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Computational Geometry

1.1 Geometry

```
const double PI=atan2(0.0.-1.0);
 template<tvpename T>
 struct point{
  T x,y;
   point(){}
   point(const T&x,const T&y):x(x),y(y){}
   point operator+(const point &b)const{
     return point(x+b.x,y+b.y); }
   point operator-(const point &b)const{
     return point(x-b.x,y-b.y); }
   point operator*(const T &b)const{
     return point(x*b,y*b); }
   point operator/(const T &b)const{
     return point(x/b,y/b); }
   bool operator==(const point &b)const{
     return x==b.x&&y==b.y; }
   T dot(const point &b)const{
     return x*b.x+y*b.y; }
   T cross(const point &b)const{
     return x*b.y-y*b.x; }
   point normal()const{//求法向量
     return point(-y,x); }
  T abs2()const{//向量長度的平方
     return dot(*this); }
   T rad(const point &b)const{//兩向量的弧度
 return fabs(atan2(fabs(cross(b)),dot(b))); }
  T getA()const{//對x軸的弧度
     T A=atan2(y,x);//超過180度會變負的
     if(A<=-PI/2)A+=PI*2;
     return A;
 template<typename T>
 struct line{
   line(){}
   point<T> p1,p2;
   T a,b,c;//ax+by+c=0
   line(const point<T>&x,const point<T>&y):p1
        (x),p2(y){}
   void pton(){//轉成一般式
     a=p1.y-p2.y;
     b=p2.x-p1.x;
     c=-a*p1.x-b*p1.y;
  T ori(const point<T> &p)const{//點和有向直
        線的關係, >0左邊、=0在線上<0右邊
     return (p2-p1).cross(p-p1);
  T btw(const point<T> &p)const{//點投影落在
        線段上<=0
                                             102
     return (p1-p).dot(p2-p);
   bool point on segment(const point<T>&p)
        const{//點是否在線段上
     return ori(p) == 0&&btw(p) <= 0;</pre>
   T dis2(const point<T> &p,bool is segment
        =0) const { // 點 跟 直 線 / 線 段 的 距 離 平 方
```

```
point<T> v=p2-p1, v1=p-p1;
  if(is segment){
                                            110
    point<T> v2=p-p2:
                                            111
    if(v.dot(v1)<=0)return v1.abs2();</pre>
                                            112
    if(v.dot(v2)>=0)return v2.abs2();
                                           113
                                            114
  T tmp=v.cross(v1);
                                            115 };
  return tmp*tmp/v.abs2():
T seg dis2(const line<T> &1)const{//兩線段
  return min({dis2(1.p1,1),dis2(1.p2,1),1. 120
       dis2(p1,1),1.dis2(p2,1)});
                                            122
point<T> projection(const point<T> &p)
     const { //點對直線的投影
                                            123
  point<T> n=(p2-p1).normal();
                                            124
  return p-n*(p-p1).dot(n)/n.abs2();
                                            125
                                            126
point<T> mirror(const point<T> &p)const{
                                            127
  //點對直線的鏡射,要先呼叫pton轉成一般式 128
  point<T> R;
  T d=a*a+b*b:
                                           130
  R.x=(b*b*p.x-a*a*p.x-2*a*b*p.y-2*a*c)/d;
  R.y=(a*a*p.y-b*b*p.y-2*a*b*p.x-2*b*c)/d;
                                           131
                                            132
  return R:
                                            133
                                            134
bool equal(const line &1)const{//直線相等
                                            135
  return ori(1.p1)==0&&ori(1.p2)==0;
bool parallel(const line &1)const{
  return (p1-p2).cross(l.p1-l.p2)==0;
                                            137
bool cross_seg(const line &1)const{
                                            138
  return (p2-p1).cross(l.p1-p1)*(p2-p1).
