

# CS186 Discussion #3

(Joins, Heap Files)

# Joins

```
SELECT <columns>
      FROM <tables>
      WHERE <predicate>
[GROUP BY <column list>
[HAVING <predicate>]];
```

```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
         year_released, genre)
```

```
SELECT *
FROM Artists, Albums
WHERE Artists.artist_id =
      Albums.artist_num;
```

```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
         year_released, genre)
```

```
SELECT *
FROM Artists A, Albums B
WHERE A.artist_id = B.artist_num;
```

```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
          year_released, genre)
```

Write a SQL expression for the following query:

The name of all songs with the genre “country” which have spent more than 2 weeks in the top 40.

```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
         year_released, genre)
```

Write a SQL expression for the following query:

The name of all songs with the genre “country” which have spent more than 2 weeks in the top 40.

```
SELECT Songs.song_name FROM Albums,
Songs WHERE Songs.album_num =
Albums.album_id AND Albums.genre =
'country' AND Songs.weeks_in_top_40 > 2;
```

```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
         year_released, genre)
```

Write a SQL expression for the following query:

The number of albums released by each artist.



```
Songs (song_id, song_name, album_num, weeks_in_top_40)
Artists (artist_id, artist_name, first_year_active)
Albums (album_id, album_name, artist_num,
         year_released, genre)
```

Write a SQL expression for the following query:

The number of albums released by each artist.

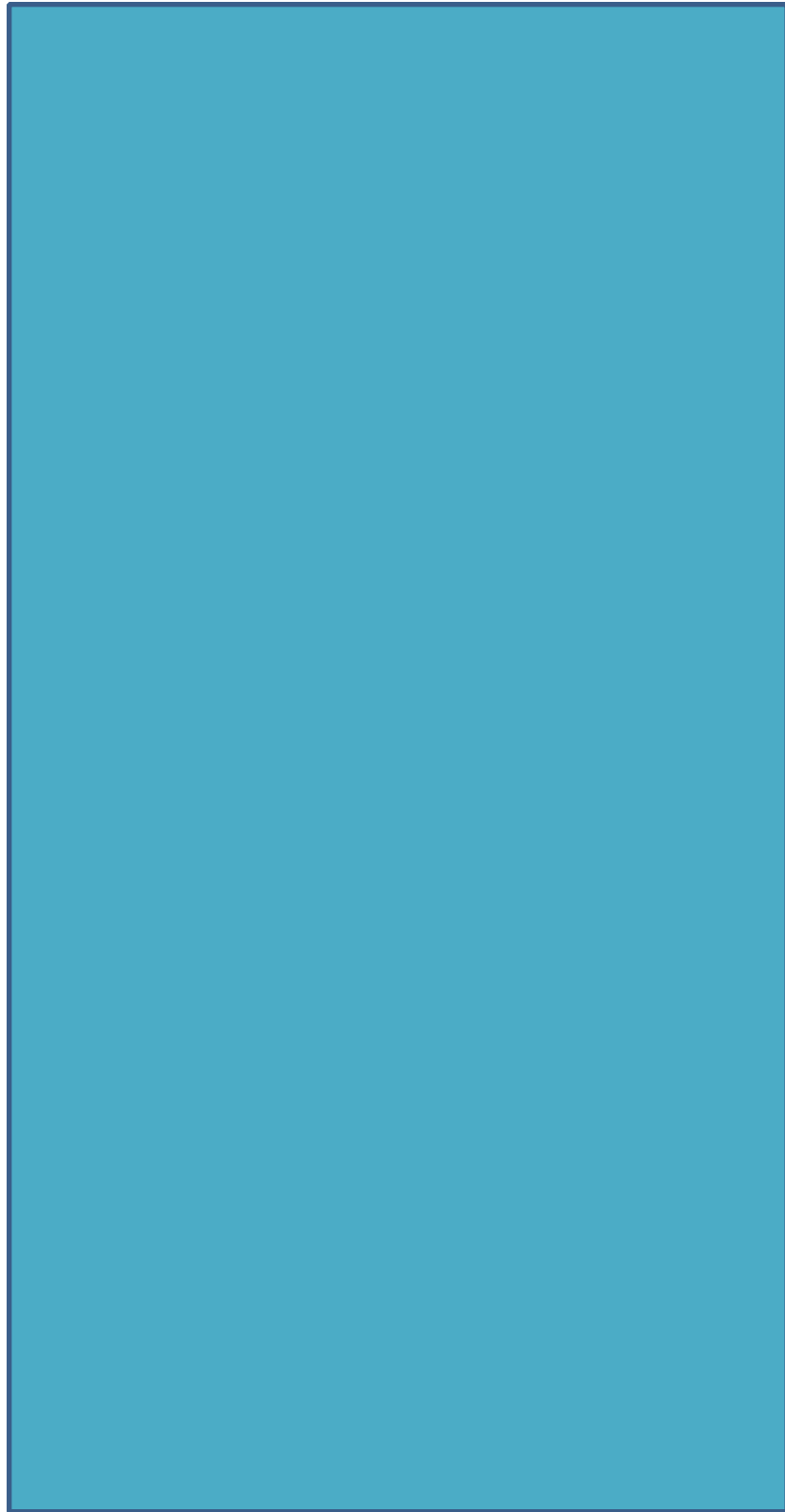
```
SELECT count(*) FROM Artists, Albums
WHERE Artists.artist_id =
Albums.artist_num GROUP BY
Artists.artist_id;
```

# Join Algorithms

```
SELECT * FROM Sailors S, Reserves R
      WHERE S.sid = R.sid;
```

# Visualizations

Sailors



# Visualizations

Sailors

Page 1

Page 2

Page 3

Page 4

# Visualizations

Sailors

Record 1
Record 2
Record 3
Record 4
Record 5

Page 2

Page 3

Page 4

# Visualizations

## Sailors

## Reserves

## Record 1

## Record 2

## Record 3

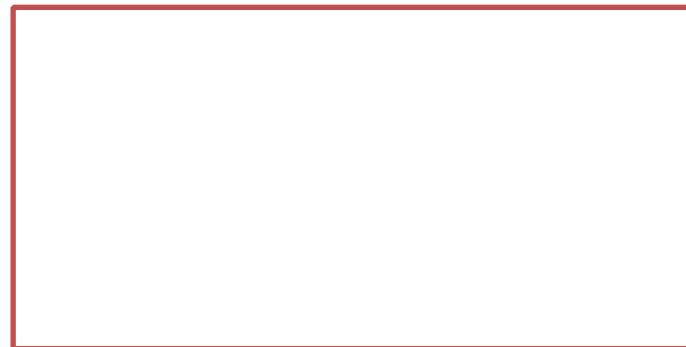
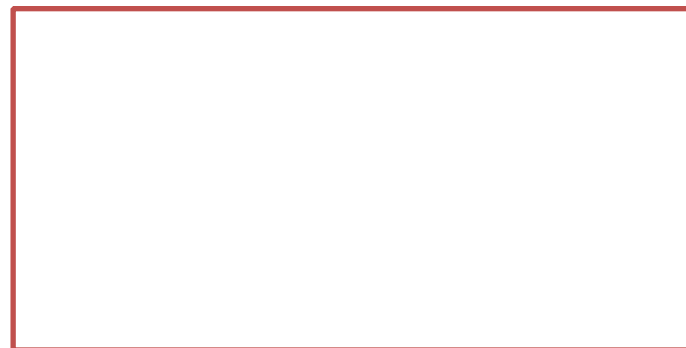
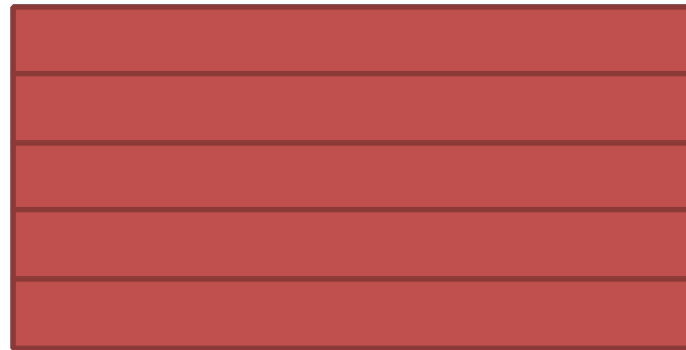
## Record 4

## Record 5

Page 2

Page 3

Page 4



# Simple Nested Loops Join

Sailors

Reserves



**Key idea:**

Take each record of S and match it with each record of R.

**Steps:**

1. Get tuple of S.
2. Iterate through each tuple in R.



# Simple Nested Loops Join

Sailors

Reserves

**Key idea:**

Take each record of S and match it with each record of R.

**Steps:**

1. Get tuple of S.
2. Iterate through each tuple in R.

(name = Bob, sid = 1)

# Simple Nested Loops Join

Sailors

Reserves

(name = Bob, sid = 1)

(sid = 3, bid = 6)
(sid = 1, bid = 4)

## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

## Output:

(name = Bob, sid = 1, bid = 4)

# Simple Nested Loops Join

Sailors

Reserves

(name = Bob, sid = 1)

(sid = 3, bid = 6)
(sid = 1, bid = 4)
(sid = 1, bid = 7)

## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

## Output:

(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)

# Simple Nested Loops Join

Sailors

Reserves

(name = Bob, sid = 1)

(sid = 3, bid = 6)
(sid = 1, bid = 4)
(sid = 1, bid = 7)

## Key idea:

Take each record of S and match it with each record of R.

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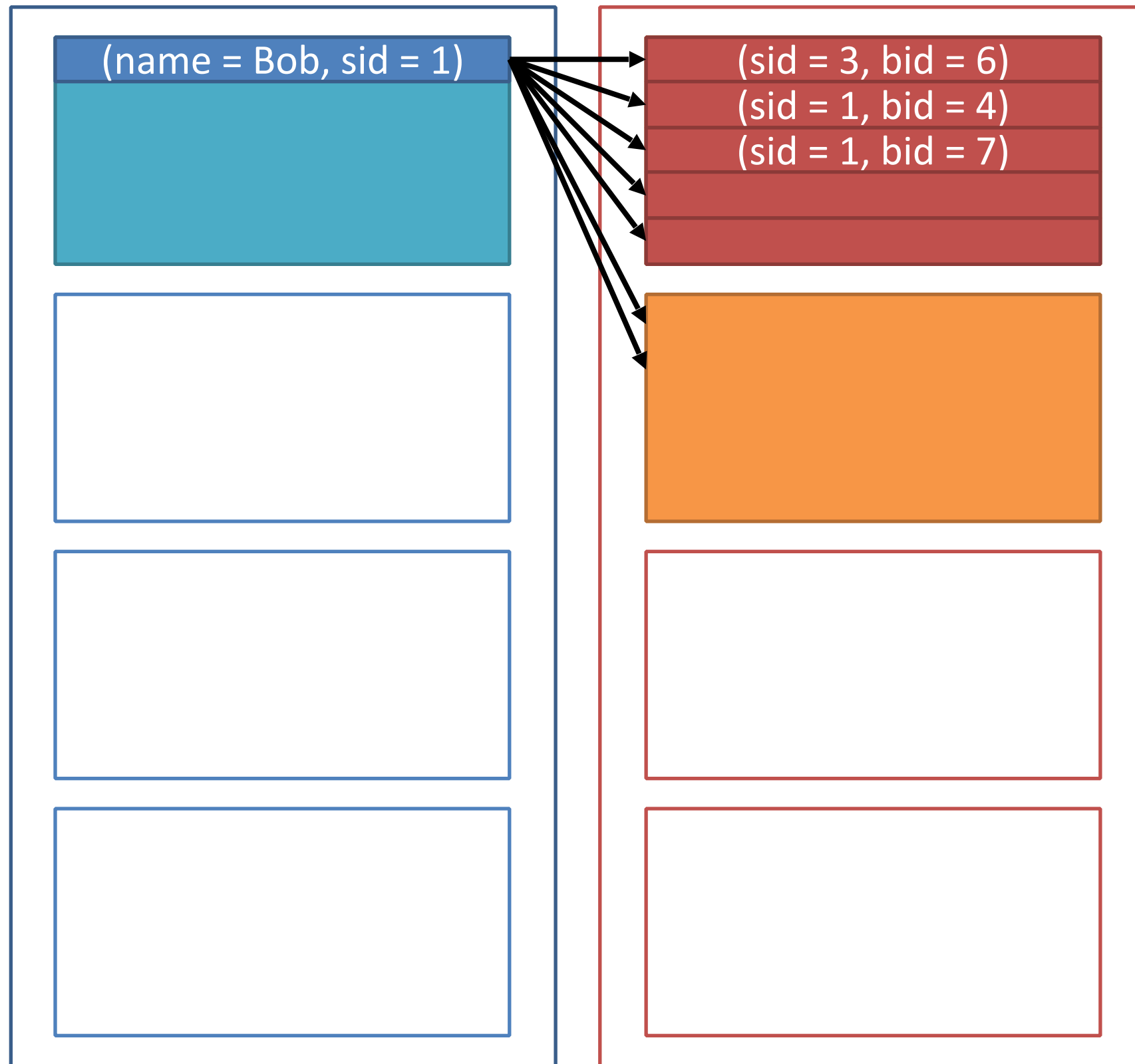
## Output:

(name = Bob, sid = 1, bid = 4)  
(name = Bob, sid = 1, bid = 7)

# Simple Nested Loops Join

Sailors

Reserves



## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

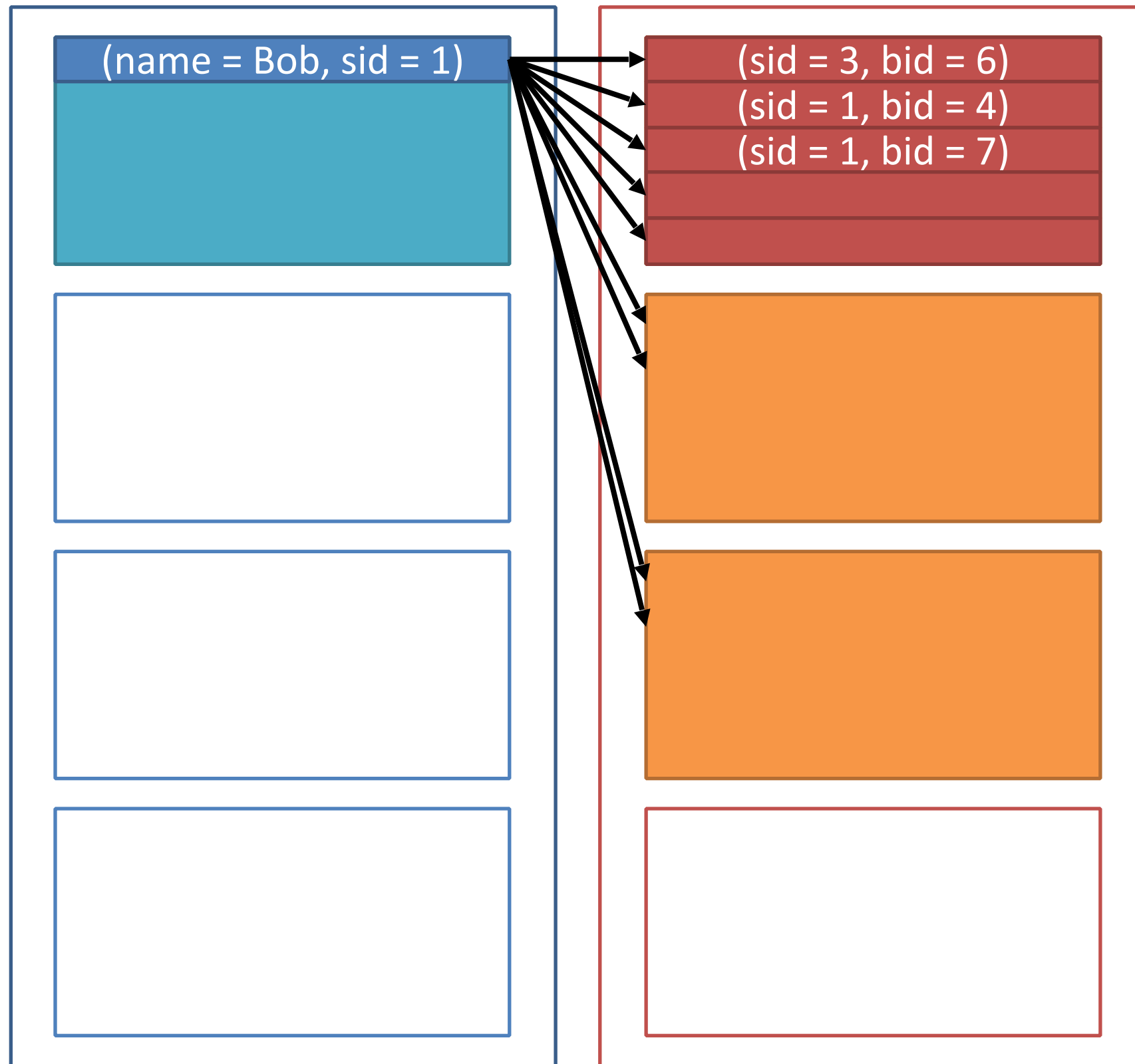
## Output:

`(name = Bob, sid = 1, bid = 4)`  
`(name = Bob, sid = 1, bid = 7)`

# Simple Nested Loops Join

Sailors

Reserves



## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

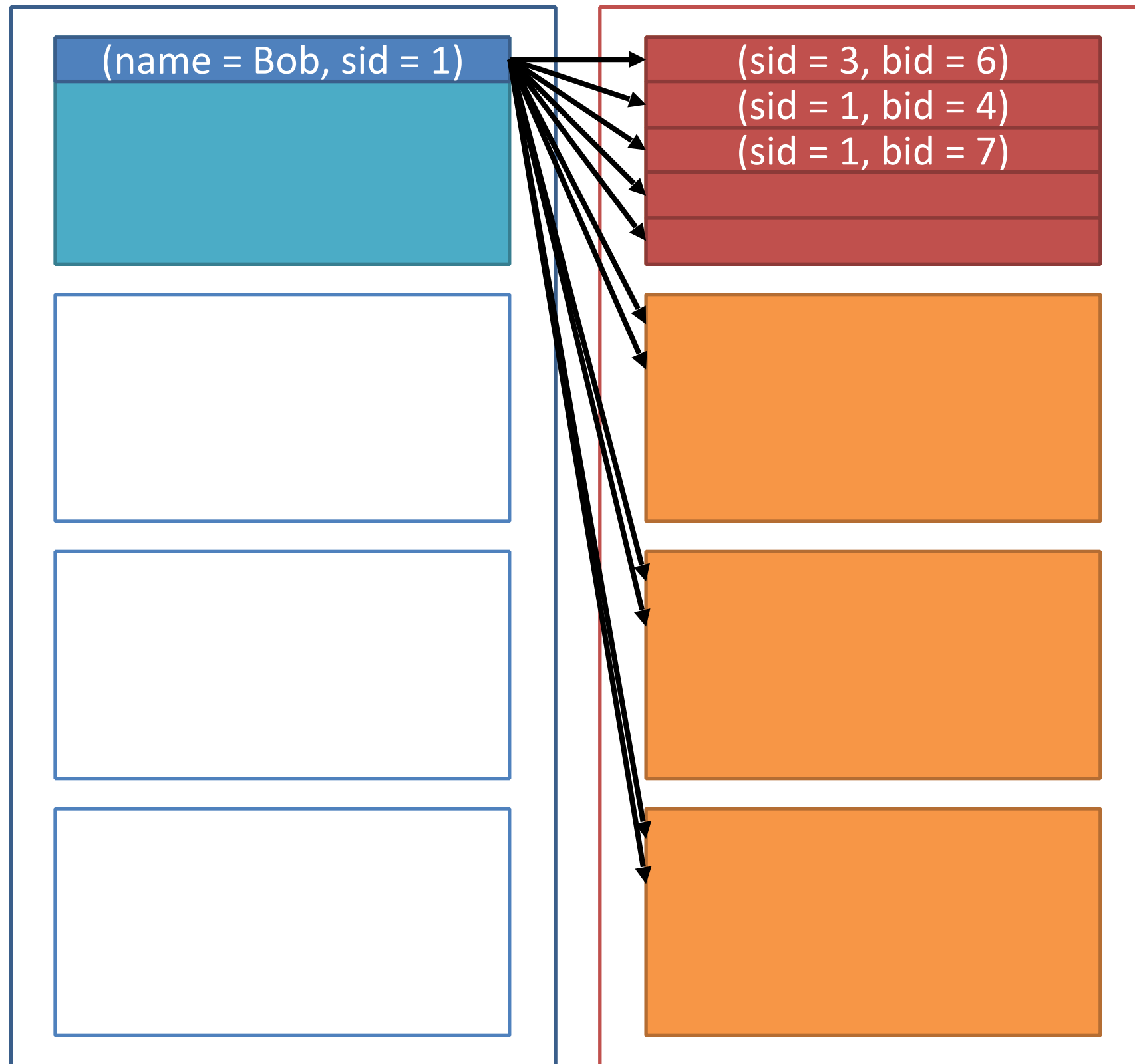
## Output:

