Explanation: The missing number from 1 to 5 is 4

Problem statement: Search an element in a sorted and rotated array

Example:

Input : arr[] = {5, 6, 7, 8, 9, 10, 1, 2, 3};

key = 3

Output : Found at index 8

Input : arr[] = {5, 6, 7, 8, 9, 10, 1, 2, 3};

key = 30

Output : Not found

Input : arr[] = {30, 40, 50, 10, 20}

key = 10

Output : Found at index 3

Problem statement 3: Program for array rotation

Write a function `rotate(ar[], d, n)` that rotates `ar[]` of size `n` by `d` elements.

Array

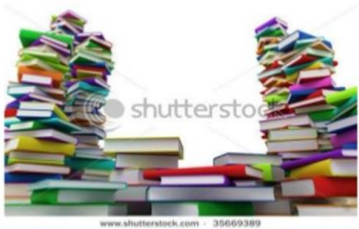
1	2	3	4	5	6	7
---	---	---	---	---	---	---

Rotation of the above array by 2 will make array

ArrayRotation1

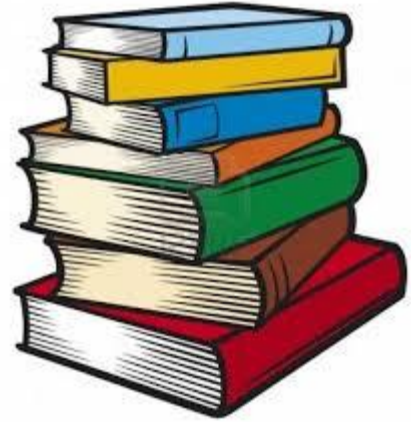
3	4	5	6	7	1	2
---	---	---	---	---	---	---