       cross(1.p2-p1)<=0;//直線是否交線段
int line_intersect(const line &1)const{// 140
     直線相交情況,-1無限多點、1交於一點、0 141
                                            142
  return parallel(1)?(ori(1.p1)==0?-1:0)
                                            143
       :1;
                                            144
                                            145
int seg_intersect(const line &l)const{
 T c1=ori(l.p1), c2=ori(l.p2);
                                            146
  T c3=1.ori(p1), c4=1.ori(p2);
  if(c1==0&&c2==0){//共線
    bool b1=btw(1.p1)>=0,b2=btw(1.p2)>=0;
    T a3=1.btw(p1),a4=1.btw(p2);
    if(b1&&b2&&a3==0&&a4>=0) return 2;
                                            149
                                            150
    if(b1&&b2&&a3>=0&&a4==0) return 3;
    if(b1&&b2&&a3>=0&&a4>=0) return 0;
                                            151
                                            152
    return -1;//無限交點
  }else if(c1*c2<=0&&c3*c4<=0)return 1;</pre>
  return 0;//不相交
                                            154
point<T> line intersection(const line &l)
                                            156
     const{/*直線交點*/
                                            157
  point<T> a=p2-p1,b=l.p2-l.p1,s=l.p1-p1;
                                            158
  //if(a.cross(b)==0)return INF;
                                            159
  return p1+a*(s.cross(b)/a.cross(b));
                                            160
point<T> seg_intersection(const line &1)
     const{//線段交點
```

```
int res=seg intersect(1);
                                                  162
       if(res<=0) assert(0);</pre>
       if(res==2) return p1;
                                                  163
       if(res==3) return p2;
                                                  164
       return line intersection(1);
                                                  165
                                                  166
116 template<typename T>
117 struct polygon{
                                                  167
    polygon(){}
                                                  168
     vector<point<T> > p;//逆時針順序
     T area()const{//面積
                                                  169
       T ans=0;
       for(int i=p.size()-1,j=0;j<(int)p.size()</pre>
            ;i=j++)
         ans+=p[i].cross(p[j]);
                                                  171
       return ans/2;
                                                  172
                                                  173
                                                  174
     point<T> center of mass()const{//重心
       T cx=0, cy=0, w=0;
                                                  175
       for(int i=p.size()-1,j=0;j<(int)p.size()</pre>
            ;i=j++){
         T a=p[i].cross(p[j]);
                                                  177
         cx+=(p[i].x+p[j].x)*a;
         cy+=(p[i].y+p[j].y)*a;
                                                  178
         w+=a:
                                                  179
       return point<T>(cx/3/w,cy/3/w);
                                                  180
                                                  181
     char ahas(const point<T>& t)const{//點是否
          在簡單多邊形內,是的話回傳1、在邊上回
          傳-1、否則回傳0
       bool c=0;
                                                  184
       for(int i=0,j=p.size()-1;i<p.size();j=i</pre>
         if(line<T>(p[i],p[j]).point_on_segment
              (t))return -1;
         else if((p[i].y>t.y)!=(p[j].y>t.y)&&
         t.x<(p[j].x-p[i].x)*(t.y-p[i].y)/(p[j]
                                                  188
              ].y-p[i].y)+p[i].x)
                                                  189
           c=!c;
                                                  190
       return c;
     char point_in_convex(const point<T>&x)
                                                  191
                                                  192
         const{
                                                  193
       int l=1,r=(int)p.size()-2;
                                                 194
       while(1<=r){//點是否在凸多邊形內,是的話
            回傳1、在邊上回傳-1、否則回傳0
                                                  195
         int mid=(1+r)/2;
                                                  196
         T a1=(p[mid]-p[0]).cross(x-p[0]);
                                                  197
         T a2=(p[mid+1]-p[0]).cross(x-p[0]);
                                                  198
         if(a1>=0&&a2<=0){
                                                  199
           T res=(p[mid+1]-p[mid]).cross(x-p[
                                                  200
                mid]);
                                                  201
           return res>0?1:(res>=0?-1:0);
                                                  202
         }else if(a1<0)r=mid-1;</pre>
                                                  203
         else l=mid+1:
                                                  204
                                                  205
       return 0;
                                                  206
     vector<T> getA()const{//凸包邊對x軸的夾角
                                                  207
       vector<T>res;//一定是遞增的
                                                  208
       for(size t i=0;i<p.size();++i)</pre>
                                                  209
```

```
res.push_back((p[(i+1)%p.size()]-p[i])
         .getA());
  return res:
bool line intersect(const vector<T>&A,
     const line<T> &1)const{//O(LogN)
  int f1=upper_bound(A.begin(),A.end(),(1.
       p1-1.p2).getA())-A.begin();
  int f2=upper bound(A.begin(),A.end(),(1.
