CS186 Discussion #3

(Joins, Heap Files)

Joins

```
SELECT <columns>
  FROM <tables>
WHERE <predicate>
[GROUP BY <column list>
 [HAVING <predicate>]];
```

Songs(song_id, song_name, album_num, weeks_in_top_40)
Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
year_released, genre)

```
SELECT *
FROM Artists, Albums
WHERE Artists.artist_id =
Albums.artist_num;
```

Songs(song_id, song_name, album_num, weeks_in_top_40)
Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
year_released, genre)

```
SELECT *
FROM Artists A, Albums B
WHERE A.artist_id = B.artist_num;
```

```
Songs(song_id, song_name, album_num, weeks_in_top_40)
Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
       year_released, genre)
```

Write a SQL expression for the following query:

The name of all songs with the genre “country” which have spent more than 2 weeks in the top 40.

```
Songs(song_id, song_name, album_num, weeks_in_top_40)
Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
       year_released, genre)
```

Write a SQL expression for the following query:

The name of all songs with the genre “country” which have spent more than 2 weeks in the top 40.

```
SELECT Songs.song_name FROM Albums,
Songs WHERE Songs.album_num =
Albums.album_id AND Albums.genre =
'country' AND Songs.weeks_in_top_40 > 2;
```

```
Songs(song_id, song_name, album_num, weeks_in_top_40)
Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
       year_released, genre)
```

Write a SQL expression for the following query:

The number of albums released by each artist.

```
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Artists(artist_id, artist_name, first_year_active)
Albums(album_id, album_name, artist_num,
       year_released, genre)
```

Write a SQL expression for the following query:

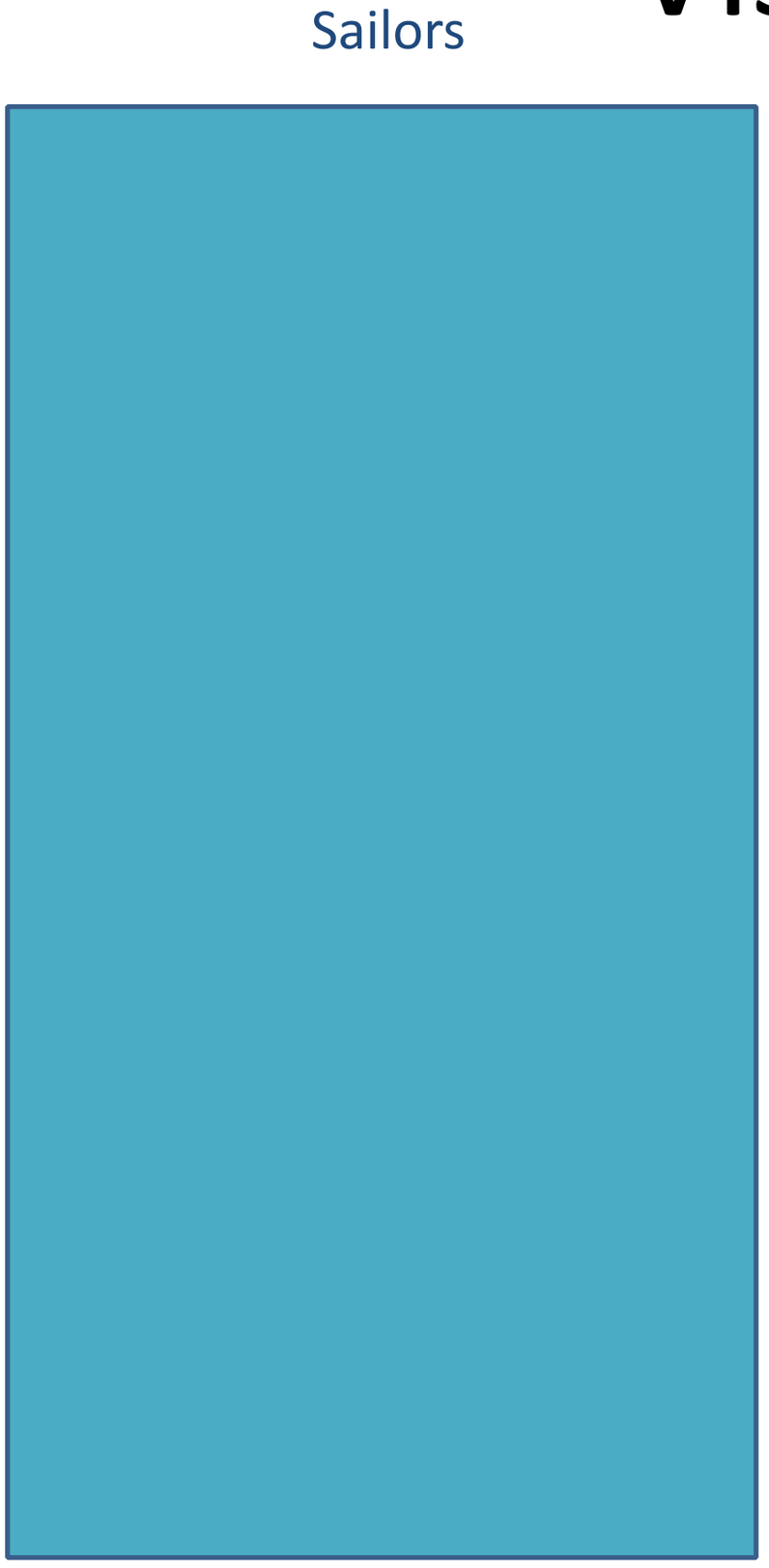
The number of albums released by each artist.

```
SELECT count(*) FROM Artists, Albums
WHERE Artists.artist_id =
Albums.artist_num GROUP BY
Artists.artist_id;
```

Join Algorithms

```
SELECT * FROM Sailors S, Reserves R
WHERE S.sid = R.sid;
```

Visualizations



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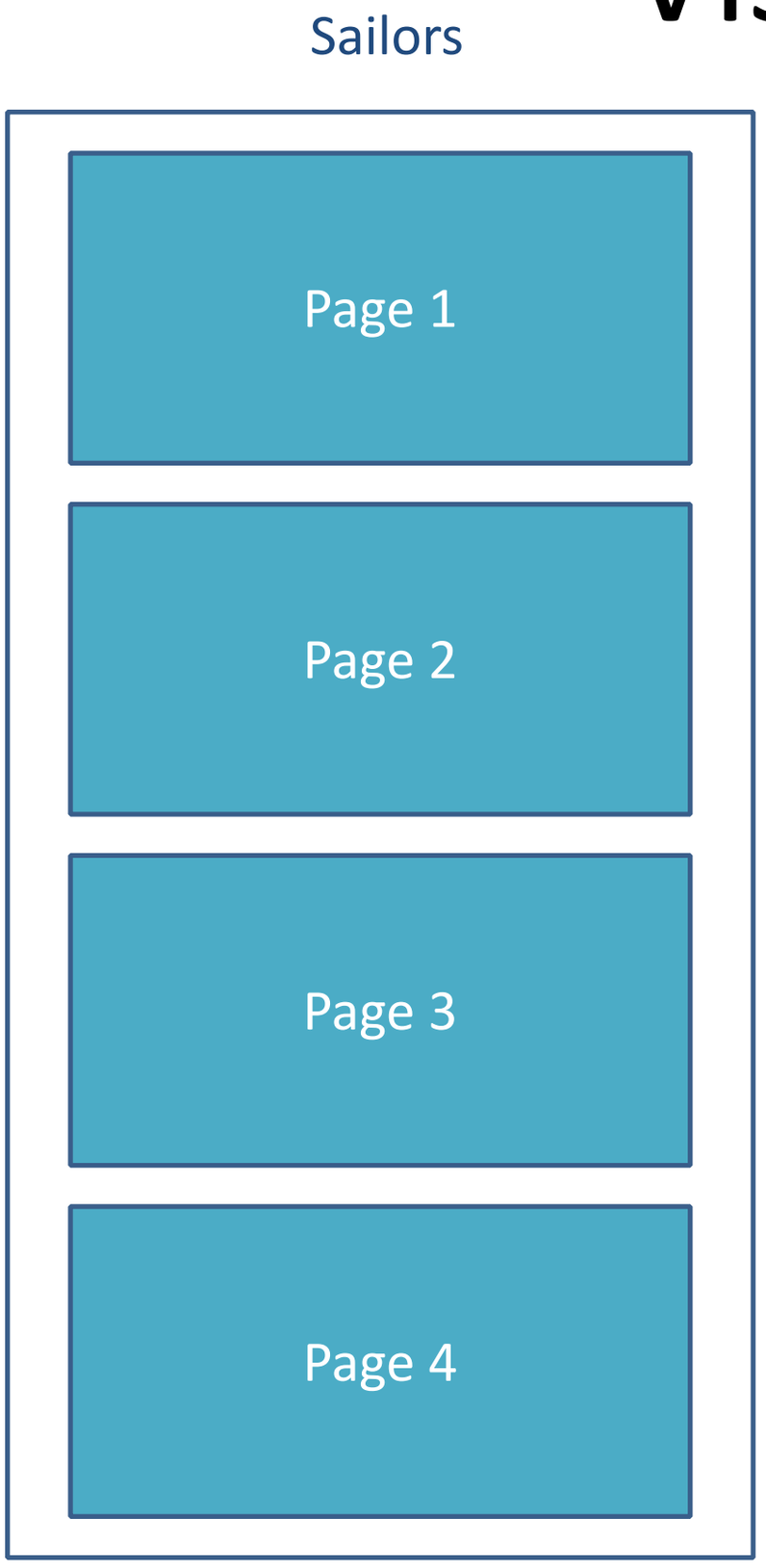

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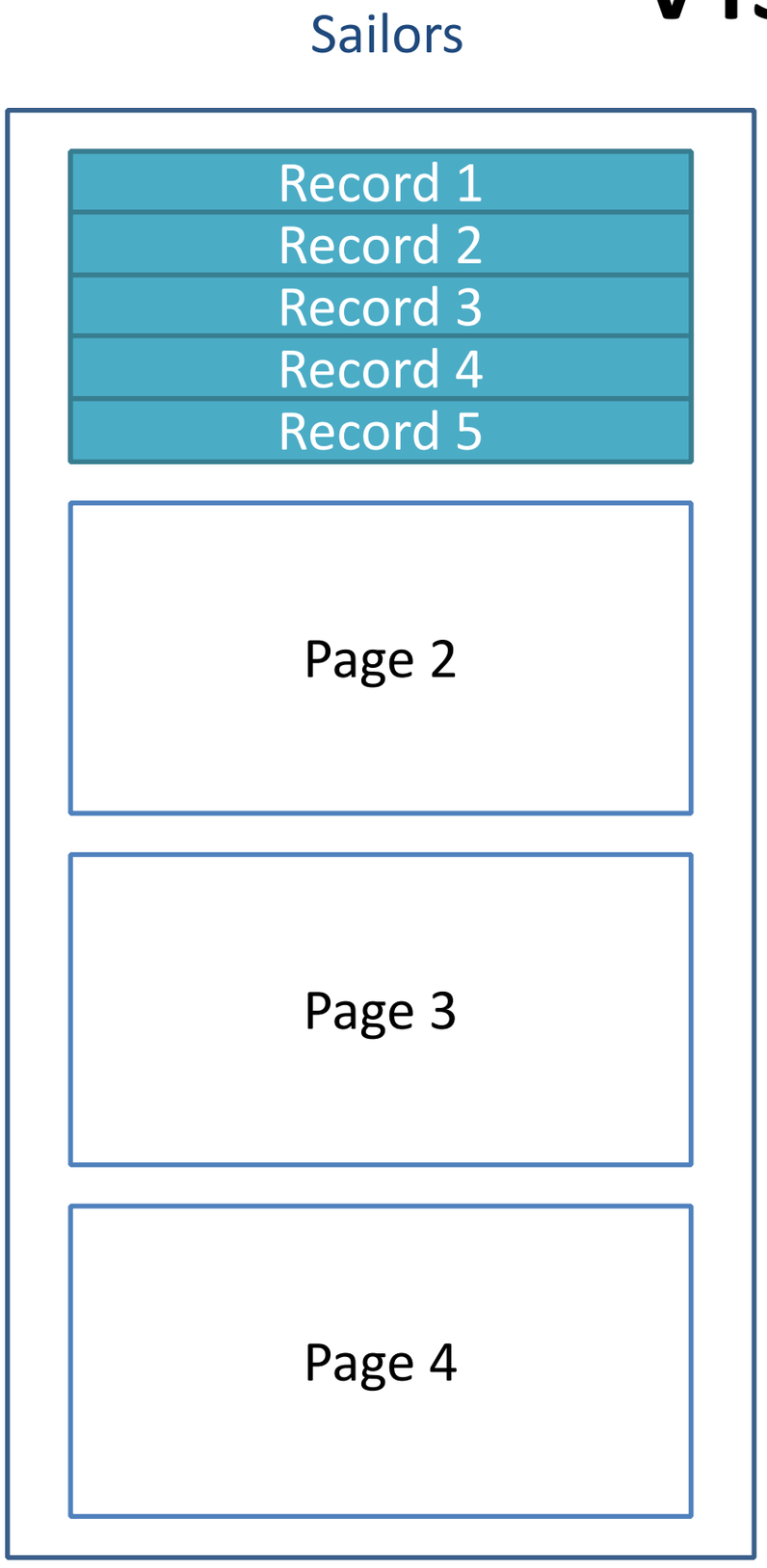

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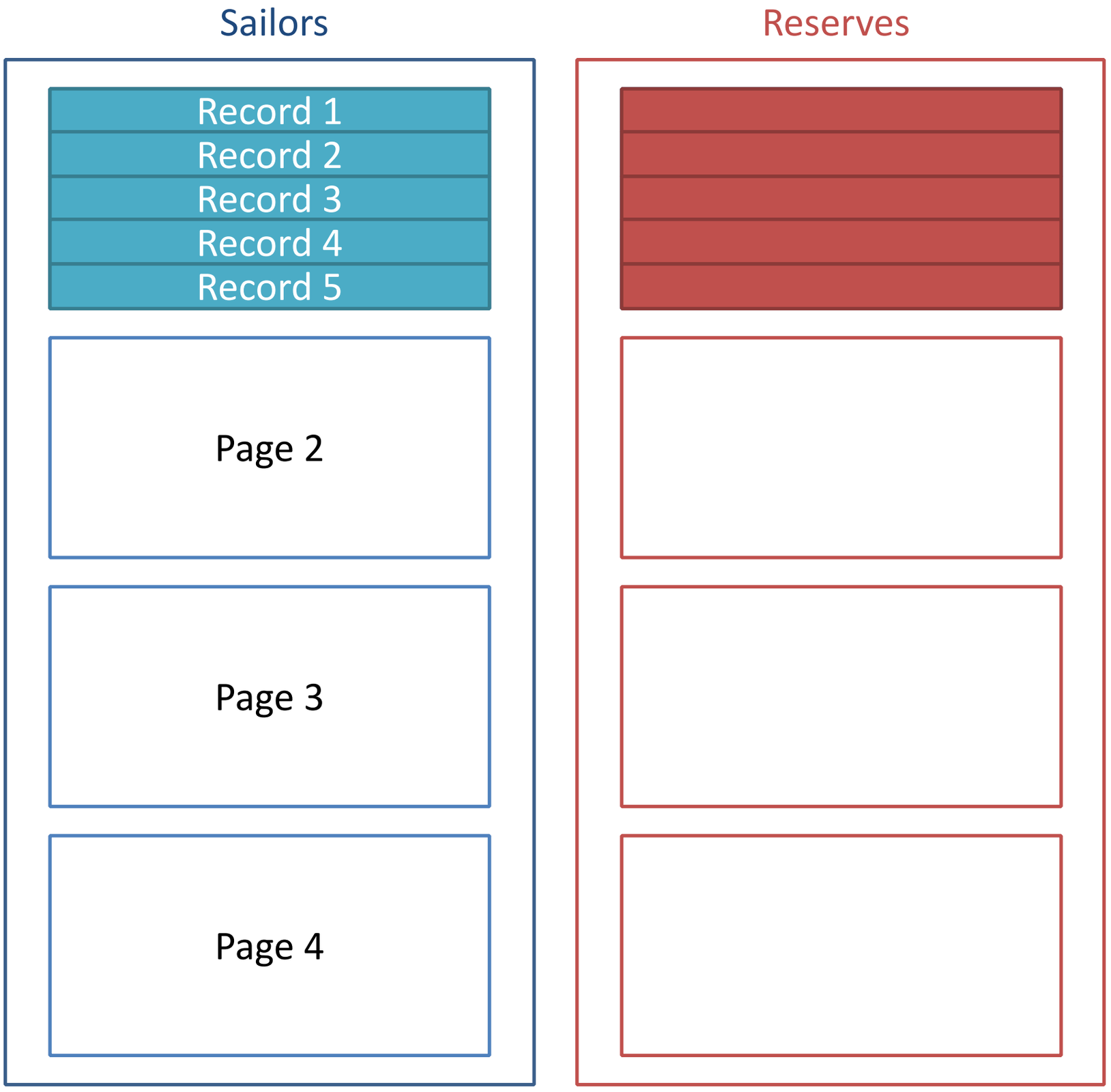
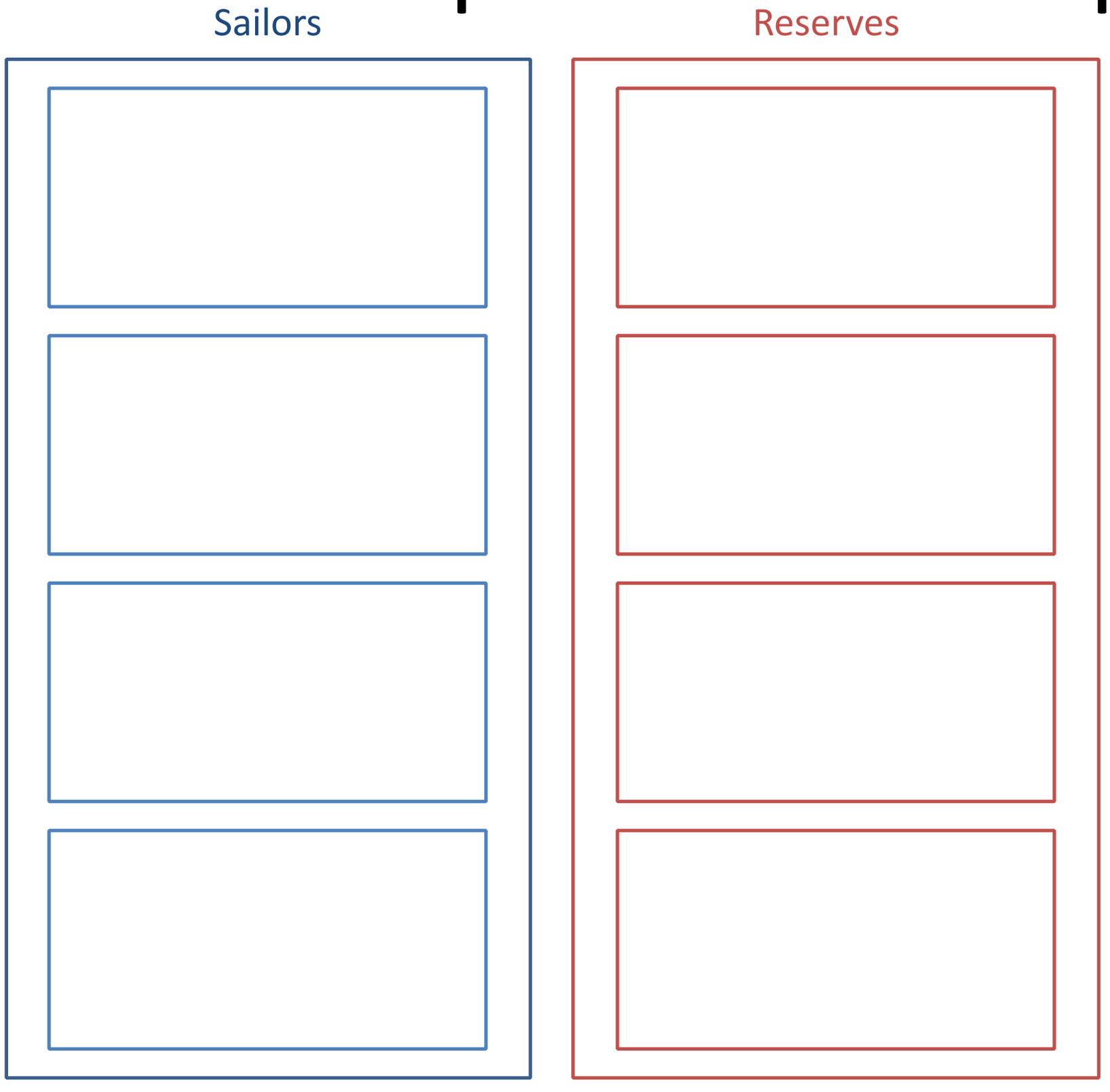


Simple Nested Loops Join



Key idea:

Take each record of S and match it with each record of R.

Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

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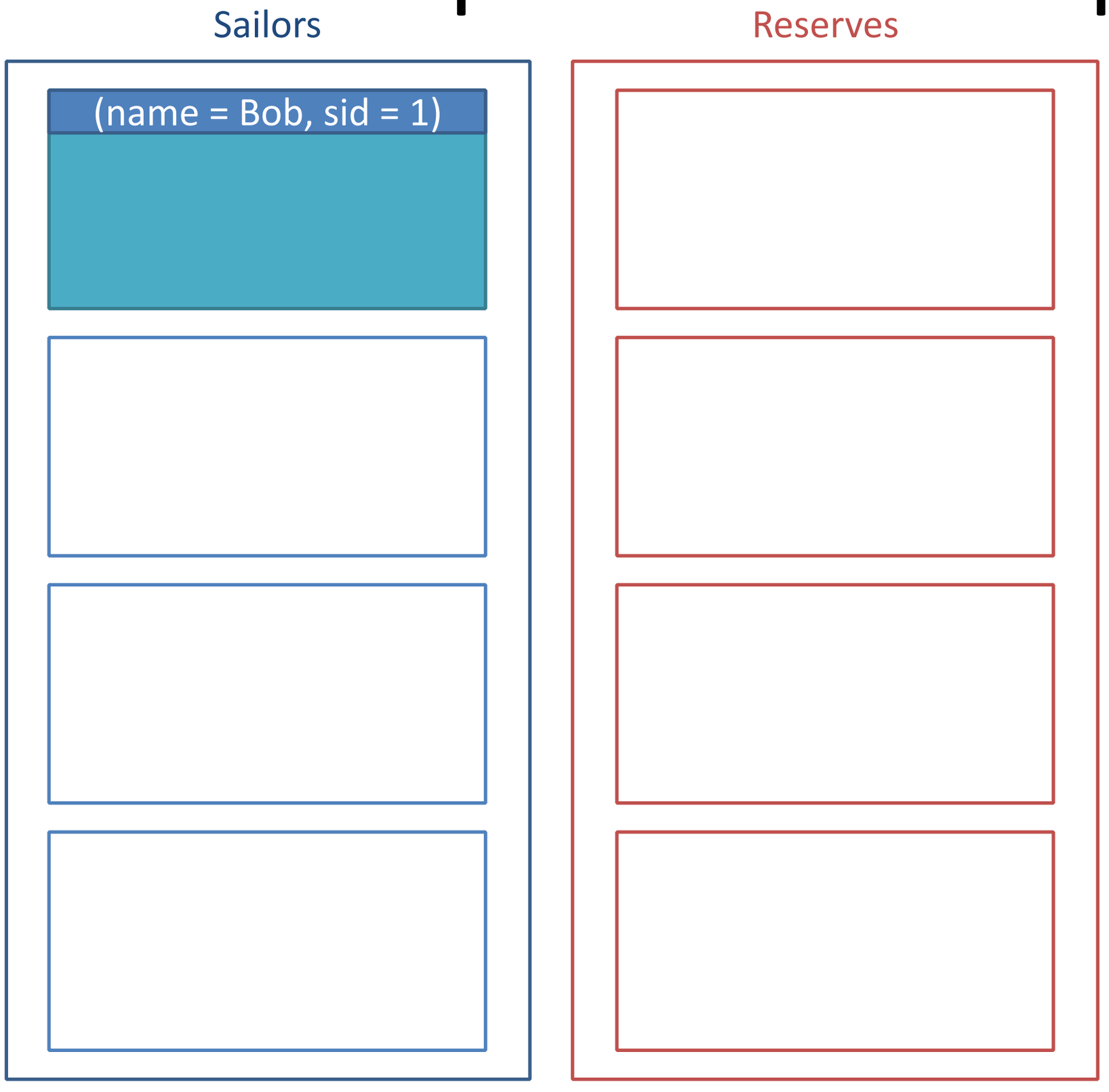

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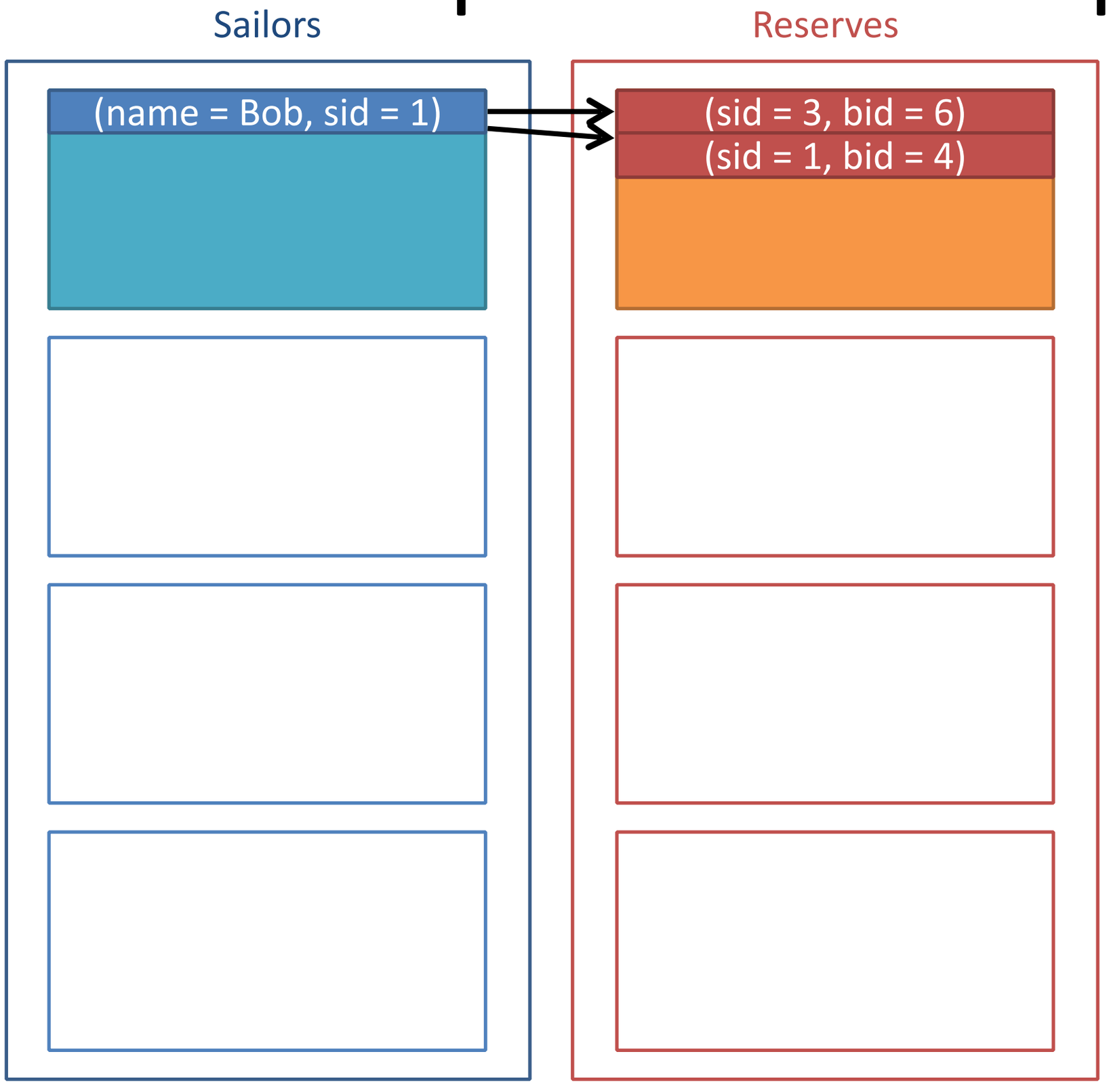