Examples of stack



Stacks

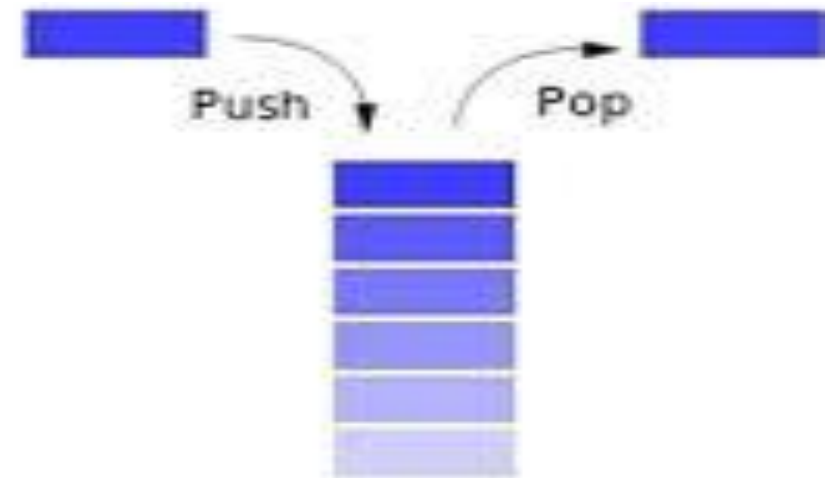
Kiran Waghmare



Stack of books



Stack of Coins



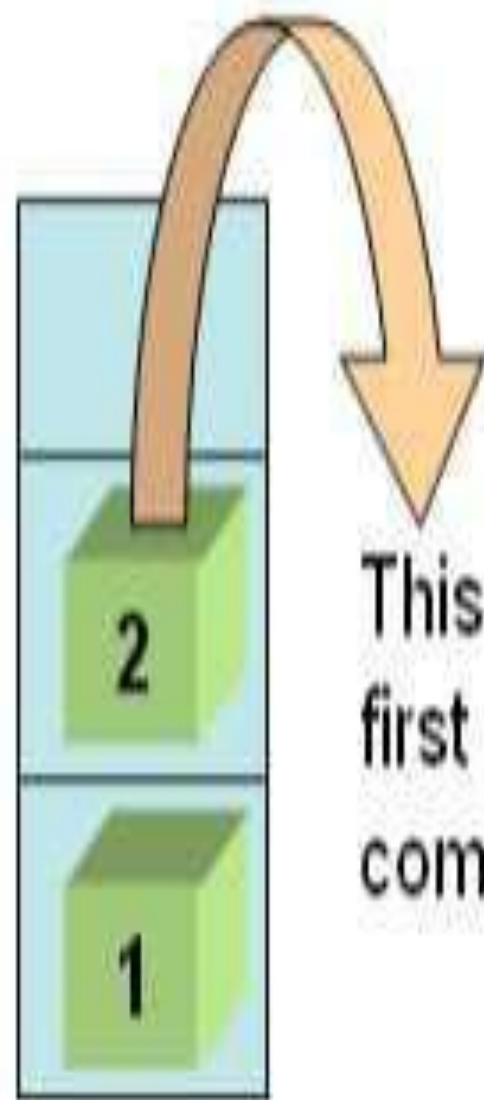
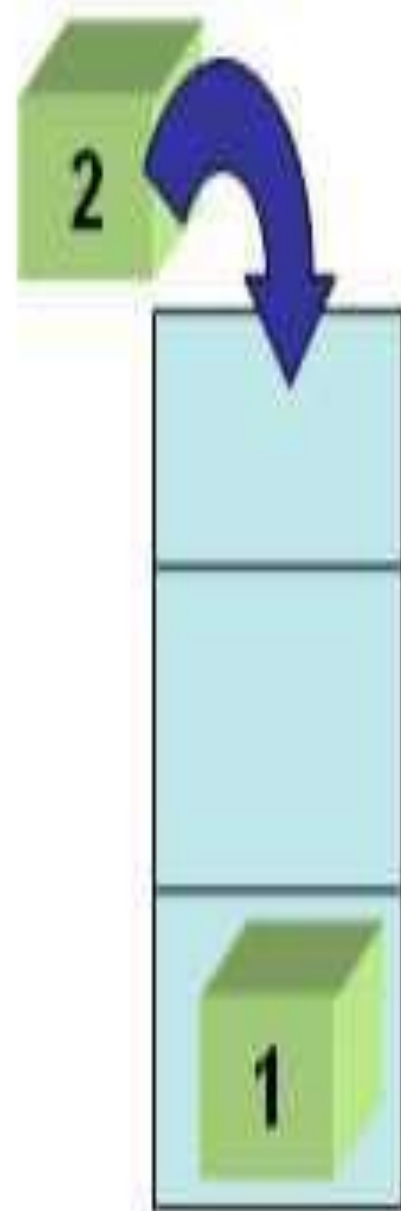
Memory stack

Standard Stack Operations

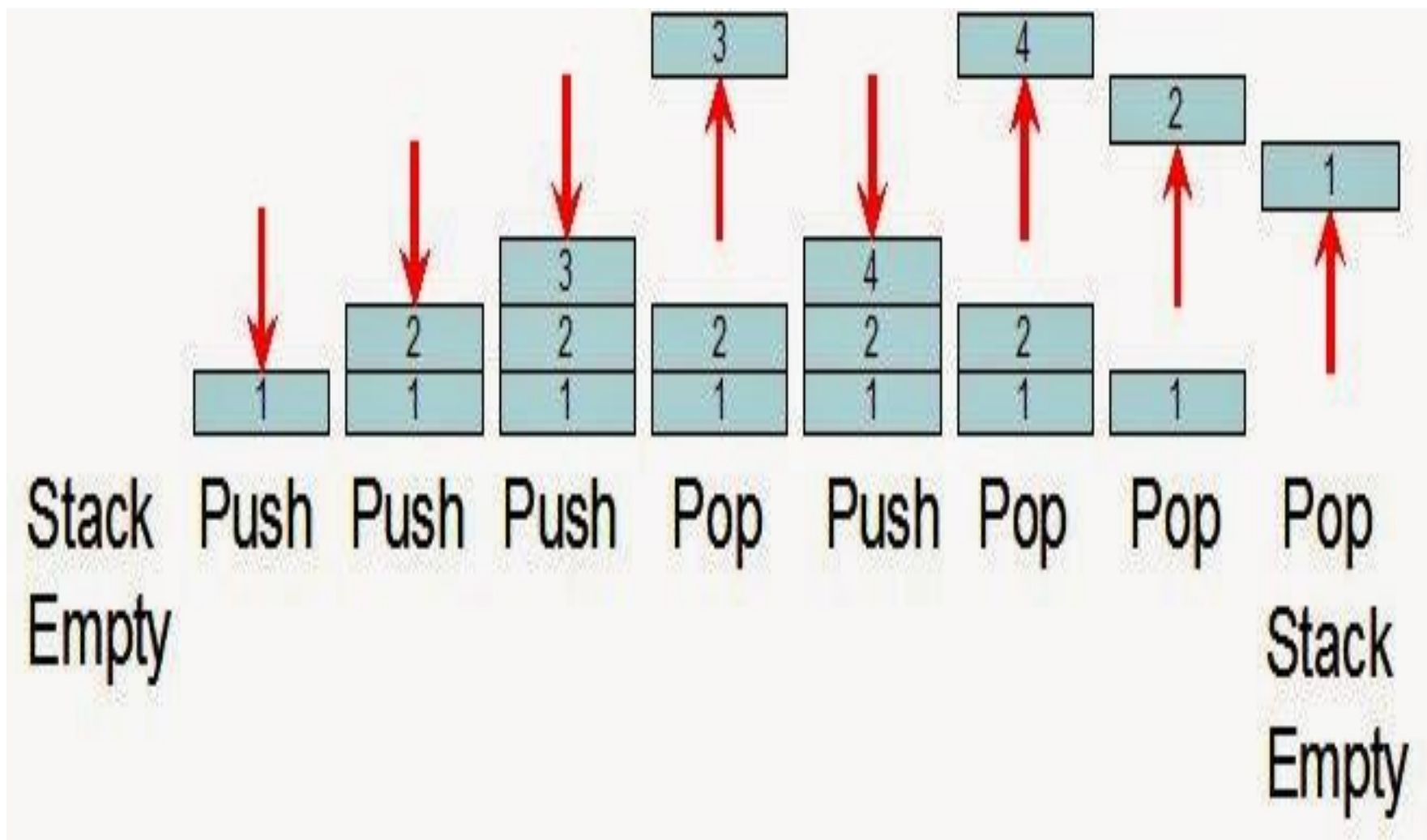
- The following are some common operations implemented on the stack:
- **push():**
 - When we insert an element in a stack then the operation is known as a push. If the stack is full then the overflow condition occurs.
- **pop():**
 - When we delete an element from the stack, the operation is known as a pop. If the stack is empty means that no element exists in the stack, this state is known as an underflow state.
- **isEmpty():**
 - It determines whether the stack is empty or not.
- **isFull():**
 - It determines whether the stack is full or not.'
- **peek():**
 - It returns the element at the given position.
- **count():**
 - It returns the total number of elements available in a stack.
- **change():**
 - It changes the element at the given position.
- **display():**
 - It prints all the elements available in the stack.