       p2-1.p1).getA())-A.begin();
  return 1.cross seg(line<T>(p[f1],p[f2]))
polygon cut(const line<T> &1)const{//△包
     對直線切割,得到直線 L左側的凸包
  polygon ans;
  for(int n=p.size(),i=n-1,j=0;j<n;i=j++){</pre>
    if(l.ori(p[i])>=0){
      ans.p.push back(p[i]);
      if(1.ori(p[j])<0)</pre>
        ans.p.push_back(1.
             line intersection(line<T>(p[i
             1,p[i])));
    }else if(l.ori(p[j])>0)
      ans.p.push back(1.line intersection(
          line<T>(p[i],p[j])));
  return ans;
static bool monotone_chain_cmp(const point
     <T>& a, const point<T>& b){//凸包排序函
  return (a.x<b.x)||(a.x==b.x&&a.y<b.y);</pre>
void monotone chain(vector<point<T> > &s){
    //凸包
  sort(s.begin(),s.end(),
       monotone chain cmp);
  p.resize(s.size()+1);
  int m=0;
  for(size t i=0;i<s.size();++i){</pre>
    while(m>=2&&(p[m-1]-p[m-2]).cross(s[i
        ]-p[m-2])<=0)--m;
    p[m++]=s[i];
  for(int i=s.size()-2,t=m+1;i>=0;--i){
    while (m>=t&&(p[m-1]-p[m-2]).cross(s[i
        ]-p[m-2])<=0)--m;
    p[m++]=s[i];
  if(s.size()>1)--m;
  p.resize(m);
T diam(){//直徑
  int n=p.size(),t=1;
  T ans=0;p.push back(p[0]);
  for(int i=0;i<n;i++){</pre>
    point<T> now=p[i+1]-p[i];
    while(now.cross(p[t+1]-p[i])>now.cross
         (p[t]-p[i]))t=(t+1)%n;
    ans=max(ans,(p[i]-p[t]).abs2());
  return p.pop_back(),ans;
T min_cover_rectangle(){//最小覆蓋矩形
```

```
int n=p.size(),t=1,r=1,l;
                                                           if(R-L<=1)return 0;</pre>
                                                           px[R]=q[R].line_intersection(q[L]);
       if(n<3)return 0;//也可以做最小周長矩形
212
                                                           for(int i=L;i<=R;++i)p.push back(px[i]); 324 struct line3D{</pre>
213
       T ans=1e99;p.push_back(p[0]);
                                                           return R-L+1;
       for(int i=0;i<n;i++){</pre>
214
215
         point<T> now=p[i+1]-p[i];
         while(now.cross(p[t+1]-p[i])>now.cross 270| };
216
                                                      template<typename T>
               (p[t]-p[i]))t=(t+1)%n;
                                                      struct triangle{
217
         while(now.dot(p[r+1]-p[i])>now.dot(p[r^{272}]
                                                        point<T> a,b,c;
               |-p[i]))r=(r+1)%n;
                                                        triangle(){}
218
         if(!i)l=r:
         while (now.dot(p[1+1]-p[i]) \le now.dot(p[275])
                                                        triangle(const point<T> &a,const point<T>
219
                                                              &b, const point<T> &c):a(a),b(b),c(c){} 331
              1]-p[i]))1=(1+1)%n;
                                                        T area()const{
220
         T d=now.abs2();
                                                          T t=(b-a).cross(c-a)/2;
         T tmp=now.cross(p[t]-p[i])*(now.dot(p[
221
                                                           return t>0?t:-t;
              r]-p[i])-now.dot(p[l]-p[i]))/d;
222
         ans=min(ans,tmp);
223
                                                        point<T> barycenter()const{//重心
                                                   280
224
       return p.pop back(),ans;
                                                   281
                                                          return (a+b+c)/3;
225
                                                   282
226
     T dis2(polygon &pl){//凸包最近距離平方
                                                        point<T> circumcenter()const{//外心
                                                   283
227
       vector<point<T> > &P=p,&Q=pl.p;
                                                   284
                                                          static line<T> u,v;
228
       int n=P.size(), m=Q.size(), l=0, r=0;
                                                           u.p1=(a+b)/2;
229
     for(int i=0;i<n;++i)if(P[i].y<P[1].y)l=i;</pre>
                                                           u.p2=point<T>(u.p1.x-a.y+b.y,u.p1.y+a.x-
                                                   286
                                                                                                      341
     for(int i=0;i<m;++i)if(Q[i].y<Q[r].y)r=i;</pre>
230
                                                               b.x);
       P.push_back(P[0]),Q.push_back(Q[0]);
231
                                                           v.p1=(a+c)/2;
232
       T ans=1e99;
                                                   288
                                                           v.p2=point<T>(v.p1.x-a.y+c.y,v.p1.y+a.x-343)
       for(int i=0;i<n;++i){</pre>
233
         while ((P[1]-P[1+1]) \cdot cross(Q[r+1]-Q[r]) 289
234
                                                           return u.line intersection(v);
              <0)r=(r+1)%m;
         ans=min(ans,line<T>(P[1],P[1+1]).