(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)

# Simple Nested Loops Join

Sailors

Reserves



## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

## Output:

(name = Bob, sid = 1, bid = 4)  
(name = Bob, sid = 1, bid = 7)

# Simple Nested Loops Join

Sailors

Reserves

## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

## Output:

(name = Bob, sid = 1, bid = 4)  
(name = Bob, sid = 1, bid = 7)



# Simple Nested Loops Join

Sailors

Reserves

(name = Bob, sid = 1)
(name = Sam, sid = 3)

(sid = 3, bid = 6)
(sid = 1, bid = 4)
(sid = 1, bid = 7)

## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

## Output:

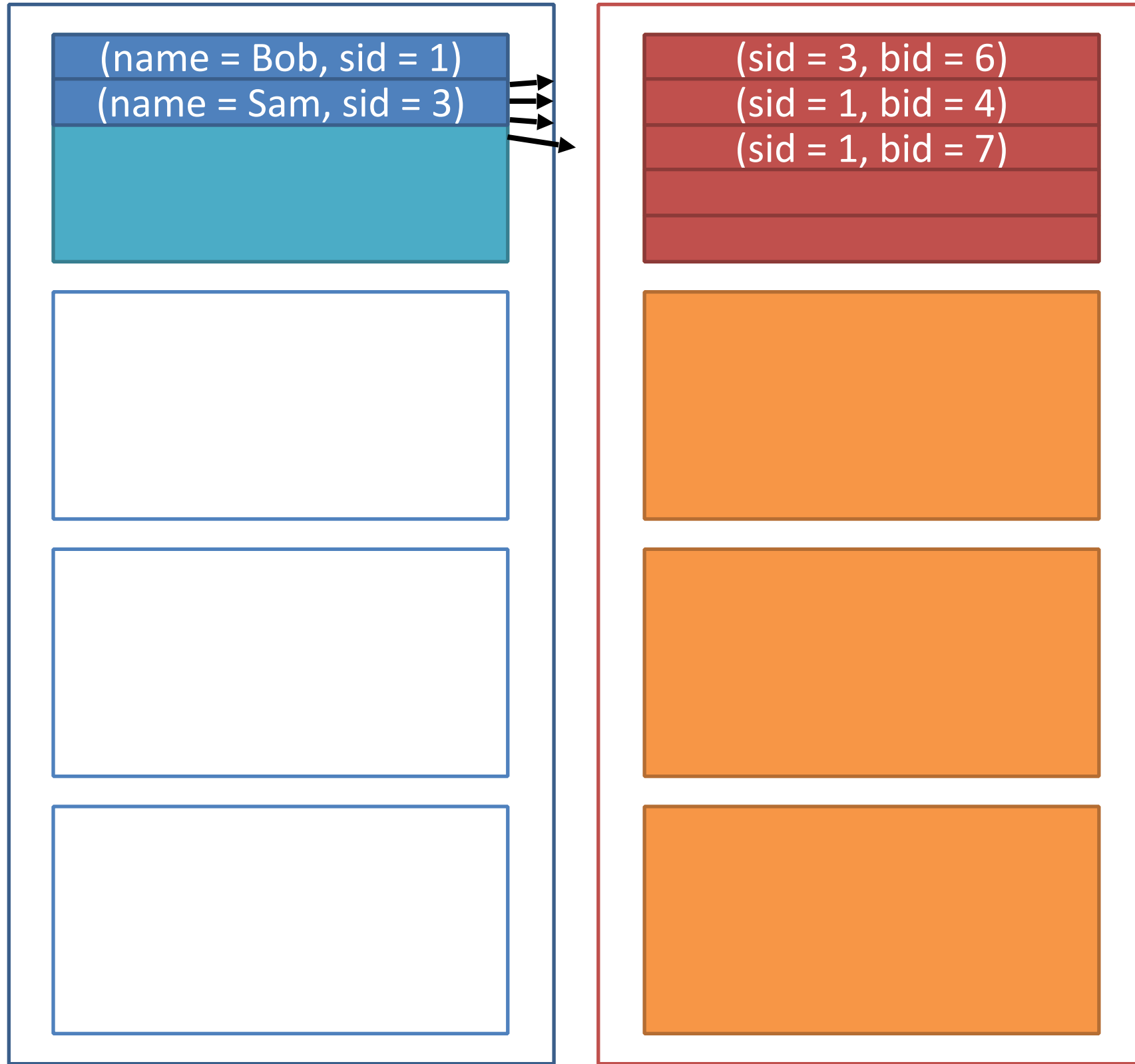
(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)
(name = Sam, sid = 3, bid = 6)

**Notation:  $[S]$  == “# pages in  $S$ ” ;  $|S|$  == “# tuples in  $S$ ”**

# Simple Nested Loops Join

# Sailors

## Reserves



## Key idea:

Take each record of S and match it with each record of R.

## Steps:

1. Get tuple of S.
2. Iterate through each tuple in R.

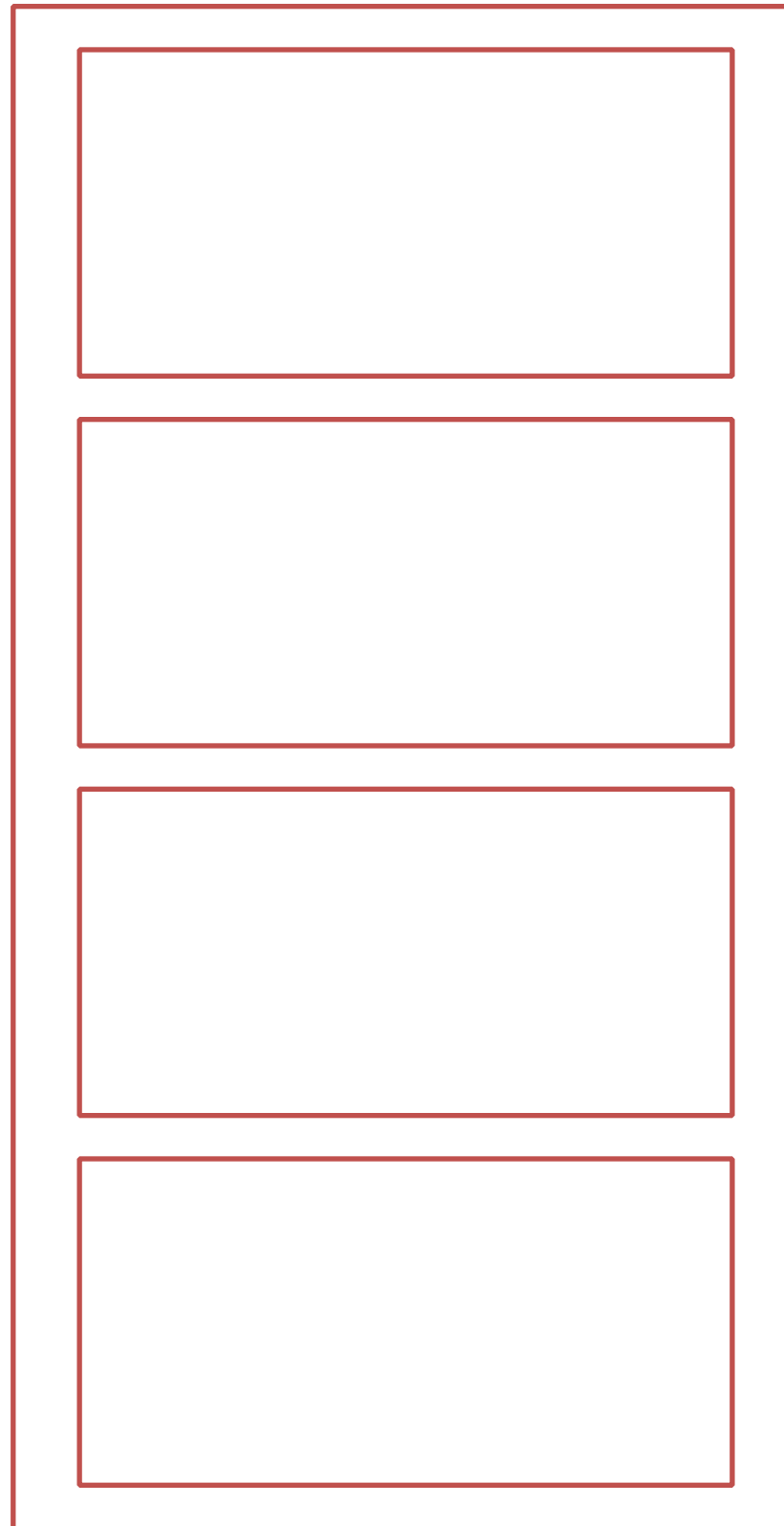
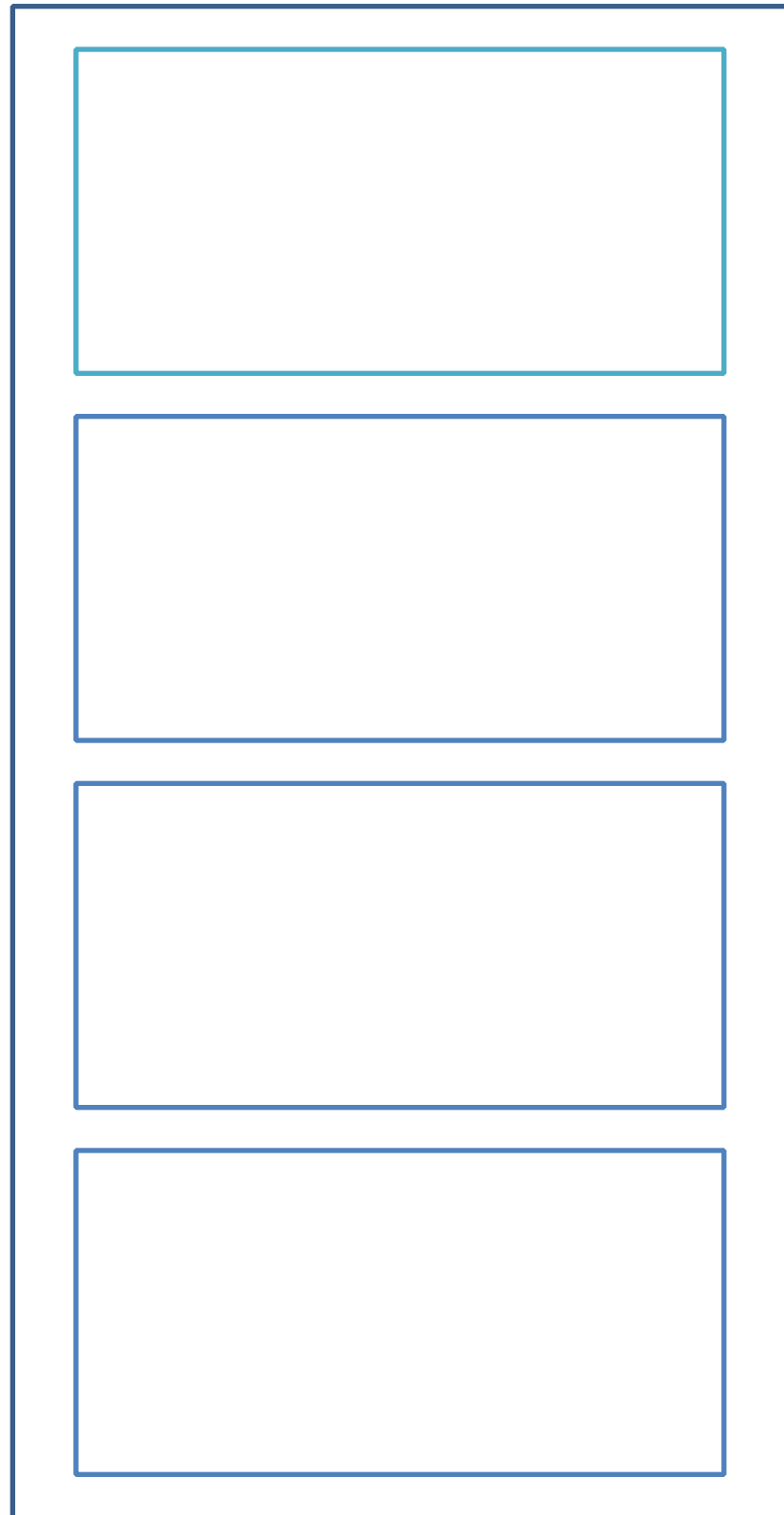
## I/Os:

$$[S] + |S| * [R]$$

# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

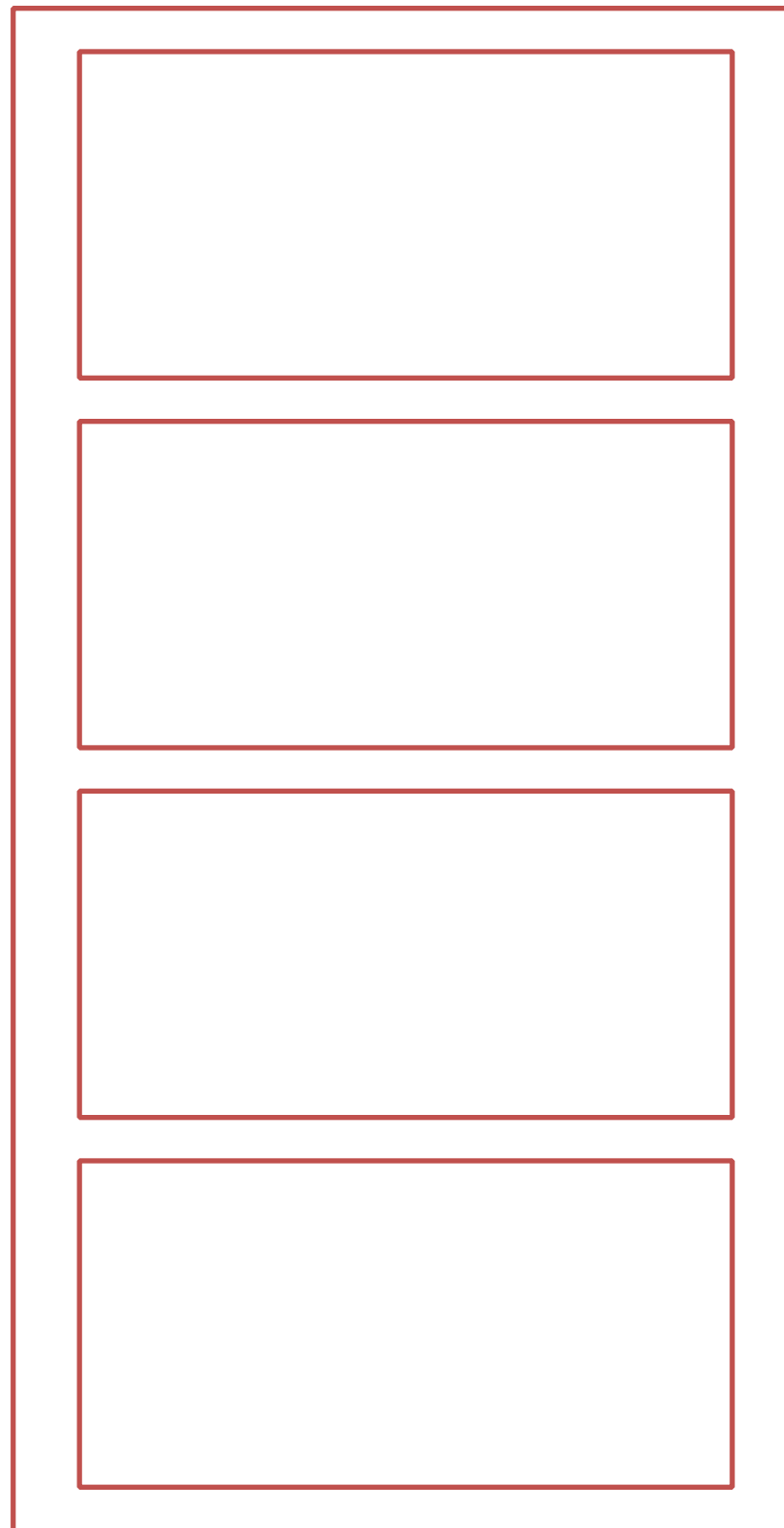
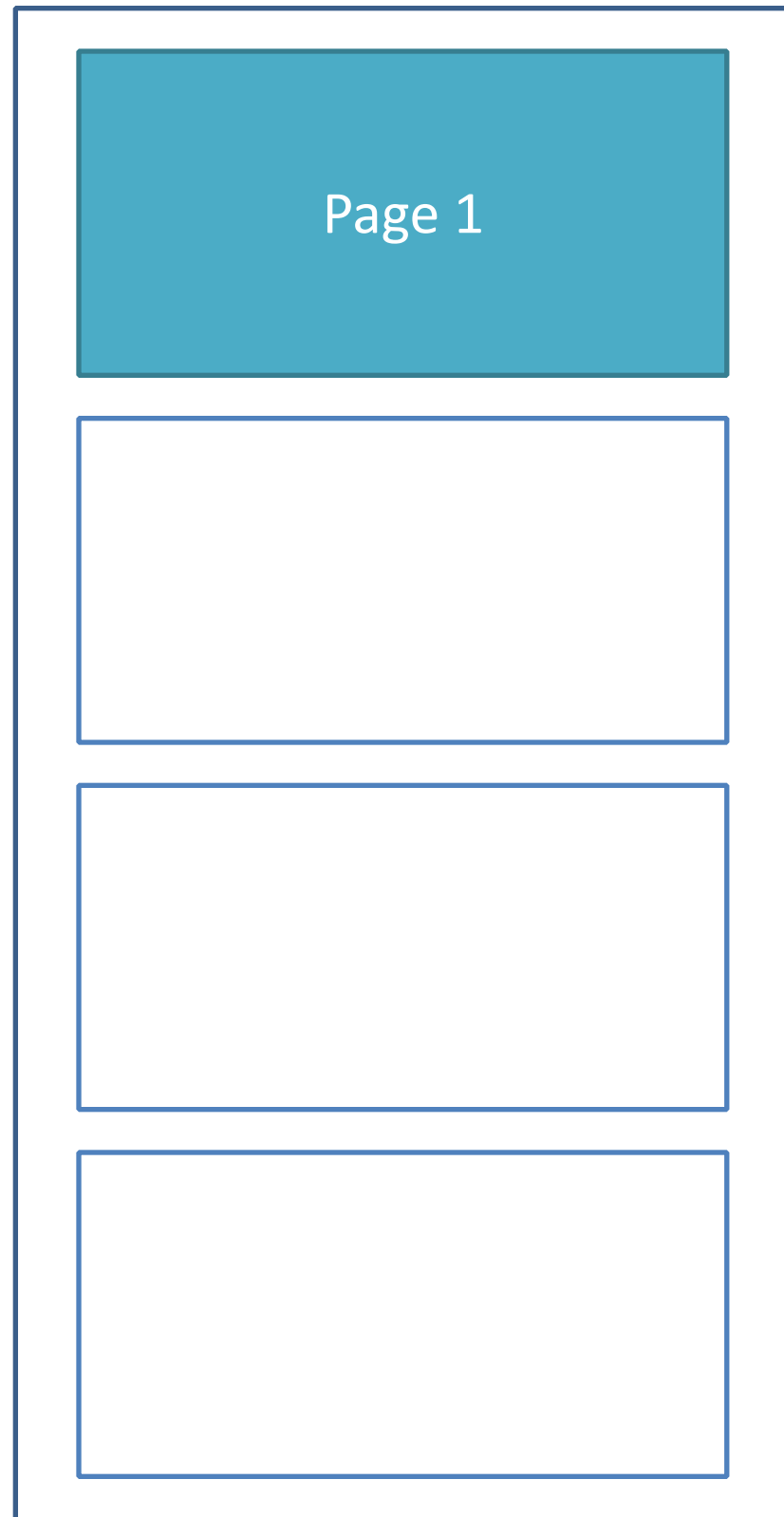
## Steps:

