CS186 Discussion #8

(Functional Dependencies)

Functional Dependencies

- X → Y reads "X determines Y"
- Used to detect redundancies and refine schema

id	rating	wage
1	3	20
2	2	10
3	3	20
4	2	10

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rating → wage

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id → rating, wage

Functional Dependencies

- Key → All attributes of relation

id	rating	wage
1	3	20
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id → rating, wage

Armstrong's Axioms

- **Reflexivity**: if $X \supseteq Y$, then $X \rightarrow Y$
  - **Examples**: $A \rightarrow A$, $AB \rightarrow A$
- **Augmentation**: if $X \rightarrow Y$, then $XZ \rightarrow YZ$ for any Z
- **Transitivity**: if $X \rightarrow Y$ and $Y \rightarrow Z$, then $X \rightarrow Z$
- Useful rules derived from Armstrong's Axioms:
  - **Union**: if $X \rightarrow Y$ and $X \rightarrow Z$, then $X \rightarrow YZ$
  - **Decomposition**: if $X \rightarrow YZ$, then $X \rightarrow Y$ and $X \rightarrow Z$

Armstrong's Axioms

- If $XA \rightarrow YA$, can you infer $X \rightarrow Y$?

Armstrong's Axioms

- If $XA \rightarrow YA$, can you infer $X \rightarrow Y$?
  - Trivial example: $A \rightarrow YA$

Worksheet: Page 2

Consider the **Works_In**(**S**sn, parking_**L**ot_num, **D**epartment_id, s**I**nce) relation. If **S** (ssn) is a key for this relationship, what is the functional dependency we can infer from that?

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- S → SLDI

If employees in the same department are given the same parking lot number, what additional functional dependency can we infer?

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- $D \rightarrow L$

Find the set of functional dependencies:

Flights(**F**light_no, **D**ate, f**R**om, **T**o, **P**lane_id),
ForeignKey(**P**lane_id)

Planes(**P**lane_id, t**Y**pe)

Seat(**S**eat_no, **P**lane_id, **L**egroom), ForeignKey(**P**lane_id)

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Seat(<u>**S**eat_no, **P**lane_id</u>, **L**egroom), ForeignKey(**P**lane_id)

- FD → RTP

- P → Y

- SP → L

Closures

- Functional dependency closure: F+
  - Set of all FDs implied by F, including trivial dependencies
  - Example: F = {A → B, B → C}
    - F+ = {A → B, B → C, A → C, A → A, A → AB, …}

Closures

- Attribute closure: X+
  - Given just X, what can we determine?
  - Example: $F = \{A \rightarrow B, B \rightarrow C\}$
    - A+ = ABC

Attribute Closures

- A methodical algorithm, given a set of FDs F:
  - Initialize X+ := X
  - Repeat until no change:
    - If U → V is in F such that U is in X+, add V to X+

- R = ABCDE
- F = {B → CD, D → E, B → A, E → C, AD → B }
- B+ = ?

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- F = {B → CD, D → E, B → A, E → C, AD → B }
- B+ = ABCDE   B is a key of R!

Boyce-Codd Normal Form (BCNF)

- Motivation: Schema design is hard, want a way to ensure a reasonable design
- BCNF ≈ “reasonable schema”

Boyce-Codd Normal Form (BCNF)

- Definition: Relation R with FDs F is in BCNF if for all $X \rightarrow A$ in F+:
  - $X \rightarrow A$ is reflexive (a trivial FD) OR
  - X is a superkey for R
    - Superkey: Key that does not need to be minimal

BCNF Decomposition

- If $X \rightarrow A$ violates BCNF, decompose R into R - A and XA
  - Repeat as necessary

- $R = ABCEG$
- $F = \{AB \rightarrow C, AC \rightarrow B, BC \rightarrow G, E \rightarrow G\}$

BCNF Decomposition

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- R = ABCEG
- F = {$AB \rightarrow C$, $AC \rightarrow B$, $BC \rightarrow G$, $E \rightarrow G$}
- ABEG, ABC

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- ABEG, ABC
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Decomposition Properties

- Lossless join: Can we reconstruct R?
  - Decomposing R into X and Y is lossless iff:
    - $X \cap Y \rightarrow X$, or
    - $X \cap Y \rightarrow Y$
  - BCNF usually produces lossless join when decomposing $X \rightarrow A$, **except** when A contains some part of X

Decomposition Properties

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  - Decomposing R into X and Y is lossless iff:
    - $X \cap Y \rightarrow X$, or
    - $X \cap Y \rightarrow Y$
- Example:
  - ABC decomposed to AB, BC
  - FDs: $A \rightarrow B$, $C \rightarrow B$
  - This is lossy! $AB \cap BC = B \rightarrow B$

Decomposition Properties

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    - $X \cap Y \rightarrow Y$
- Example:
  - ABC decomposed to AC, BC
  - FDs: $A \rightarrow B$, $C \rightarrow A$
  - This is lossless! $AC \cap BC = C \rightarrow AC$

Decomposition Properties

- Dependency-preserving: Can we verify all FDs without joins?

- Example: ABE, EG, ABC
  - $F = \{AB \rightarrow C, AC \rightarrow B, BC \rightarrow G, E \rightarrow G\}$

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Decomposition Properties

- Dependency-preserving: Can we verify all FDs without joins?

- Example: ABE, EG, ABC
  - F = {AB → C, AC → B, BC → G, E → G}
  - This dependency was not preserved!
  - We can fix this by adding BCG, but this may break BCNF.