                                                        point<T> incenter()const{//內心
                                                   291
              seg dis2(line\langle T \rangle (Q[r],Q[r+1])));
                                                          T A=sqrt((b-c).abs2()),B=sqrt((a-c).abs2
                                                                                                      347
236
         l=(l+1)%n;
                                                                ()),C=sqrt((a-b).abs2());
237
                                                           return point<T>(A*a.x+B*b.x+C*c.x,A*a.y+
238
       return P.pop back(),Q.pop back(),ans;
                                                               B*b.y+C*c.y)/(A+B+C);
239
                                                   294
     static char sign(const point<T>&t){
                                                        point<T> perpencenter()const{//垂心
241
       return (t.y==0?t.x:t.y)<0;</pre>
                                                           return barvcenter()*3-circumcenter()*2:
                                                   296
242
                                                   297
     static bool angle cmp(const line<T>& A,
          const line<T>& B){
                                                      template<typename T>
244
       point<T> a=A.p2-A.p1,b=B.p2-B.p1;
                                                      struct point3D{
                                                                                                      355
245
       return sign(a)<sign(b) | | (sign(a) == sign(b)</pre>
                                                        T x, y, z;
            )&&a.cross(b)>0);
                                                        point3D(){}
                                                        point3D(const T&x,const T&y,const T&z):x(x 357
     int halfplane intersection(vector<line<T>
                                                             ),y(y),z(z){}
          > &s){//半平面交
                                                        point3D operator+(const point3D &b)const{
                                                                                                      358
                                                           return point3D(x+b.x,y+b.y,z+b.z);}
248
       sort(s.begin(),s.end(),angle cmp);//線段
            左側為該線段半平面
                                                        point3D operator-(const point3D &b)const{
                                                           return point3D(x-b.x,y-b.y,z-b.z);}
249
       int L,R,n=s.size();
                                                        point3D operator*(const T &b)const{
250
       vector<point<T> > px(n);
                                                           return point3D(x*b,y*b,z*b);}
       vector < line < T > > q(n);
251
                                                        point3D operator/(const T &b)const{
252
       q[L=R=0]=s[0];
                                                          return point3D(x/b,y/b,z/b);}
       for(int i=1;i<n;++i){</pre>
                                                        bool operator == (const point3D &b)const{
         while(L<R&&s[i].ori(px[R-1])<=0)--R;
254
                                                           return x==b.x&&y==b.y&&z==b.z;}
255
         while(L<R&&s[i].ori(px[L])<=0)++L;</pre>
                                                   314
                                                        T dot(const point3D &b)const{
256
         a[++R]=s[i];
                                                   315
                                                           return x*b.x+y*b.y+z*b.z;}
257
         if(q[R].parallel(q[R-1])){
                                                        point3D cross(const point3D &b)const{
258
                                                           return point3D(y*b.z-z*b.y,z*b.x-x*b.z,x
259
           if(q[R].ori(s[i].p1)>0)q[R]=s[i];
                                                                *b.y-y*b.x);}
260
261
         if(L<R)px[R-1]=q[R-1].