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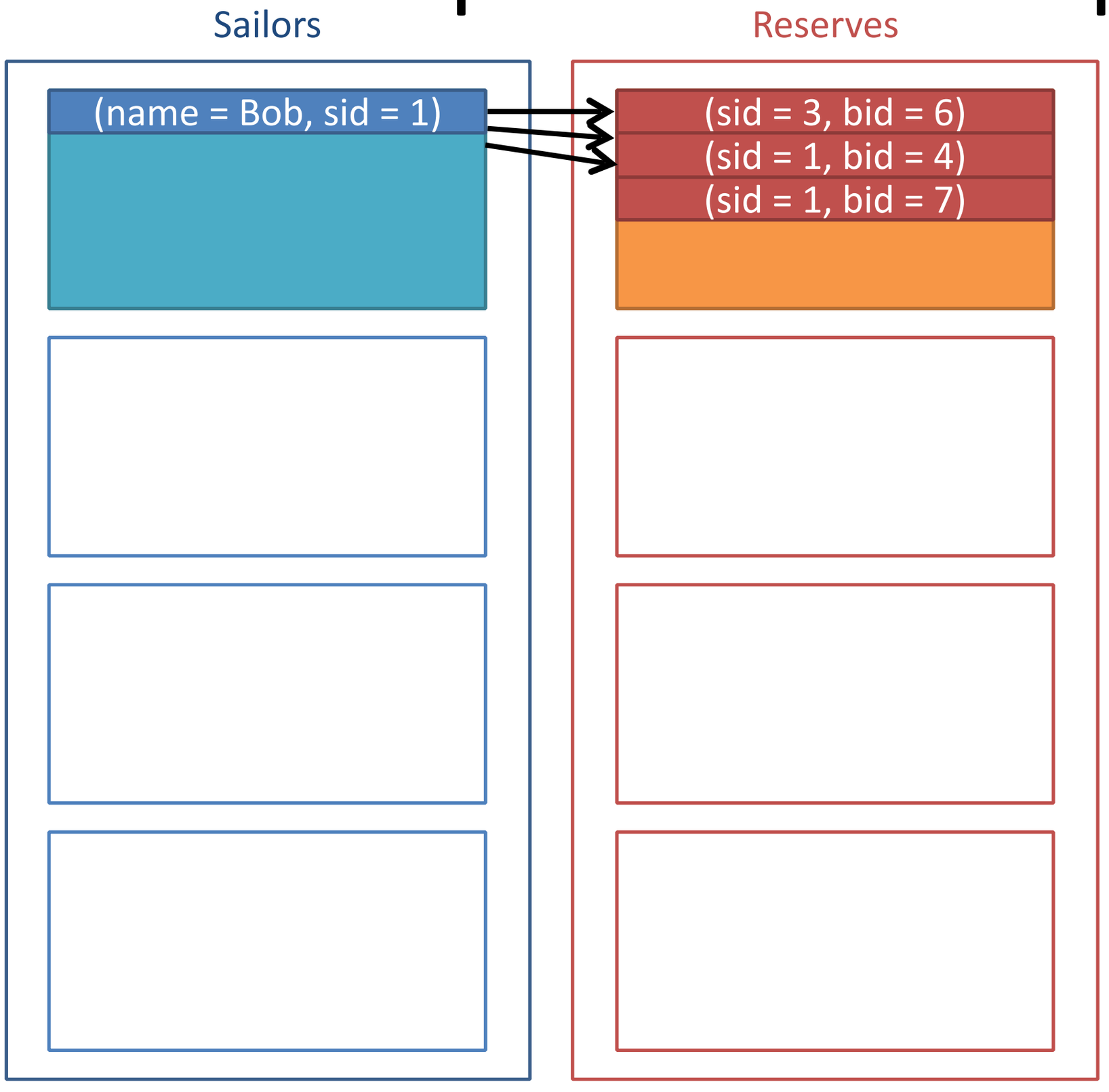

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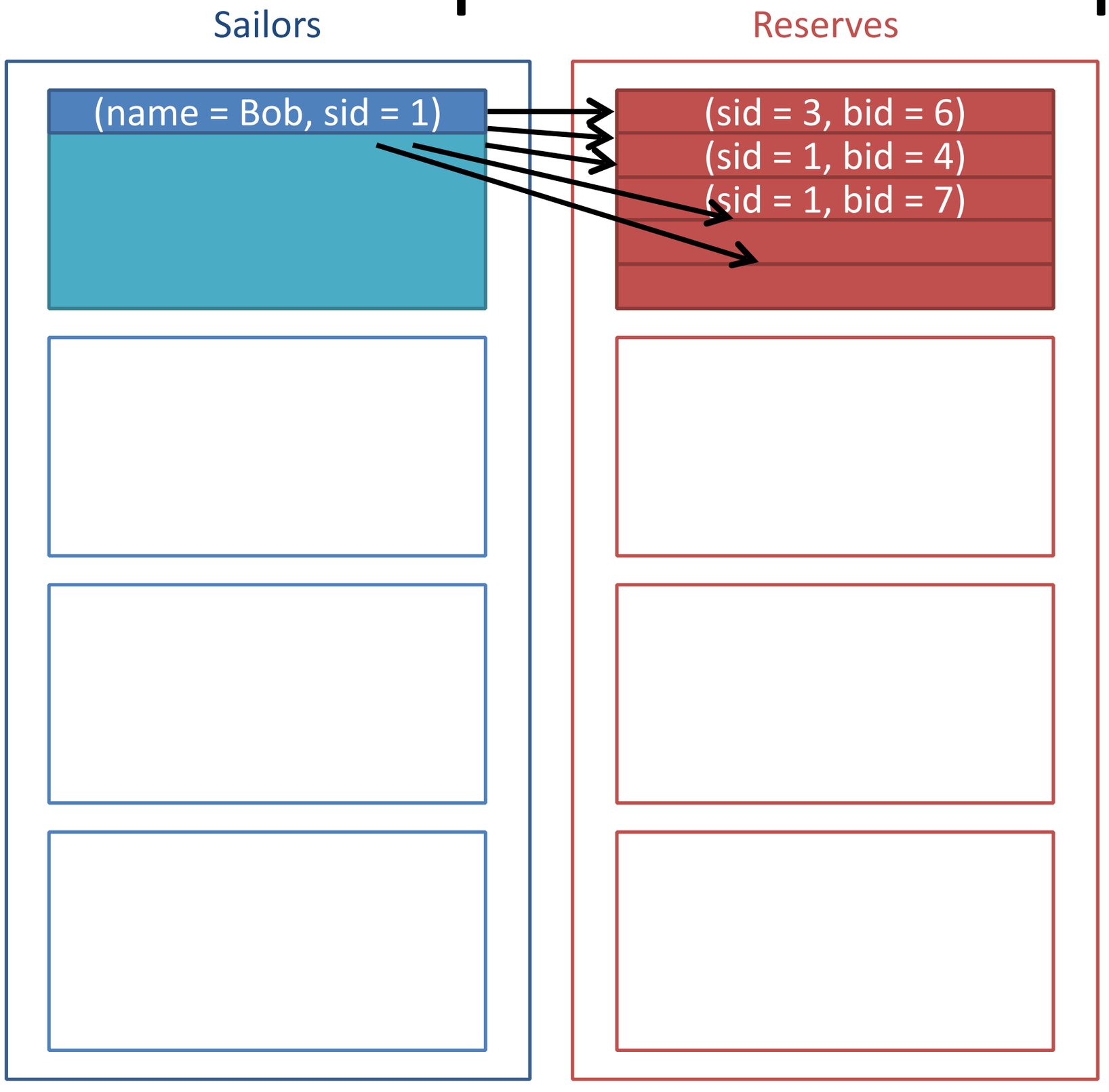

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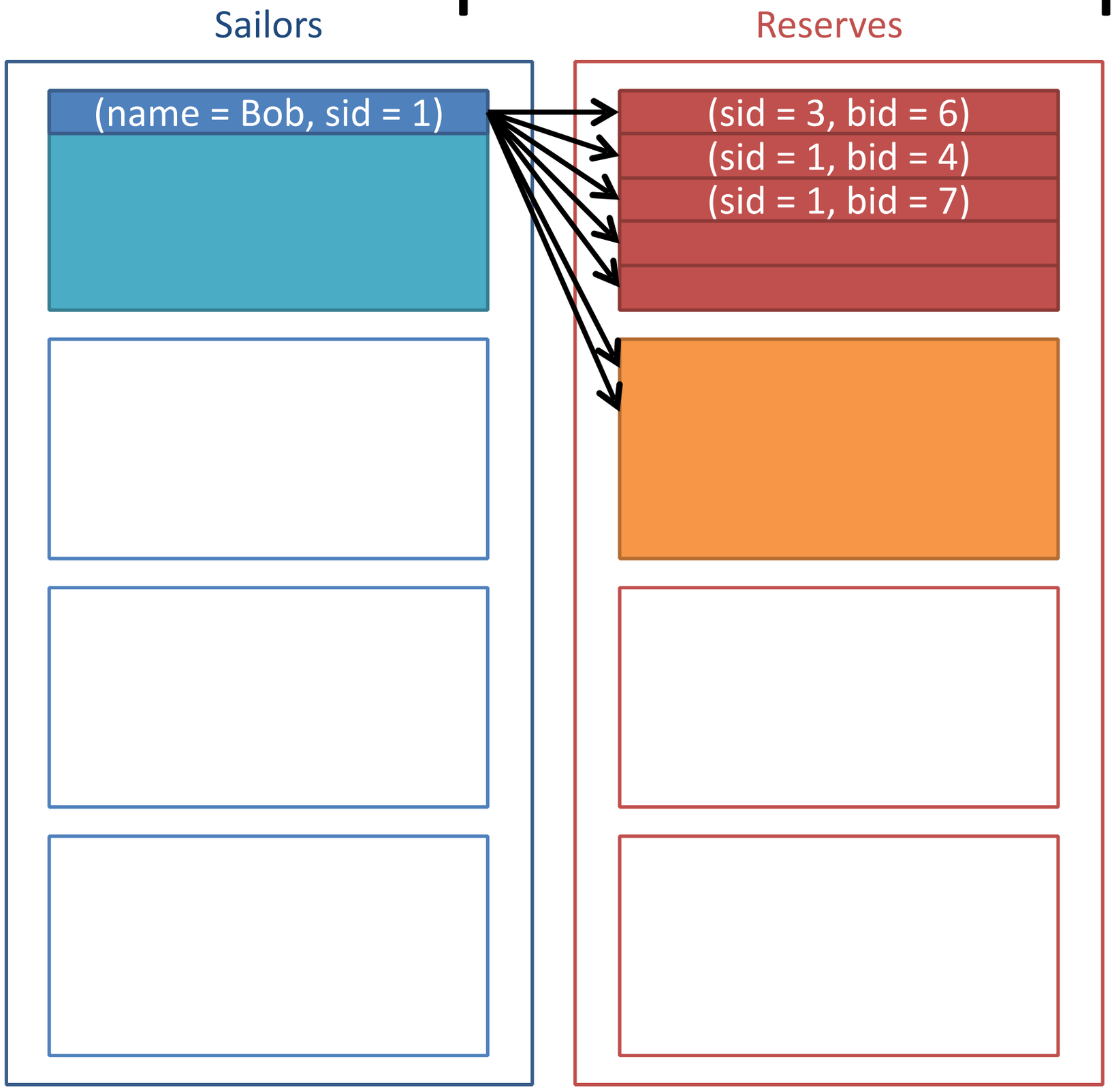

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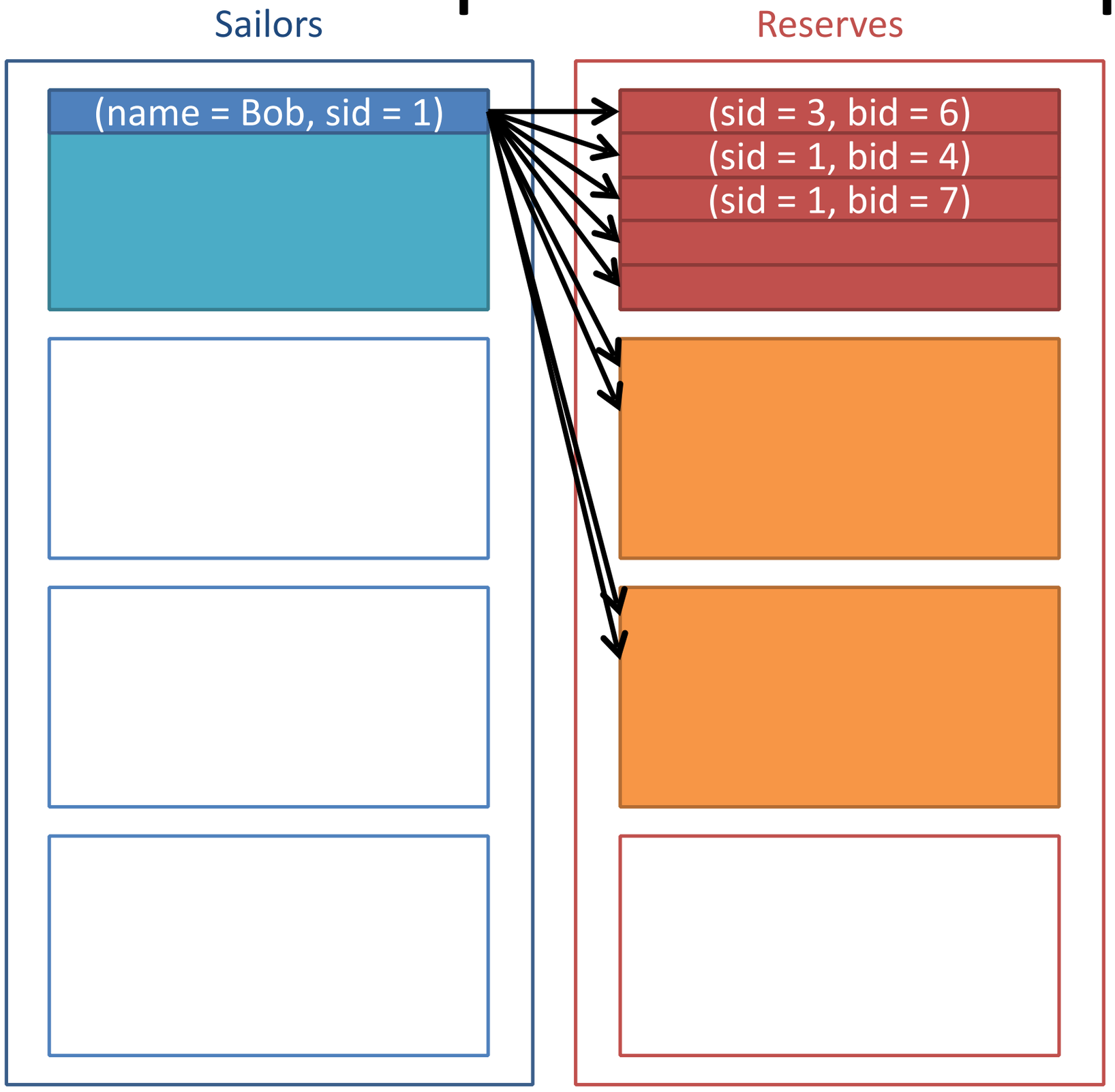

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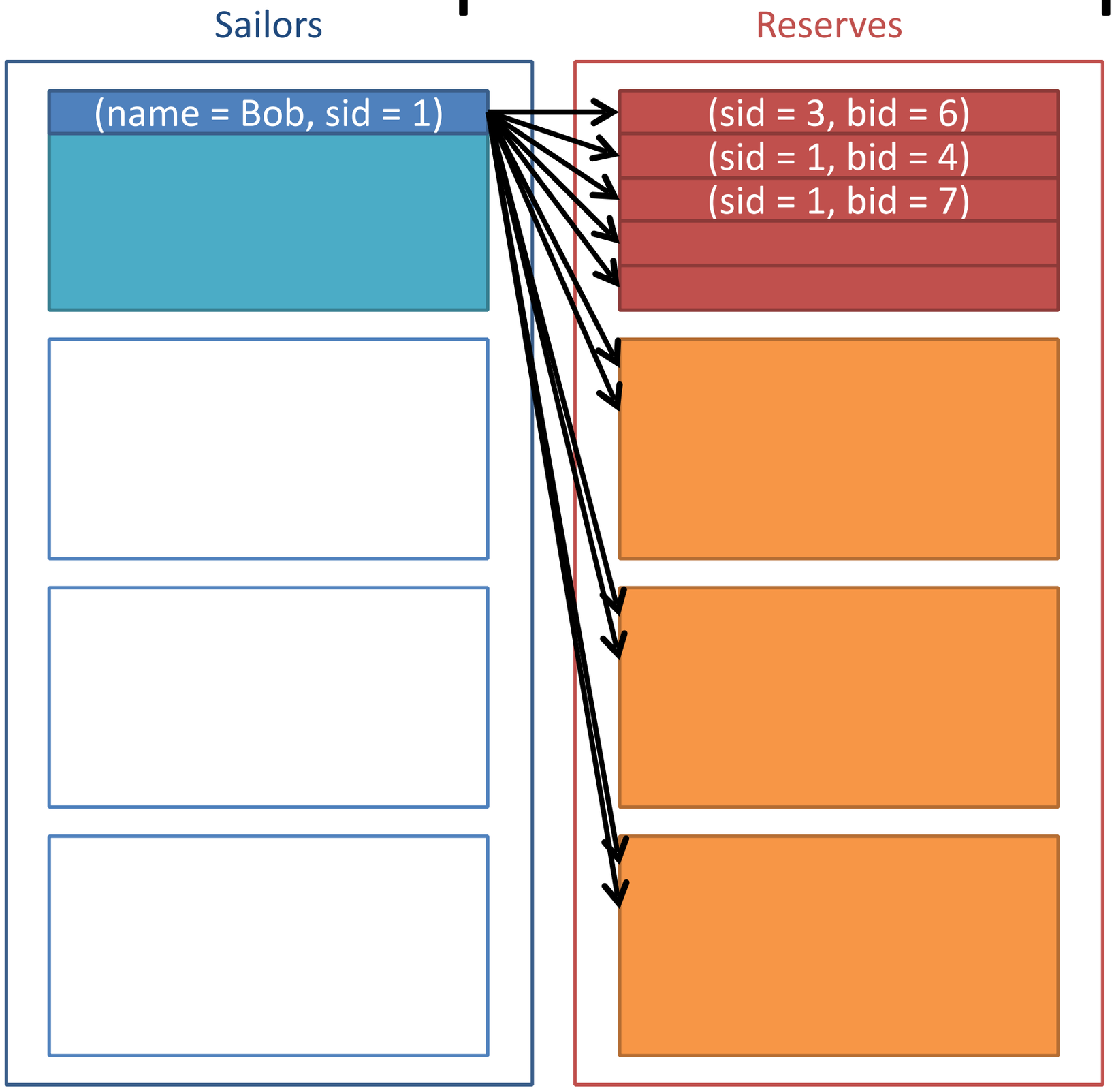

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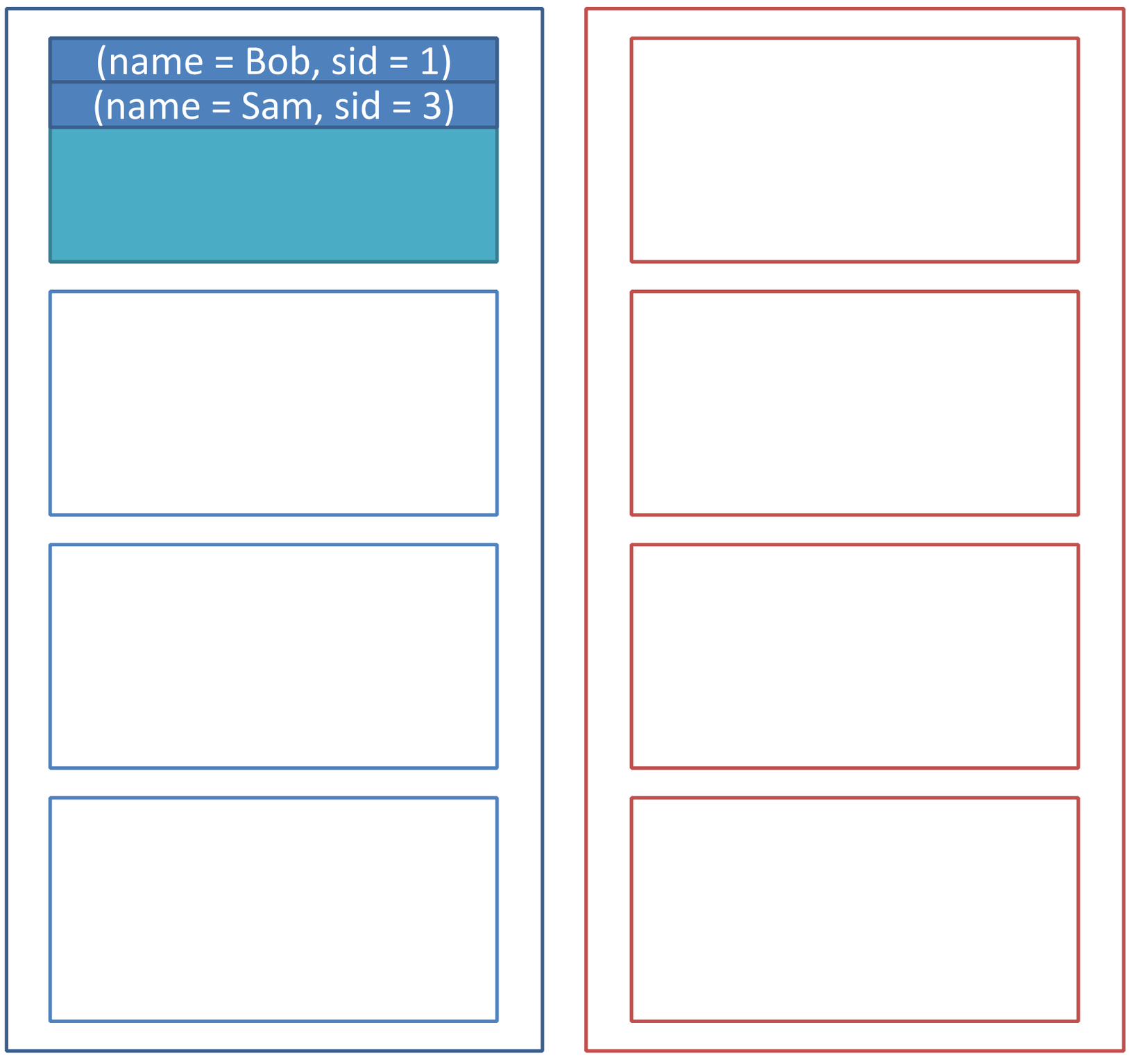

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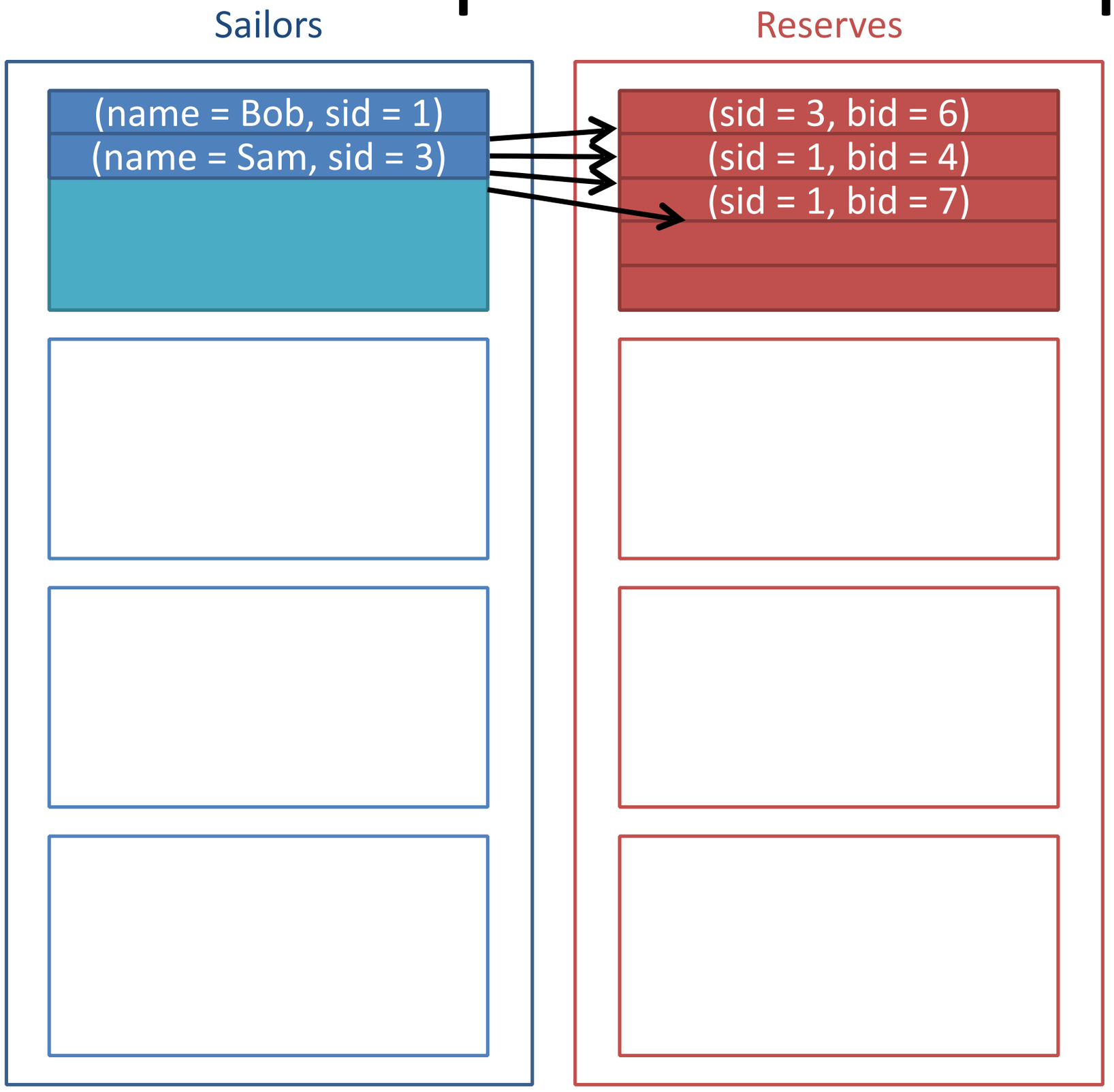

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Notation: [S] == "# pages in S" ; |S| == "# tuples in S"