1. Get page of S.
2. Iterate through each page in R.
3. Compare tuples in each.

# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

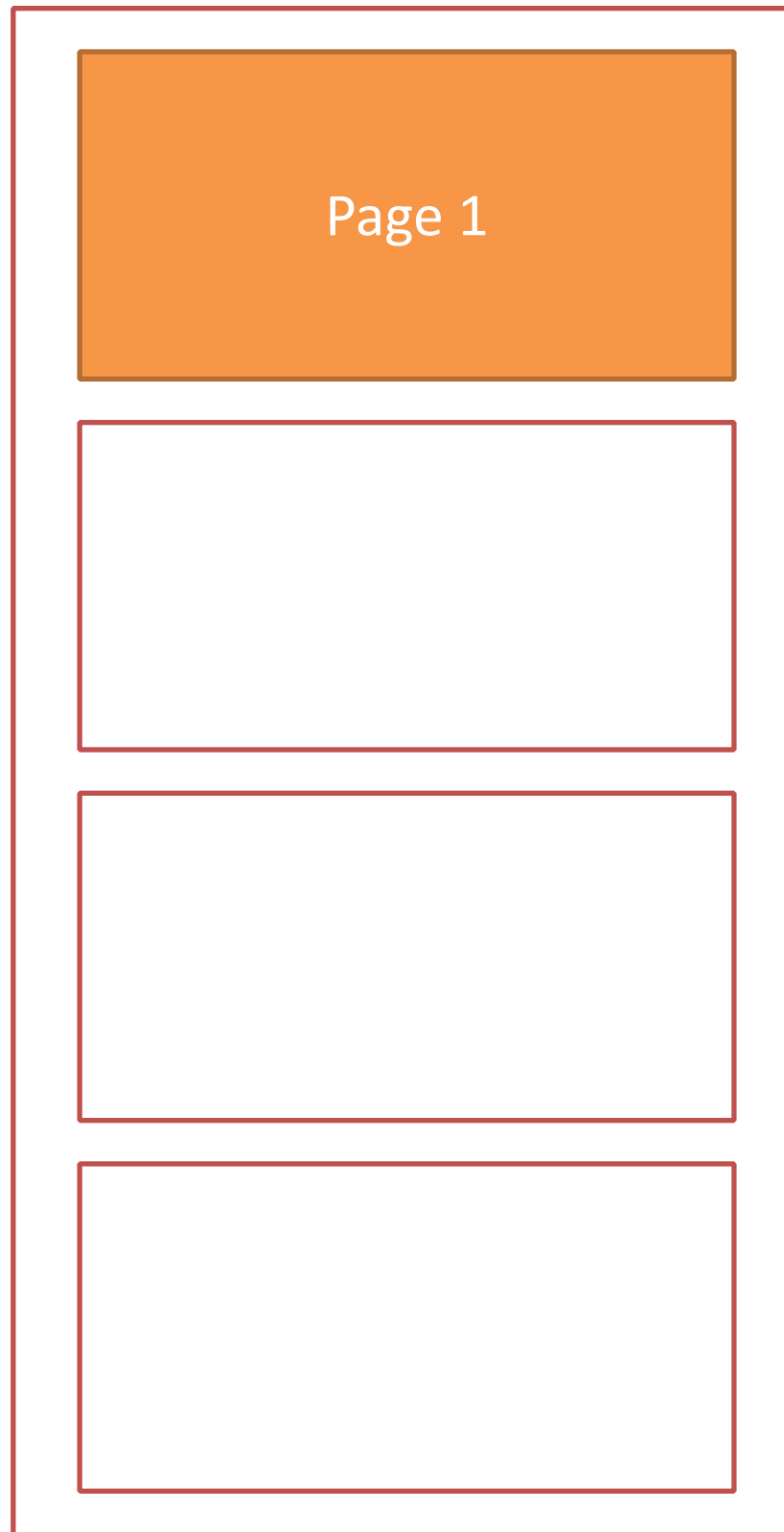
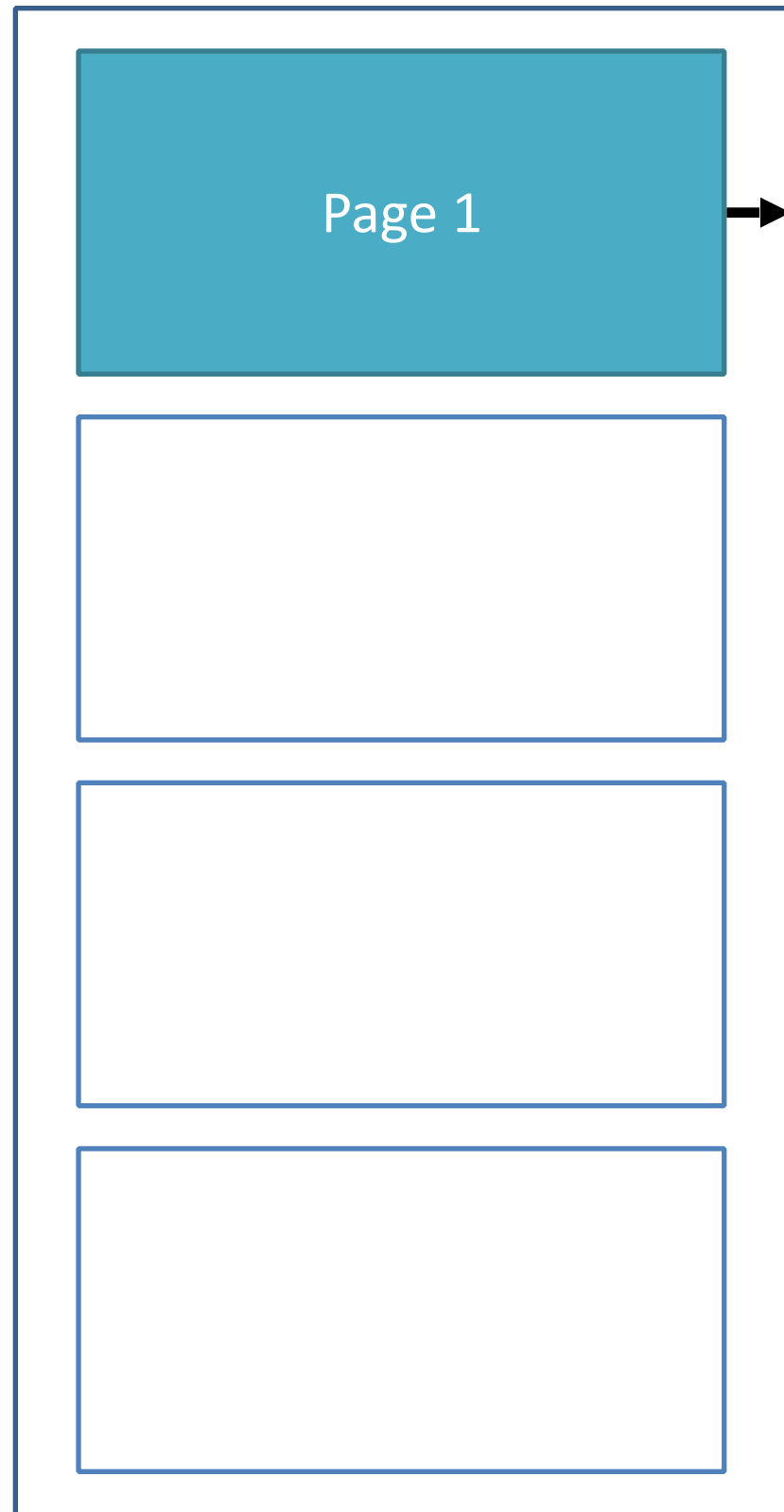
## Steps:

1. Get page of S.
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# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

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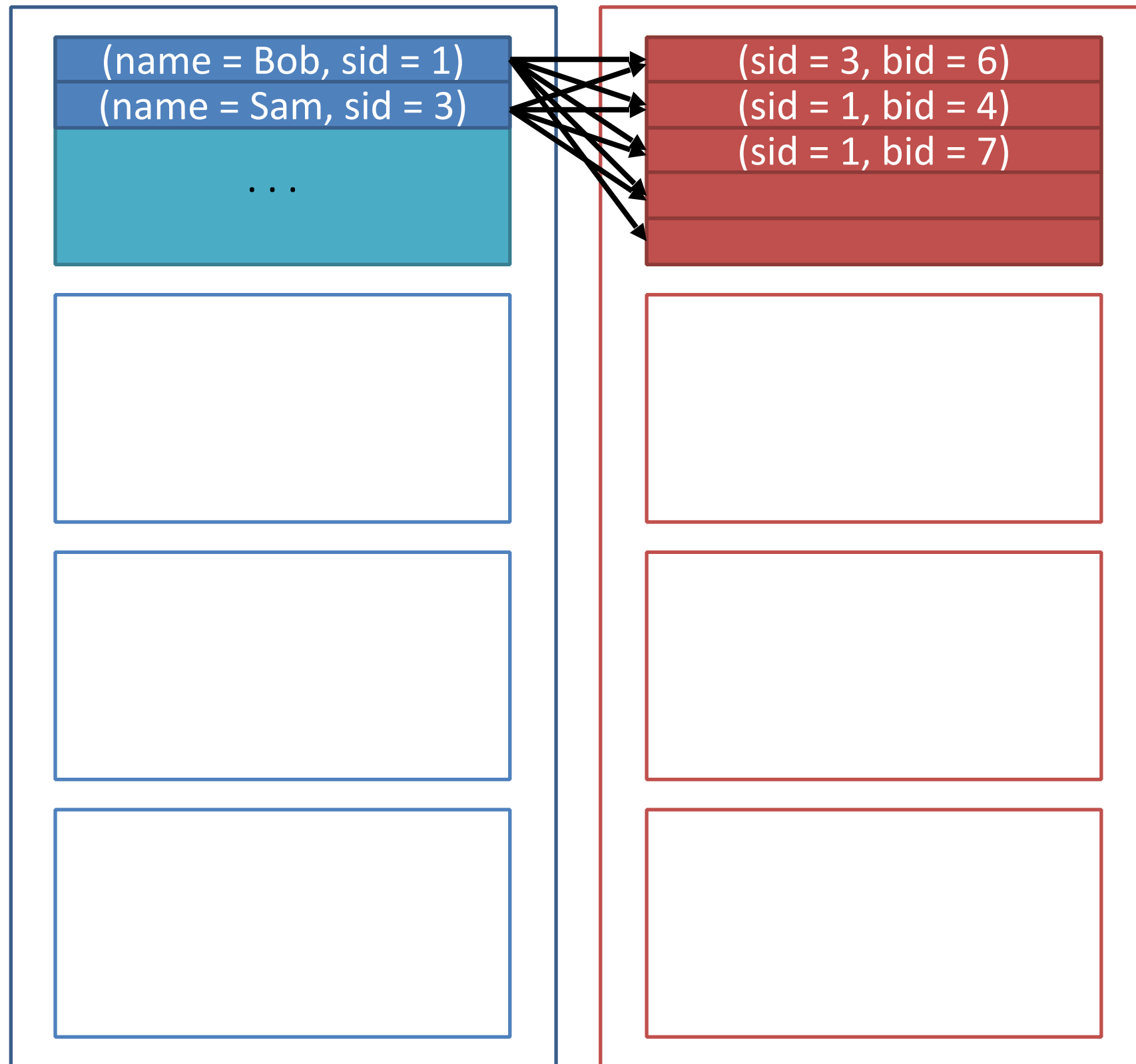
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# Page-Oriented Nested Loops Join

Sailors

Reserves



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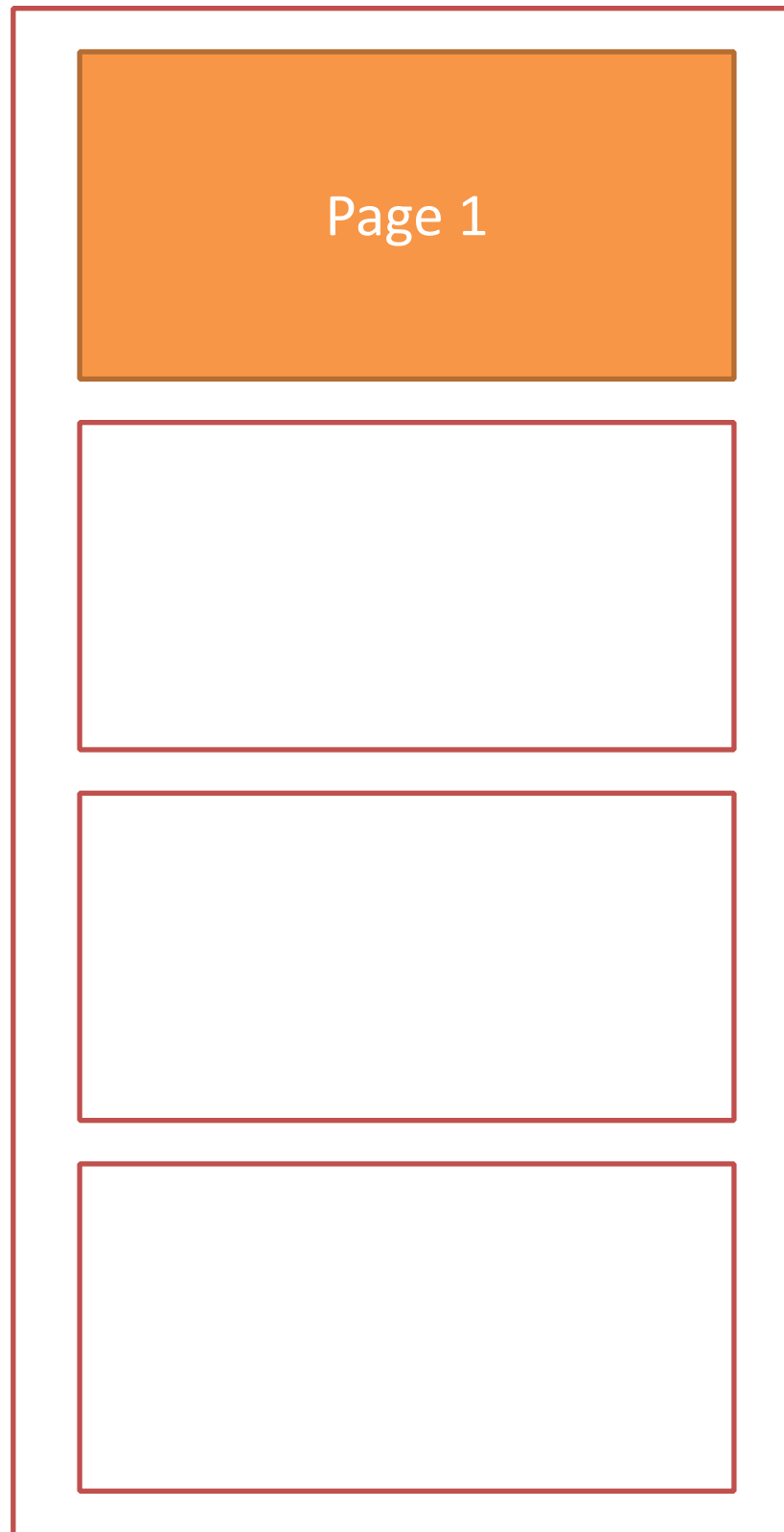
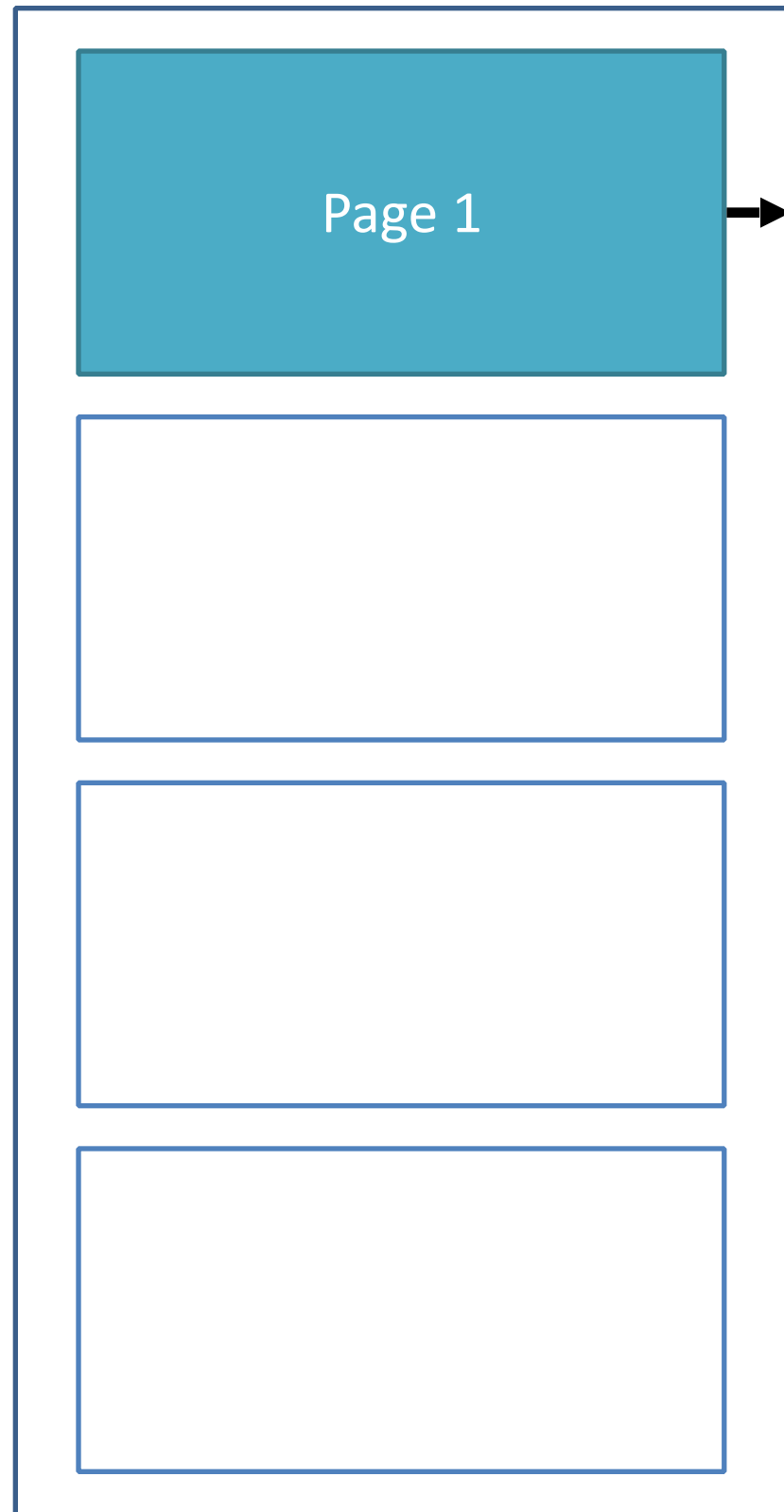
## Output:

`(name = Bob, sid = 1, bid = 4)`  
`(name = Bob, sid = 1, bid = 7)`  
`(name = Sam, sid = 3, bid = 6)`

# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

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1. Get page of S.
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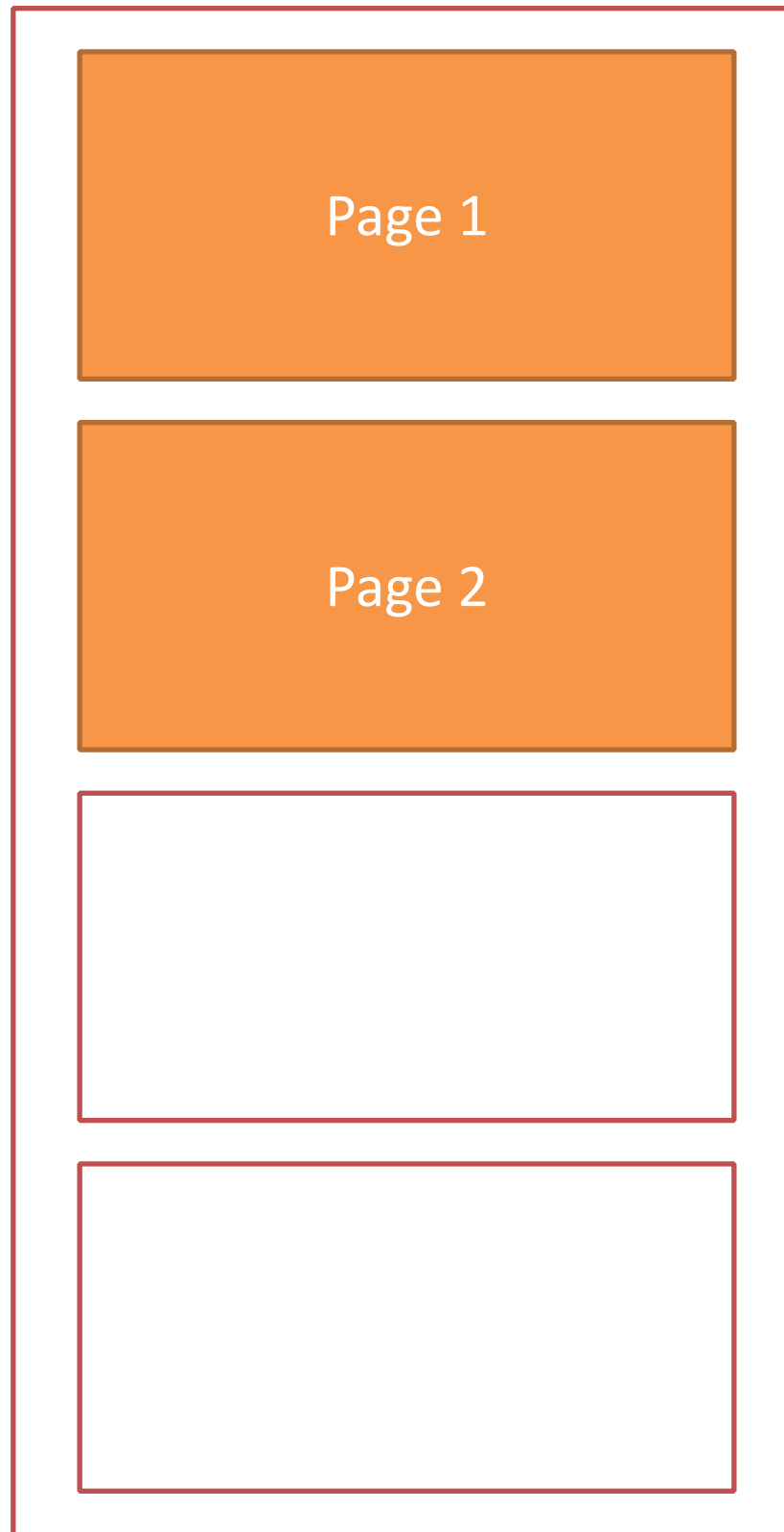
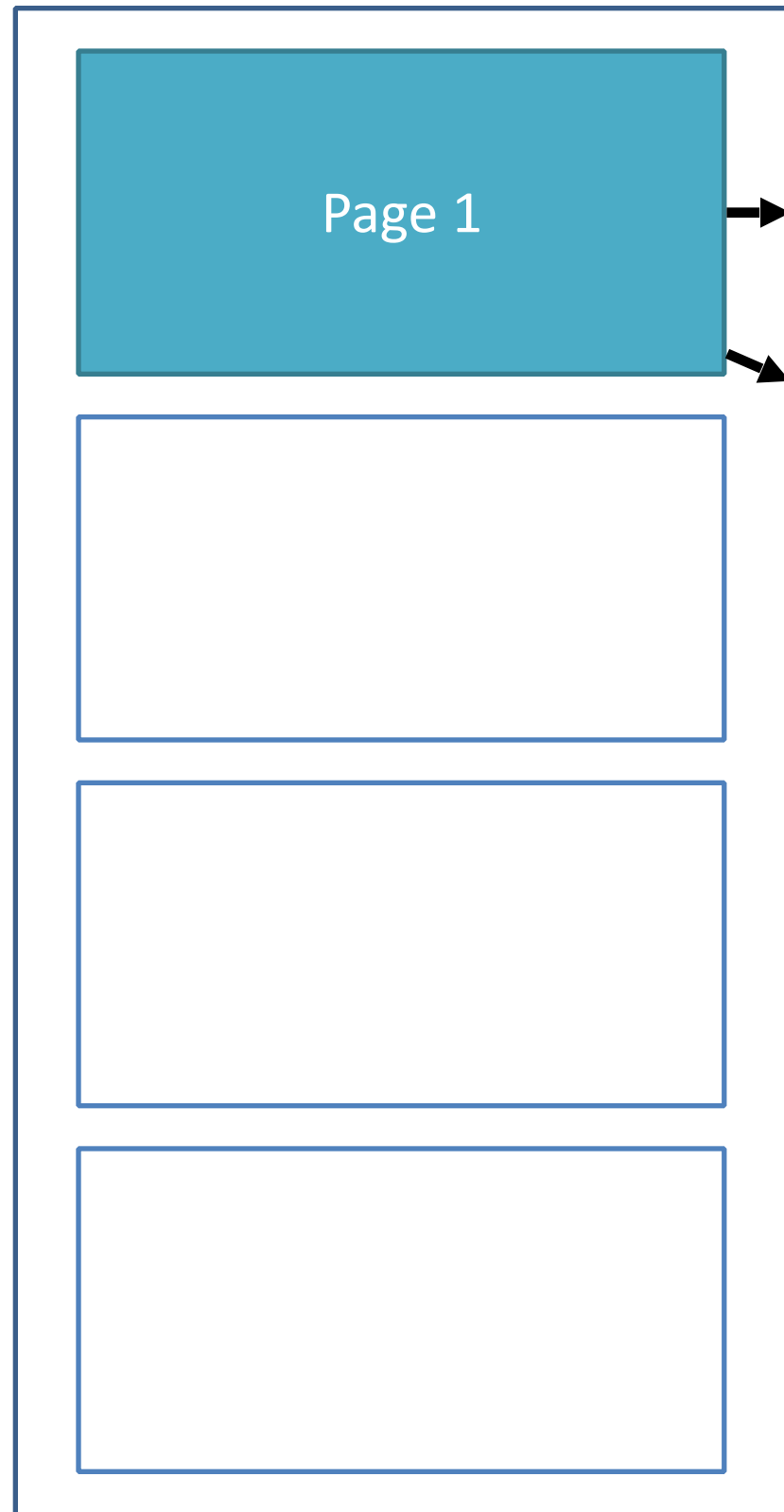
## Output:

(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)
(name = Sam, sid = 3, bid = 6)

# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

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## Output:

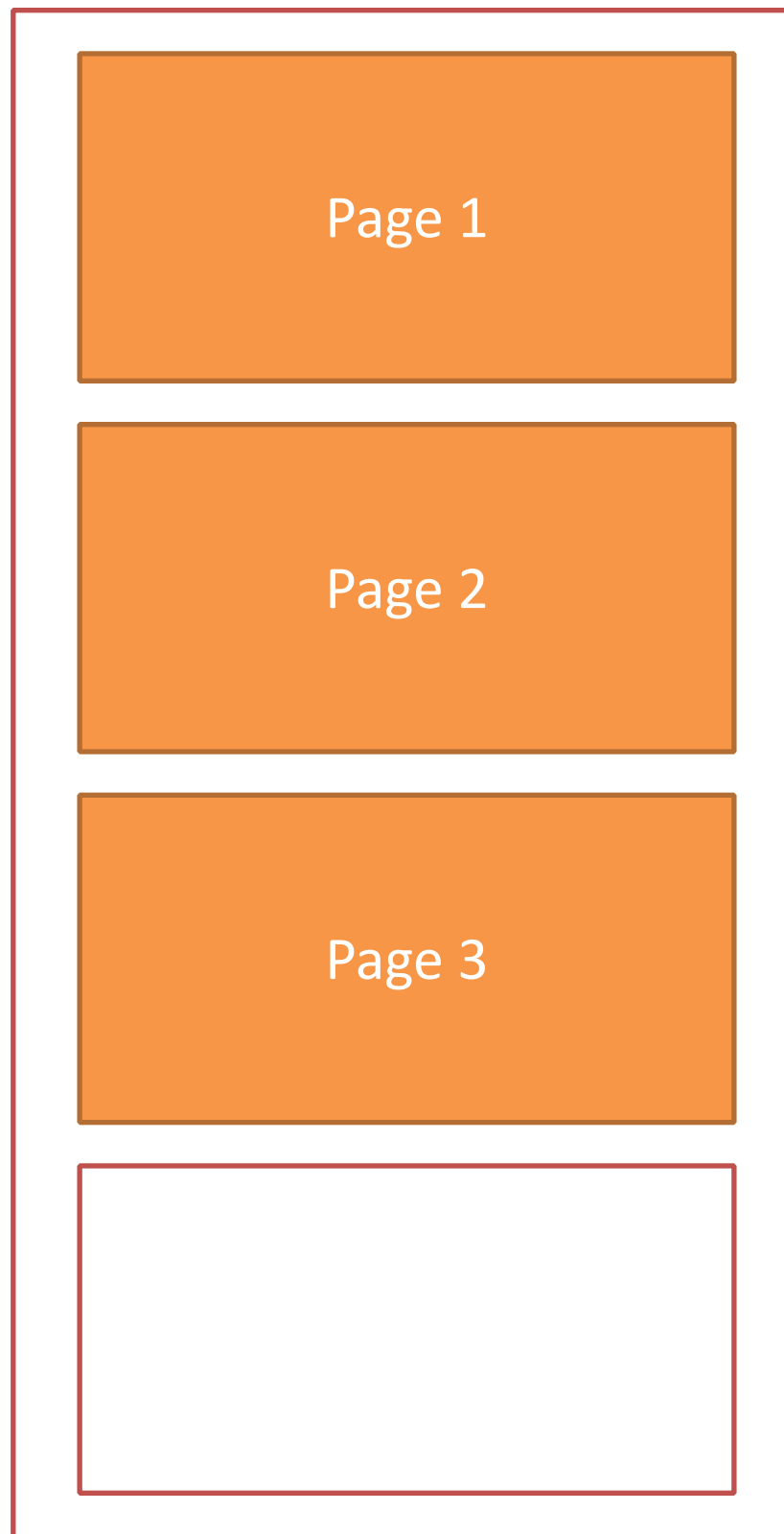
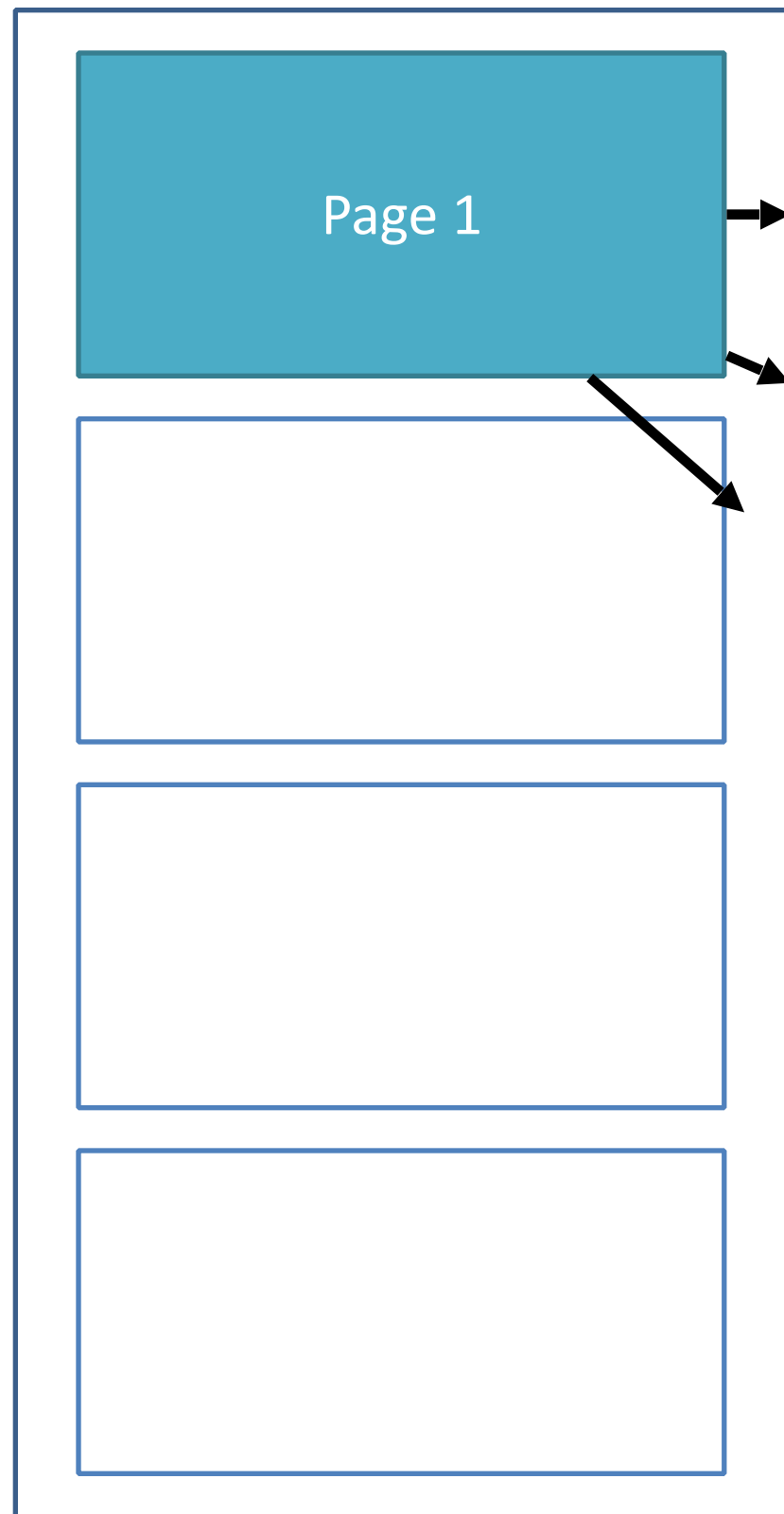
(name = Bob, sid = 1, bid = 4)
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# Page-Oriented Nested Loops Join

Sailors

Reserves



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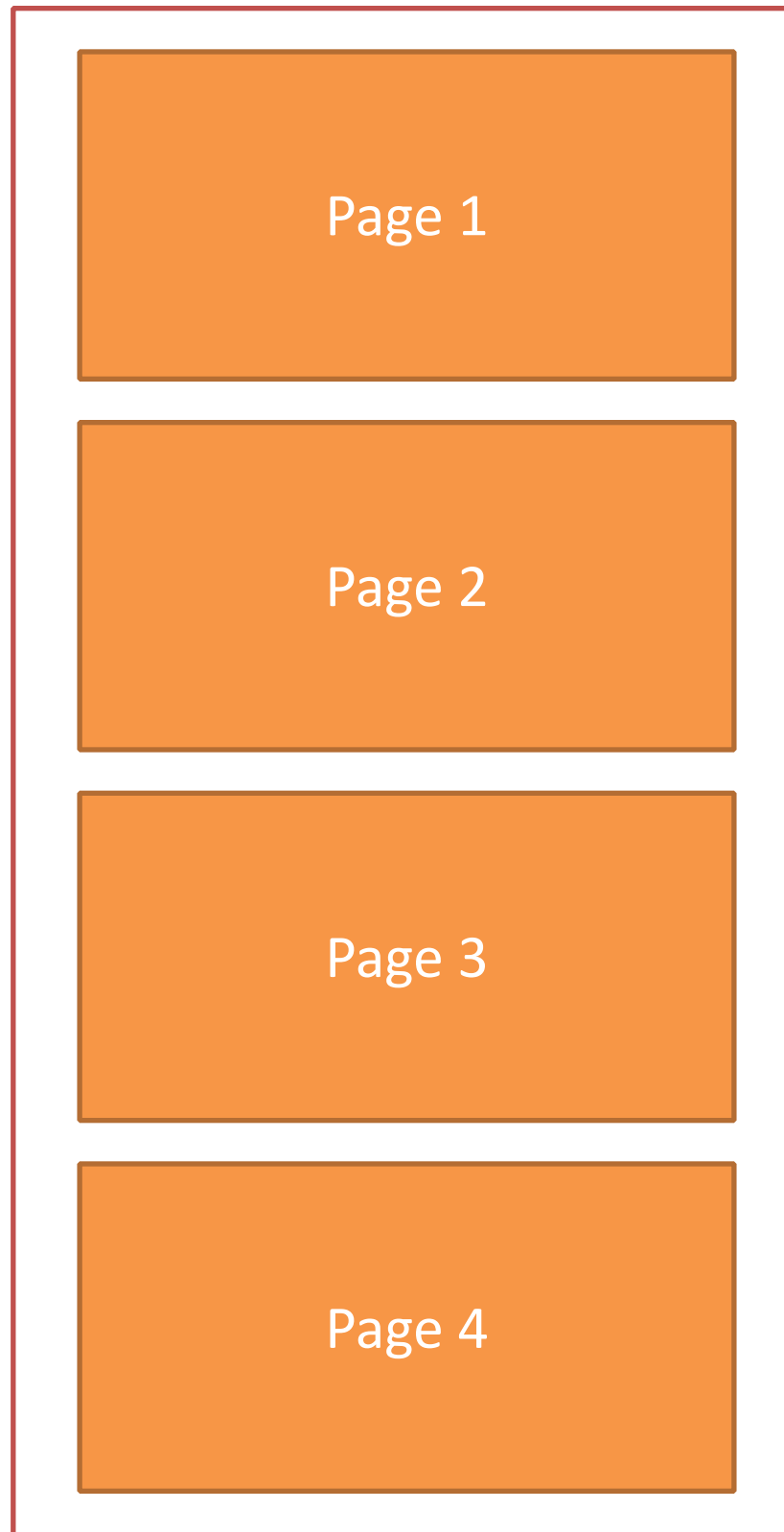
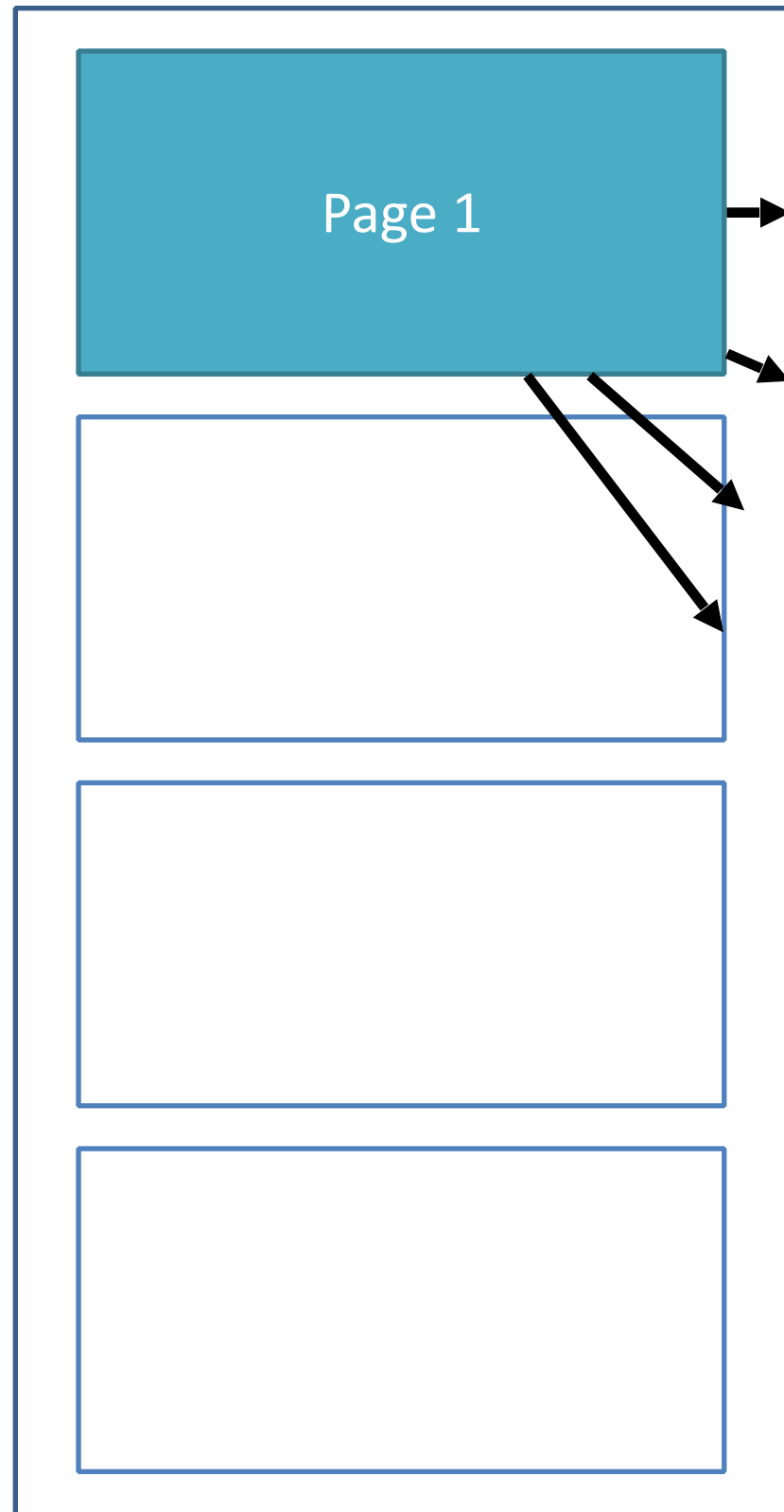
## Output:

(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)
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# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

## Steps:

1. Get page of S.
2. Iterate through each page in R.
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## Output:

(name = Bob, sid = 1, bid = 4)
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# Page-Oriented Nested Loops Join

Sailors

Reserves

Page 1

Page 2

## Key idea:

Take each page of S  
and match with each  
page of R.