                                                        T abs2()const{//向量長度的平方
              line_intersection(q[R]);
                                                           return dot(*this);}
262
                                                        T area2(const point3D &b)const{//和b、原點
263
       while(L<R&&g[L].ori(px[R-1])<=0)--R;
                                                              圍成面積的平方
       p.clear():
                                                           return cross(b).abs2()/4;}
```

```
371
323 template<typename T>
                                                  372
    point3D<T> p1,p2;
                                                  373
     line3D(){}
     line3D(const point3D<T> &p1,const point3D<</pre>
         T> &p2):p1(p1),p2(p2){}
     T dis2(const point3D<T> &p, bool is segment 377
          =0) const { // 點 跟 直 線 / 線 段 的 距 離 平 方
                                                  378
                                                  379
       point3D < T > v = p2 - p1, v1 = p - p1;
       if(is_segment){
         point3D<T> v2=p-p2;
         if(v.dot(v1)<=0)return v1.abs2();</pre>
         if(v.dot(v2)>=0)return v2.abs2();
                                                  381
       point3D<T> tmp=v.cross(v1);
       return tmp.abs2()/v.abs2();
                                                  382
     pair<point3D<T>,point3D<T> > closest pair(
          const line3D<T> &1)const{
       point3D < T > v1 = (p1 - p2), v2 = (1.p1 - 1.p2);
       point3D<T> N=v1.cross(v2),ab(p1-l.p1);
                                                  386
       //if(N.abs2()==0)return NULL;平行或重合
                                                  387
       T tmp=N.dot(ab),ans=tmp*tmp/N.abs2();//
                                                  388
            最近點對距離
       point3D<T> d1=p2-p1,d2=l.p2-l.p1,D=d1.
           cross(d2),G=1.p1-p1;
       T t1=(G.cross(d2)).dot(D)/D.abs2();
                                                  389
       T t2=(G.cross(d1)).dot(D)/D.abs2();
                                                  390
       return make pair(p1+d1*t1,1.p1+d2*t2);
                                                  391
                                                  392
     bool same side(const point3D<T> &a,const
                                                  393
         point3D<T> &b)const{
                                                  394
       return (p2-p1).cross(a-p1).dot((p2-p1).
                                                  395
           cross(b-p1))>0;
                                                  396
  };
                                                  397
352 template<typename T>
                                                  398
353 struct plane{
    point3D<T> p0,n;//平面上的點和法向量
     plane(){}
     plane(const point3D<T> &p0, const point3D<T
                                                  402
         > &n):p0(p0),n(n){}
                                                  404
    T dis2(const point3D<T> &p)const{//點到平
                                                  405
          面距離的平方
                                                  406
       T tmp=(p-p0).dot(n);
                                                  407
       return tmp*tmp/n.abs2();
                                                  408
     point3D<T> projection(const point3D<T> &p)
                                                  410
                                                  411
       return p-n*(p-p0).dot(n)/n.abs2();
                                                  412
                                                 413
     point3D<T> line intersection(const line3D
                                                  414
         T> &1)const{
                                                  415
      T tmp=n.dot(1.p2-1.p1);//等於0表示平行或
                                                  416
            重合該平面
                                                  417
       return 1.p1+(1.p2-1.p1)*(n.dot(p0-1.p1)/
                                                  418
                                                  419
     line3D<T> plane intersection(const plane &
                                                  420