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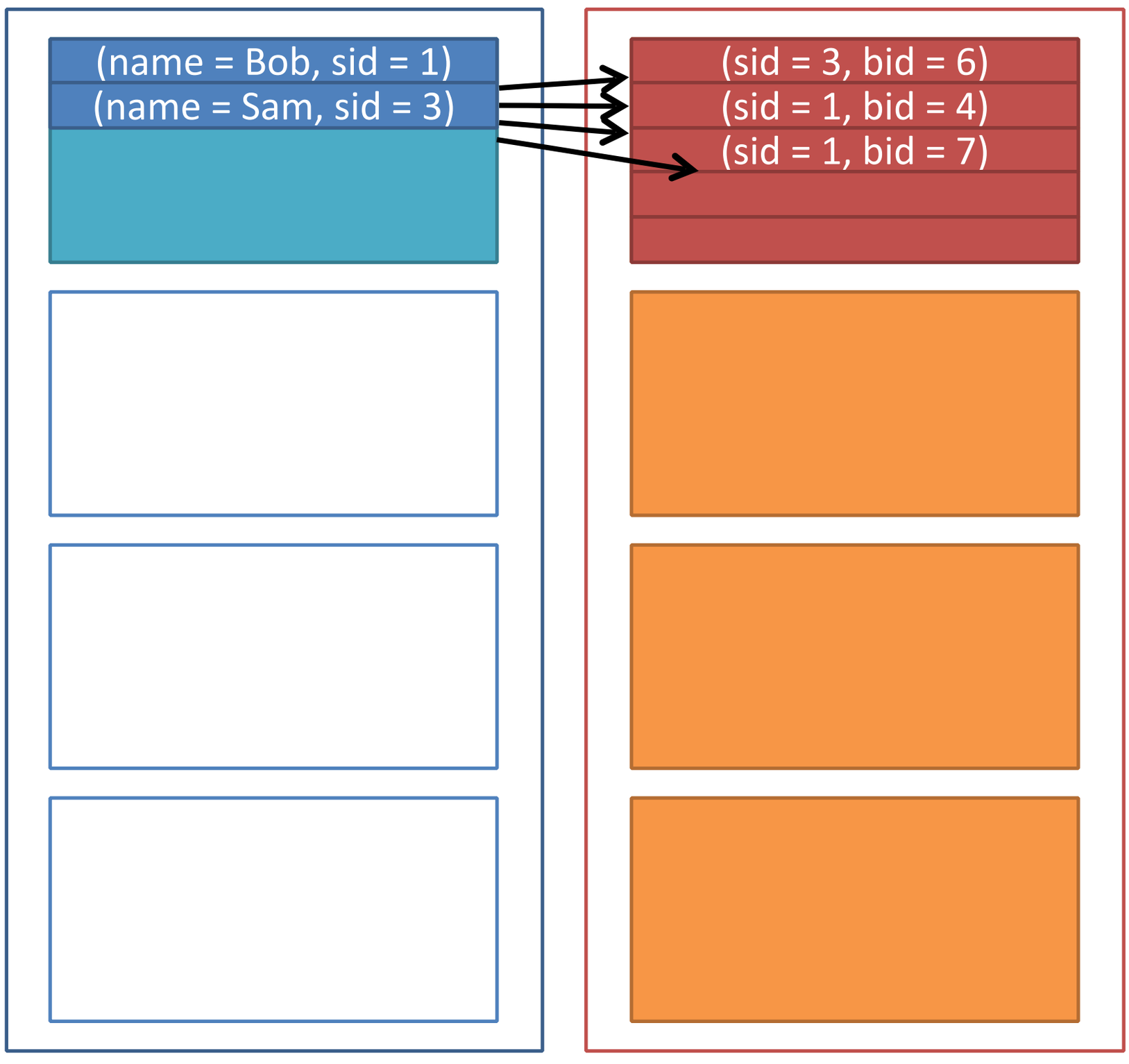

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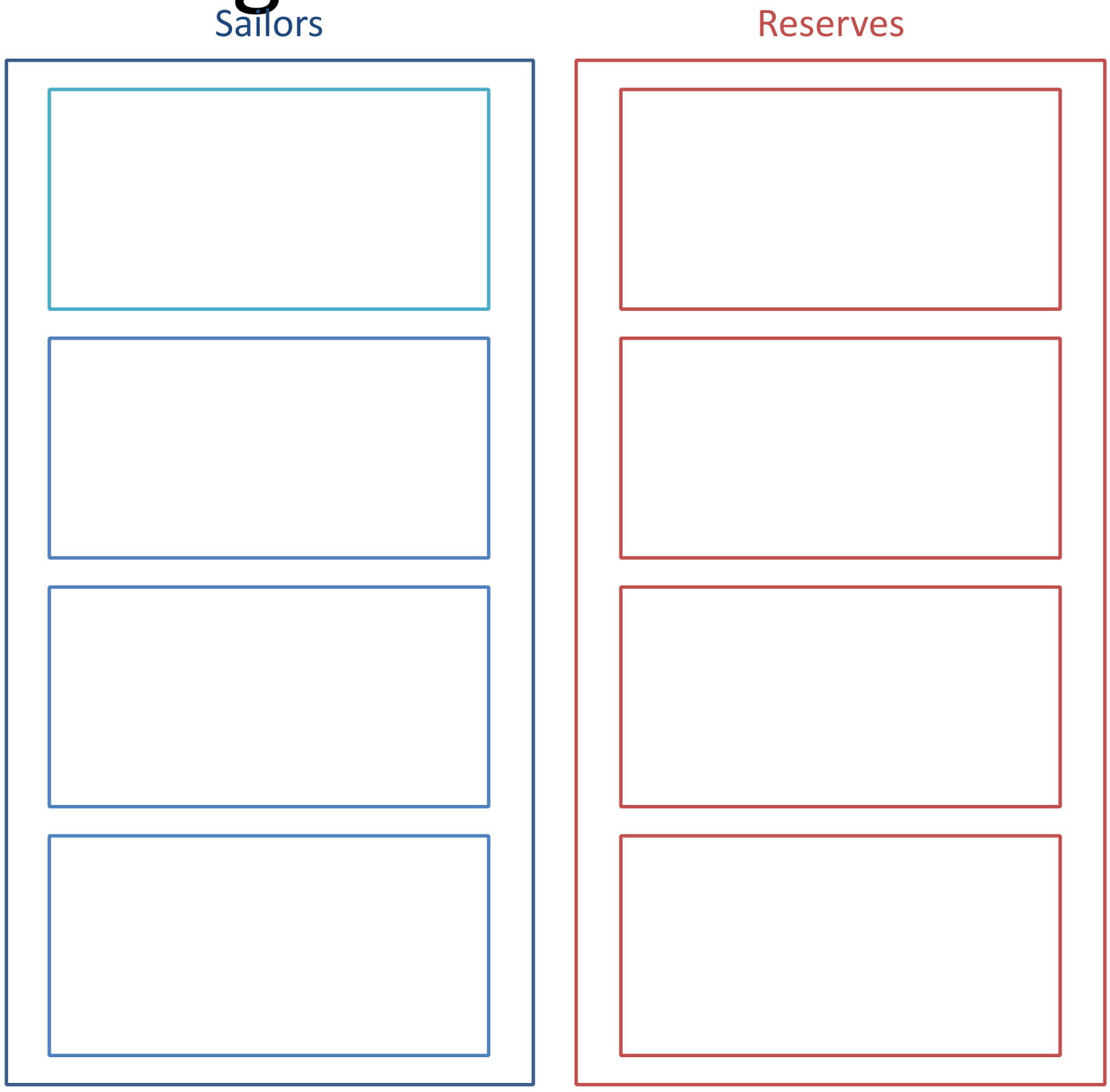
I/Os:

[S] + |S|*[R]

Page-Oriented Nested Loops Join



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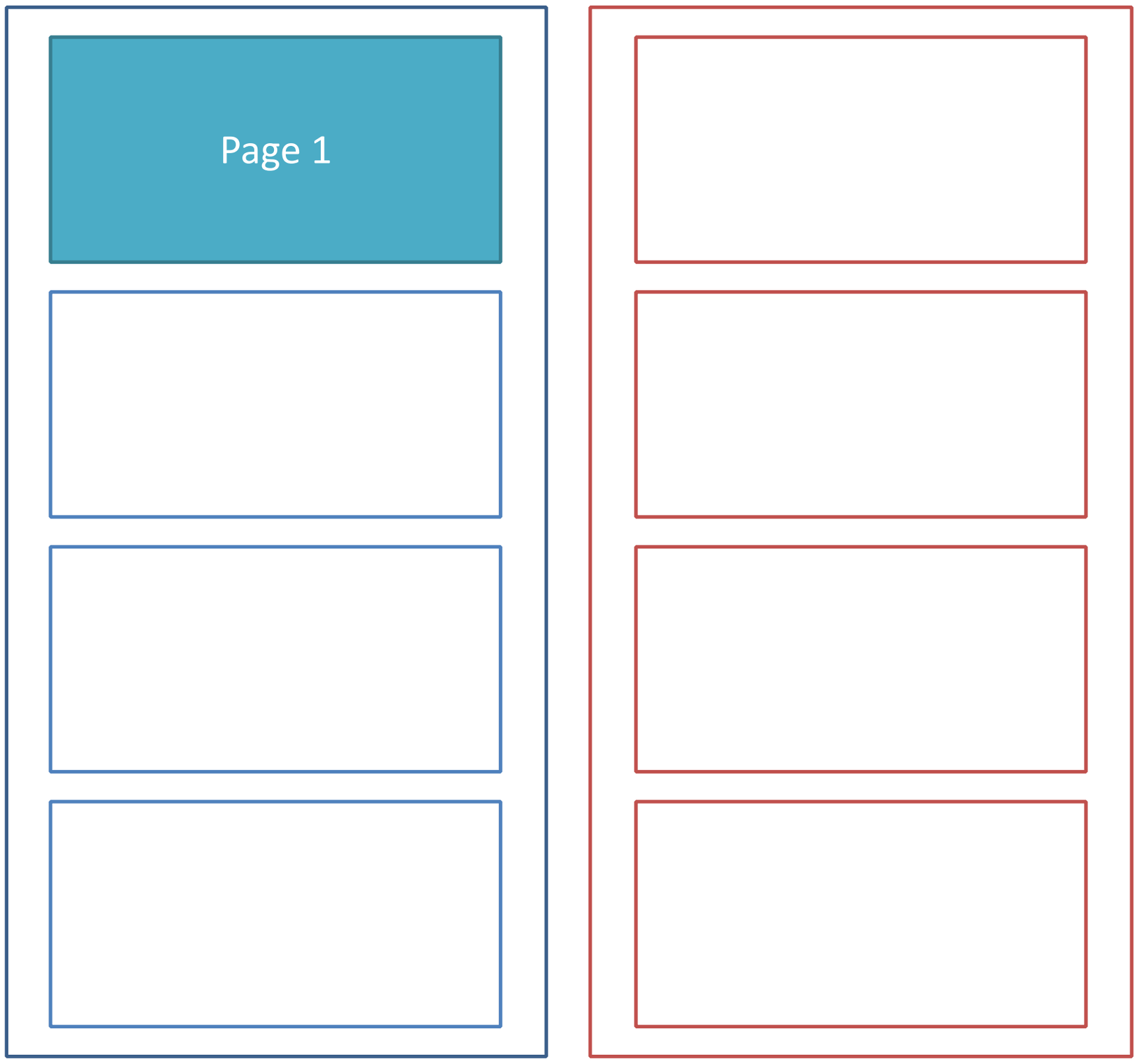

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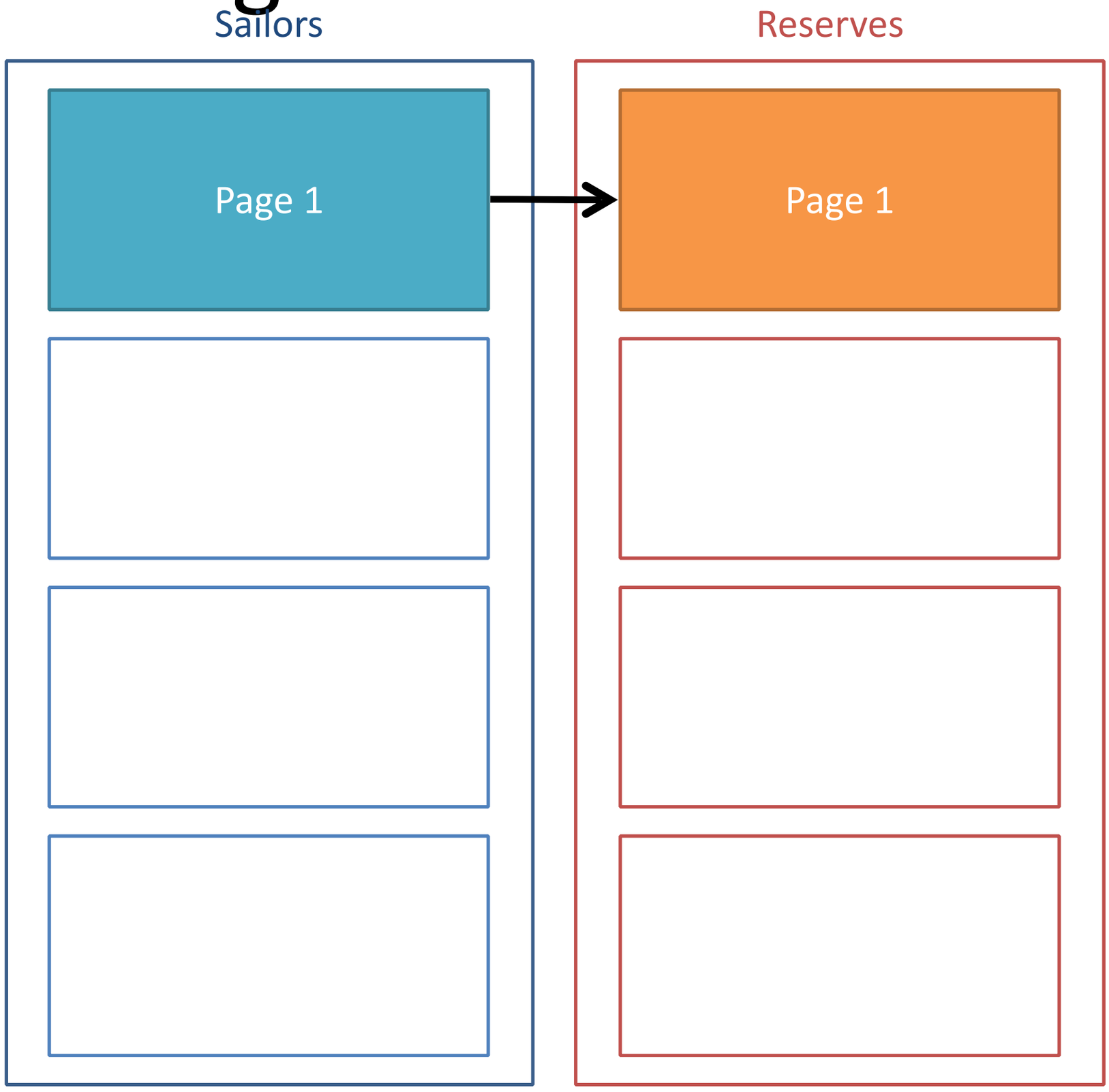

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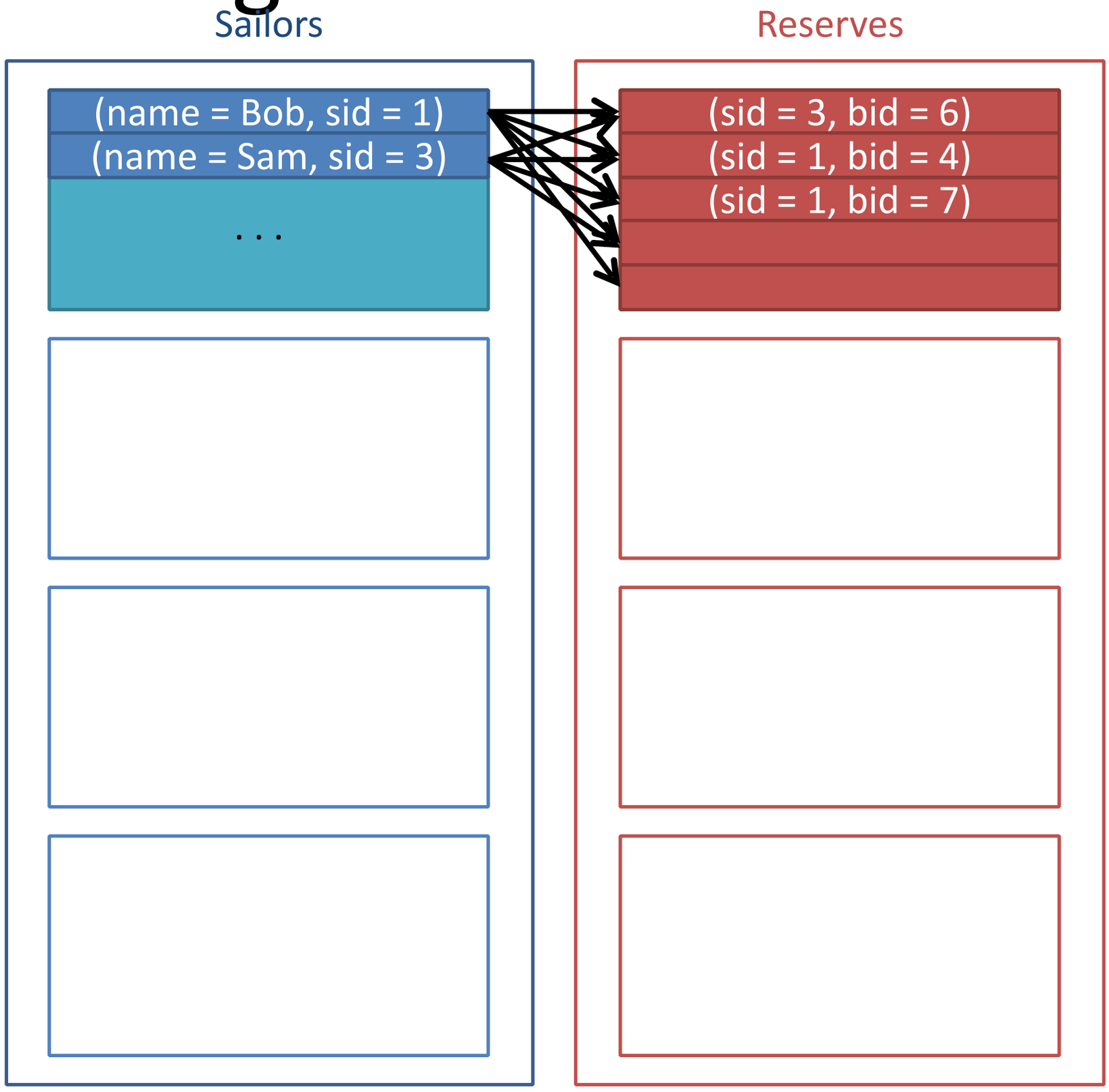

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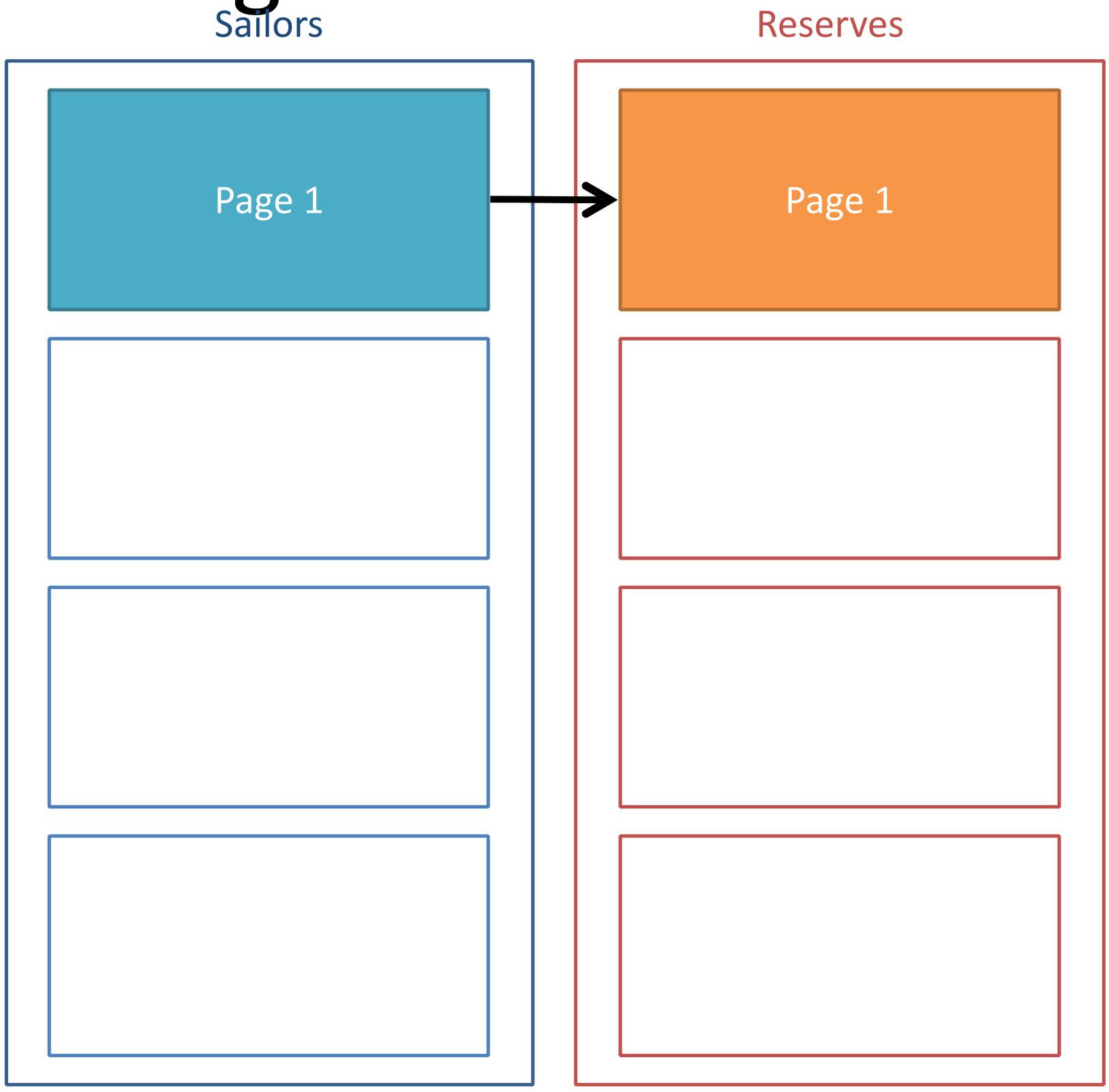

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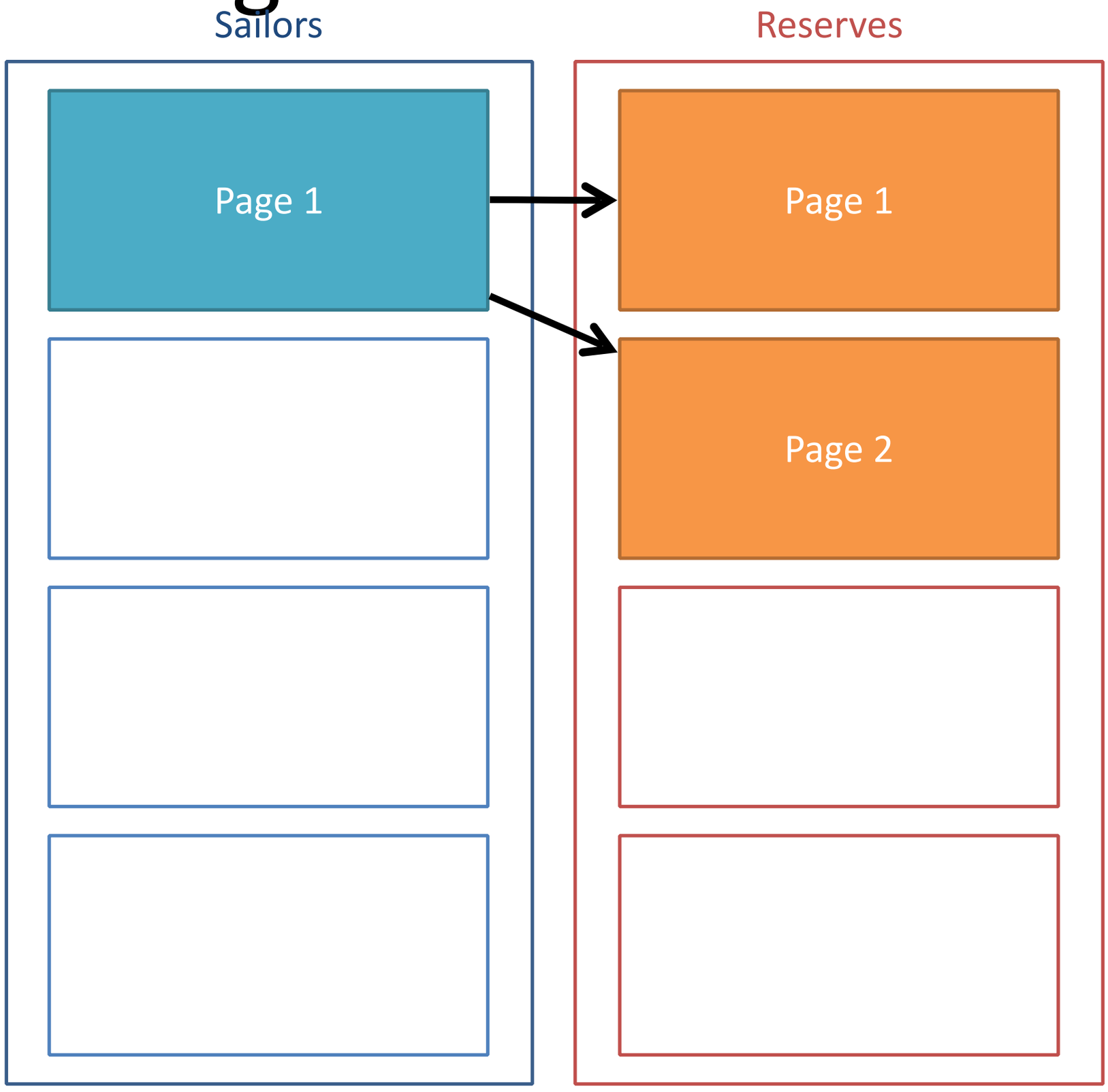

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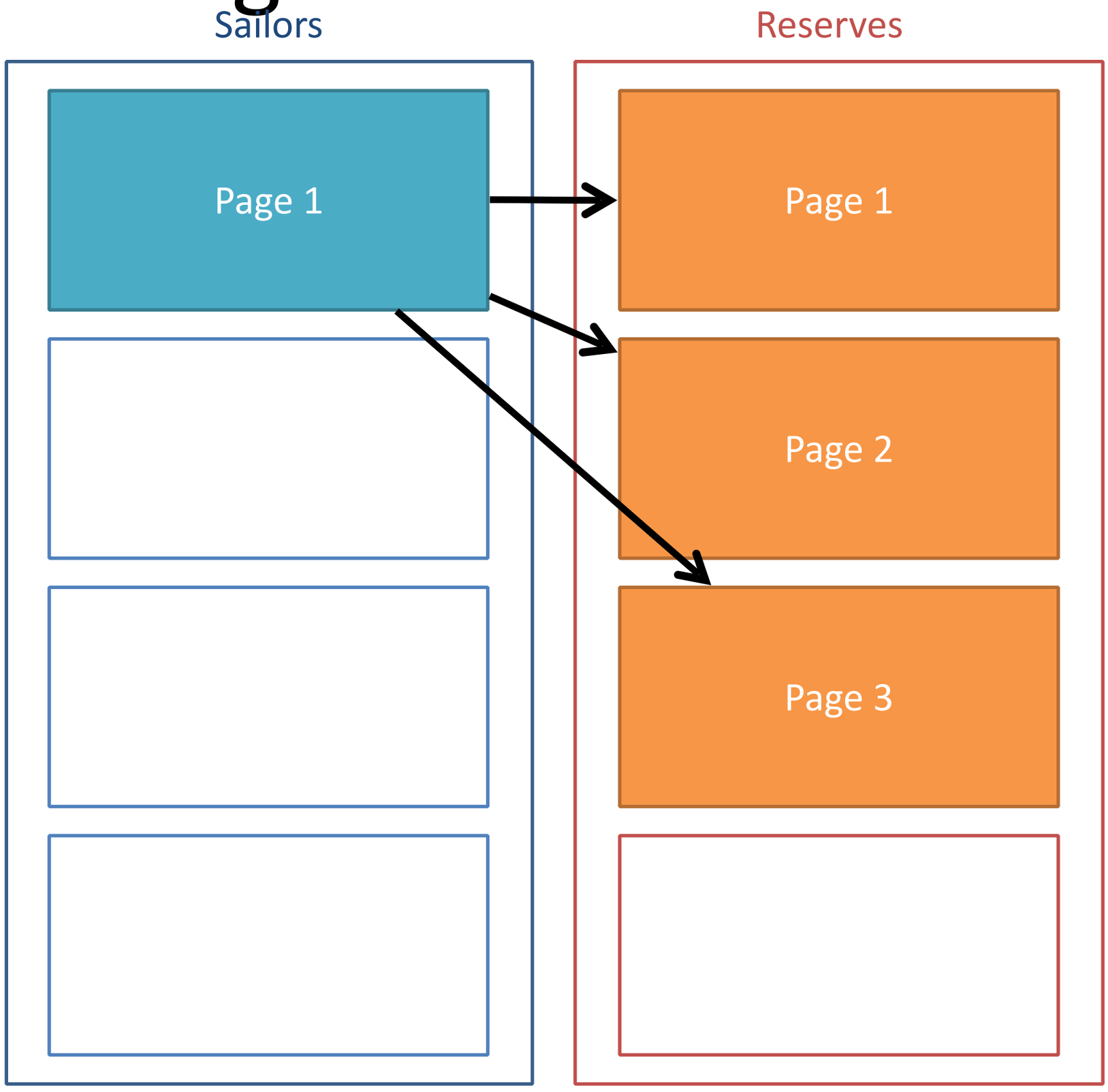