## Steps:

1. Get page of S.
2. Iterate through each page in R.
3. Compare tuples in each.

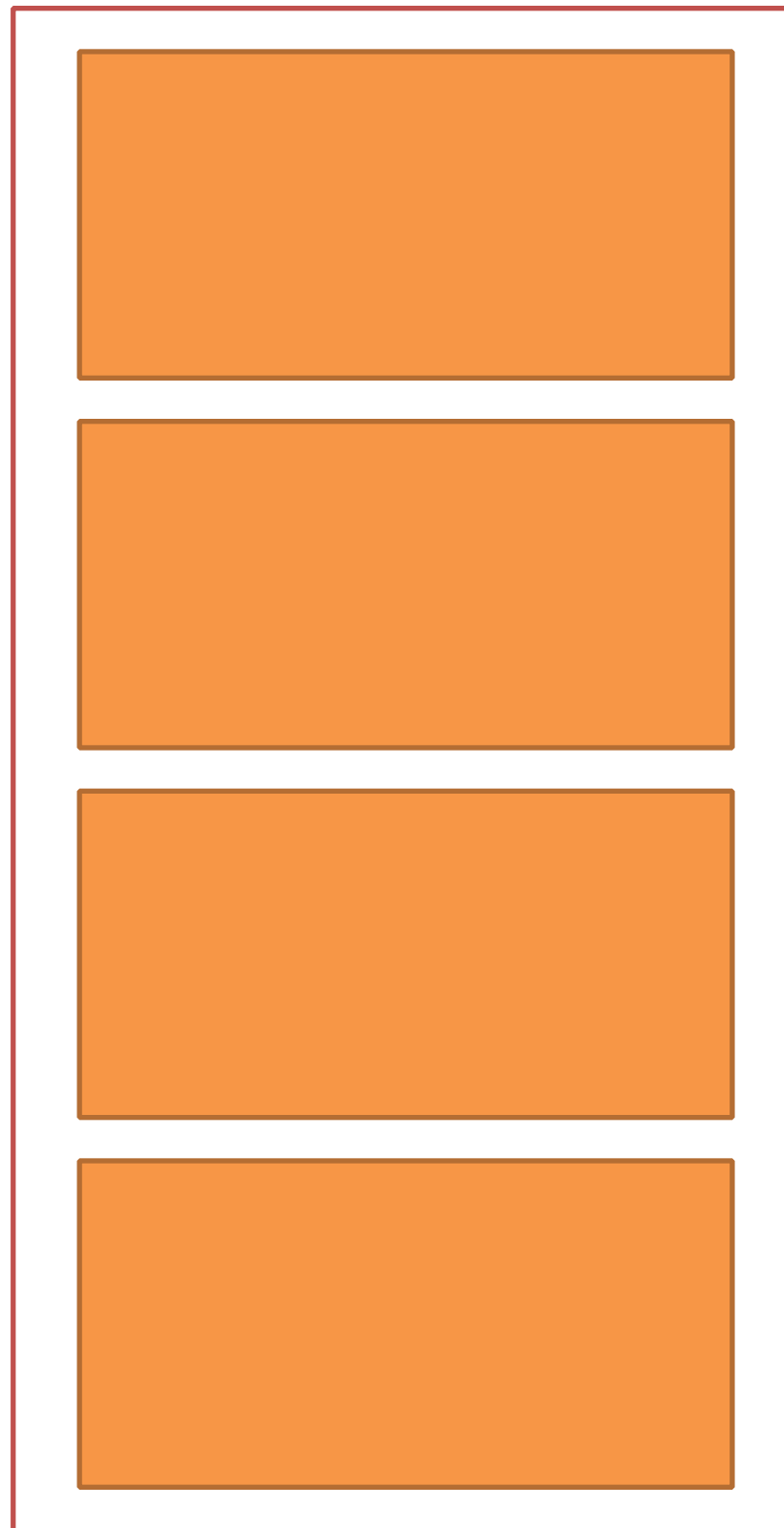
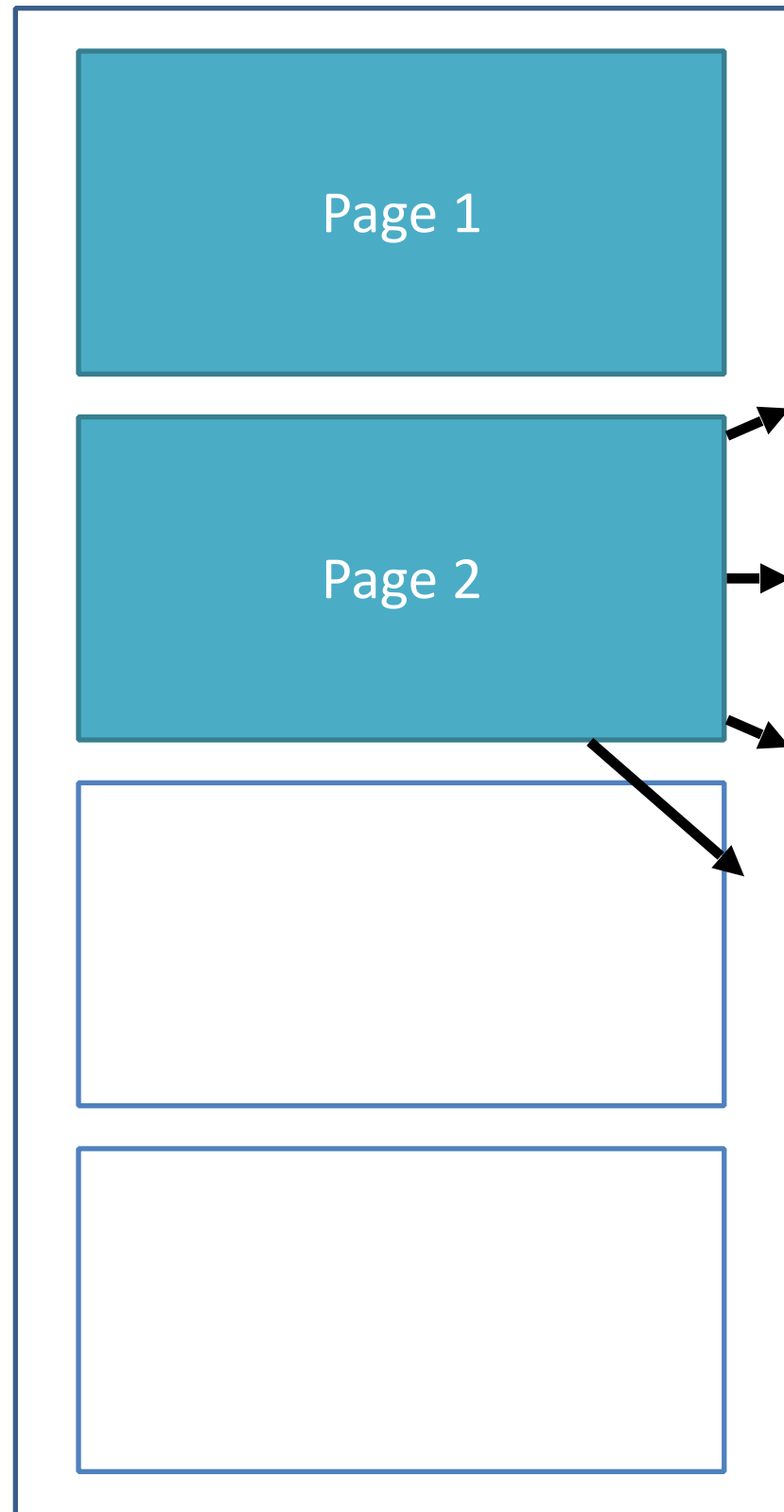
## Output:

(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)
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# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

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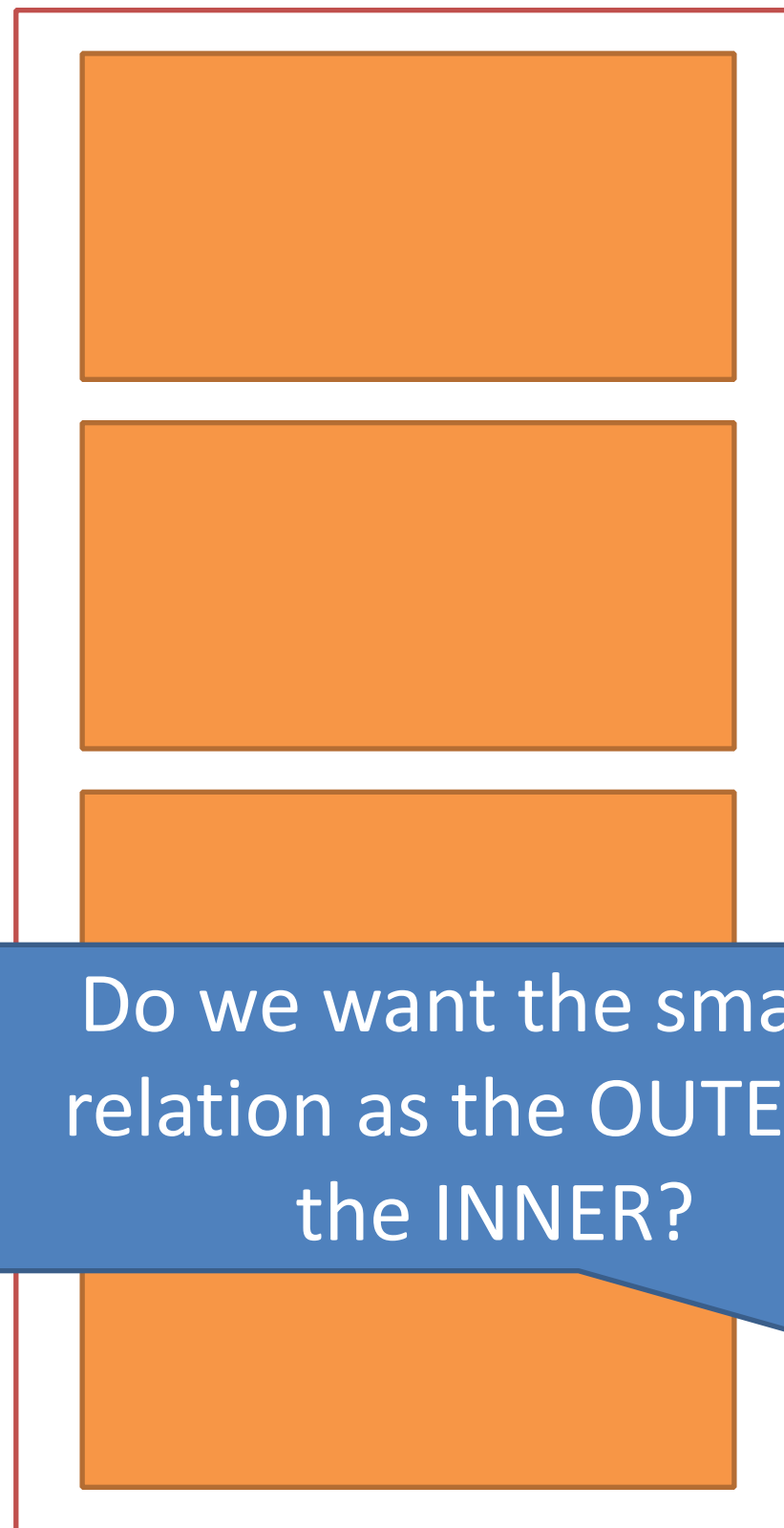
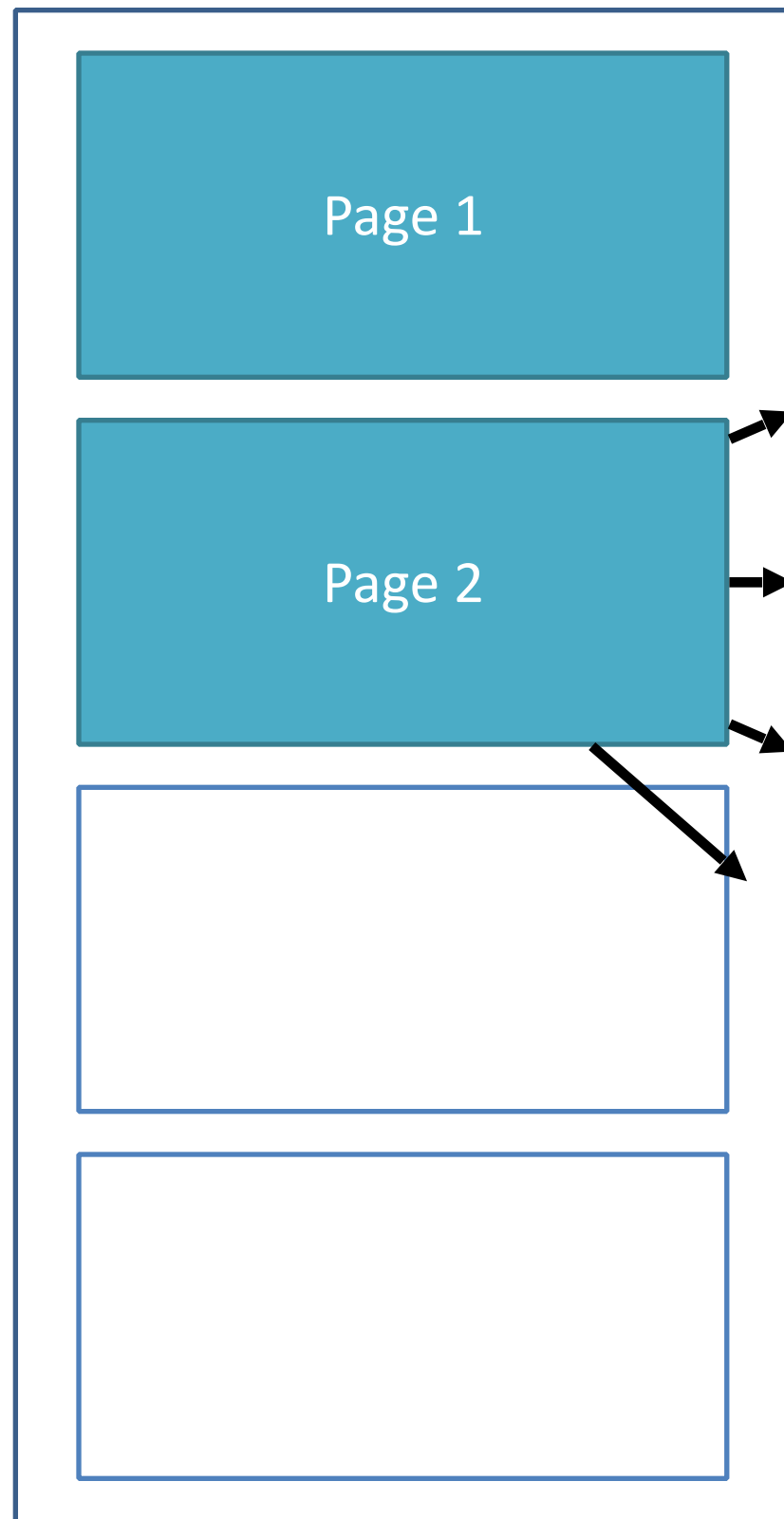
(name = Bob, sid = 1, bid = 4)
(name = Bob, sid = 1, bid = 7)
(name = Sam, sid = 3, bid = 6)

**Notation:  $[S]$  == “# pages in S” ;  $|S|$  == “# tuples in S”**

# Page-Oriented Nested Loops Join

Sailors

Reserves



## Key idea:

Take each page of S and match with each page of R.

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1. Get page of S.
2. Iterate through each page in R.

Compare tuples in each.

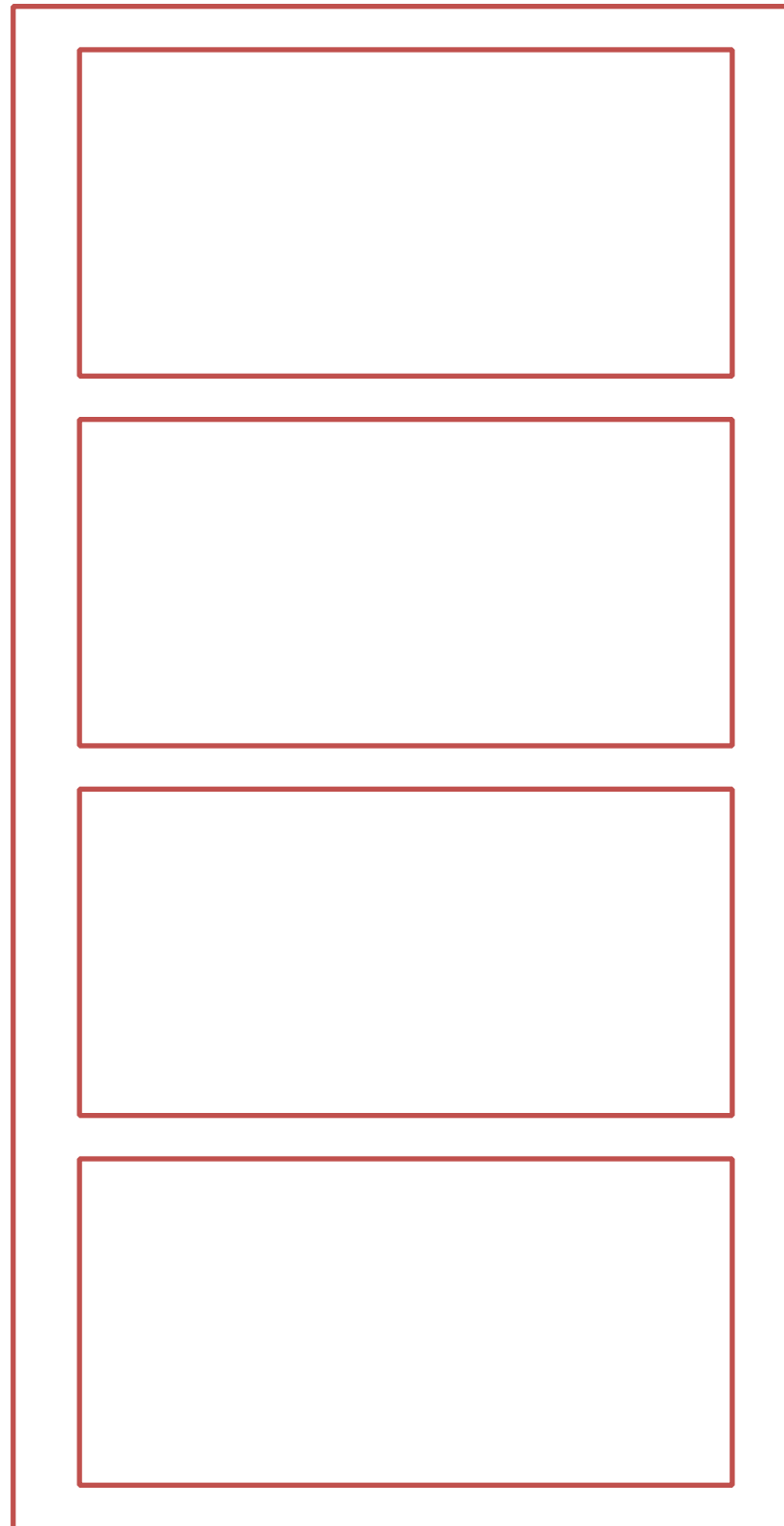
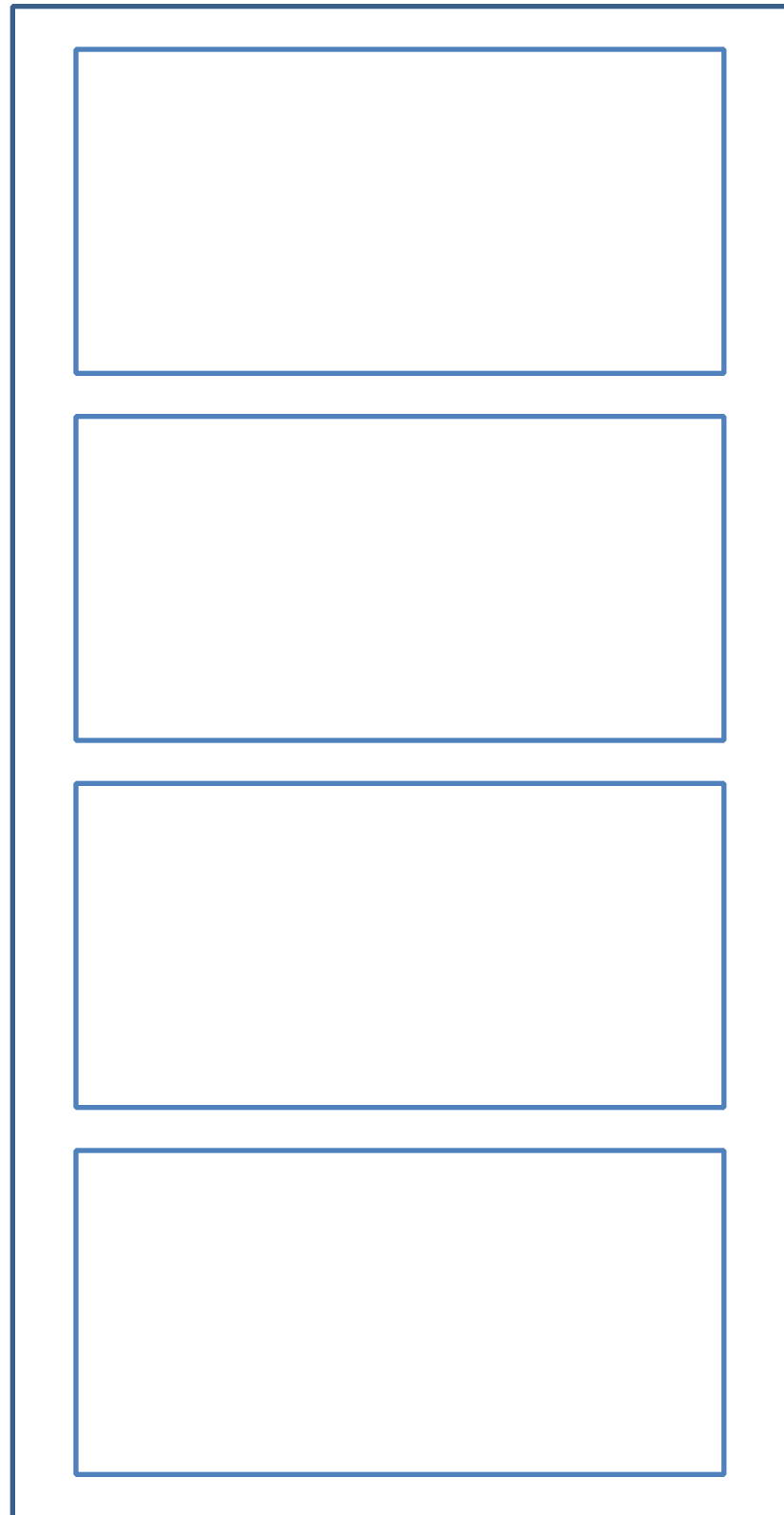
Do we want the smaller relation as the OUTER or the INNER?

$$[S] + [S] * [R]$$

# Chunk Nested Loops Join

Sailors

Reserves



**Key idea:**

Take **k pages** of S  
and match with  
each page of R.