       point3D<T> e=n.cross(pl.n),v=n.cross(e);
       T tmp=pl.n.dot(v);//等於0表示平行或重合
            該平面
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```
point3D<T> q=p0+(v*(pl.n.dot(pl.p0-p0))/
       return line3D<T>(q,q+e);
374 };
375 template<tvpename T>
376 struct triangle3D{
     point3D<T> a.b.c:
     triangle3D(){}
     triangle3D(const point3D<T> &a,const
          point3D<T> &b, const point3D<T> &c):a(a
          ),b(b),c(c){}
     bool point_in(const point3D<T> &p)const{//
          點在該平面上的投影在三角形中
        return line3D<T>(b,c).same side(p,a)&&
            line3D<T>(a,c).same side(p,b)&&
            line3D<T>(a,b).same side(p,c);
383 };
384 template<typename T>
385 struct tetrahedron{//四面體
     point3D<T> a,b,c,d;
     tetrahedron(){}
     tetrahedron(const point3D<T> &a,const
          point3D<T> &b, const point3D<T> &c,
          const point3D<T> &d):a(a),b(b),c(c),d(
          d){}
     T volume6()const{//體積的六倍
       return (d-a).dot((b-a).cross(c-a));
     point3D<T> centroid()const{
       return (a+b+c+d)/4;
     bool point_in(const point3D<T> &p)const{
       return triangle3D<T>(a,b,c).point_in(p)
            &&triangle3D<T>(c,d,a).point_in(p);
   };
399 template<typename T>
400 struct convexhull3D{
     static const int MAXN=1005;
     struct face{
       int a,b,c;
       face(int a,int b,int c):a(a),b(b),c(c){}
     vector<point3D<T>> pt;
     vector<face> ans;
     int fid[MAXN][MAXN];
     void build(){
       int n=pt.size();
       ans.clear();
       memset(fid,0,sizeof(fid));
       ans.emplace back(0,1,2);//注意不能共線
       ans.emplace back(2,1,0);
       int ftop = 0;
       for(int i=3, ftop=1; i<n; ++i,++ftop){</pre>
         vector<face> next;
         for(auto &f:ans){
           T d=(pt[i]-pt[f.a]).dot((pt[f.b]-pt[
                f.a]).cross(pt[f.c]-pt[f.a]));
           if(d<=0) next.push back(f);</pre>
           int ff=0;
           if(d>0) ff=ftop;
           else if(d<0) ff=-ftop;</pre>
```

```
fid[f.a][f.b]=fid[f.c]=fid[f.c
                ][f.a]=ff;
425
         for(auto &f:ans){
426
427
           if(fid[f.a][f.b]>0 && fid[f.a][f.b
               ]!=fid[f.b][f.a])
             next.emplace back(f.a,f.b,i);
           if(fid[f.b][f.c]>0 && fid[f.b][f.c
               ]!=fid[f.c][f.b])
             next.emplace_back(f.b,f.c,i);
           if(fid[f.c][f.a]>0 && fid[f.c][f.a
                ]!=fid[f.a][f.c])
             next.emplace_back(f.c,f.a,i);
433
434
         ans=next;
435
436
    point3D<T> centroid()const{
      point3D<T> res(0,0,0);
439
       T vol=0:
       for(auto &f:ans){
         T`tmp=pt[f.a].dot(pt[f.b].cross(pt[f.c
         res=res+(pt[f.a]+pt[f.b]+pt[f.c])*tmp;
         vol+=tmp:
       return res/(vol*4);
446
447 };
```

1.2 SmallestCircle

```
using PT=point<T>; using CPT=const PT;
2 PT circumcenter(CPT &a,CPT &b,CPT &c){
   PT u=b-a, v=c-a;
  T c1=u.abs2()/2,c2=v.abs2()/2;
   T d=u.cross(v);
   return PT(a.x+(v.y*c1-u.y*c2)/d,a.y+(u.x*
        c2-v.x*c1)/d);
8 void solve(PT p[],int n,PT &c,T &r2){
   random shuffle(p,p+n);
   c=p[0]; r2=0; // c,r2 = 圓心,半徑平方
 for(int i=1;i<n;i++)if((p[i]-c).abs2()>r2){
     c=p[i]; r2=0;
 for(int j=0;j<i;j++)if((p[j]-c).abs2()>r2){
       c.x=(p[i].x+p[j].x)/2;
       c.y=(p[i].y+p[j].y)/2;
       r2=(p[j]-c).abs2();
 for(int k=0;k<j;k++)if((p[k]-c).abs2()>r2){
         c=circumcenter(p[i],p[j],p[k]);
         r2=(p[i]-c).abs2();
```

1.3 最近點對

```
1 template < typename IT = point < T > * >
  T cloest pair( IT L, IT R){
    if(R-L <= 1) return INF;</pre>
    IT mid = L+(R-L)/2;
    T x = mid -> x;
    T d = min(cloest pair(L,mid),cloest pair(
    inplace merge(L, mid, R, ycmp);
    static vector<point> b; b.clear();
    for(auto u=L;u<R;++u){</pre>
      if((u->x-x)*(u->x-x)>=d) continue;
      for(auto v=b.rbegin();v!=b.rend();++v){
        T dx=u\rightarrow x-v\rightarrow x, dy=u\rightarrow y-v\rightarrow y;
        if(dy*dy>=d) break;
        d=min(d,dx*dx+dy*dy);
      b.push back(*u);
    return d:
  T closest_pair(vector<point<T>> &v){
    sort(v.begin(),v.end(),xcmp);
    return closest_pair(v.begin(), v.end());
```

28

29

34

updated[node] = false;

if (start > end or start > r or end < 1)</pre>

2 Data Structure

2.1 segment tree

```
template <typename T>