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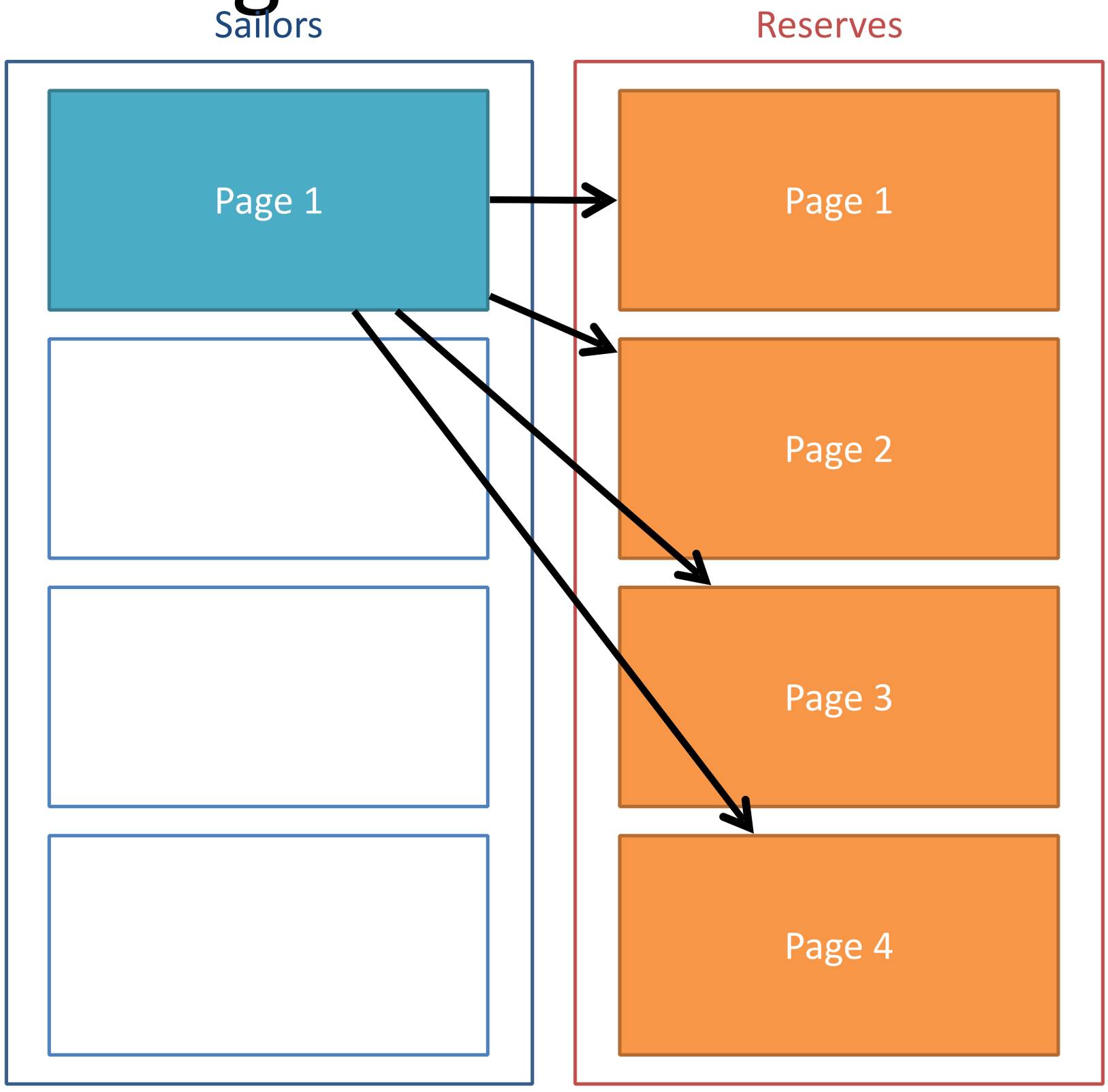

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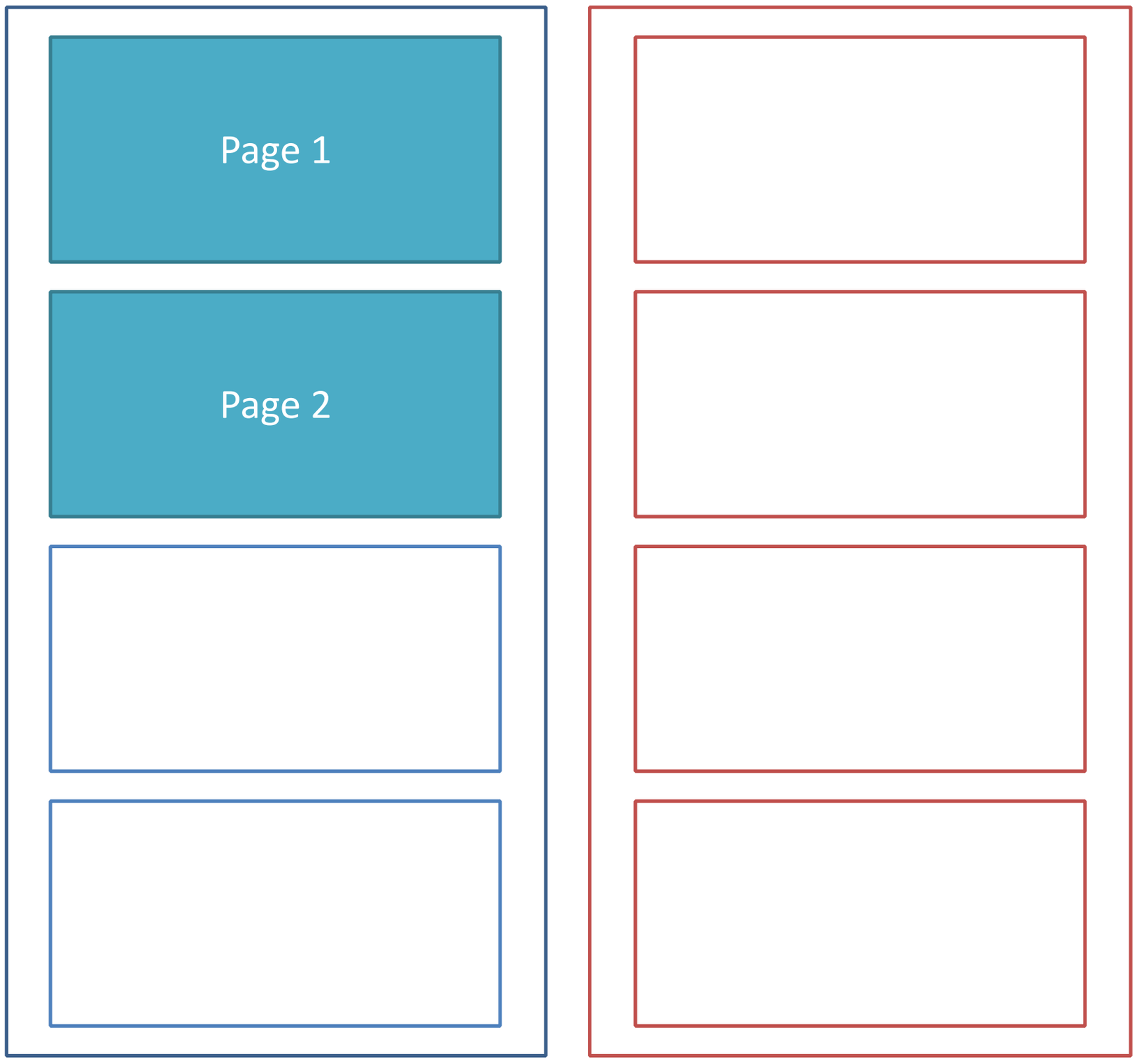

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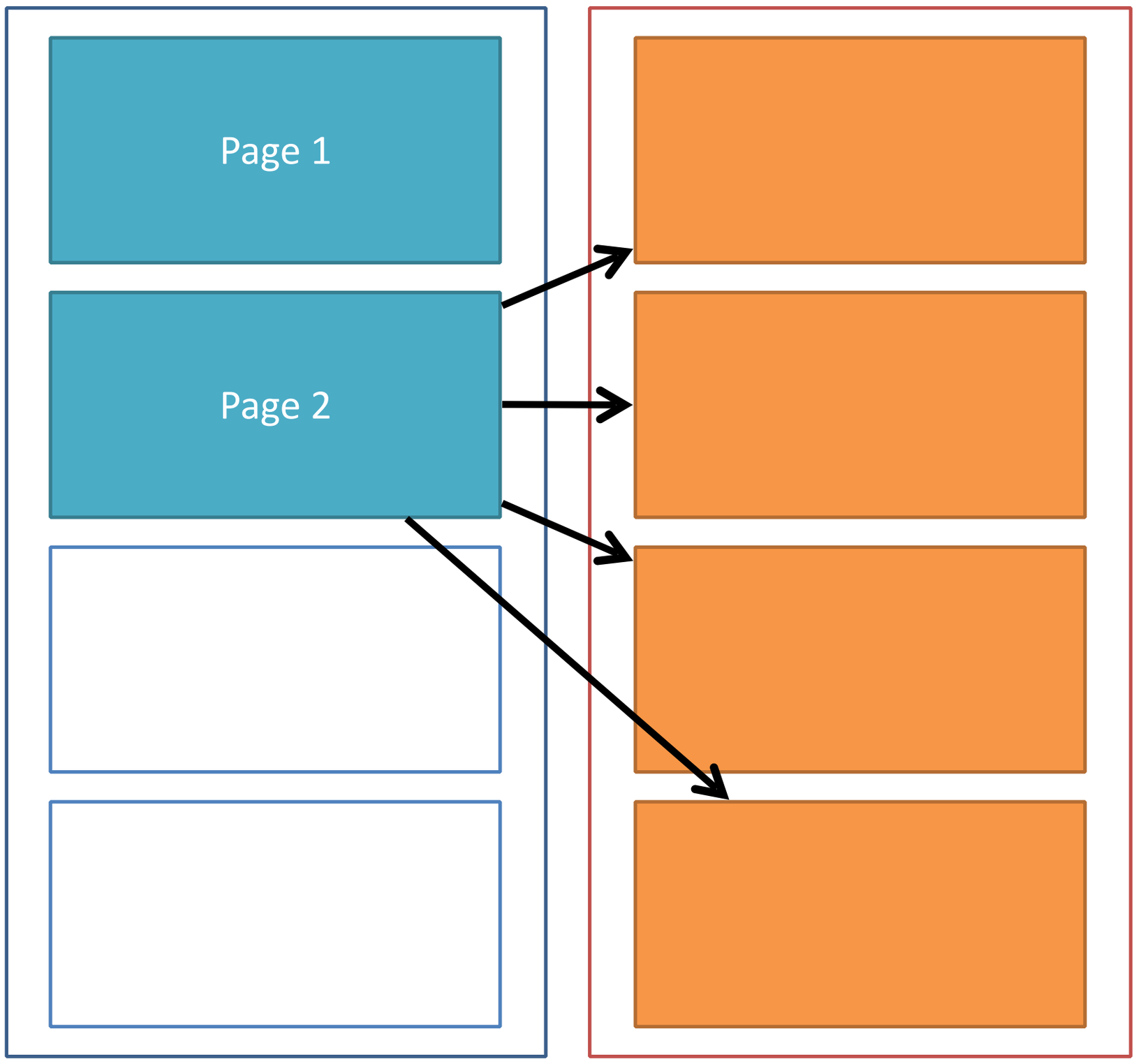

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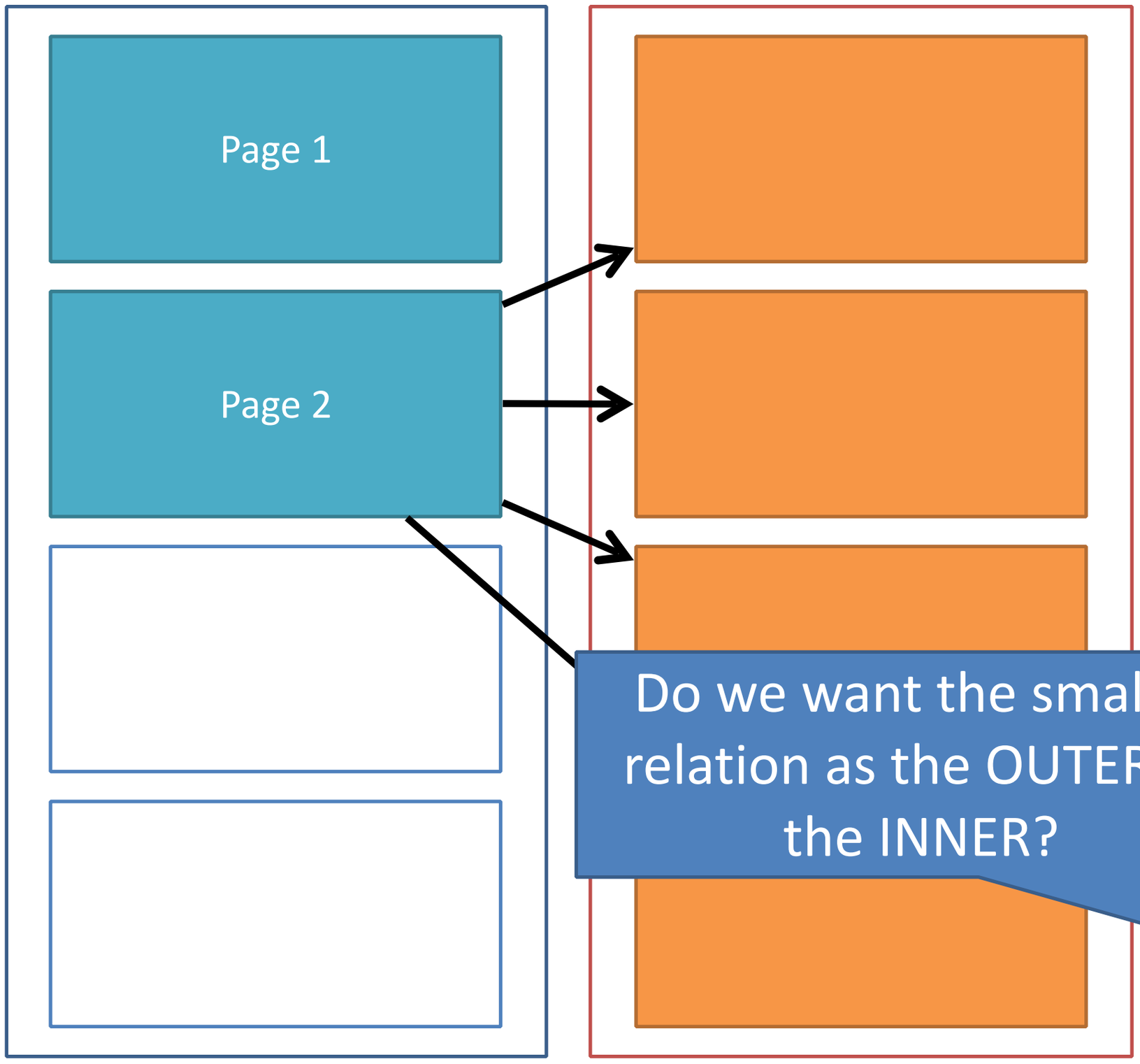

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:

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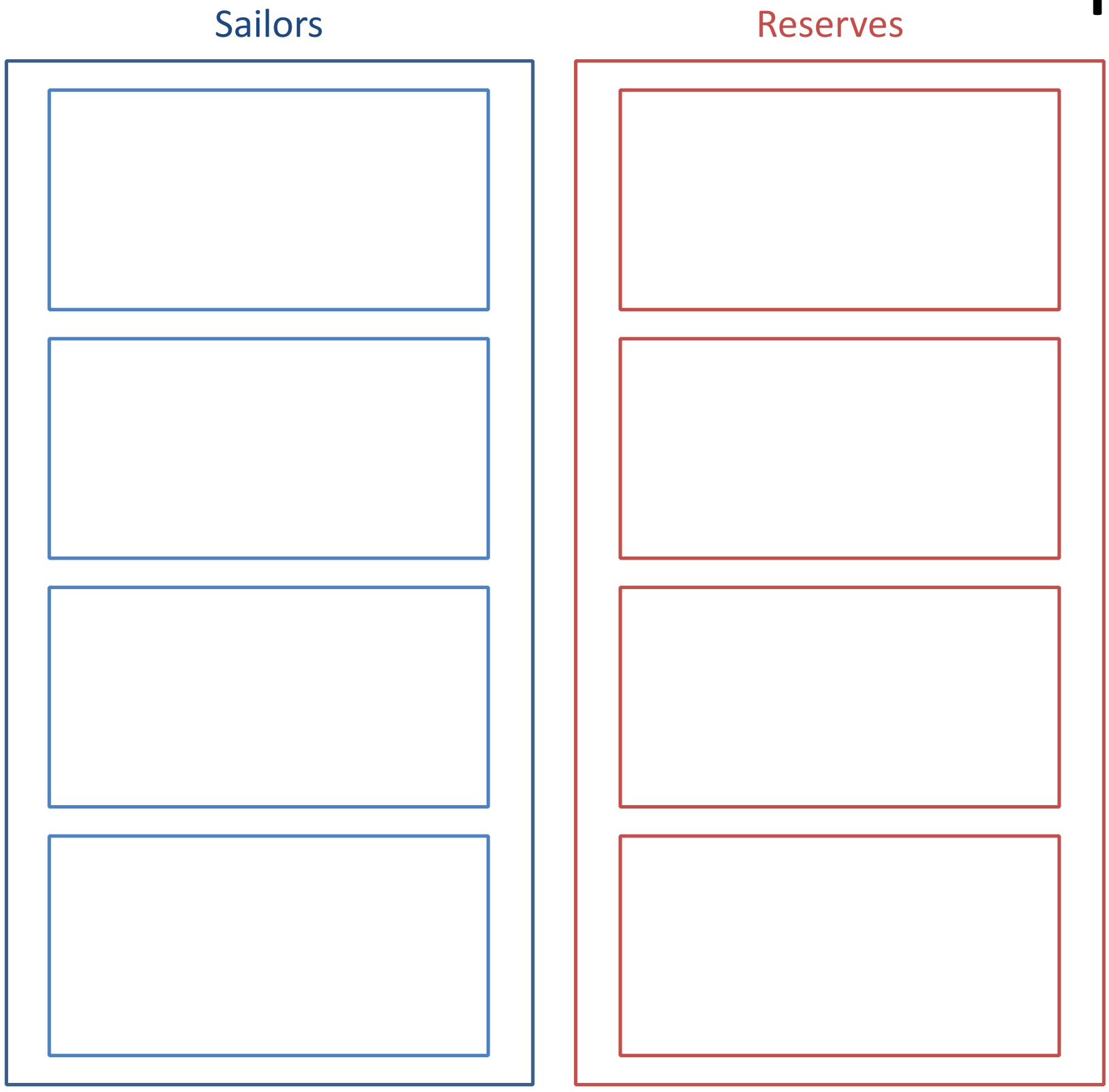

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Take **k pages** of S and match with each page of R.

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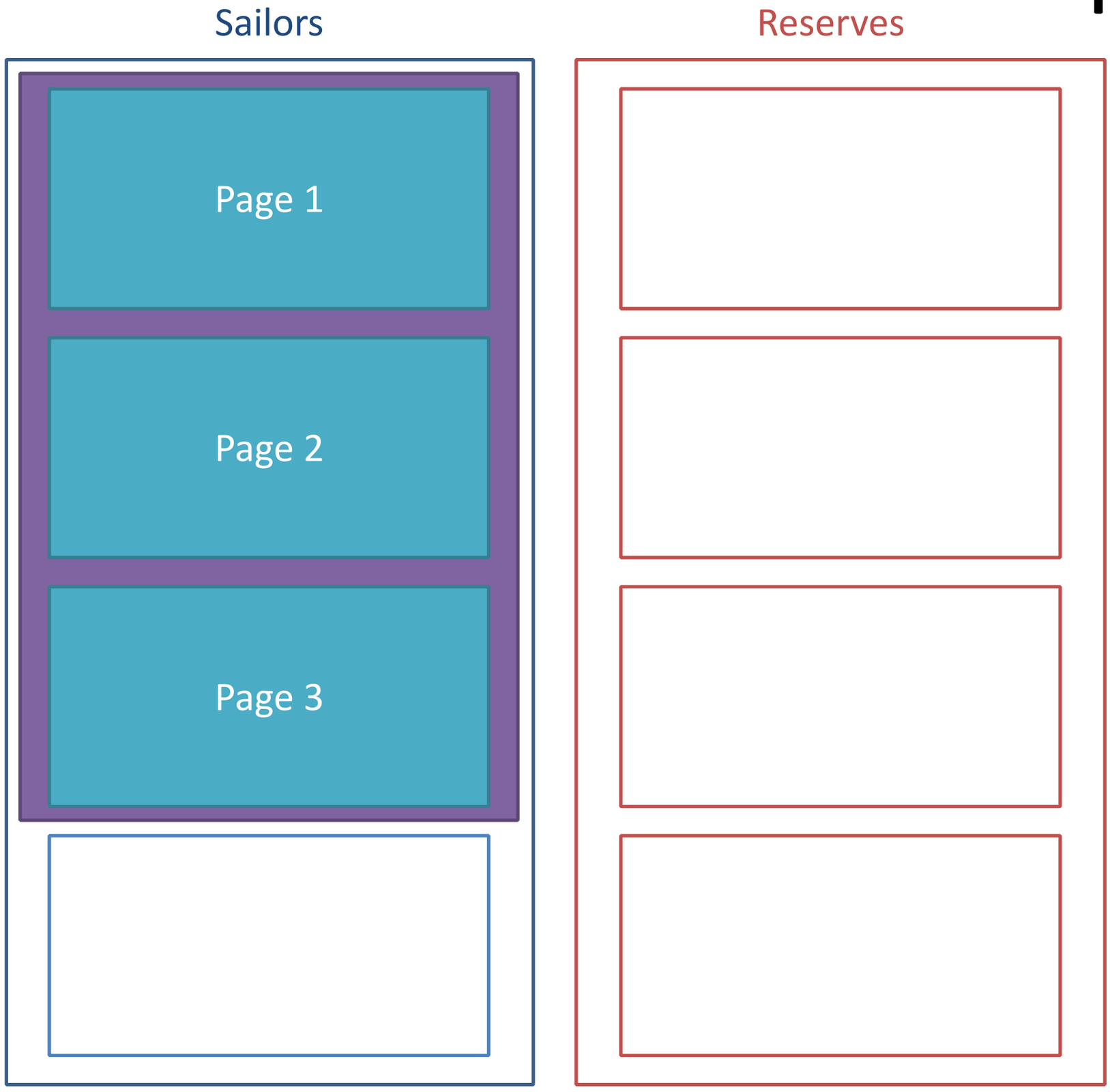

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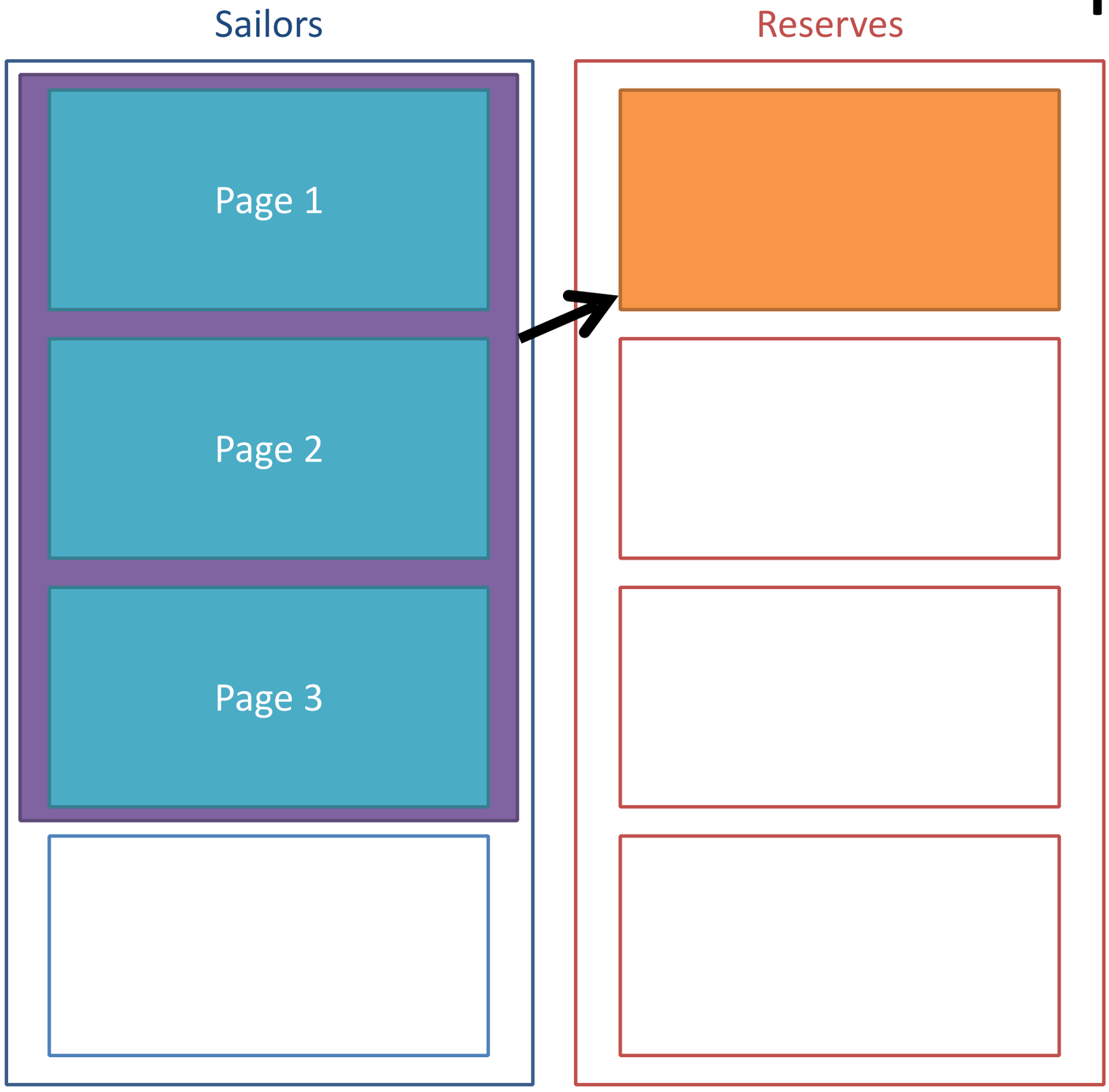