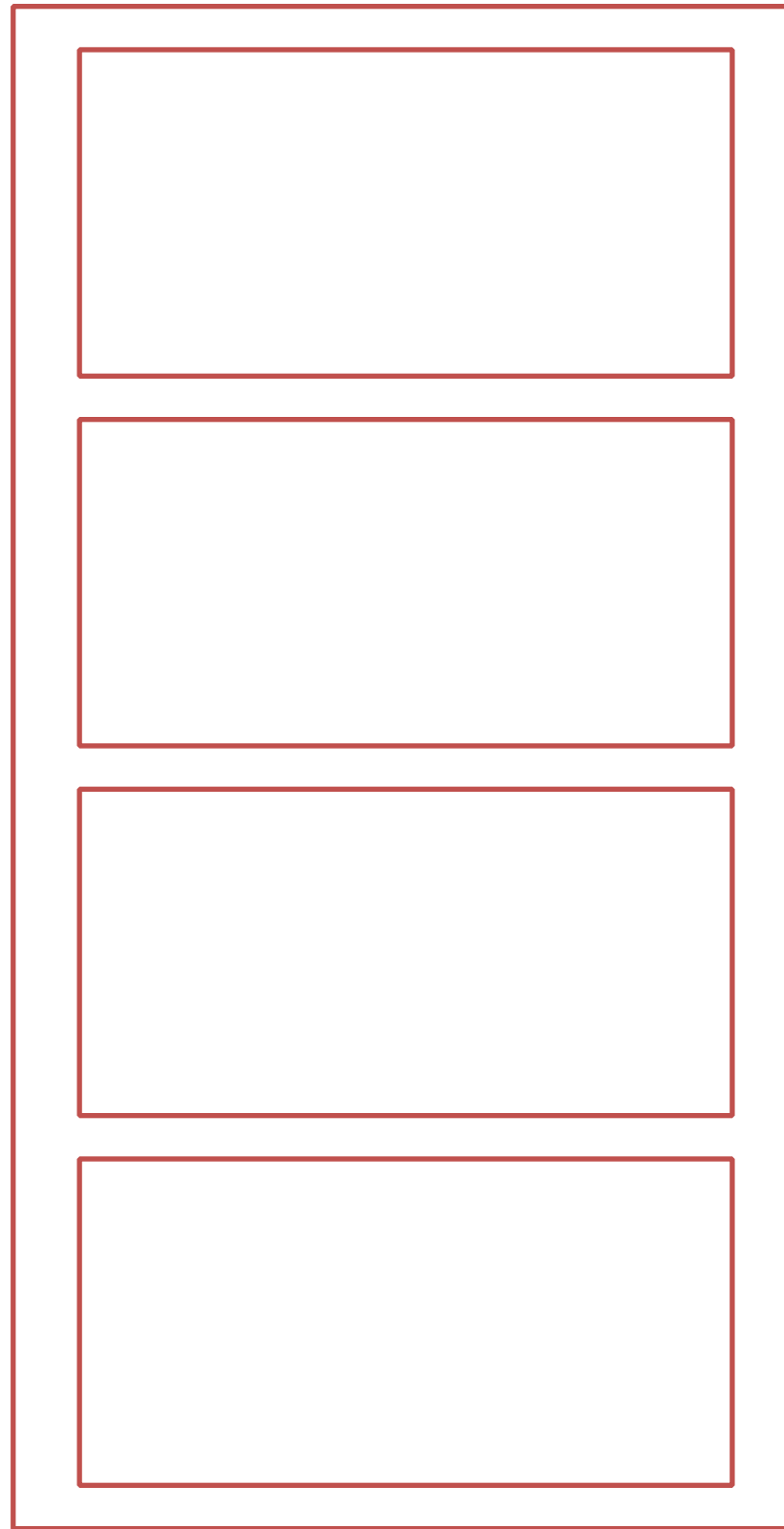
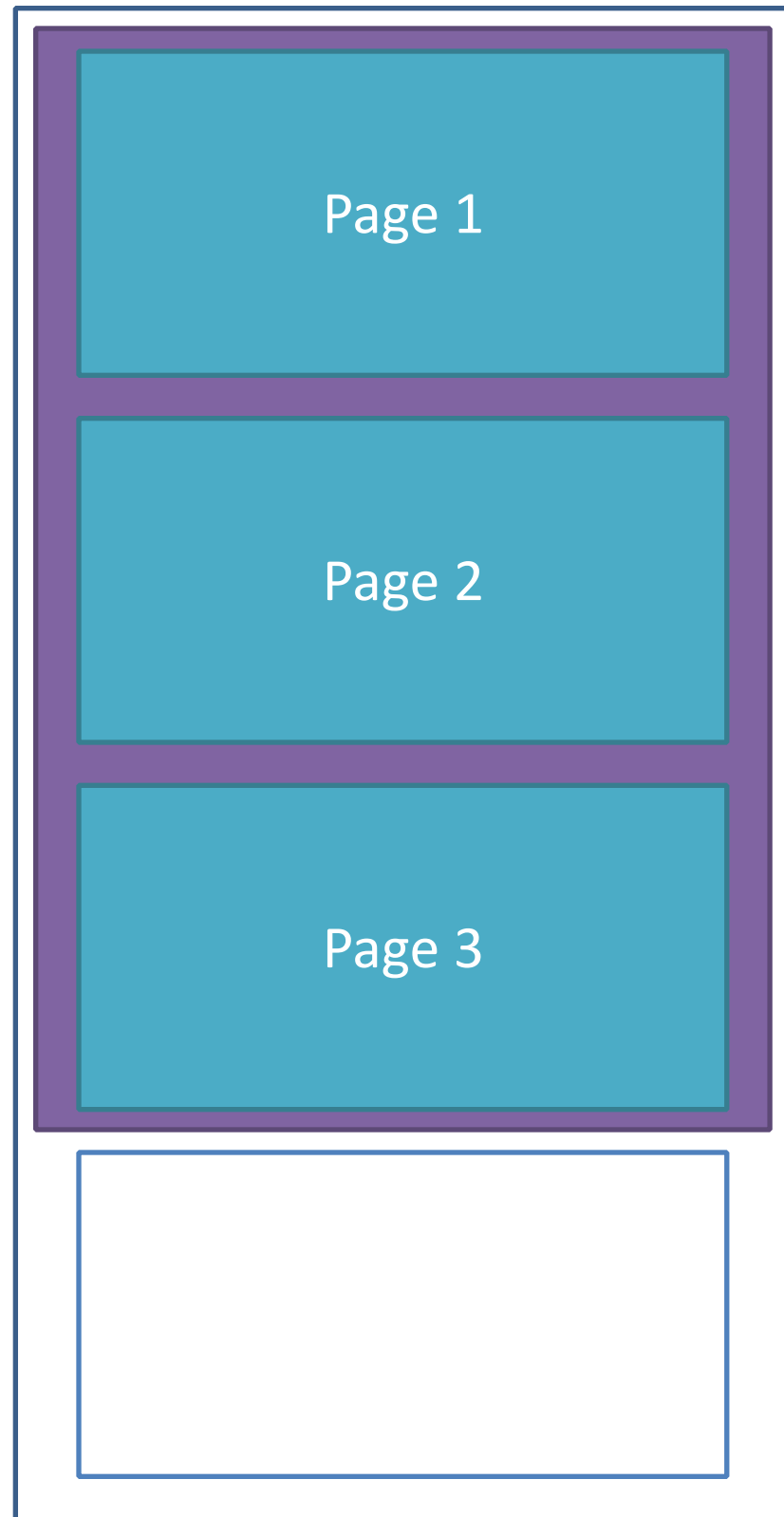
**Steps:**

1. Get **k** pages of S.
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each page in R.
3. Compare tuples in  
each.

# Chunk Nested Loops Join

Sailors

Reserves



**Key idea:**

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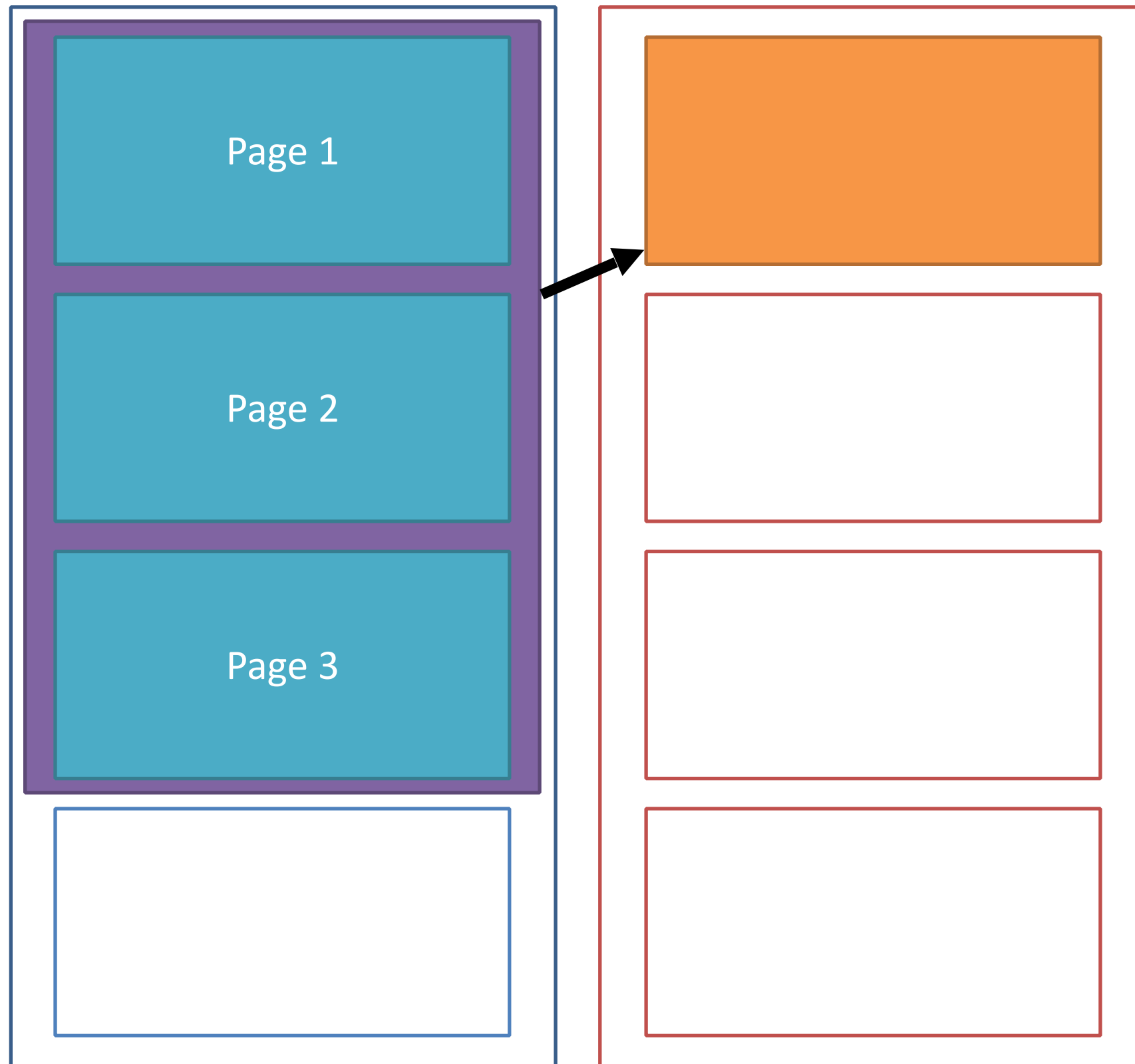
**Steps:**

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# Chunk Nested Loops Join

Sailors

Reserves



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**Steps:**

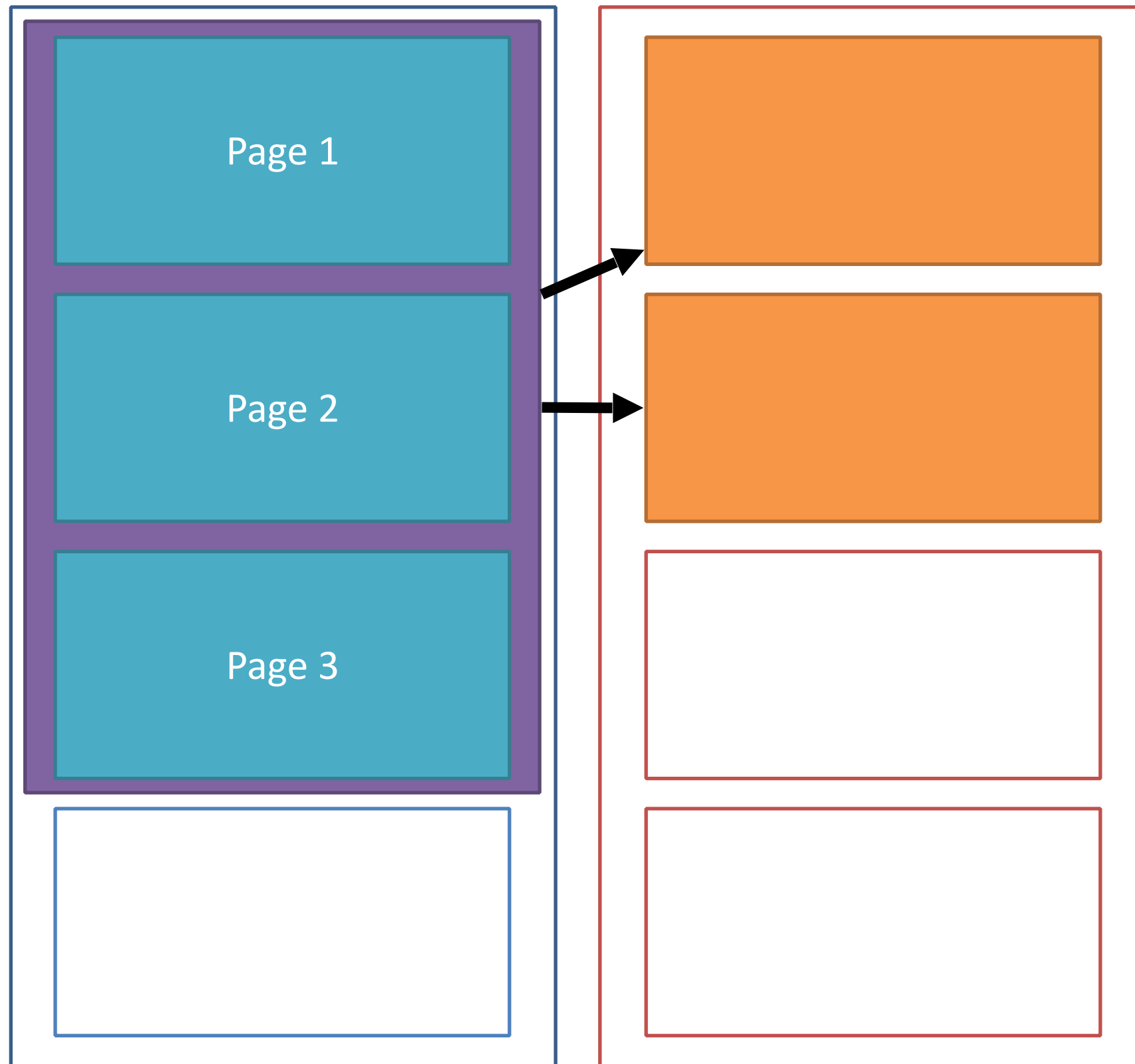
1. Get **k** pages of S.
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# Chunk Nested Loops Join

Sailors

Reserves



**Key idea:**

Take **k pages** of S  
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each page of R.

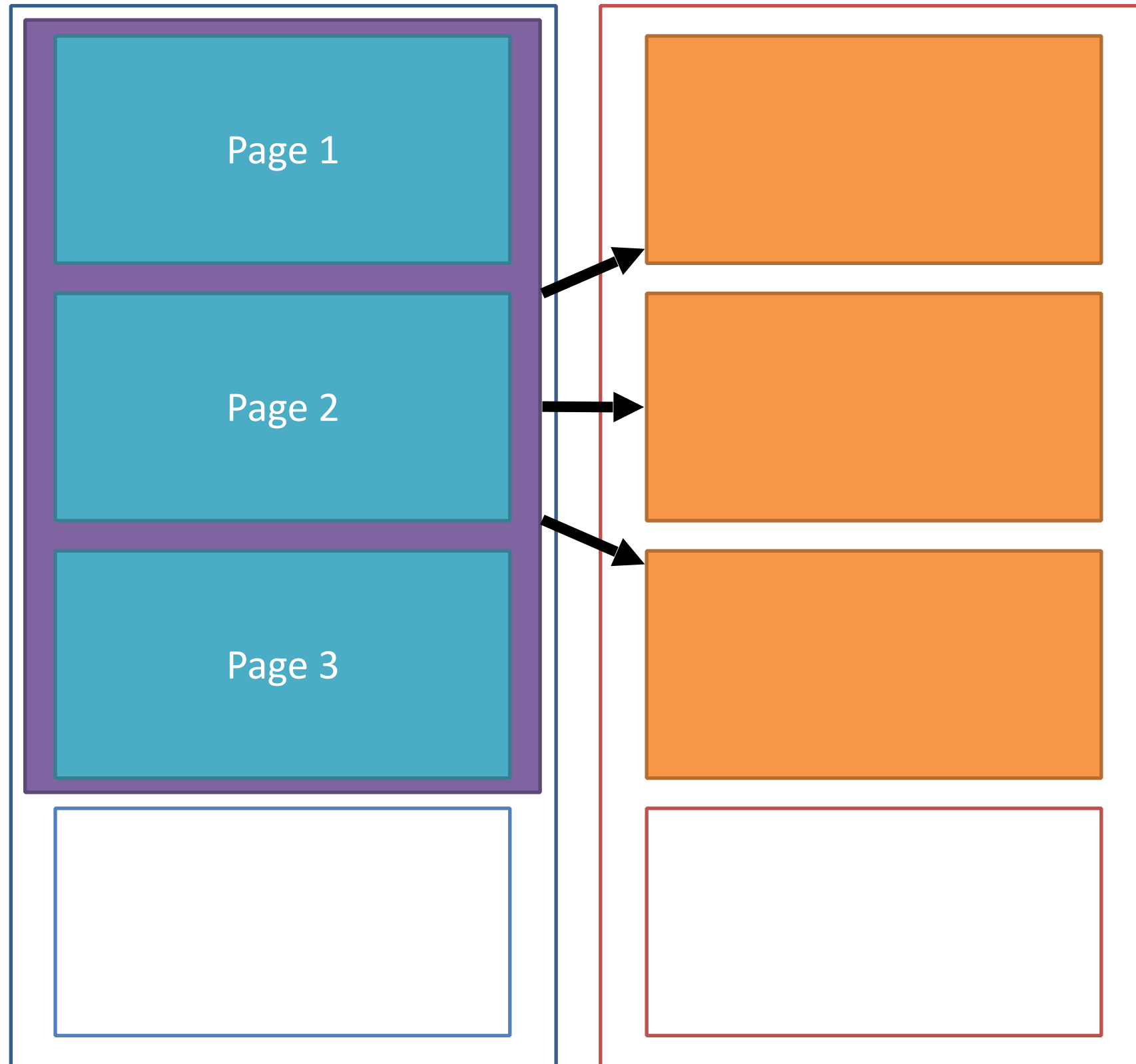
**Steps:**

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2. Iterate through  
each page in R.
3. Compare tuples in  
each.

