

Condition Variables Solutions

Condition Variable Scenario

- In this scenario, the processing thread creates a `unique_lock` instance to lock the mutex, but the fetching thread creates a `lock_guard`
- Why do the two threads use different lock types?
 - The processing thread needs to be able to unlock the mutex. This is done inside the `wait()` call. In some applications, the waiting thread may need to directly unlock the mutex as well.
 - Therefore the mutex must support `unlock()`
 - Hence `unique_lock` is used

Condition Variable Scenario

- The fetching thread only needs the mutex to protect the critical section. The extra flexibility and overhead of `unique_lock` is not needed here, so we use the simpler `lock_guard`
- We put the `lock_guard` and the critical section in their own scope
- The mutex will be automatically unlocked by the `lock_guard` destructor

Condition Variable Example

- Now reverse the order of the threads, so that the writer thread is started first. Add a sleep (say, half a second) before starting the reader thread. What happens?
 - The writing thread completes before the reading thread runs
 - The reading thread has not called `wait()` on the condition variable
 - The condition variable does not have any waiting threads to notify
 - The notification is "lost"
 - By the time the reading thread calls `wait()`, the writing thread has completed
 - The condition variable does not receive any more notifications
 - The reading thread blocks for ever