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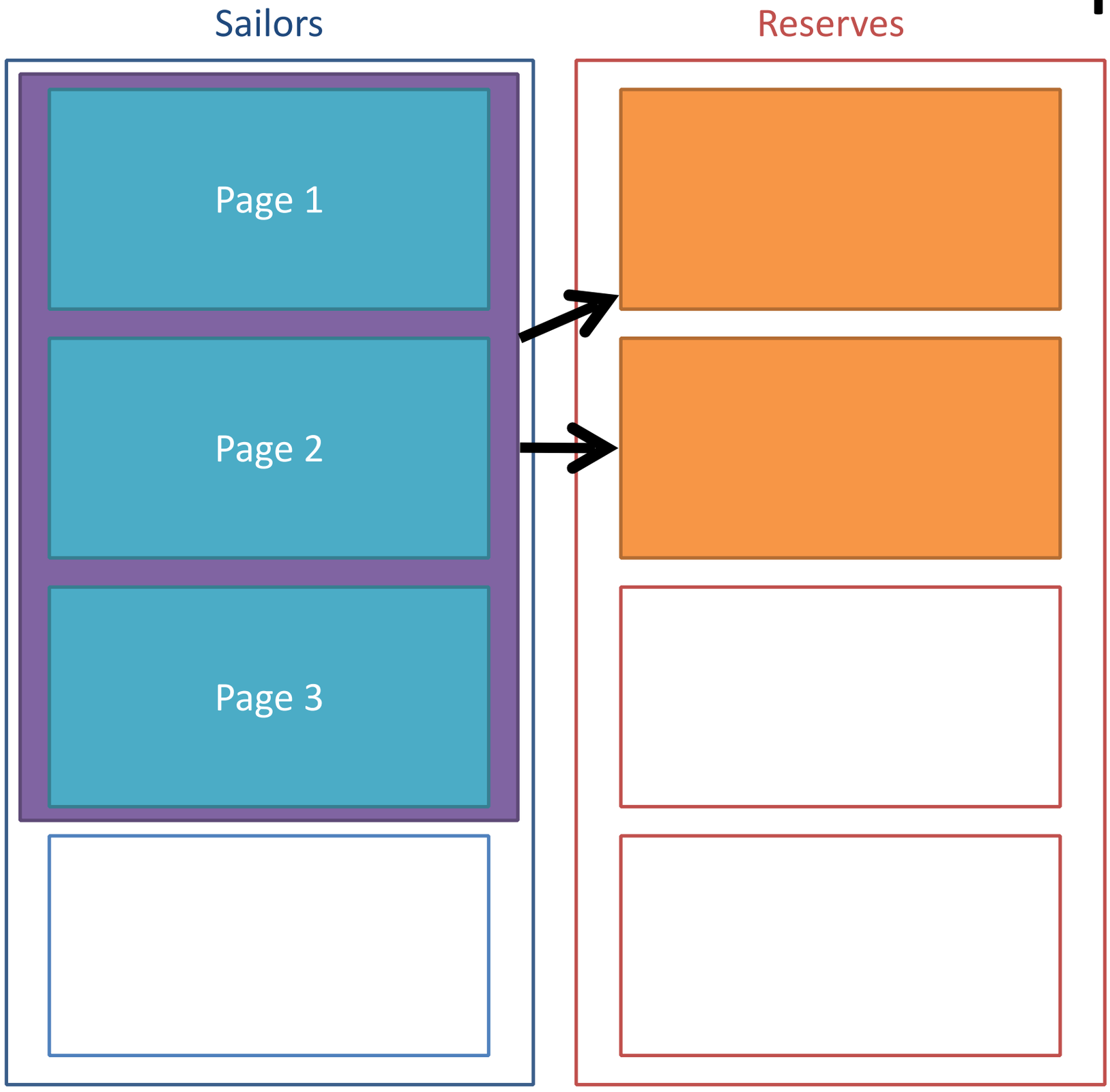

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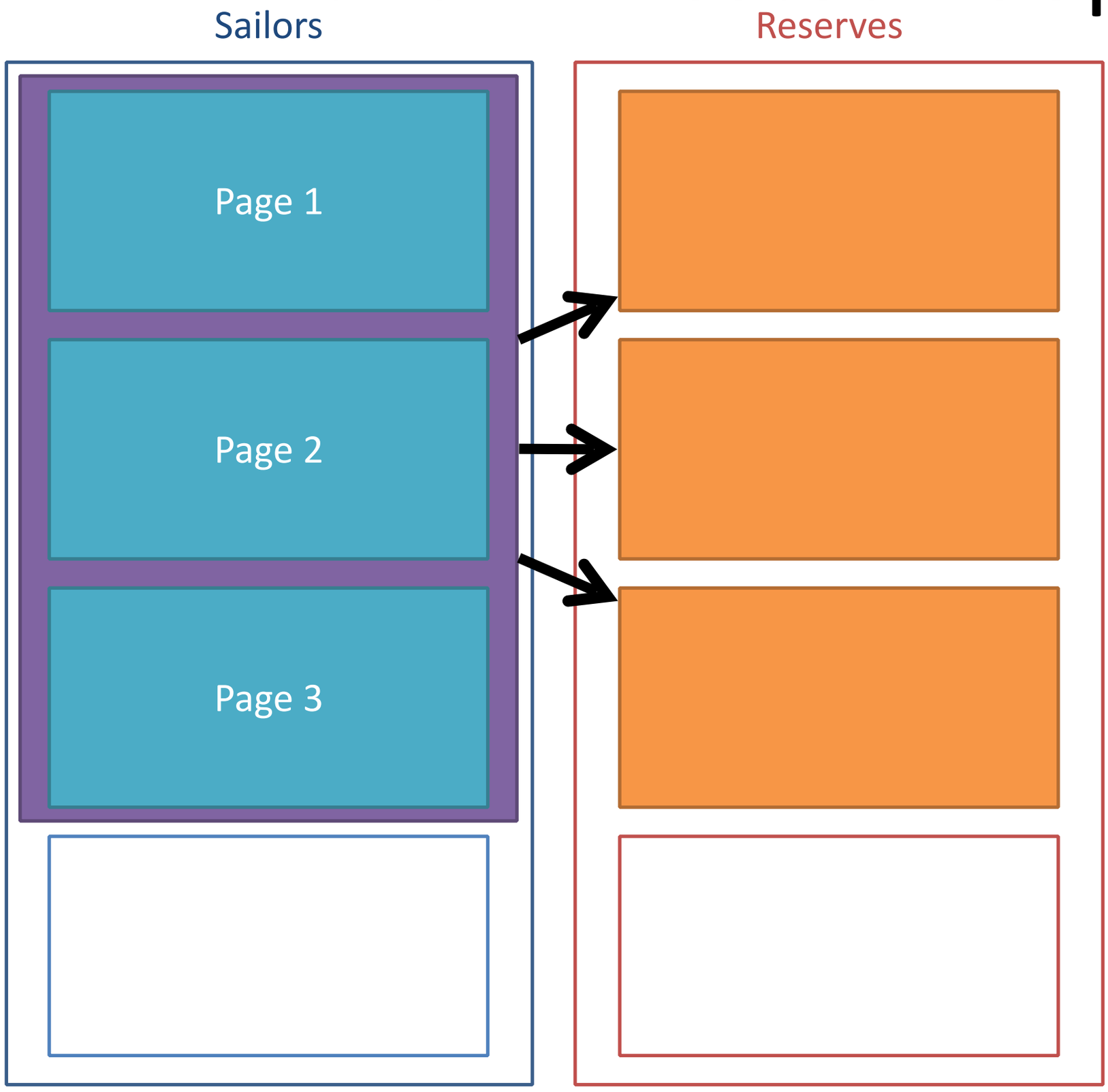

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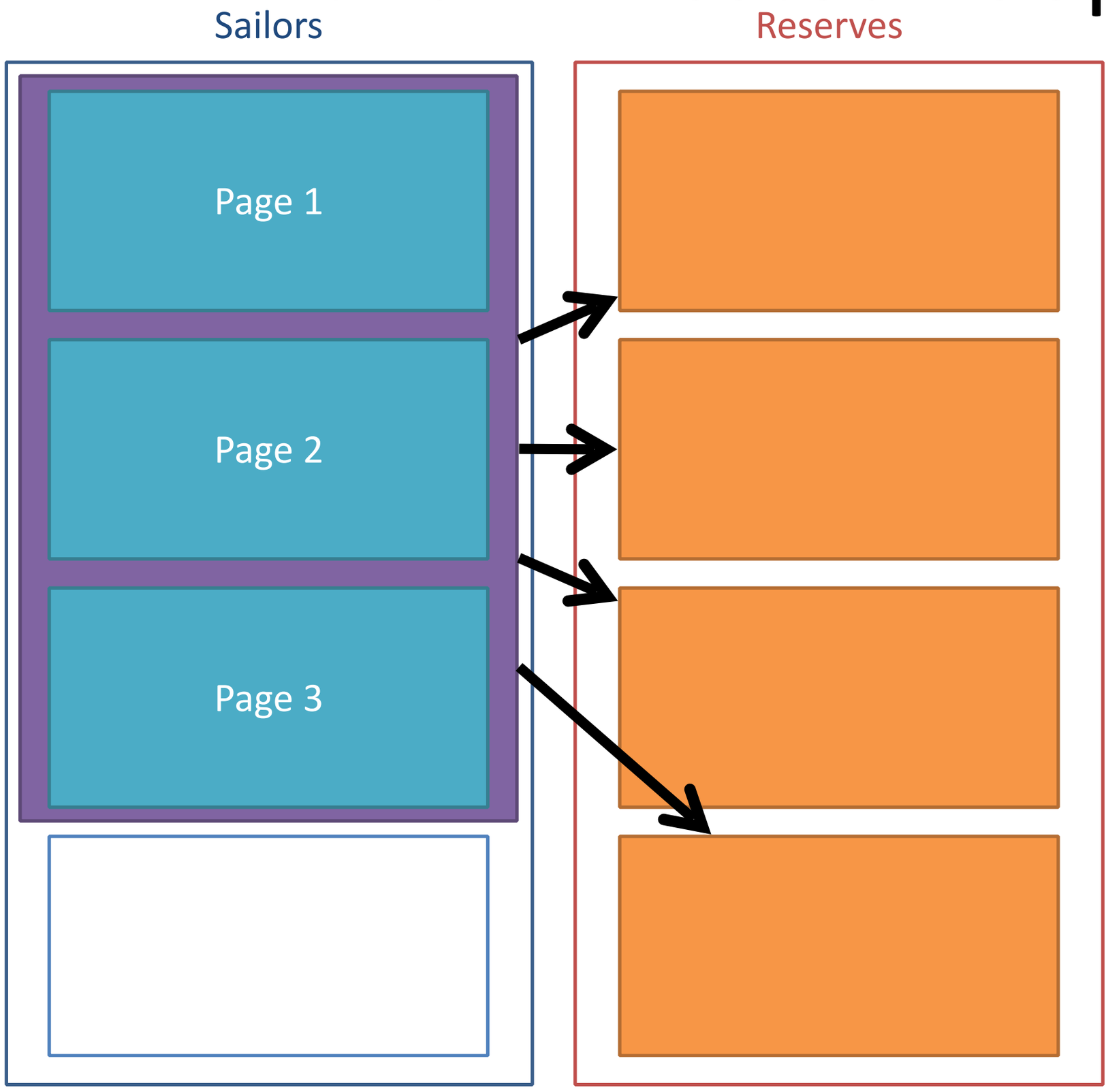

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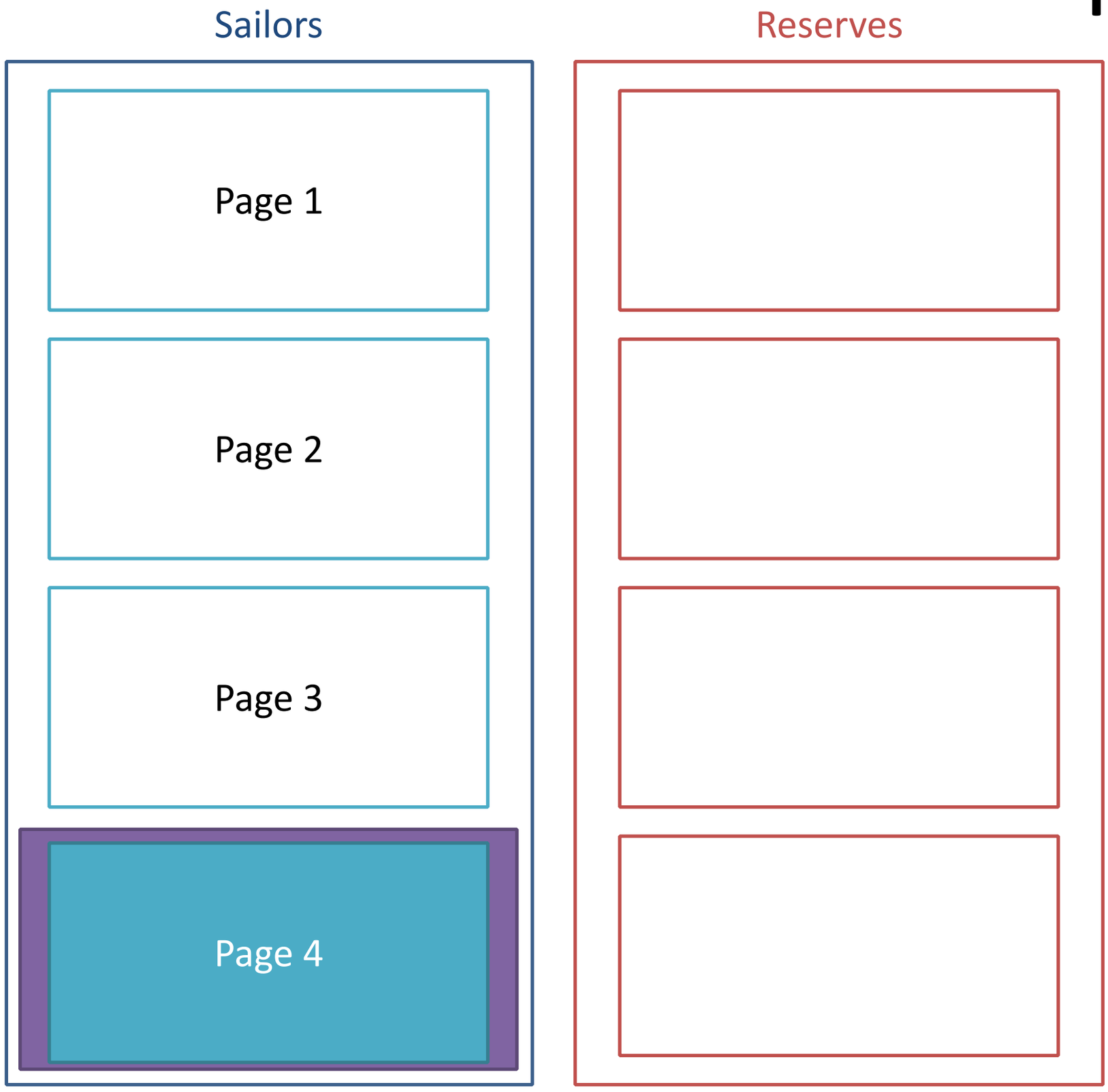

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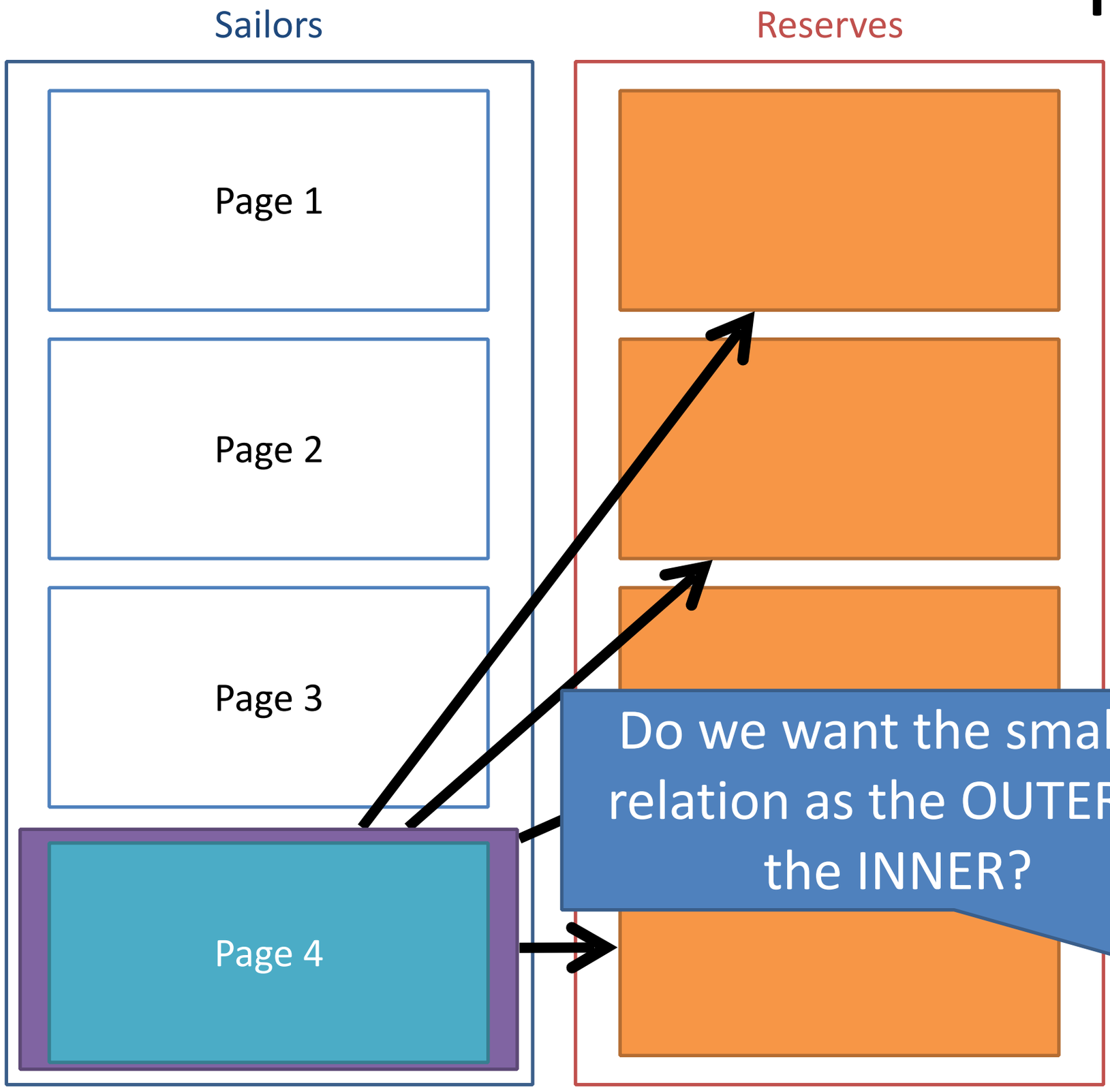

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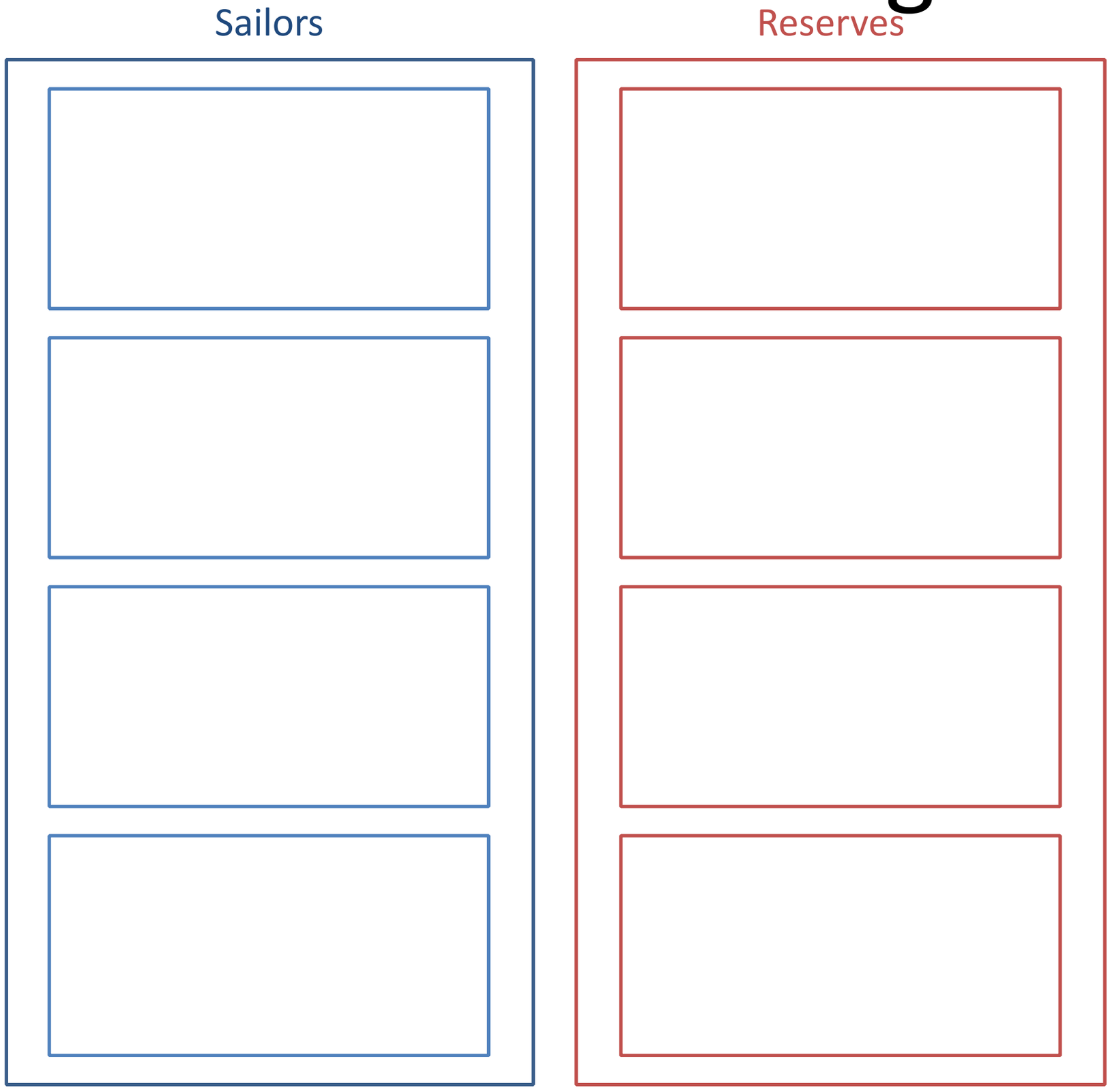
Cost:

[S] + ([S] / k)*[R]

Sort-Merge Join



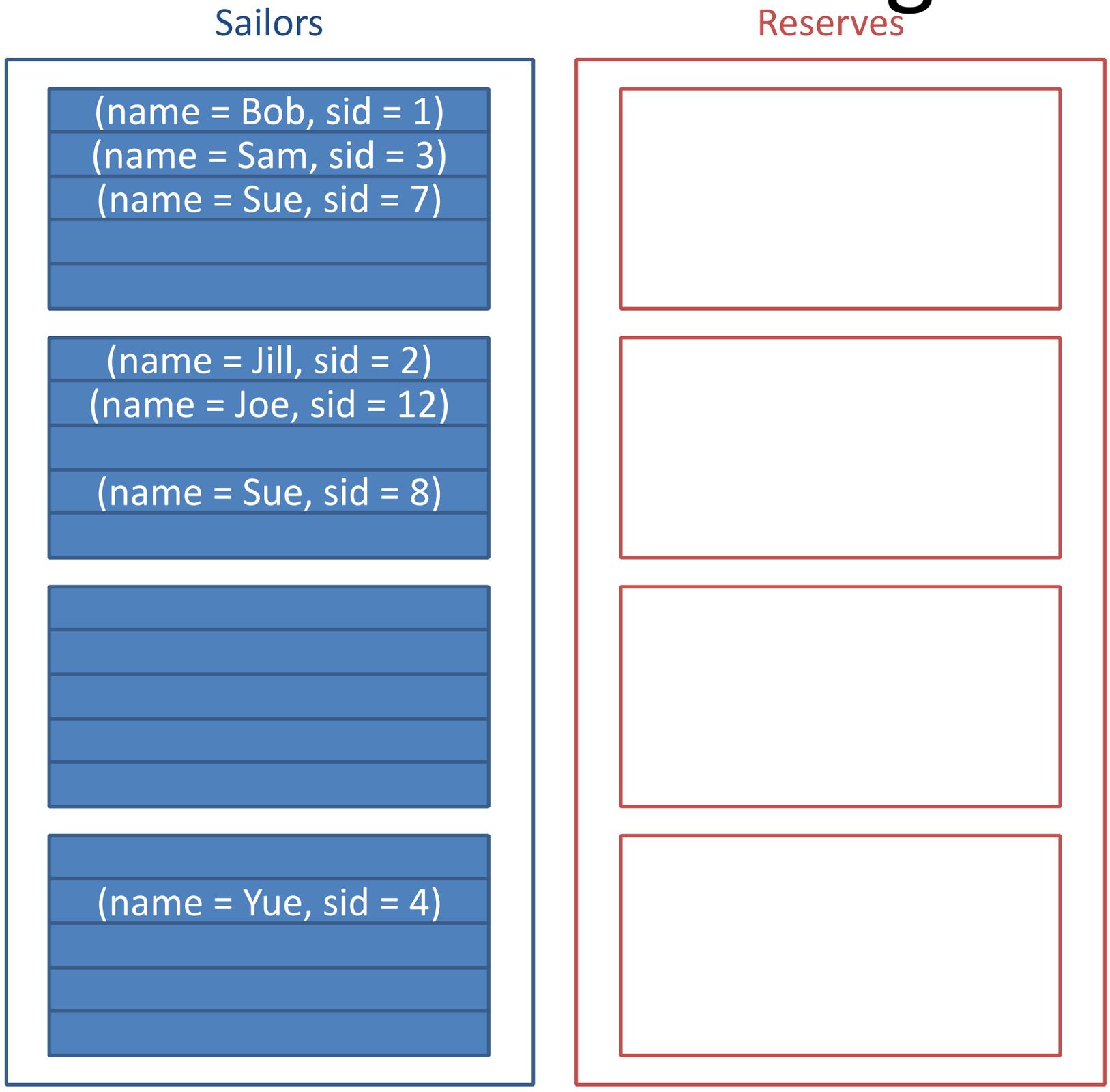
Key idea:

Sort S and R, then merge them!