# Chunk Nested Loops Join

Sailors

Reserves



**Key idea:**

Take **k pages** of S and match with each page of R.

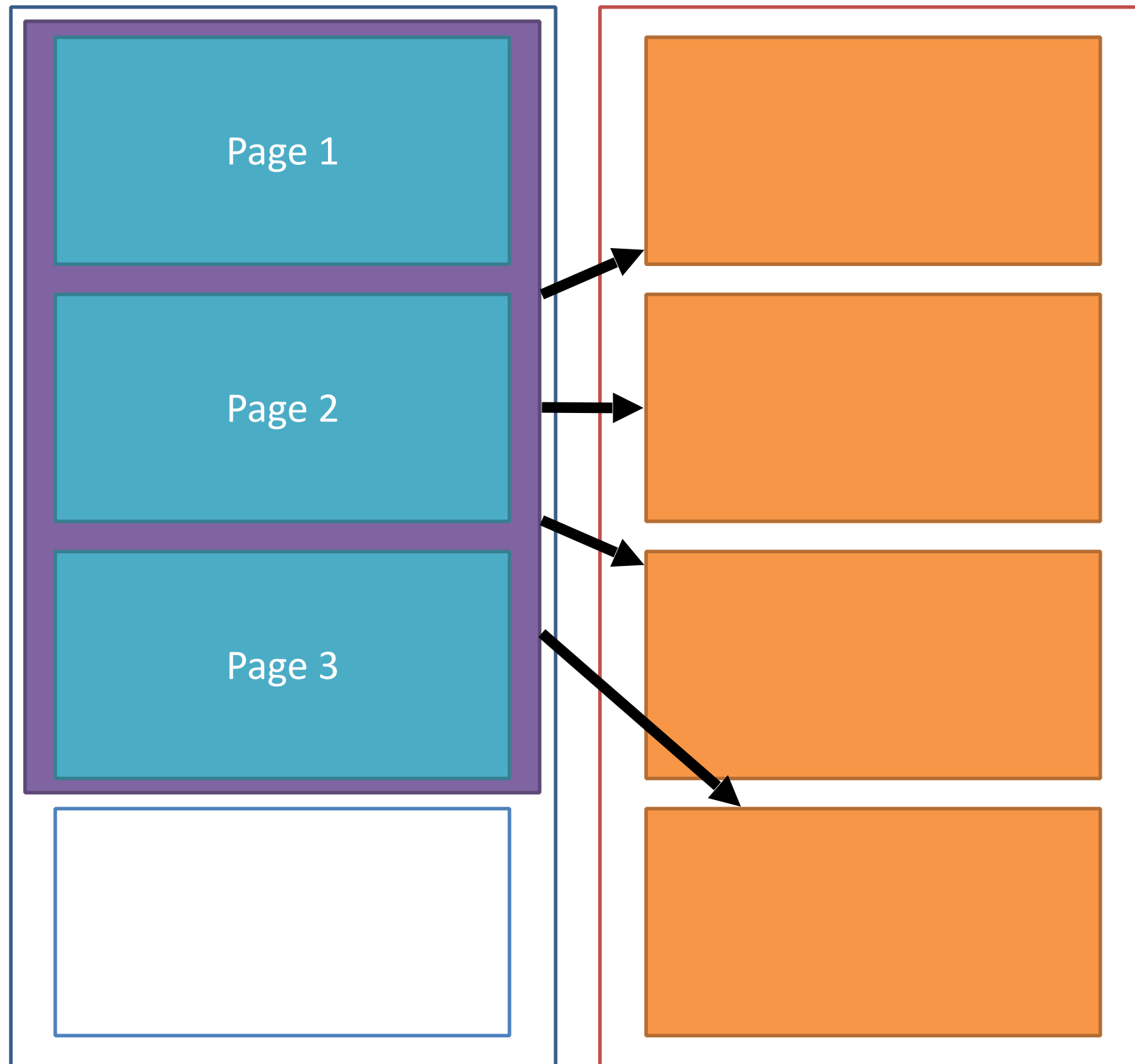
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# Chunk Nested Loops Join

Sailors

Reserves



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# Chunk Nested Loops Join

Sailors

Reserves

Page 1

Page 2

Page 3

Page 4

**Key idea:**

Take **k pages** of S  
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**Steps:**

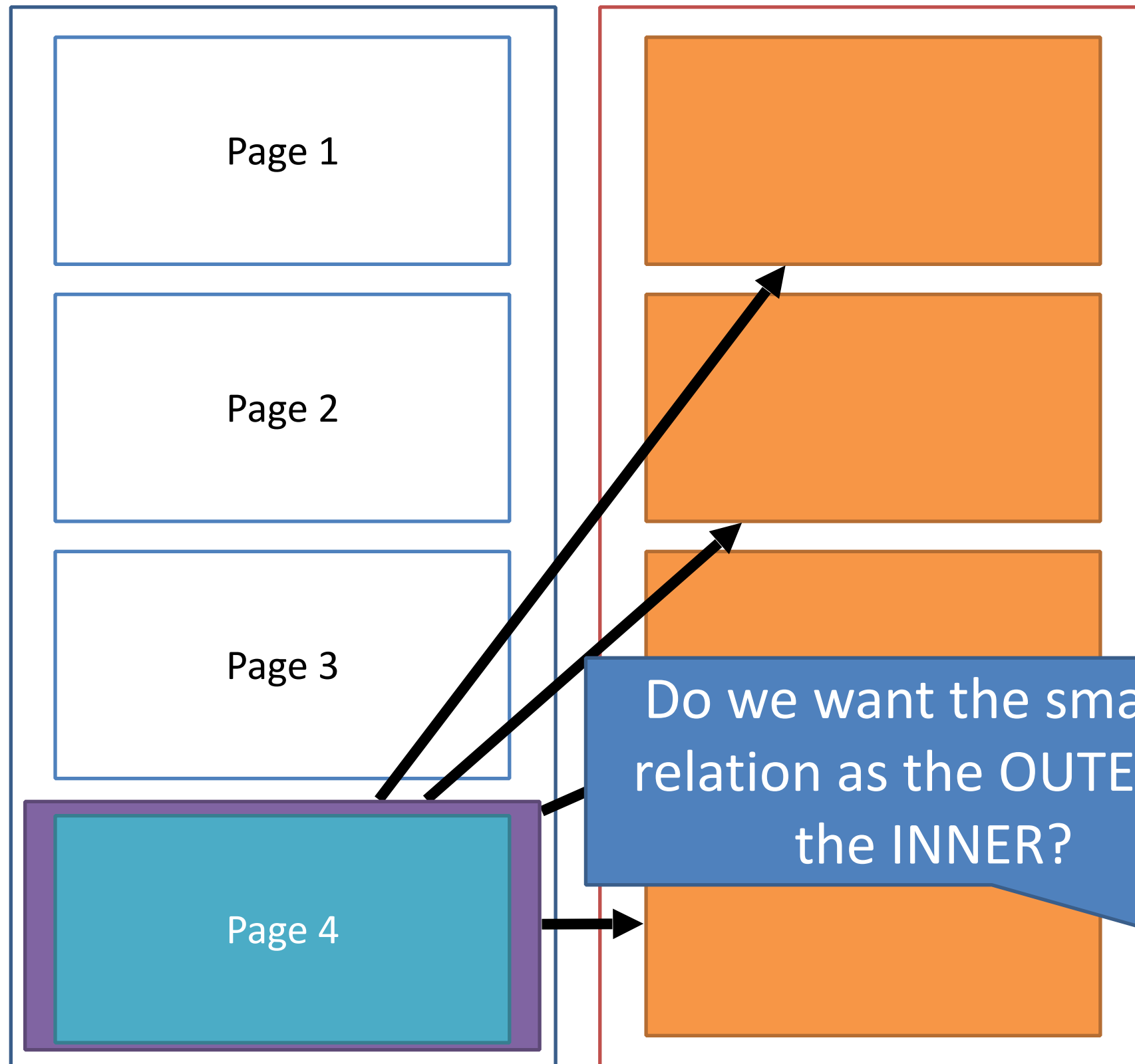
1. Get **k** pages of S.
2. Iterate through  
each page in R.
3. Compare tuples in  
each.

**Notation:  $[S]$  == “# pages in S” ;  $|S|$  == “# tuples in S”**

# Chunk Nested Loops Join

Sailors

Reserves



**Key idea:**

Take **k pages** of S and match with each page of R.

**Steps:**

1. Get **k** pages of S.
2. Iterate through each page in R.

Compare tuples in each.

:

$$[S] + ([S] / k) * [R]$$

# Sort-Merge Join

Sailors


Reserves


**Key idea:**

Sort S and R, then  
merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

# Sort-Merge Join

Sailors

Reserves

**Key idea:**

Sort S and R **on join column**, then merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

(name = Bob, sid = 1)

(name = Sam, sid = 3)

(name = Sue, sid = 7)

(name = Jill, sid = 2)

(name = Joe, sid = 12)

(name = Sue, sid = 8)

(name = Yue, sid = 4)

# Sort-Merge Join

Sailors

Reserves

**Key idea:**

Sort S and R **on join column**, then merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

(name = Bob, sid = 1)

(name = Jill, sid = 2)

(name = Sam, sid = 3)

(name = Yue, sid = 4)

(name = Sue, sid = 7)

(name = Sue, sid = 8)

(name = Joe, sid = 12)

...



# Sort-Merge Join

Sailors

Reserves

**Key idea:**

Sort S and R **on join column**, then merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

(name = Bob, sid = 1)
(name = Jill, sid = 2)
(name = Sam, sid = 3)
(name = Yue, sid = 4)
(name = Sue, sid = 7)

(name = Sue, sid = 8)
(name = Joe, sid = 12)
...

(sid = 1, bid = 4)
(sid = 1, bid = 7)
(sid = 3, bid = 6)
(sid = 4, bid = 3)
(sid = 8, bid = 1)

(sid = 8, bid = 13)
(sid = 8, bid = 15)
(sid = 12, bid = 1)
...

# Sort-Merge Join

Sailors

Reserves

**Key idea:**

Sort S and R **on join column**, then merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

**Output:**

(name = Bob, sid = 1, bid = 4)

(name = Bob, sid = 1)

(name = Jill, sid = 2)

(name = Sam, sid = 3)

(name = Yue, sid = 4)

(name = Sue, sid = 7)

(name = Sue, sid = 8)

(name = Joe, sid = 12)

...

(sid = 1, bid = 4)

(sid = 1, bid = 7)

(sid = 3, bid = 6)

(sid = 4, bid = 3)

(sid = 8, bid = 1)

(sid = 8, bid = 13)

(sid = 8, bid = 15)

(sid = 12, bid = 1)

...

# Sort-Merge Join

Sailors

Reserves

**Key idea:**

Sort S and R **on join column**, then merge them!

**Steps:**

1. Sort S and R.
2. “Zip” or merge.

**Output:**

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...

**Notation:  $[S]$  == “# pages in  $S$ ” ;  $|S|$  == “# tuples in  $S$ ”**

# Sort-Merge Join

## Sailors

## Reserves

```
graph TD
    L1["(name = Bob, sid = 1)  
(name = Jill, sid = 2)  
(name = Sam, sid = 3)  
(name = Yue, sid = 4)  
(name = Sue, sid = 7)"]
    L2["(name = Sue, sid = 8)  
(name = Joe, sid = 12)  
...  
  
"]
    L3["  
  
  
  
"]
    L1 --- L2
    L2 --- L3
```

(name = Bob, sid = 1)  
(name = Jill, sid = 2)  
(name = Sam, sid = 3)  
(name = Yue, sid = 4)  
(name = Sue, sid = 7)

(name = Sue, sid = 8)  
(name = Joe, sid = 12)  
...

[illegible]

## Key idea:

# Sort S and R **on join column**, then merge them!

## Steps:

1. Sort S and R.
2. “Zip” or merge.

## I/Os:

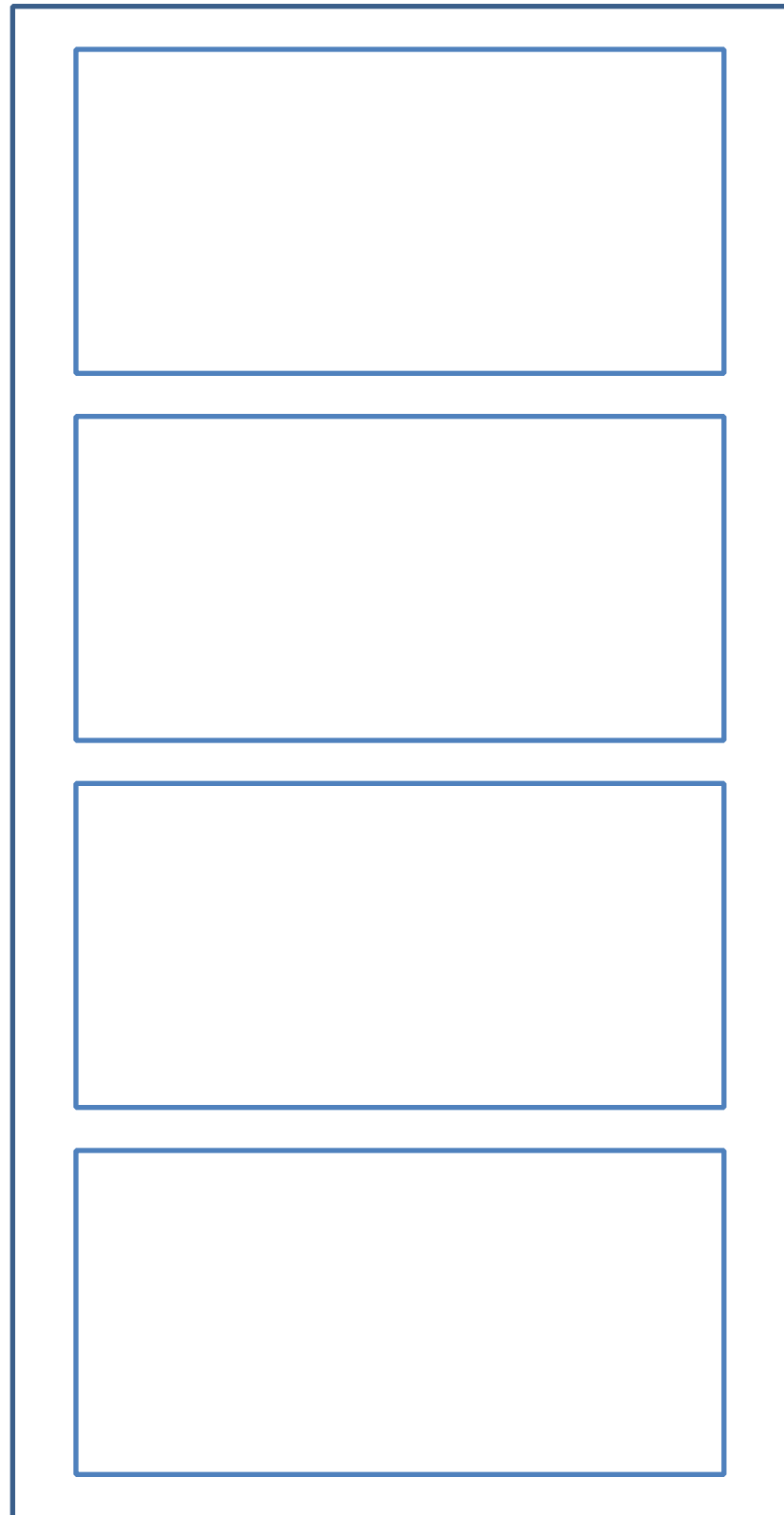
$$\sim 5([S] + [R])$$

## Sorting: 4([S]+[R])

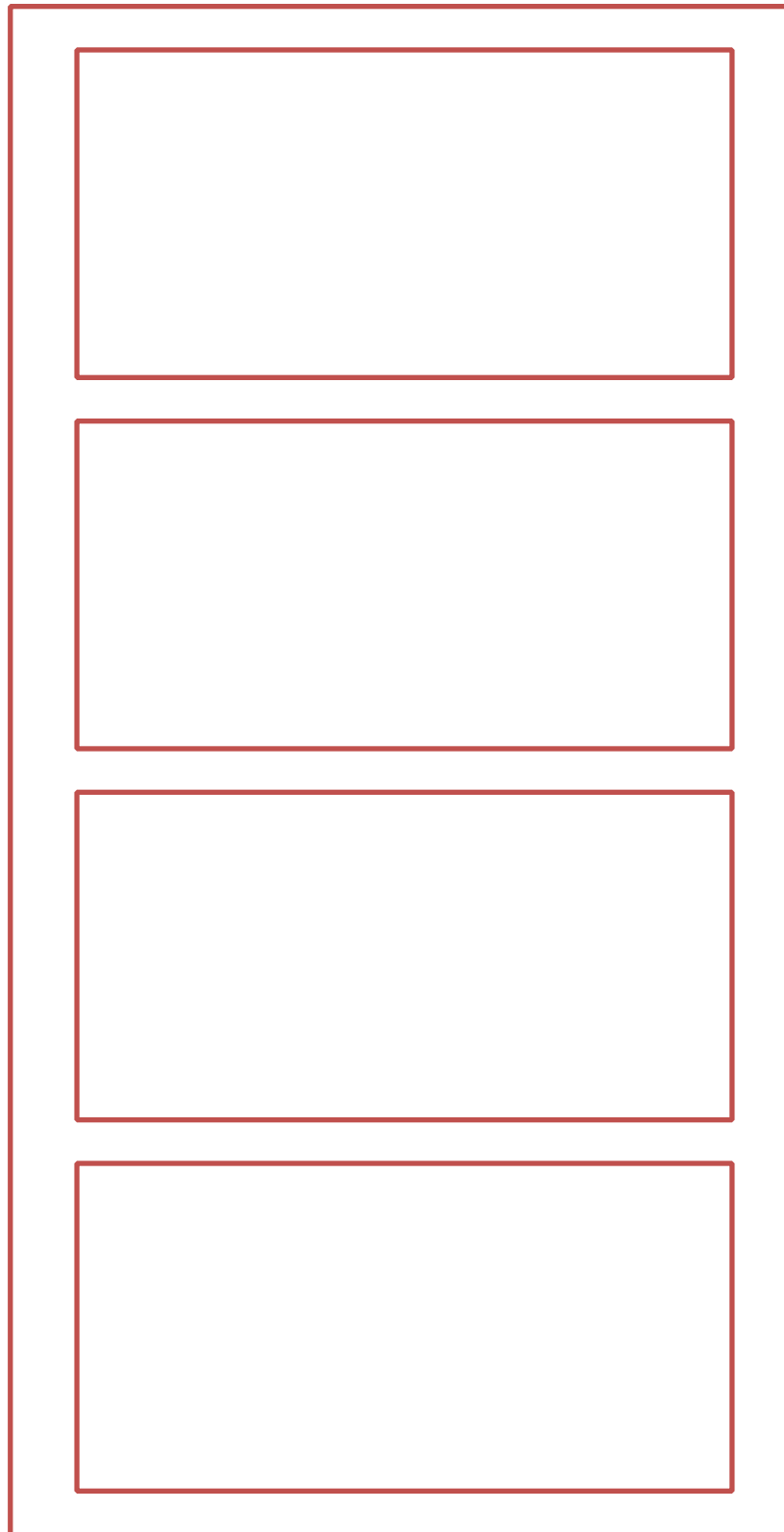
## Merging: [S]+[R]

# Hash-Join

Sailors



Reserves



**Key idea:**

Partition S and R  
using same hash fn,  
then collect same  
partitions

**Steps:**

1. Partition S and R
2. Re-Hash, collect

# Hash-Join

Sailors

Reserves

## Key idea:

Partition S and R  
using same hash fn,  
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partitions

## Steps:

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(name = Bob, sid = 1)

(name = Sam, sid = 3)

(name = Sue, sid = 7)

(name = Jill, sid = 2)

(name = Joe, sid = 12)

(name = Sue, sid = 8)

(name = Yue, sid = 4)

# Hash function: $\text{sid mod 4}$ Hash-Join

Sailors

Reserves

**Key idea:**

Partition S and R  
using same hash fn,  
then collect same  
partitions

**Steps:**

1. Partition S and R
2. Re-Hash, collect

(name = Joe, sid = 12)
(name = Sue, sid = 8)
(name = Yue, sid = 4)
...