Empty Stack

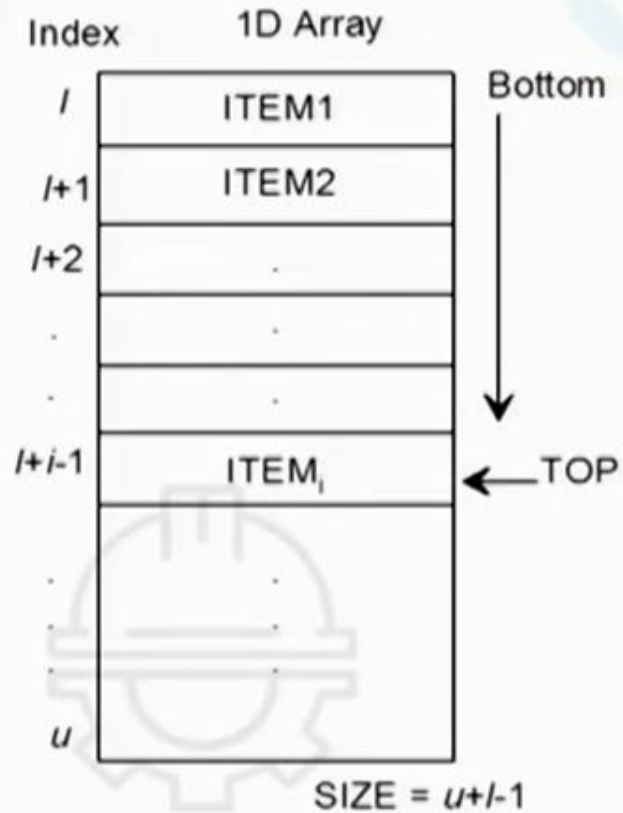


This will be the
first object to
come out.