Steps:

1. Sort S and R.
2. "Zip" or merge.

Sort-Merge Join



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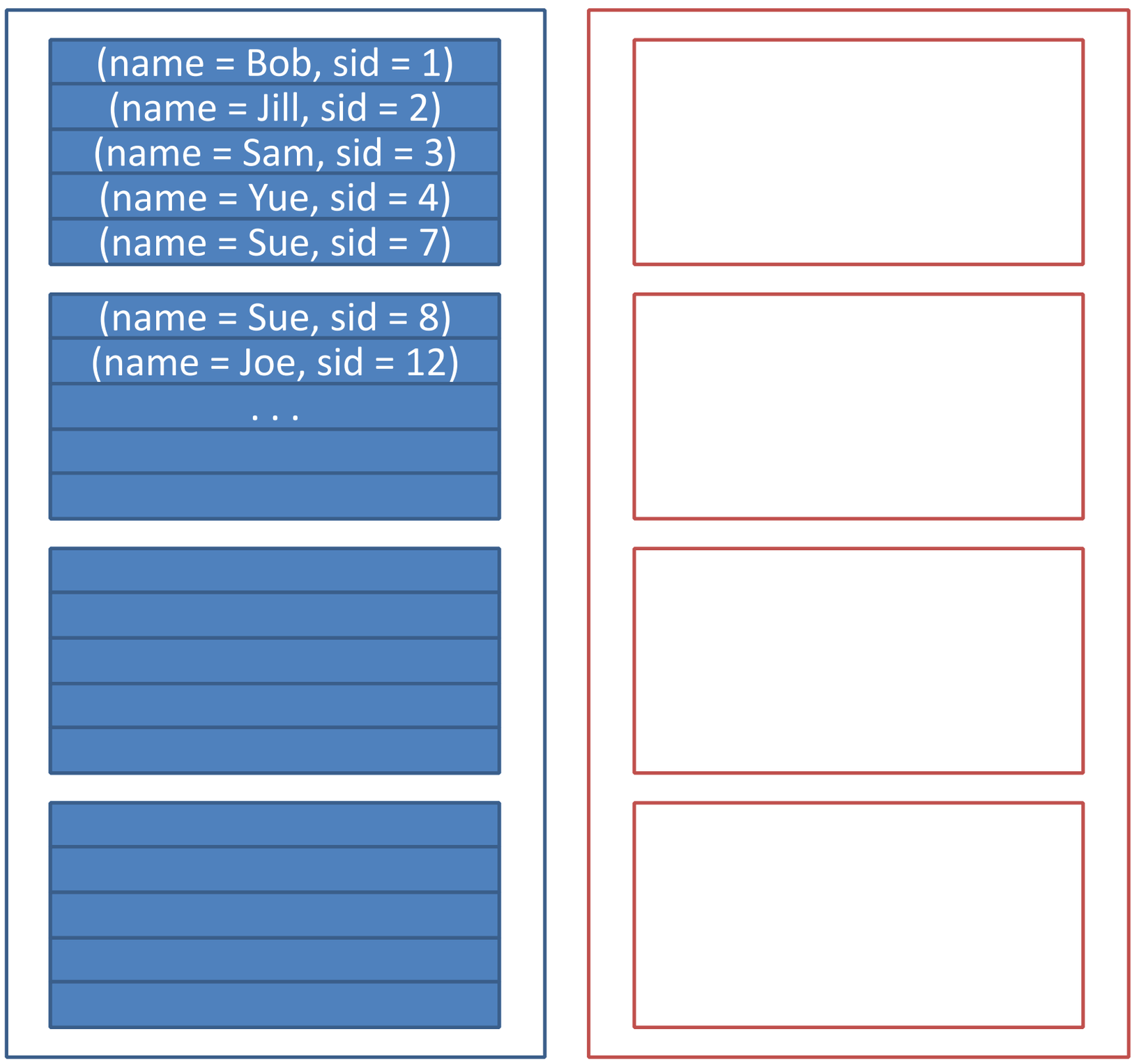

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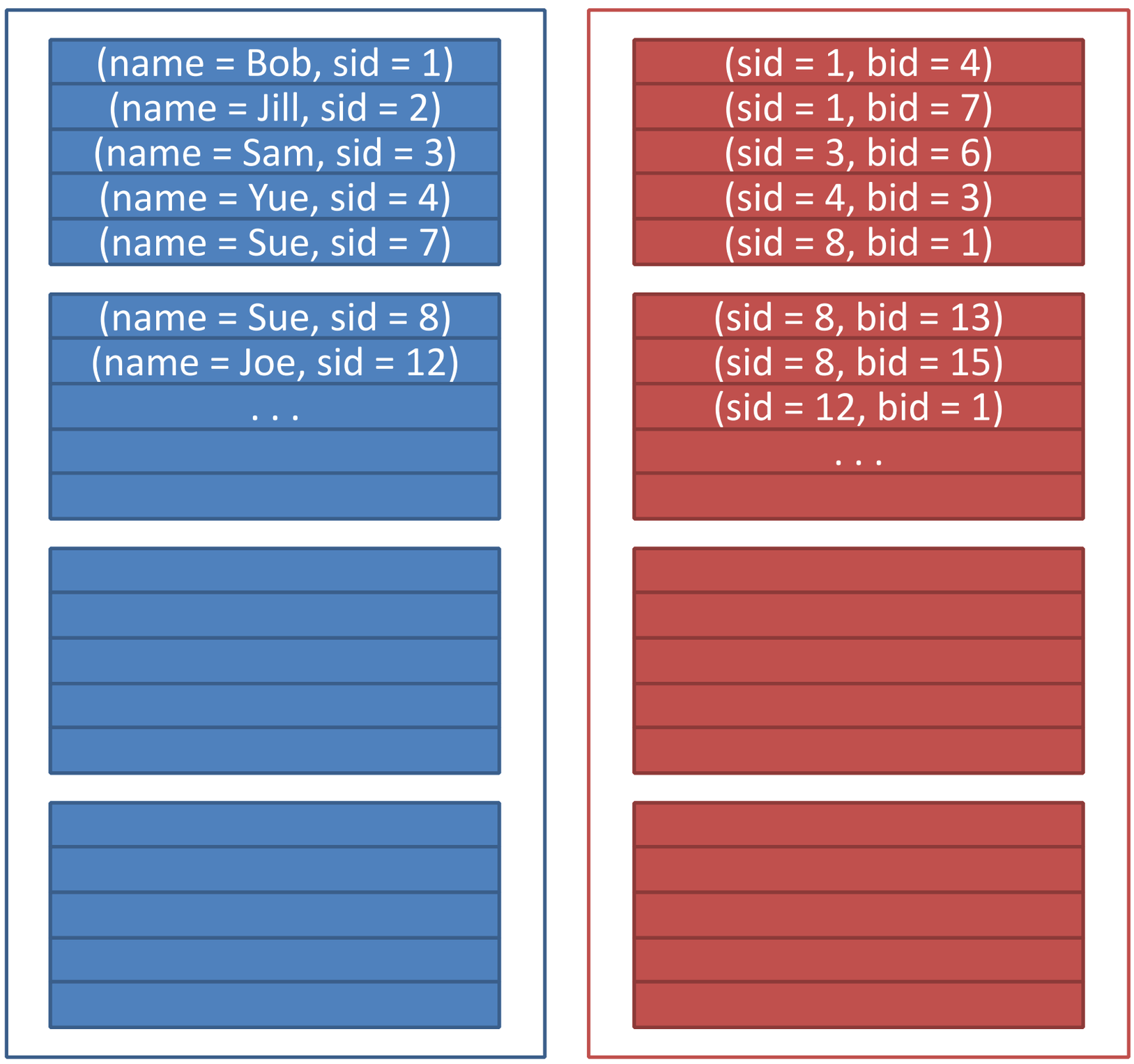

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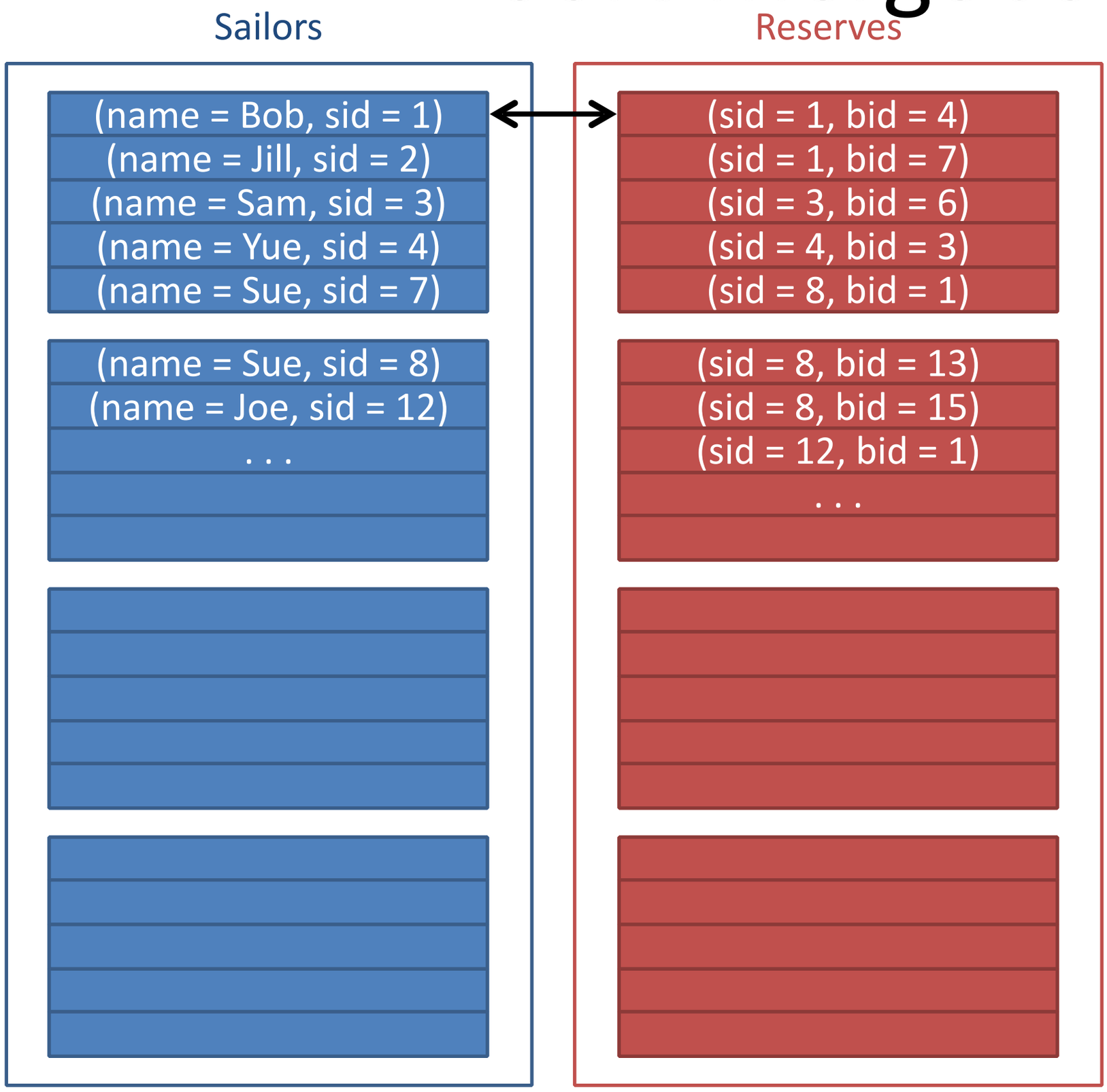

Key idea:

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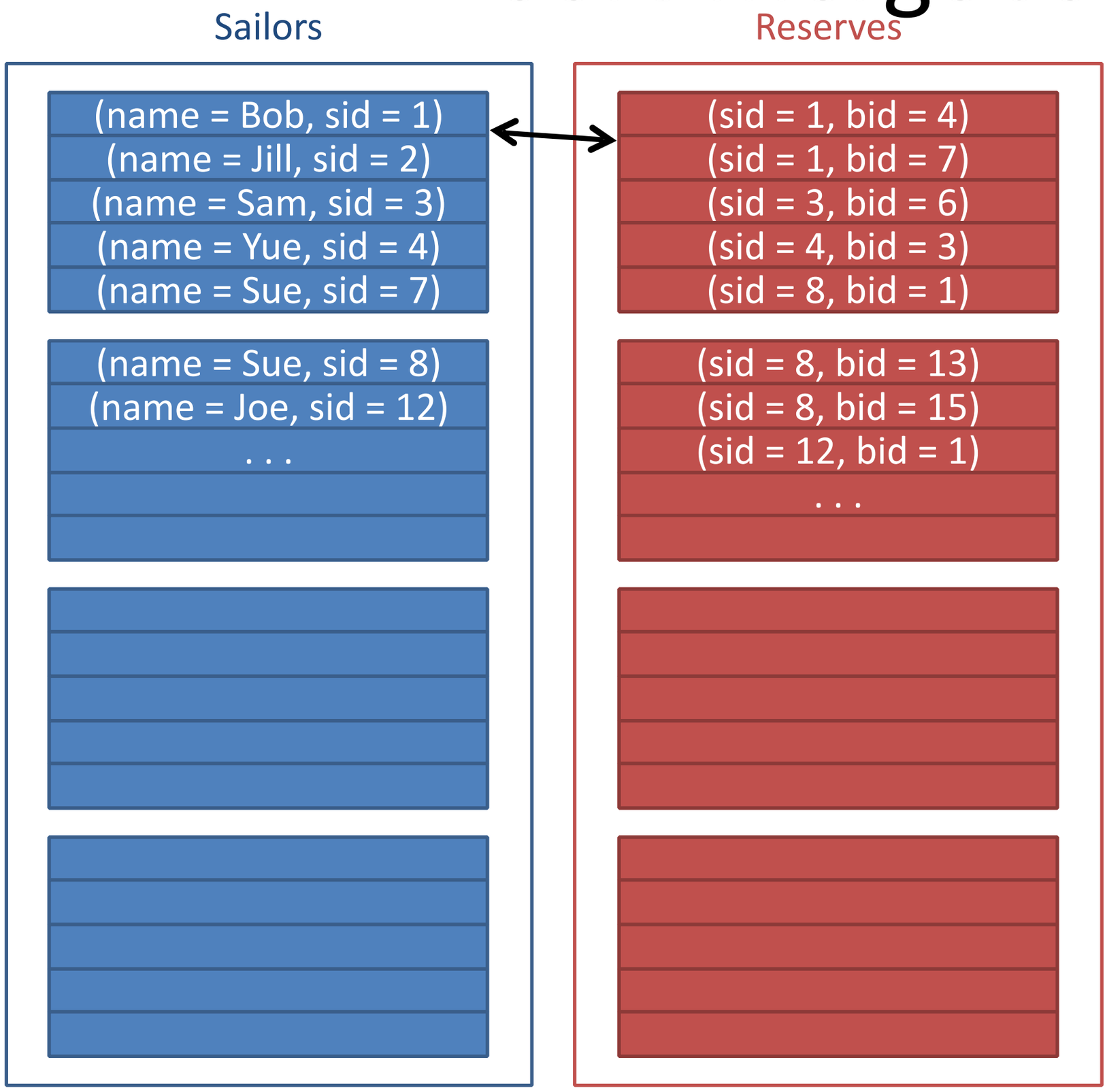
Steps:

1. Sort S and R.
2. "Zip" or merge.

Output:

(name = Bob, sid = 1, bid = 4)

Sort-Merge Join



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1. Sort S and R.
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Output:

(name = Bob, sid = 1, bid = 4)
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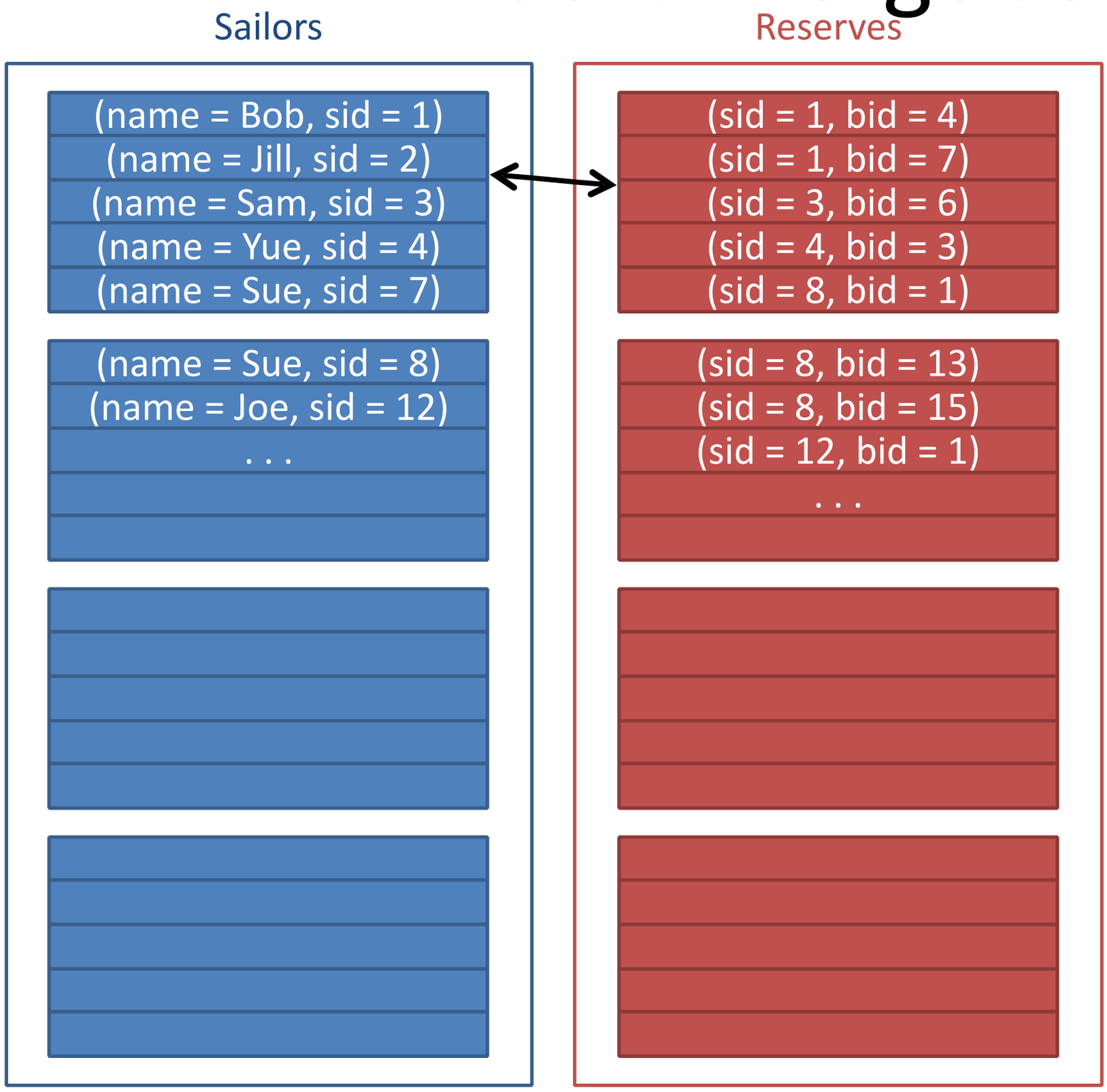

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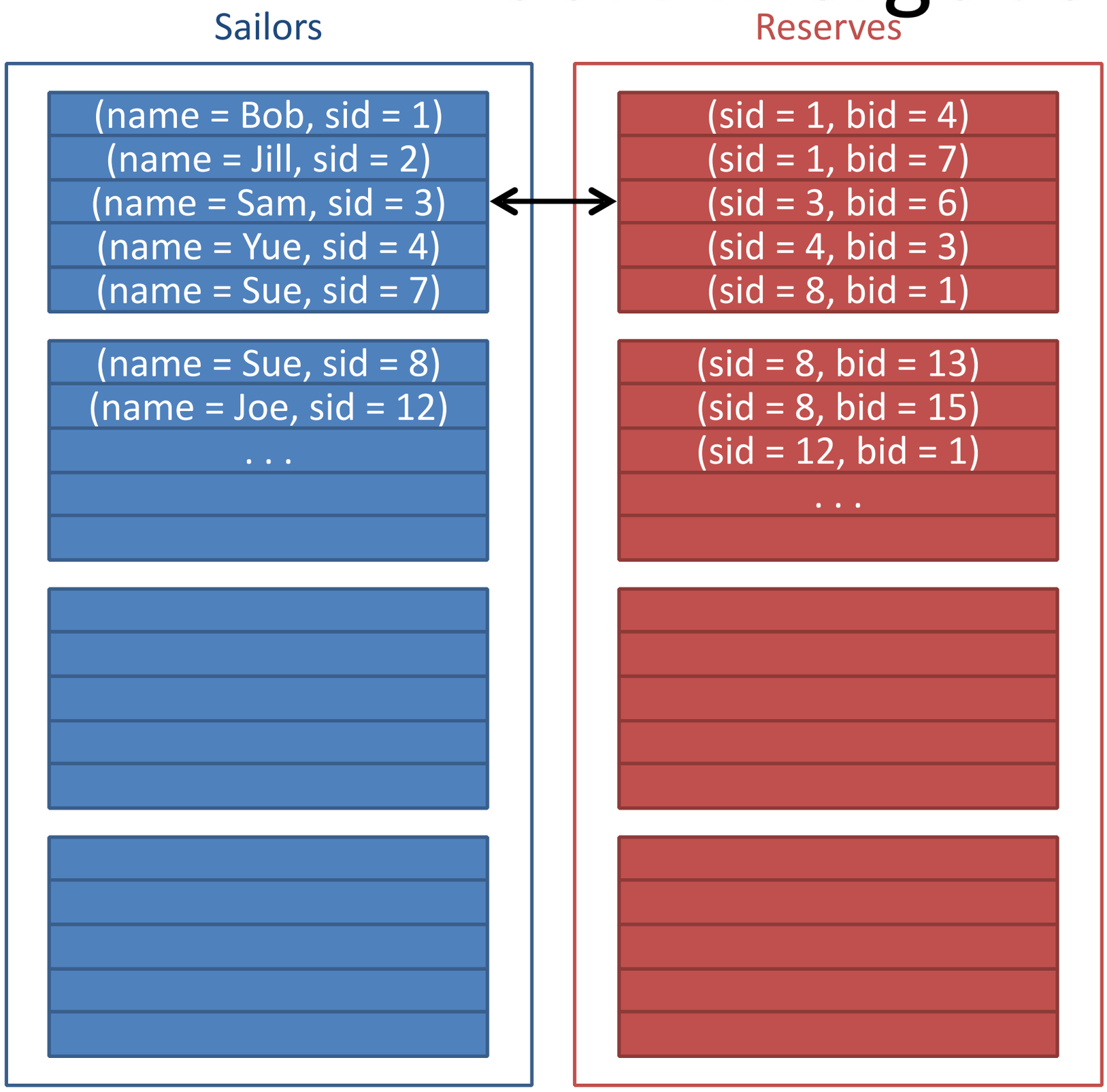

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(name = Bob, sid = 1, bid = 4)
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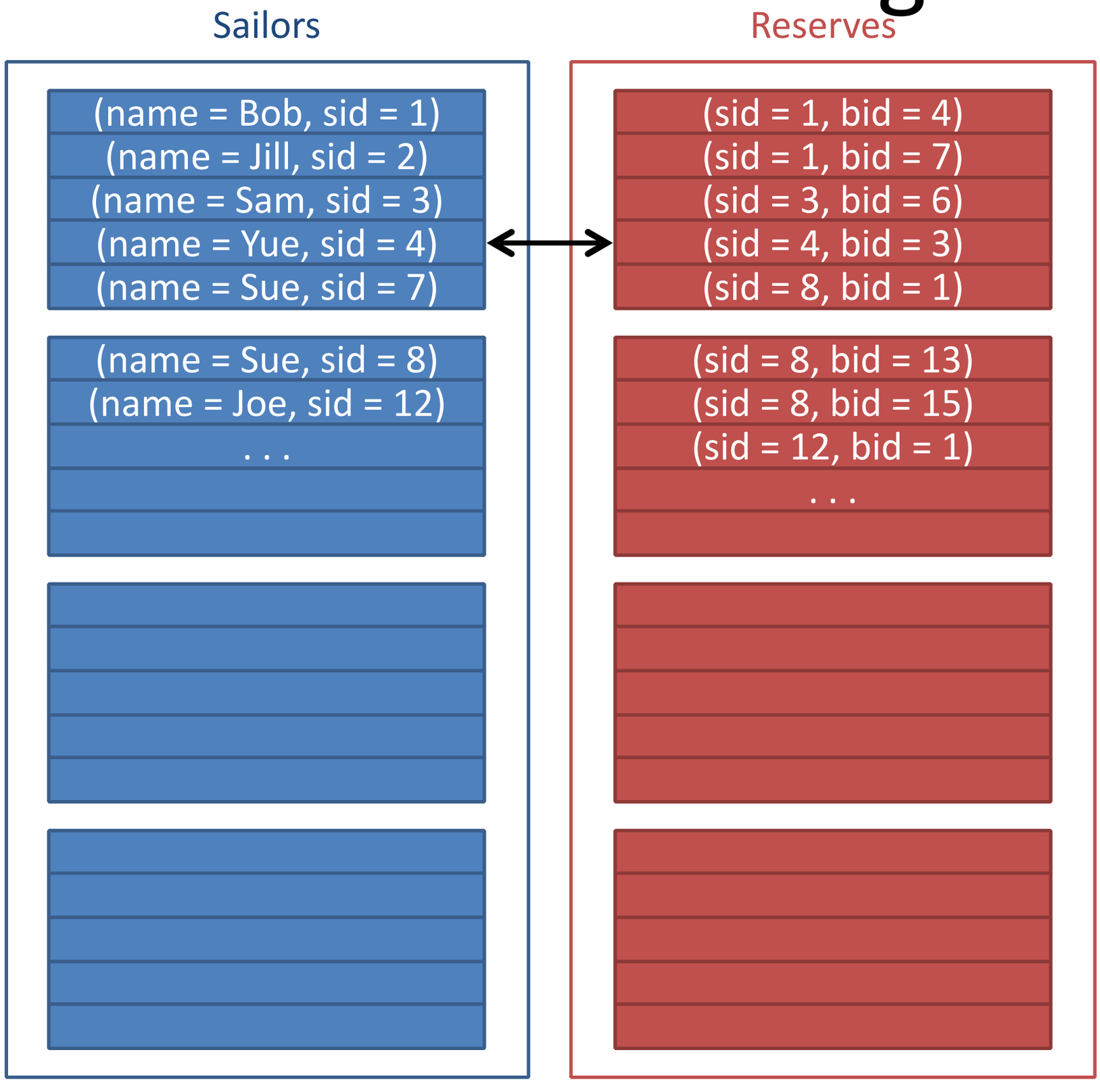

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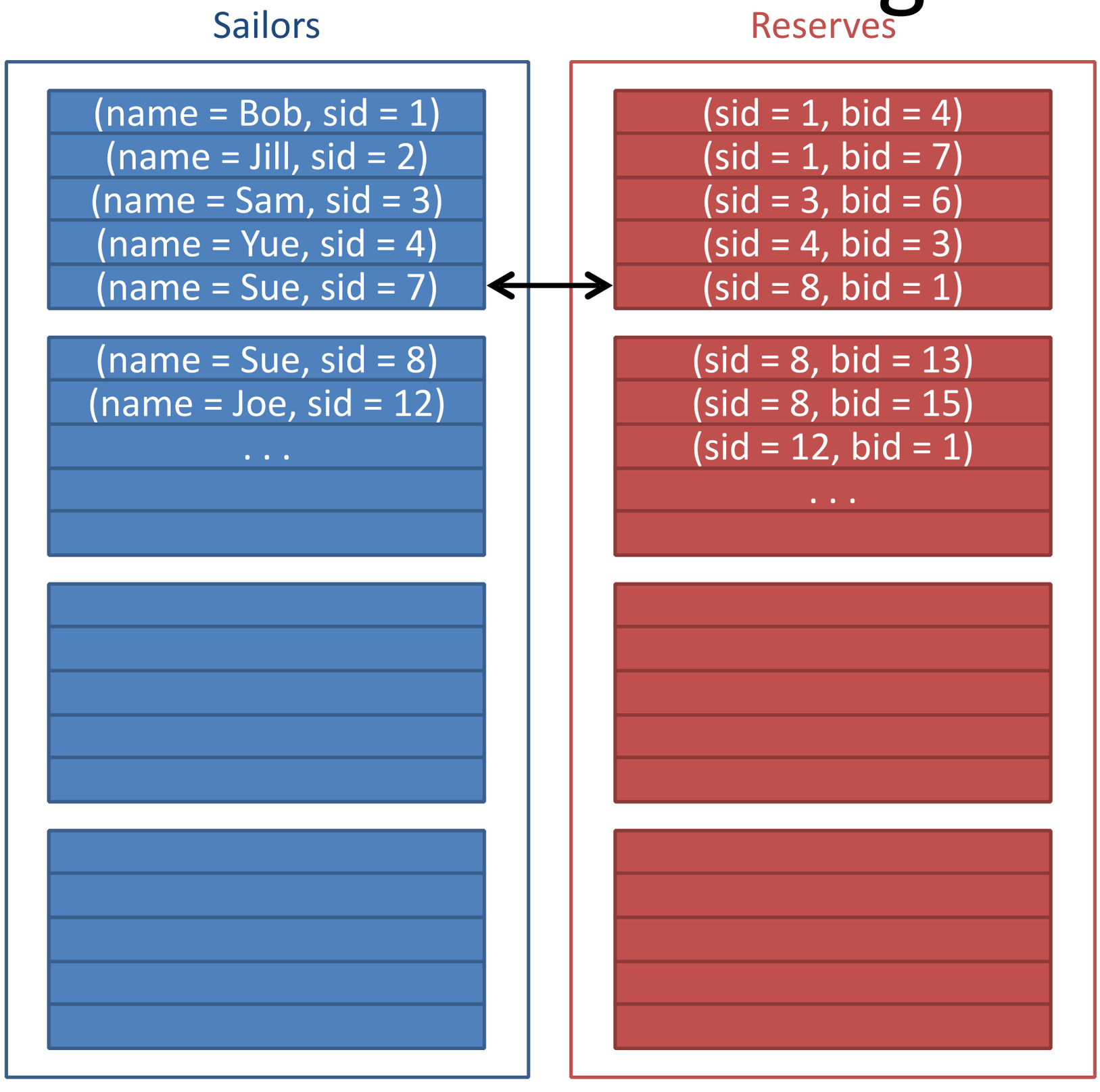