(name = Bob, sid = 1)
...

(name = Jill, sid = 2)
...

(name = Sue, sid = 7)
(name = Sam, sid = 3)
...

(sid = 12, bid = 1)
(sid = 8, bid = 13)
(sid = 8, bid = 15)
(sid = 4, bid = 3)
(sid = 8, bid = 1)

(sid = 1, bid = 4)
(sid = 1, bid = 7)
...

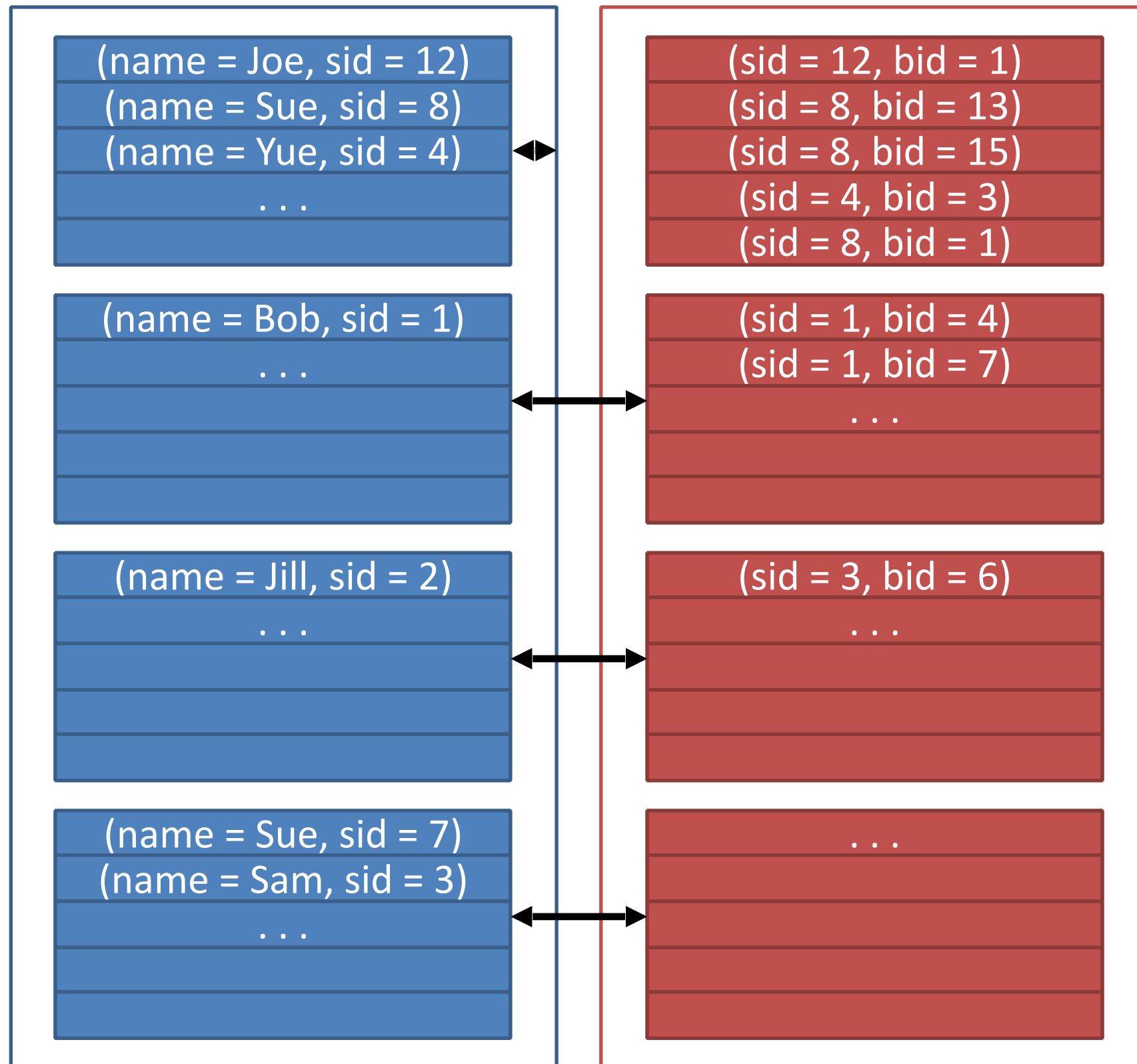
(sid = 3, bid = 6)
...

...
-----

# Hash-Join

Sailors

Reserves



## Key idea:

Partition S and R  
using same hash fn,  
then collect same  
partitions

## Steps:

1. Partition S and R
2. Re-Hash, collect



**Notation:  $[S]$  == “# pages in  $S$ ” ;  $|S|$  == “# tuples in  $S$ ”**

# Hash-Join

Sailors

Reserves

**Key idea:**

Partition  $S$  and  $R$   
using same hash fn,  
then collect same  
partitions

**Steps:**

1. Partition  $S$  and  $R$   
using same hash fn, collect

3( $[S] + [R]$ )

Partition: 2( $[S] + [R]$ )

Re-Hash:  $[S] + [R]$

NOTE: This is no different  
from what we previously  
assumed about hashing.

NOTE: What are we  
assuming about the size of  
partitions?

(name = Joe, sid = 12)  
(name = Sue, sid = 8)  
(name = Yue, sid = 4)  
...

(name = Bob, ...)

(name = Jill, sid = 2)  
...

(name = Sue, sid = 7)  
(name = Sam, sid = 3)  
...

(sid = 12, bid = 1)  
(sid = 8, bid = 13)  
(sid = 8, bid = 15)  
(sid = 4, bid = 3)  
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...

...

...

# Join Cheatsheet

**Notation:**  $[S]$  == “# pages in  $S$ ” ;

$|S|$  == “# tuples in  $S$ ”

- Chunk nested loop join
  - Take  $k$  pages of  $S$  and match with each page of  $R$ .
  - Total Cost:  $[S] + ([S] / k) * [R]$
- Sort merge join
  - Sort  $S$  and  $R$  **on join column**, then merge them!
  - Total Cost:  $\sim 5([S] + [R])$
- Hash join
  - Partition  $S$  and  $R$  using same hash fn, then collect same partitions
  - Total Cost:  $\sim 3([S] + [R])$ 
    - Assuming  $\text{len}(\text{partition}) \leq B$  pages

When is a chunk-nested loops join the best?

# When is a chunk-nested loops join the best?

- Not using an equality predicate
- Join is just a cross product

When is a sort-merge join the best?

# When is a sort-merge join the best?

- Skewed input data
- Small memory size
- Want sorted output/have sorted input

When is a hash-join the best?

# When is a hash-join the best?

- One partition large, the other small (can keep in memory)



We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How many disk reads are needed to perform Chunk Nested Loops Join?

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How many disk reads are needed to perform Chunk Nested Loops Join?

$(\# \text{ of pages in smaller relation}) + ((\# \text{ of pages in smaller relation}) / (\# \text{ of pages in memory} - 2 \text{ for I/O})) * (\# \text{ of pages in larger relation})$

$$= 50 + (50/10) * (100) = 550$$

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How about a Hash Join? (Assume no recursive partitioning)

We have 12 pages of memory, and we want to join two tables [R] and [S] where [R] is 100 pages and [S] is 50 pages.

How about a Hash Join? (Assume no recursive partitioning)

(# of pages in both relations) \* (1 read before hashing + 1 read after hashing)

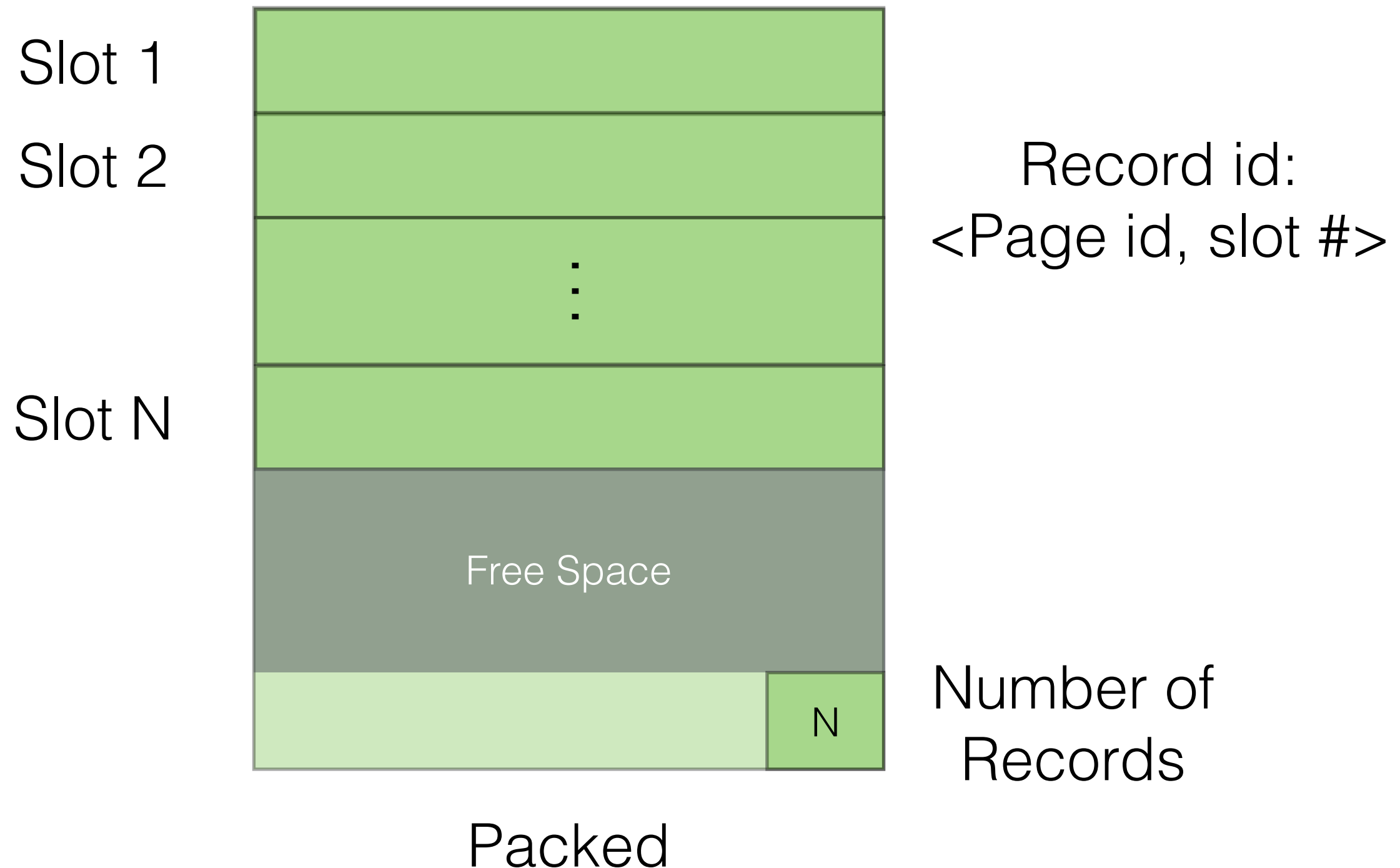
$$=(100+50) * 2$$

$$= 300$$

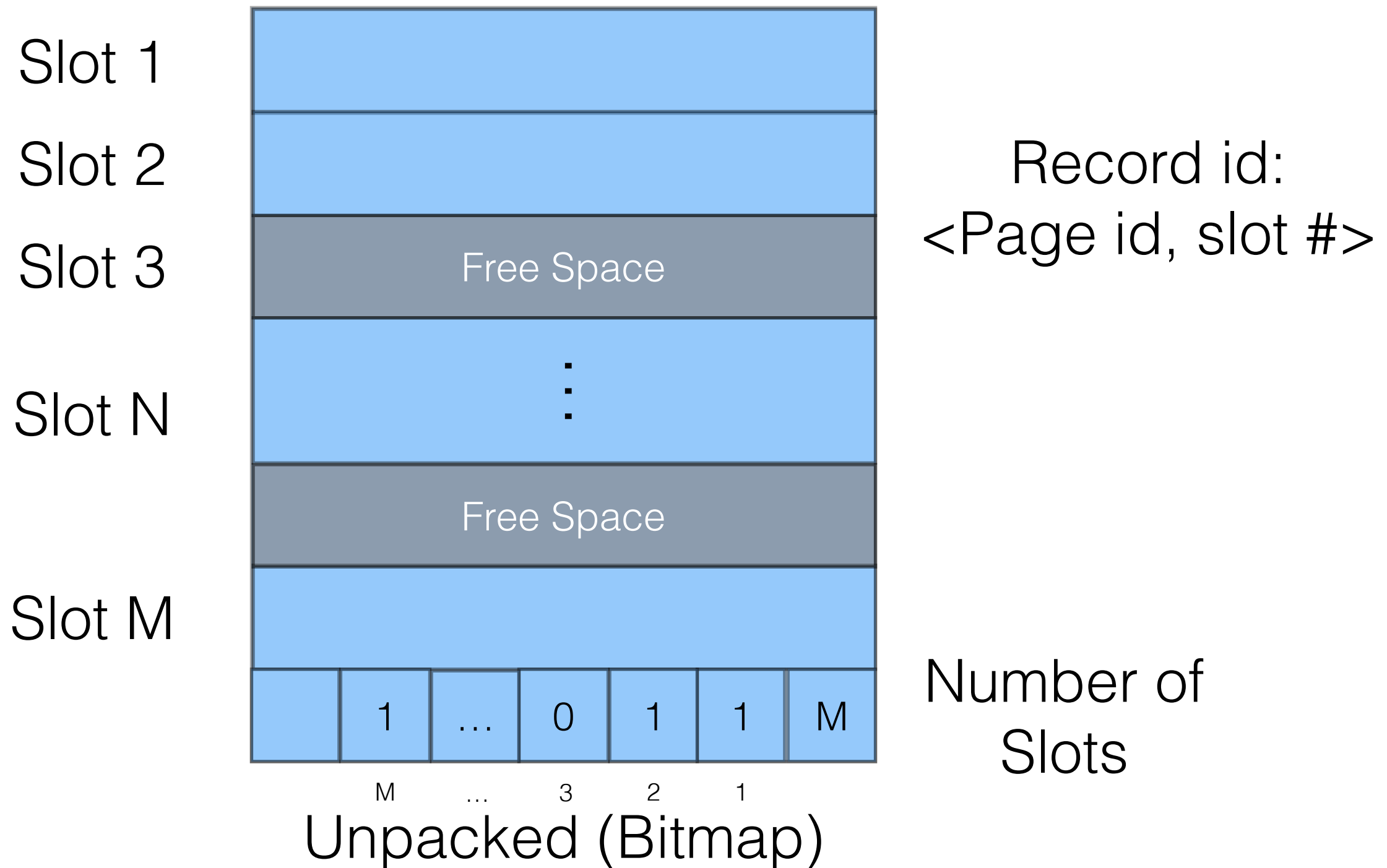
# Heap Files

(Page Formats)

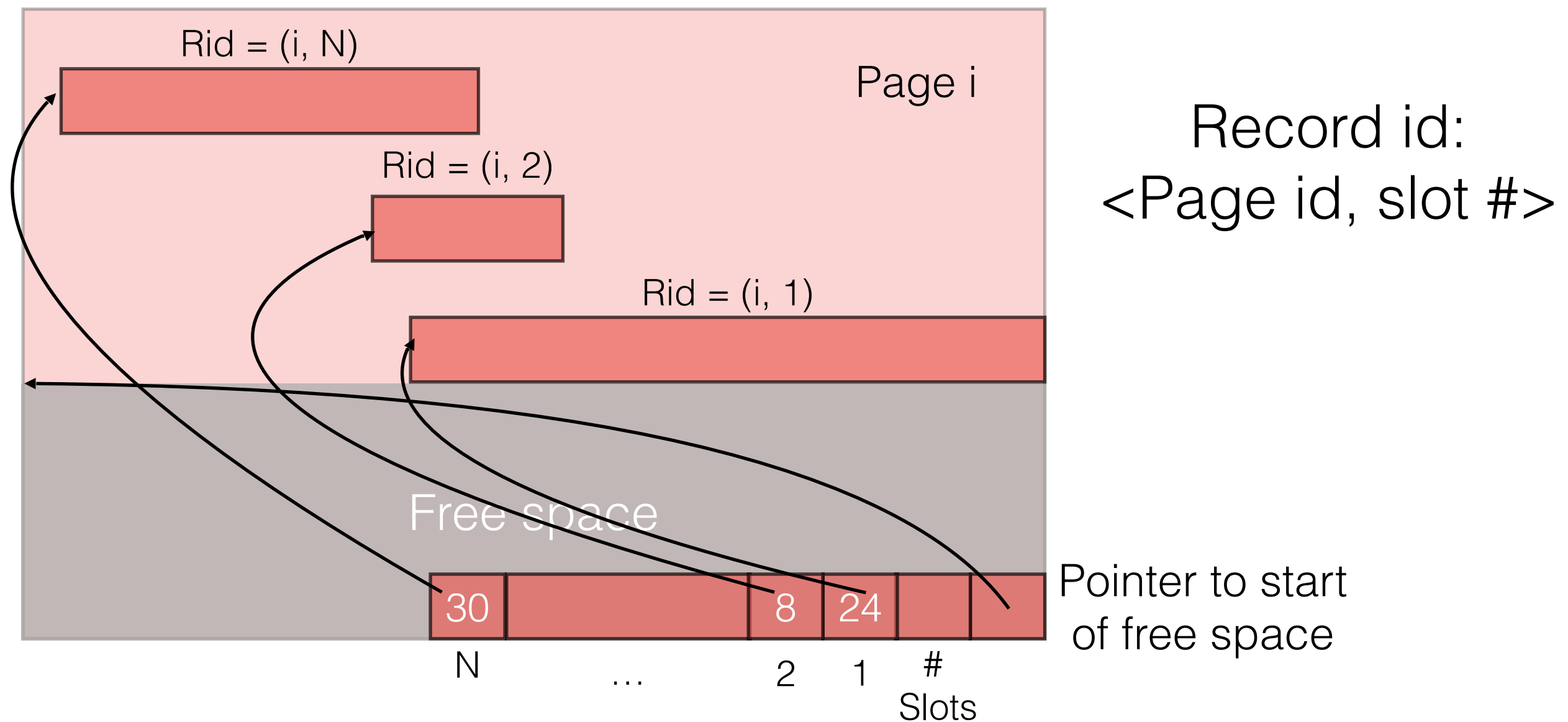
# Fixed-Length Records



# Fixed-Length Records



# Variable Length Records





What are the advantages and disadvantages of using slotted pages or bitmaps over just tightly packing records together?

What are the advantages and disadvantages of using slotted pages or bitmaps over just tightly packing records together?

- Allow movement of records without changing record ID
- Slotted pages support variable-length records

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

What's the size of the largest record you can insert?

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

What's the size of the largest record you can insert?

Need 4 bytes for the entry, so  $(80 - 4) = 76$  bytes

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

At most, how many 1-byte large records can you insert?

You have a slotted page with 80 bytes of free space, and it costs 4 bytes to store a directory entry.

At most, how many 1-byte large records can you insert?

- Amount of space taken up by  $x$  1-byte records  
= (1 byte for record + 4 for directory entry)  
= (5 bytes / record)
- Free space / (amount of space per record)  
=  $80 / 5 = 16$  records