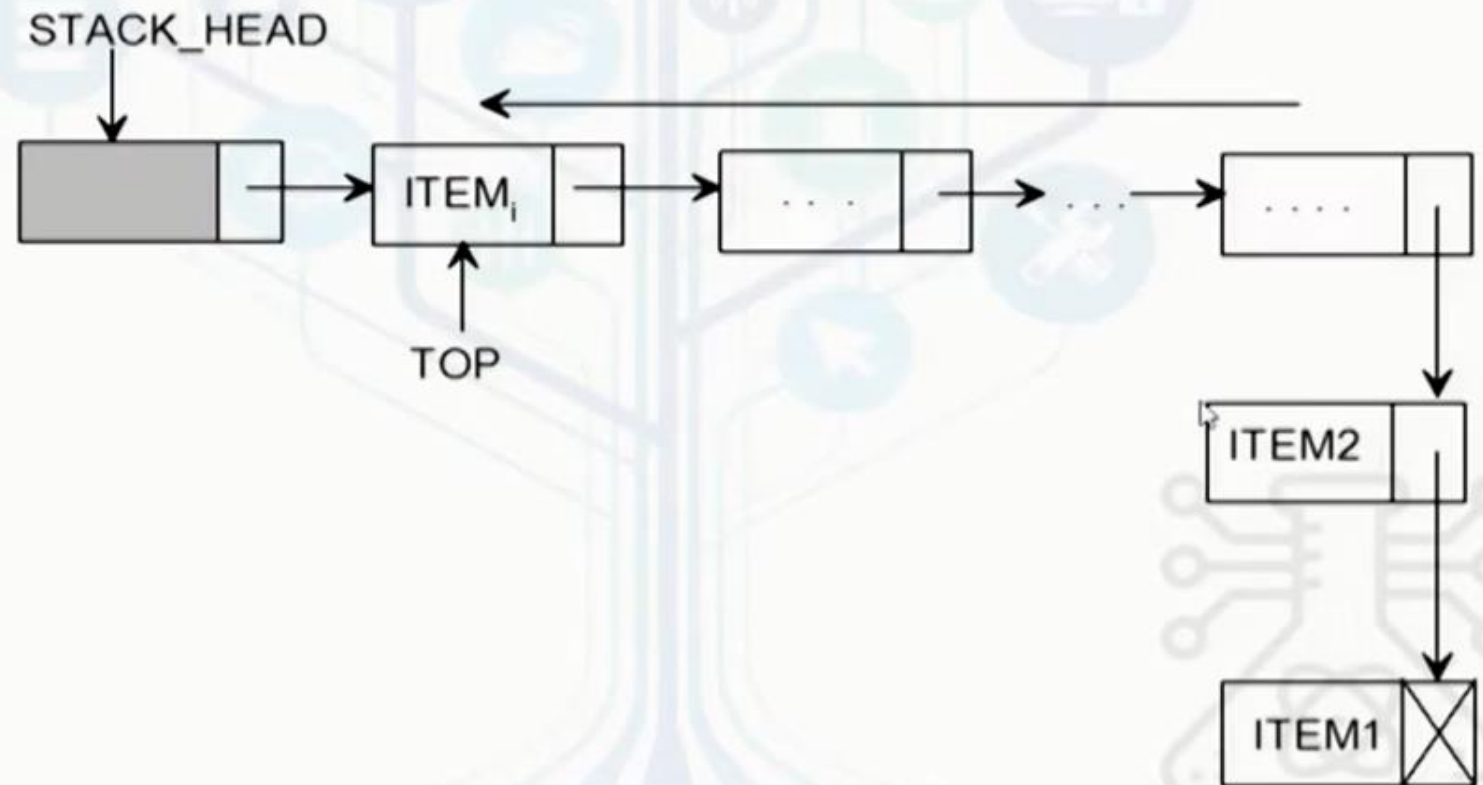


Memory representations

Array representation



Linked list representation




```

S = new int[size];
top = -1;
}

```

```

public void push(int j)
{
    S[++top] = j;
}

```

```

public int pop()
{
    return S[top--];
}

```

```

public int peek()
{
    return S[top];
}

```

```

isEmpty()

```

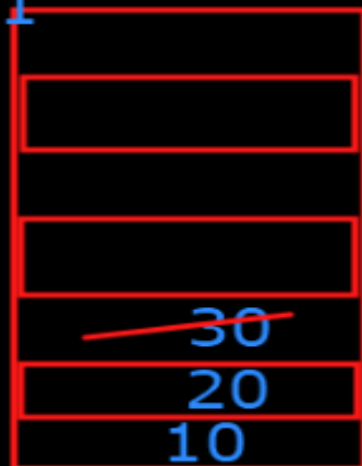
```

Push(10);
Push (20);
Push(30);
pop();

```

size=7

0 to 6

~~-1~~

TOP

TOP

s[0]

TOP = ~~-1~~ 0
STACK

```

class StackApp
{

```