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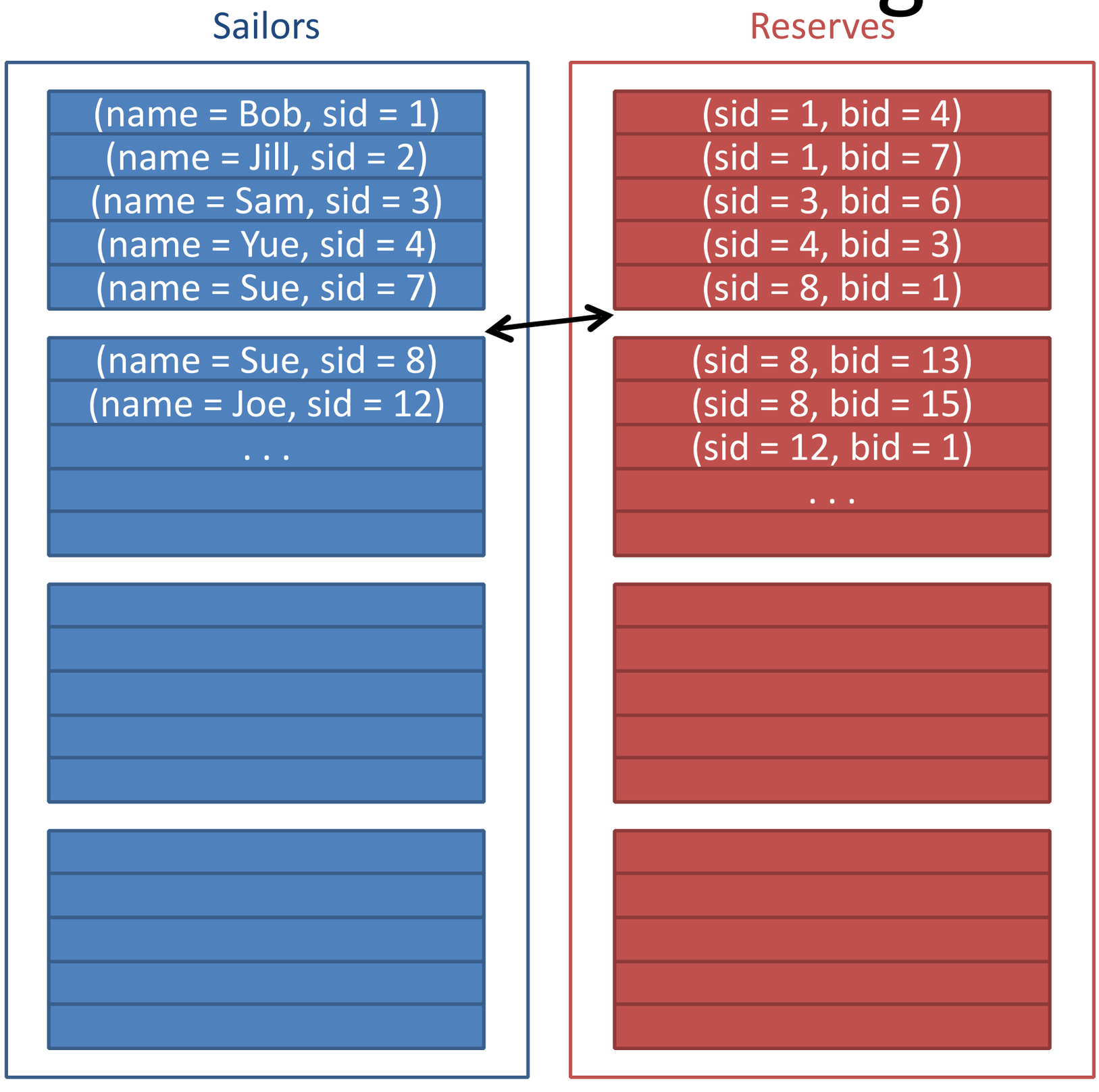

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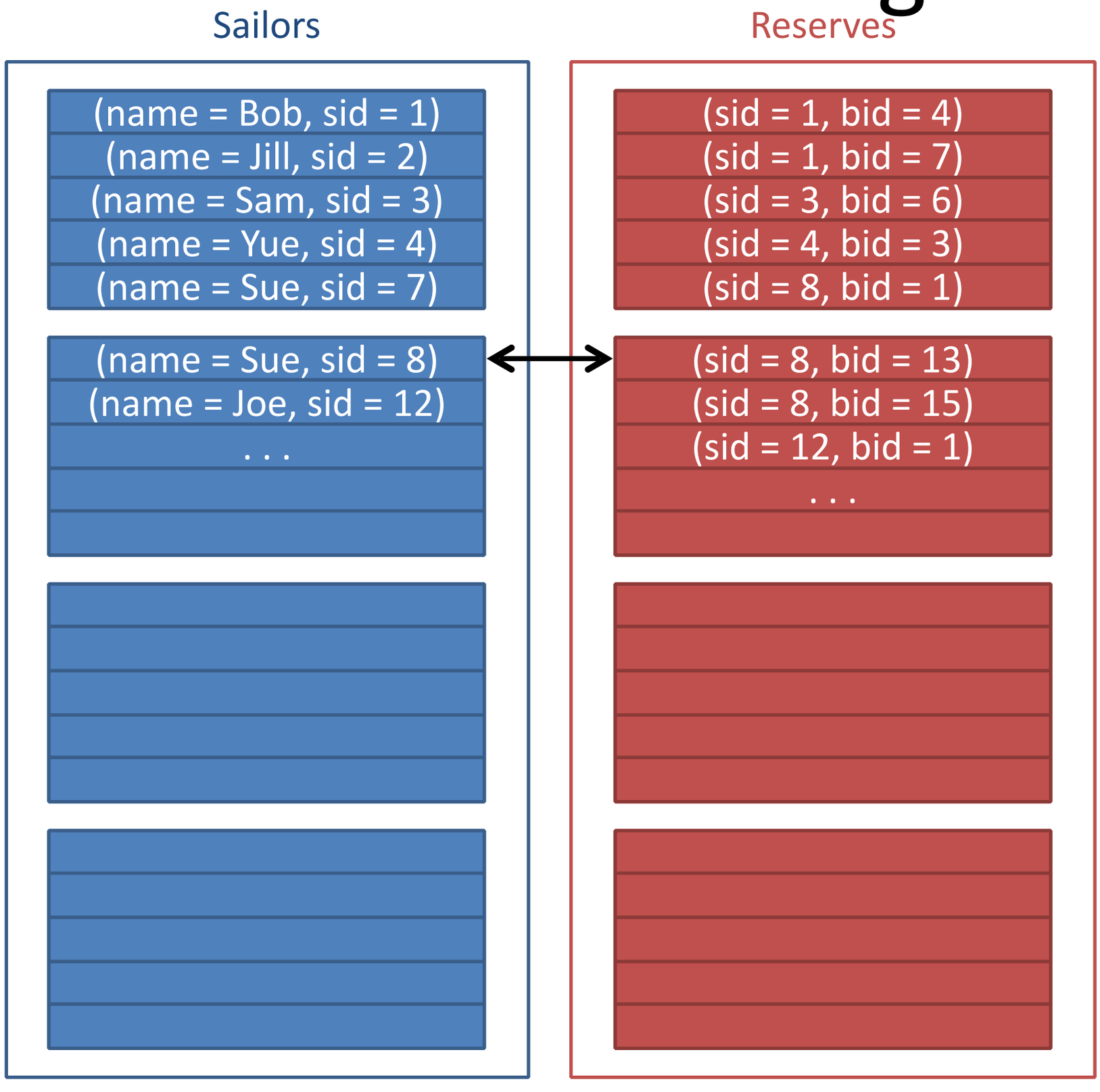

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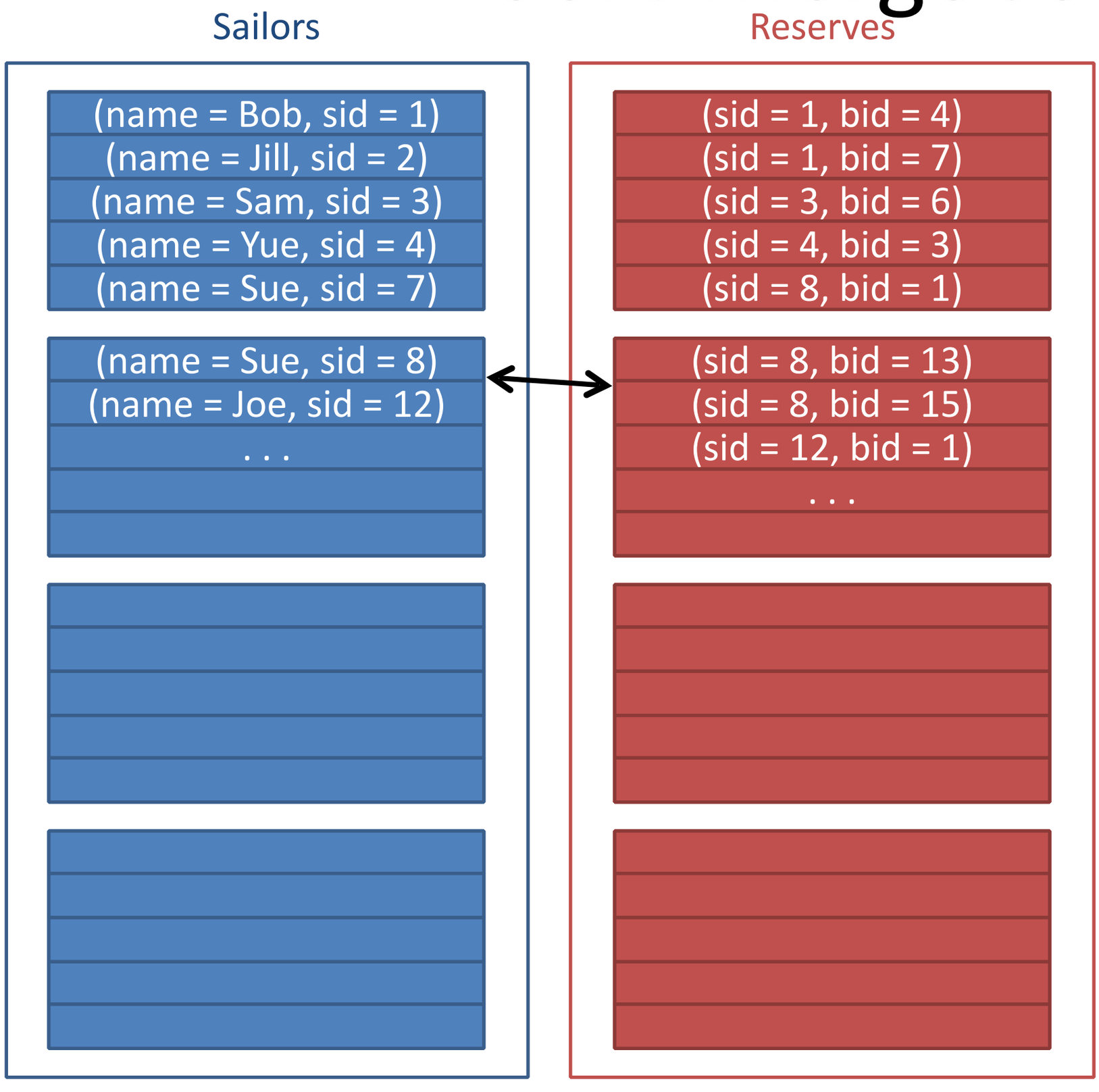

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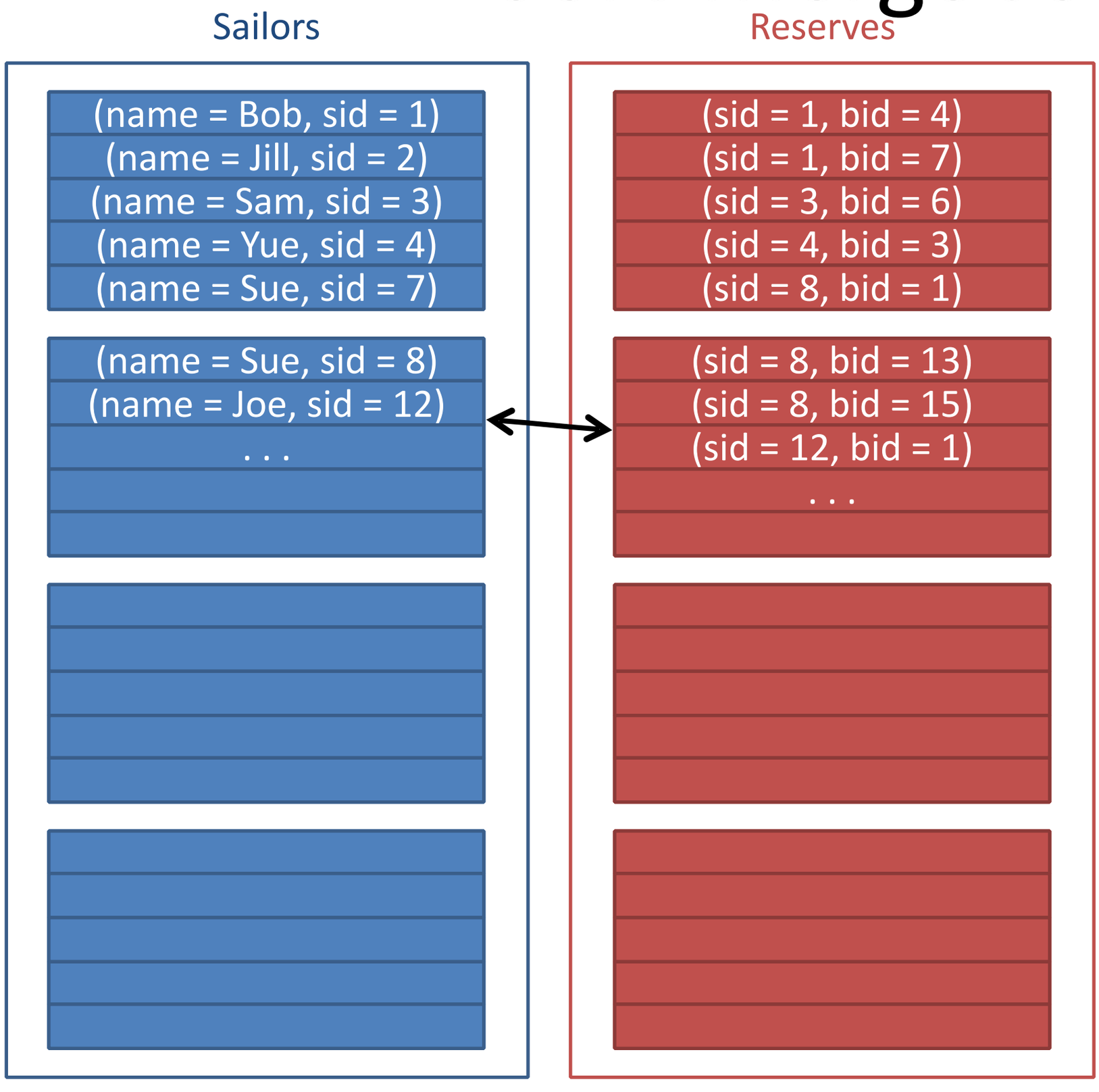

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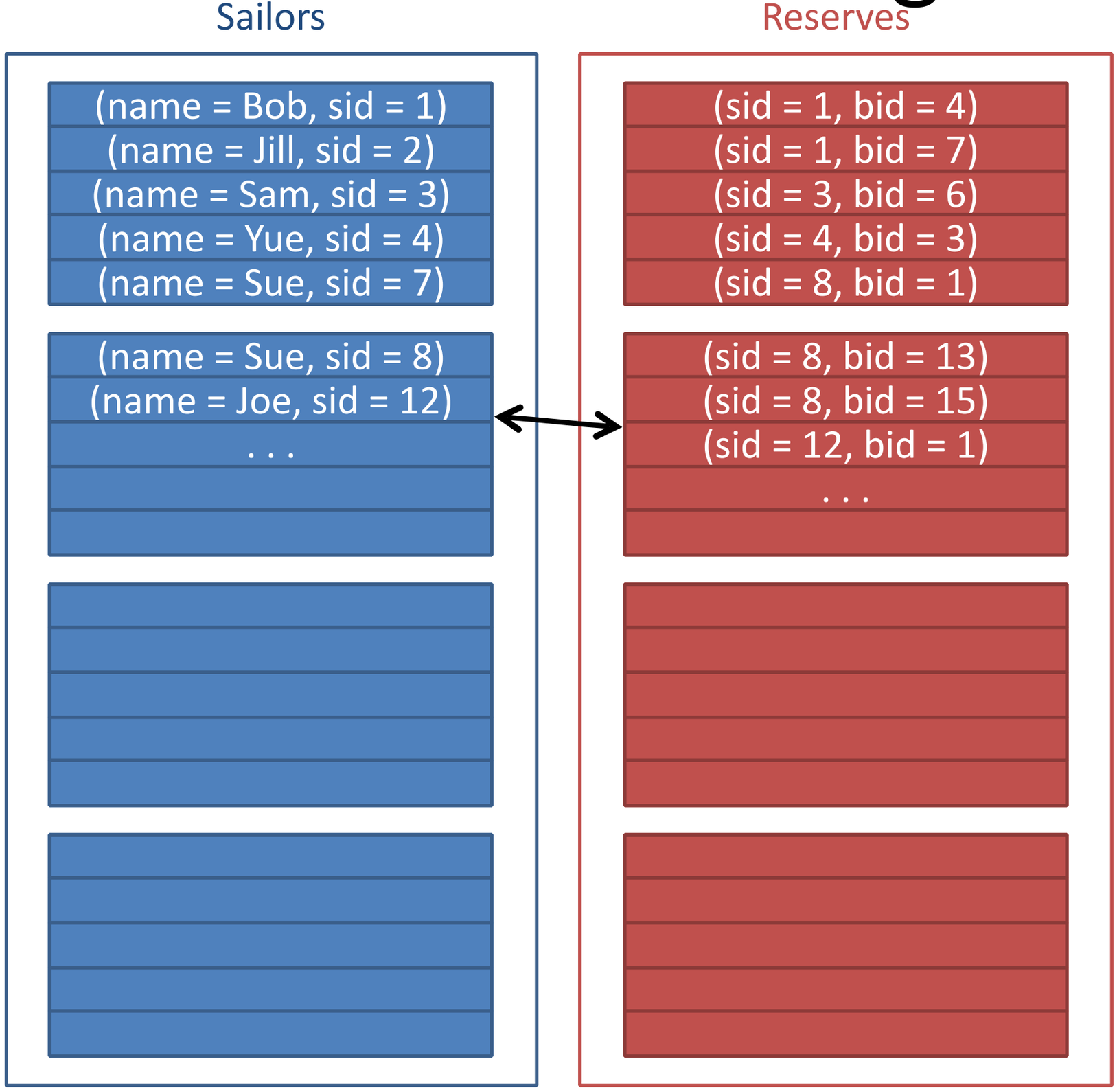
Notation: [S] == "# pages in S" ; |S| == "# tuples in S"

Sort-Merge Join



Key idea:

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Steps:

1. Sort S and R.
2. "Zip" or merge.

I/Os:

~5([S] + [R])

Sorting: 4([S]+[R])

Merging: [S]+[R]

Optimizing Sort-Merge Join

Sailors

(name = Bob, sid = 1)
(name = Jill, sid = 2)
(name = Sam, sid = 3)
(name = Yue, sid = 4)
(name = Sue, sid = 7)

(name = Sue, sid = 8)
(name = Joe, sid = 12)
. . .

Reserves

(sid = 1, bid = 4)
(sid = 1, bid = 7)
(sid = 3, bid = 6)
(sid = 4, bid = 3)
(sid = 8, bid = 1)

(sid = 8, bid = 13)
(sid = 8, bid = 15)
(sid = 12, bid = 1)
. . .

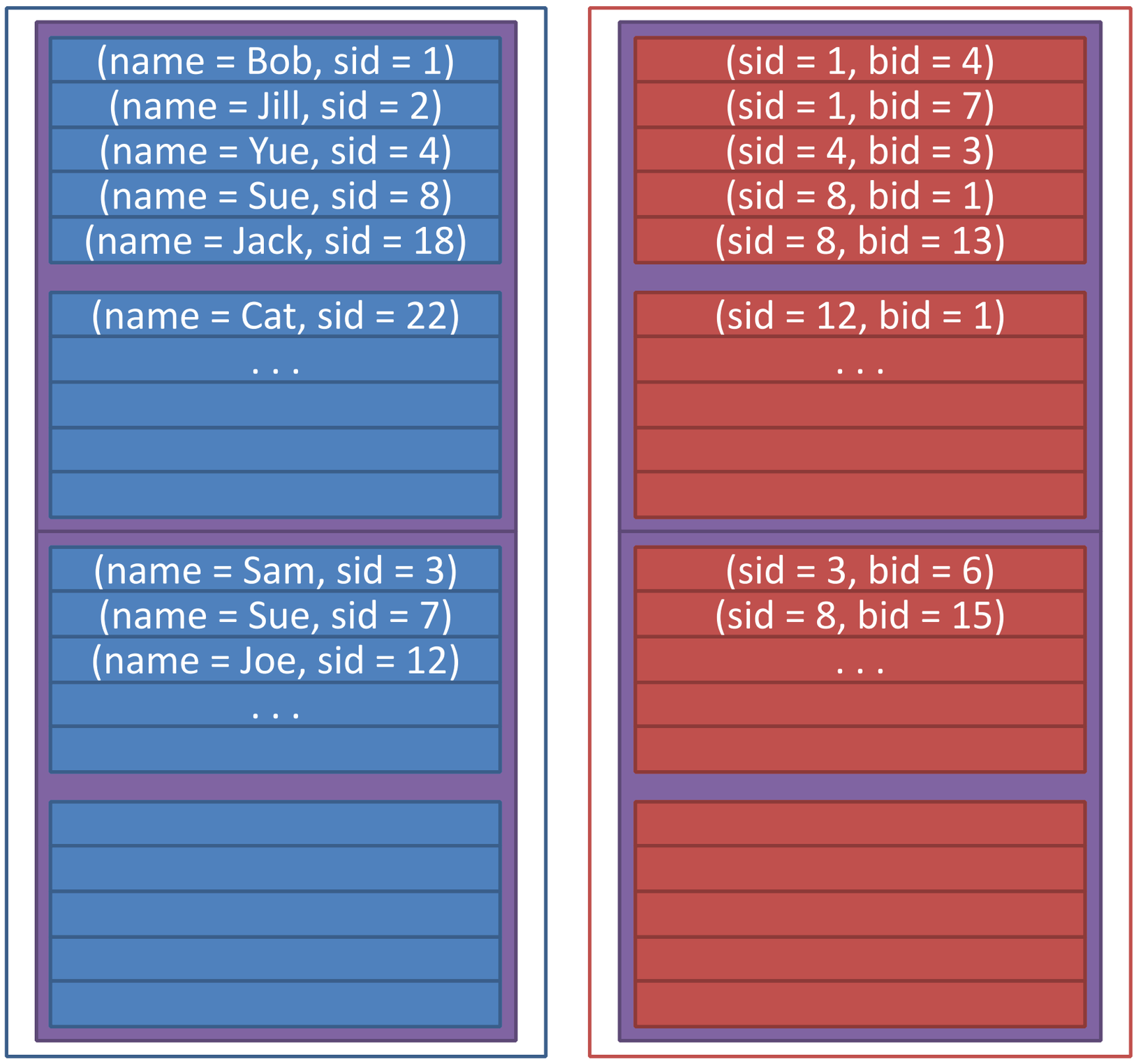
Key idea:

Internal Sort on both. Perform merge on all runs!

Steps:

1. Internal sort S and R. (Pass 0)
2. Merge all runs.

Optimizing Sort-Merge Join



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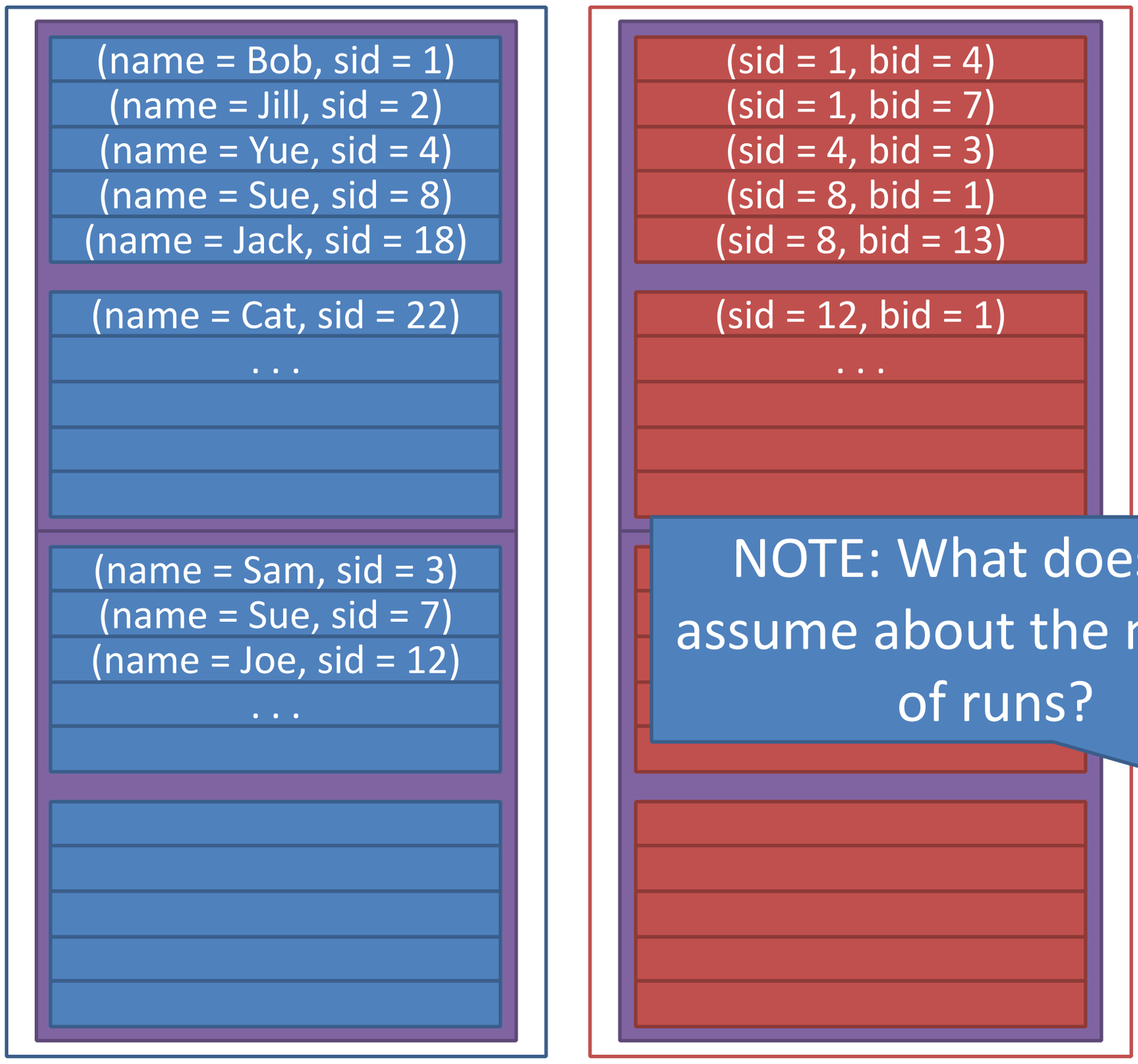

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Internal Sort on both. Perform merge on all runs!