Minimal Cover

- G for a set of FDs F
  - Closure of G = closure of F
  - Right hand side of each FD is G is a single attribute
- Implies lossless join and dependency preserving decomposition

- Example: $A \rightarrow B$, $ABCD \rightarrow E$, $EF \rightarrow GH$, $ACDF \rightarrow EG$

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- Example: $A \rightarrow B$, $ABCD \rightarrow E$, $EF \rightarrow GH$, $ACDF \rightarrow EG$
- $A \rightarrow B$

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- Example: A → B, ABCD → E, EF → GH, ACDF → EG
- A → B, ACD → E

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- Example: A → B, ABCD → E, EF → GH, ACDF → EG
- A → B, ACD → E, EF → G, EF → H

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- Example: A → B, ABCD → E, EF → GH, ACDF → EG
- A → B, ACD → E, EF → G, EF → H

Worksheet: Page 3

Now consider the attribute set R = ABCDE and the FD set F = {AB → C, A → D, D → E, AC → B}. Compute the closure for the following attributes.

- A:
- AB:
- B:
- D:

Now consider the attribute set R = ABCDE and the FD set F = {AB → C, A → D, D → E, AC → B}. Compute the closure for the following attributes.

- A: ADE
- AB: ABCDE
- B: B
- D: DE

Consider the relation with attributes ABCDE and FD's: $AB \rightarrow ABCDE$, $D \rightarrow A$, $E \rightarrow B$

Is this in BCNF?

Consider the relation with attributes ABCDE and FD's: AB → ABCDE, D → A, E → B
Is this in BCNF?

- **This is not BCNF**
  - D → A and E → B are not trivial, and neither D nor E are keys.
  - This is 3NF because both A and B are parts of a key of the relation (see the definition of 3NF).

Decompose R = ABCDEFG into BCNF, given the FD set:

$F = \{AB \rightarrow CD,\ C \rightarrow EF,\ G \rightarrow A,\ G \rightarrow F,\ CE \rightarrow F\}$.

Decompose R = ABCDEFG into BCNF, given the FD set:
F = {AB → CD, C → EF, G → A, G → F, CE → F}.

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Decompose R = ABCDEFG into BCNF, given the FD set:
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- ABEFG, ABCD
- BEFG, ABCD, AG

Decompose R = ABCDEFG into BCNF, given the FD set:
F = {AB → CD, C → EF, G → A, G → F, CE → F}.

- ABEFG, ABCD
- BEFG, ABCD, AG
- BEG, FG, ABCD, AG

Decompose R = ABCDEFG into BCNF, given the FD set:
F = {AB → CD, C → EF, G → A, G → F, CE → F}.

- ABEFG, ABCD
- BEFG, ABCD, AG
- BEG, FG, ABCD, AG

Does BEG, FG, ABCD, AG preserve dependencies?

$F = \{AB \rightarrow CD, C \rightarrow EF, G \rightarrow A, G \rightarrow F, CE \rightarrow F\}.$

Does BEG, FG, ABCD, AG preserve dependencies?
$F = \{AB \rightarrow CD, C \rightarrow EF, G \rightarrow A, G \rightarrow F, CE \rightarrow F\}$.

- No, $C \rightarrow EF$ and $CE \rightarrow F$ are not preserved.

Give a minimal cover for:

$F = \{AB \rightarrow CD,\ C \rightarrow EF,\ G \rightarrow A,\ G \rightarrow F,\ CE \rightarrow F\}.$

Give a minimal cover for:

$F = \{AB \to CD, C \to EF, G \to A, G \to F, CE \to F\}$.

- $AB \to C$
- $AB \to D$

Give a minimal cover for:

$F = \{AB \rightarrow CD, C \rightarrow EF, G \rightarrow A, G \rightarrow F, CE \rightarrow F\}$.

- $AB \rightarrow C$
- $AB \rightarrow D$
- $C \rightarrow F$
- $C \rightarrow E$

Give a minimal cover for:

$F = \{AB \rightarrow CD,\ C \rightarrow EF,\ G \rightarrow A,\ G \rightarrow F,\ CE \rightarrow F\}.$

- $AB \rightarrow C$
- $AB \rightarrow D$
- $C \rightarrow F$
- $C \rightarrow E$
- $G \rightarrow A$

Give a minimal cover for:

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Give a minimal cover for:

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- $AB \rightarrow D$
- $C \rightarrow F$
- $C \rightarrow E$
- $G \rightarrow A$
- $G \rightarrow F$

More lossless practice:

F = {AB -> CDE, BE -> X, A -> E}

Is ABC, BCDEX a lossless decomposition?

More lossless practice:
F = {AB -> CDE, BE -> X, A -> E}
Is ABC, BCDEX a lossless decomposition?

- No, it is lossy. ABC ∩ BCDEX = BC, which is not a superkey of ABC nor BCDEX.

R: ABCDE

Given FD $=\{AE \rightarrow BC, AC \rightarrow D, CD \rightarrow BE, D \rightarrow E\}$

Give three candidate keys.

R: ABCDE

Given FD ={AE → BC, AC → D, CD → BE, D → E}

Give three candidate keys.

- AE, AC and AD are candidate keys, as each of their attribute closures include all attributes and no subset of them is a super key by itself.

R: ABCDE

Given FD ={$AE \to BC$, $AC \to D$, $CD \to BE$, $D \to E$}

Is R already in BCNF?

R: ABCDE

Given FD ={AE → BC, AC → D, CD → BE, D → E}

Is R already in BCNF?

- No, because both CD → BE and D → E violate BCNF.

R: ABCD

Given FD $=\{A \rightarrow B, B \rightarrow D, C \rightarrow D\}$

Decomposed to AB, CD, AC. Is this lossless?

R: ABCD

Given FD ={A → B, B → D, C → D}

Decomposed to AB, CD, AC. Is this lossless?

- Yes, a lossless decomposition would be: ABC CD which is lossless because C is a key for CD and then a further decomposition of ABC into AB and AC which is lossless because A is a key for AB.