Steps:

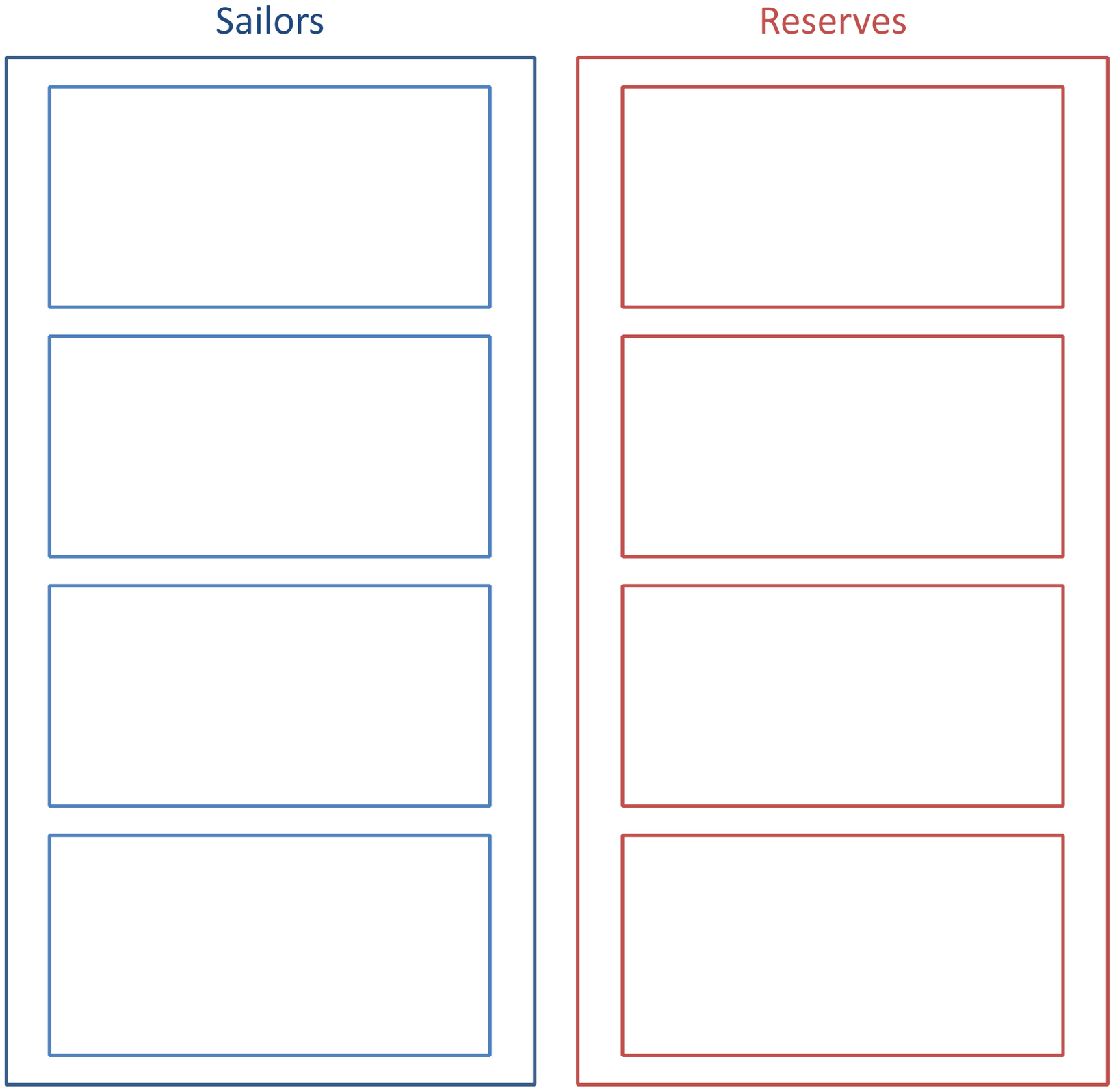
1. Internal sort S and R. (Pass 0)
2. ge all runs.

*3([S] + [R])

Pass 0: 2([S]+[R])

Merging: [S]+[R]

Hash-Join



Key idea:

Partition S and R using same hash fn, then collect same partitions

Steps:

1. Partition S and R
2. Re-Hash, collect

Hash-Join

Sailors

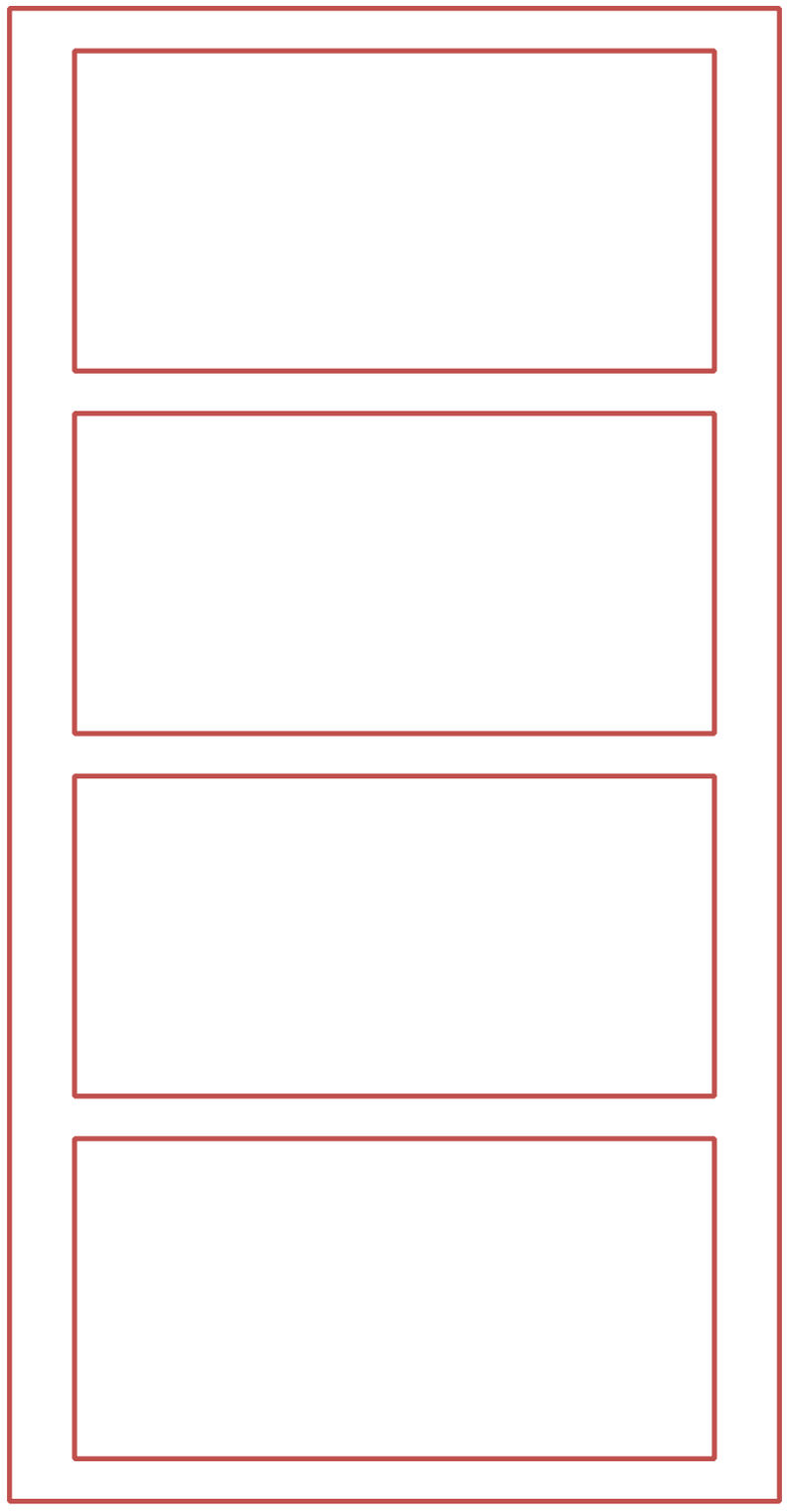
(name = Bob, sid = 1)
(name = Sam, sid = 3)
(name = Sue, sid = 7)

(name = Jill, sid = 2)
(name = Joe, sid = 12)

(name = Sue, sid = 8)

(name = Yue, sid = 4)

Reserves



Key idea:

Partition S and R using same hash fn, then collect same partitions

Steps:

1. Partition S and R
2. Re-Hash, collect

Hash-Join

Hash function: **sid mod 4**

Sailors

(name = Joe, sid = 12)
(name = Sue, sid = 8)
(name = Yue, sid = 4)
. . .

(name = Bob, sid = 1)
. . .

(name = Jill, sid = 2)
. . .

(name = Sue, sid = 7)
(name = Sam, sid = 3)
. . .

Reserves

(sid = 12, bid = 1)
(sid = 8, bid = 13)
(sid = 8, bid = 15)
(sid = 4, bid = 3)
(sid = 8, bid = 1)

(sid = 1, bid = 4)
(sid = 1, bid = 7)
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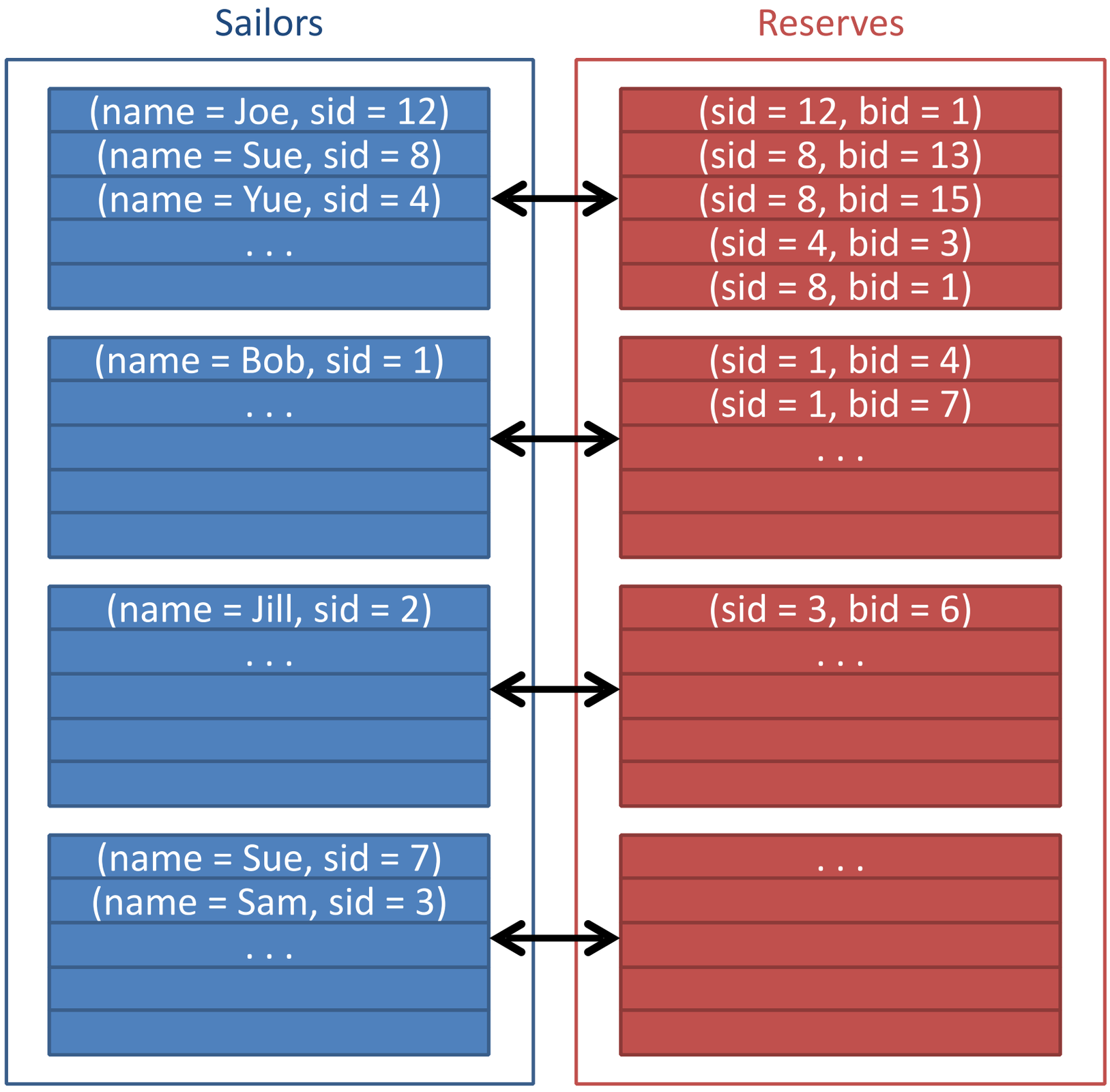

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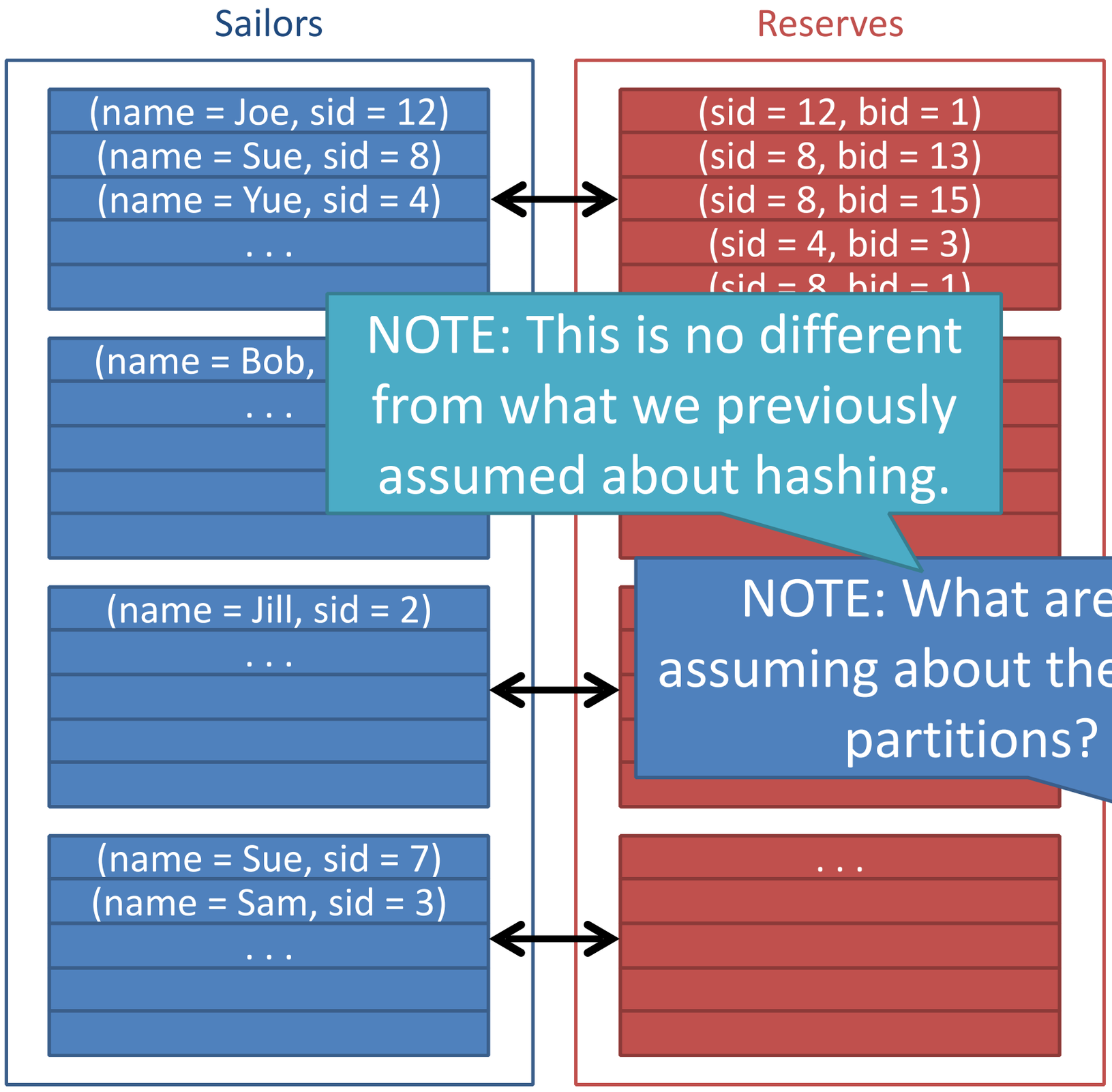
Notation: [S] == "# pages in S" ; |S| == "# tuples in S"

Hash-Join



Key idea:

Partition S and R using same hash fn, then collect same partitions

Steps:

1. Partition S and R
2. Hash, collect

3([S] + [R])

Partition: 2([S]+[R])

Re-Hash: [S]+[R]

Join Cheatsheet

Notation: [S] == "# pages in S" ;
|S| == "# tuples in S"

- Chunk nested loop join
  - Take **k pages** of S and match with each page of R.
  - Total Cost: [S] + ([S] / k)*[R]
- Sort merge join
  - Sort S and R **on join column**, then merge them!
  - Total Cost: ~5([S] + [R])
- Hash join
  - Partition S and R using same hash fn, then collect same partitions
  - Total Cost: ~3([S] + [R])
    - Assuming len(partition) ≤ B pages

When is a chunk-nested loops join the best?

When is a chunk-nested loops join the best?

- Not using an equality predicate
- Join is just a cross product

When is a sort-merge join the best?

When is a sort-merge join the best?

- Skewed input data
- Small memory size
- Want sorted output/have sorted input

When is a hash-join the best?

When is a hash-join the best?

- One partition large, the other small (can keep in memory)

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How many disk reads are needed to perform Chunk Nested Loops Join?

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How many disk reads are needed to perform Chunk Nested Loops Join?

(# of pages in smaller relation) + ((# of pages in smaller relation) / (# of pages in memory - 2 for I/O)) * (# of pages in larger relation)

= 50 + (50/10) * (100) = 550

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How about a Hash Join? (Assume no recursive partitioning)

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How about a Hash Join? (Assume no recursive partitioning)

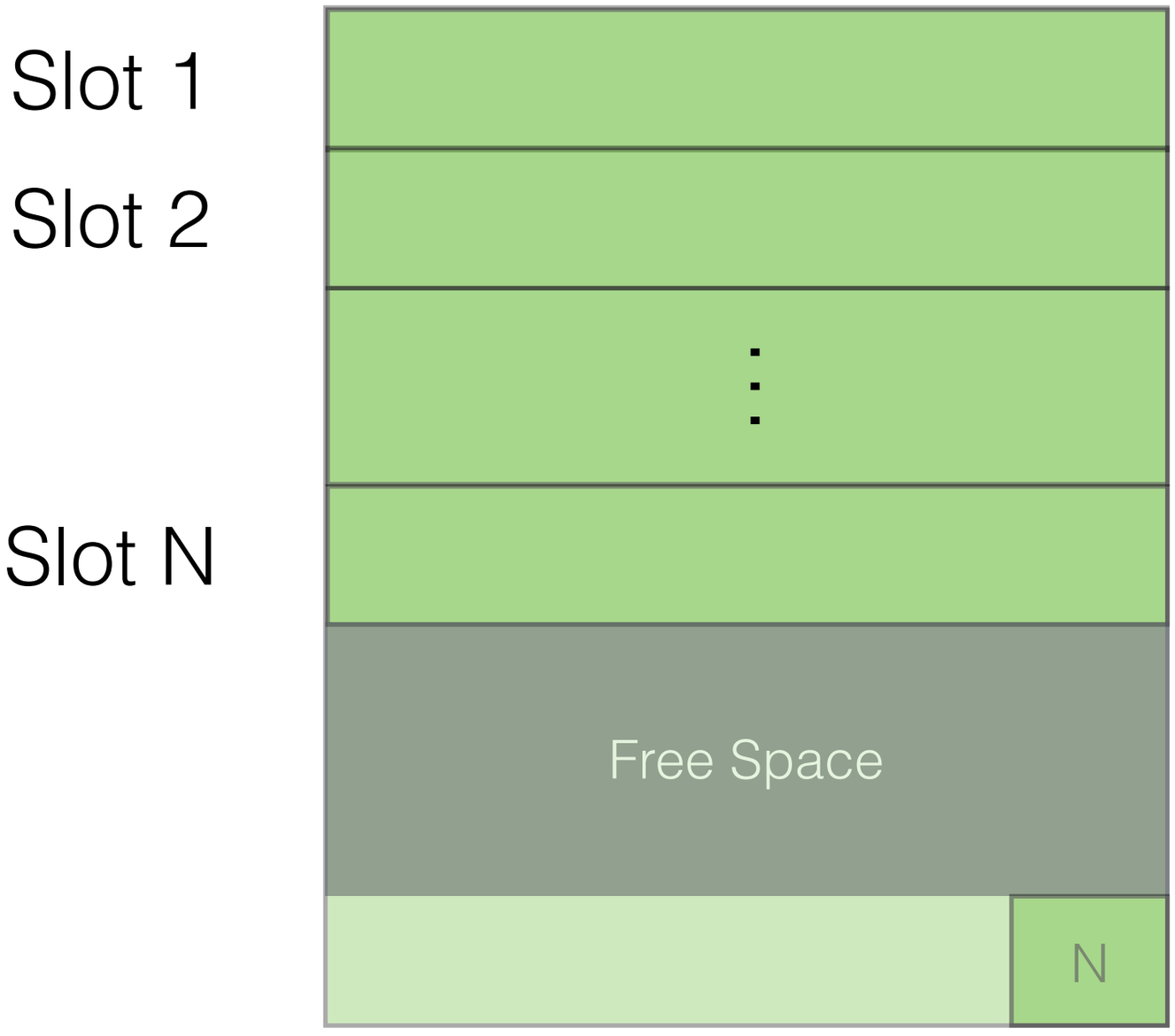
(# of pages in both relations) * (1 read before hashing + 1 read after hashing)

=(100+50) * 2
= 300

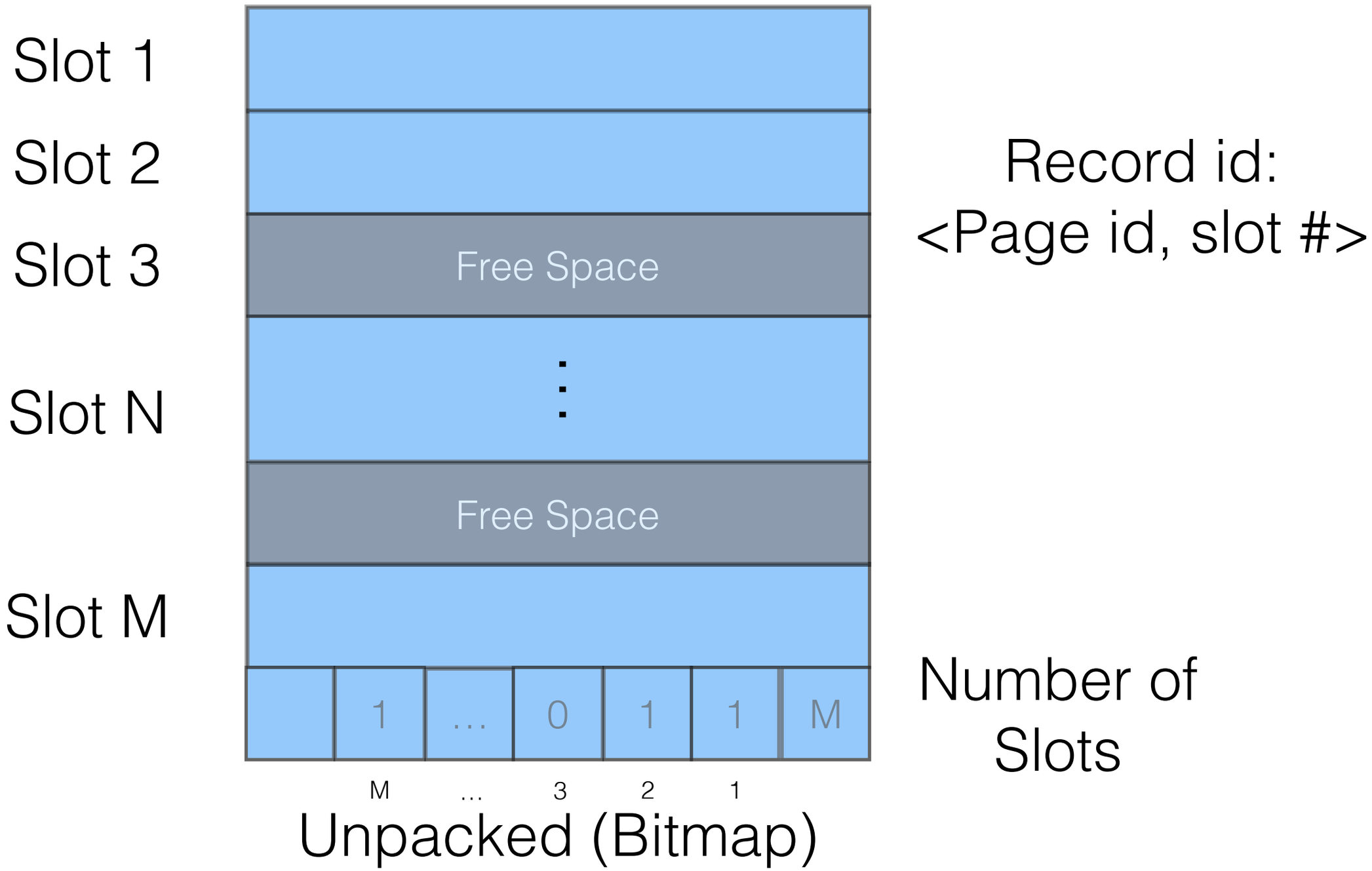
Heap Files

(Page Formats)

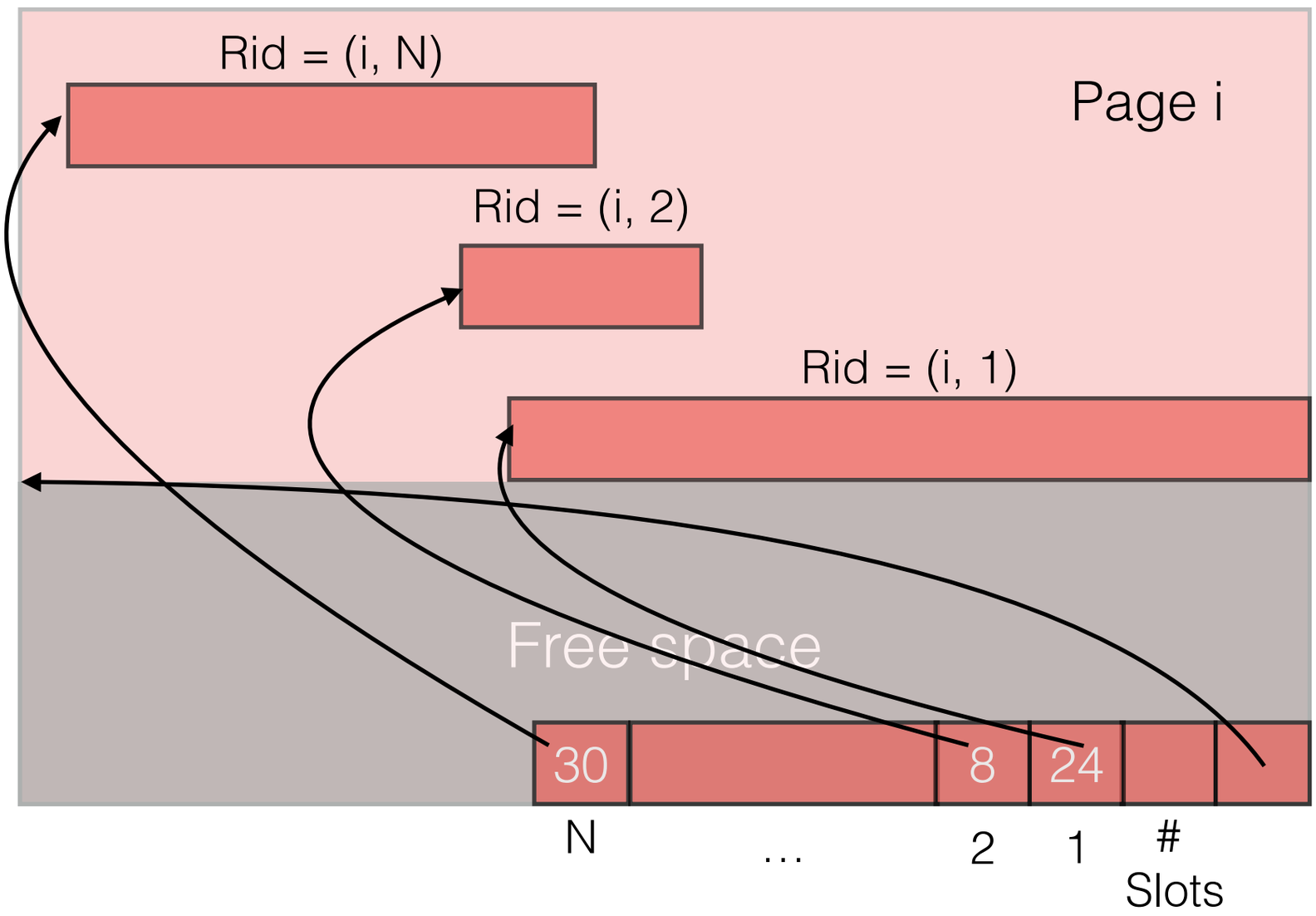
Fixed-Length Records



Fixed-Length Records



Variable Length Records



Record id:
<Page id, slot #>

Slotted Page

What are the advantages and disadvantages of using slotted pages or bitmaps over just tightly packing records together?

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- Allow movement of records without changing record ID
- Slotted pages support variable-length records

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

What's the size of the largest record you can insert?

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

What's the size of the largest record you can insert?

Need 4 bytes for the entry, so (80 - 4) = 76 bytes

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

At most, how many 1-byte large records can you insert?

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

At most, how many 1-byte large records can you insert?

- Amount of space taken up by x 1-byte records
= (1 byte for record + 4 for directory entry)
= (5 bytes / record)

- Free space / (amount of space per record)
= 80 / 5 = 16 records