class SegmentTree {
private:
  std::vector<T> tree, lazy;
  std::vector<bool> updated;
  int n:
  void build(int node, int start, int end,
       const std::vector<T>& arr) {
    if (start == end) {
      tree[node] = arr[start];
    } else {
      int mid = (start + end) / 2;
      build(2 * node, start, mid, arr);
      build(2 * node + 1, mid + 1, end, arr)
      tree[node] = tree[2 * node] + tree[2 *
            node + 1];
  void updateRange(int node, int start, int
       end, int 1, int r, T val) {
    if (updated[node]) {
                                               73
      tree[node] = lazy[node] * (end - start
           + 1);
      if (start != end) {
        lazy[node * 2] = lazy[node];
        lazy[node * 2 + 1] = lazy[node];
        updated[node * 2] = updated[node];
        updated[node * 2 + 1] = updated[node 78] };
            ];
```

```
if (start >= 1 and end <= r) {</pre>
    tree[node] = val * (end - start + 1);
    if (start != end) {
      lazy[node * 2] = val;
      lazy[node * 2 + 1] = val;
      updated[node * 2] = true;
      updated[node * 2 + 1] = true;
    return;
  int mid = (start + end) / 2;
  updateRange(node * 2, start, mid, 1, r,
  updateRange(node * 2 + 1, mid + 1, end,
      1, r, val);
  tree[node] = tree[node * 2] + tree[node
       * 2 + 1]:
T queryRange(int node, int start, int end,
      int 1, int r) {
  if (start > end or start > r or end < 1) 22</pre>
        return 0:
  if (updated[node]) {
    tree[node] = lazy[node] * (end - start 25
          + 1);
    if (start != end) {
      lazy[node * 2] = lazy[node];
      lazy[node * 2 + 1] = lazy[node];
      updated[node * 2] = updated[node];
      updated[node * 2 + 1] = updated[node
                                             32
    updated[node] = false;
                                             33
  if (start >= 1 and end <= r) return tree 35
       [node];
  int mid = (start + end) / 2;
  T p1 = queryRange(node * 2, start, mid,
      1, r);
  T p2 = queryRange(node * 2 + 1, mid + 1,
        end, 1, r);
  return p1 + p2;
                                             42
SegmentTree(const std::vector<T>& arr) {
 n = arr.size();
  tree.resize(4 * n);
  lazy.resize(4 * n);
  updated.resize(4 * n, false);
  build(1, 0, n - 1, arr);
void updateRange(int 1, int r, T val) {
                                             52
    updateRange(1, 0, n - 1, l, r, val); }
                                             53
T queryRange(int 1, int r) { return
                                             55
     queryRange(1, 0, n - 1, 1, r); }
                                             56
                                             57
```

2.2 DisjointSet

```
1 struct DisjointSetUnion {
    public:
      DisjointSetUnion()
        : n(0) {}
       explicit DisjointSetUnion(int n)
        : n(n)
        , parent or size(n, -1) {}
      int merge(int a, int b) {
        assert(0 <= a && a < n);
        assert(0 <= b && b < _n);
        int x = leader(a), y = leader(b);
        if (x == y) return x;
        if (-parent_or_size[x] < -</pre>
             parent or size[y])
           std::swap(x, y);
        parent_or_size[x] += parent_or_size[y
        parent_or_size[y] = x;
        return x;
       bool same(int a, int b) {
        assert(0 <= a && a < n);
        assert(0 <= b && b < _n);
        return leader(a) == leader(b);
      int leader(int a) {
        assert(0 <= a && a < _n);
        if (parent or size[a] < 0) return a;</pre>
        return parent or size[a] = leader(
             parent or size[a]);
      int size(int a) {
        assert(0 <= a && a < n);
        return -parent or size[leader(a)];
       std::vector<std::vector<int>> groups() {
        std::vector<int> leader buf( n),
              group_size(_n);
        for (int i = 0; i < _n; i++) {</pre>
          leader buf[i] = leader(i):
          group_size[leader_buf[i]]++;
         std::vector<std::vector<int>> result(
        for (int i = 0; i < _n; i++) {
          result[i].reserve(group size[i]);
        for (int i = 0; i < n; i++) {</pre>
          result[leader_buf[i]].push_back(i);
         result.erase(
          std::remove if(
           result.begin(), result.end(),
           [&](const std::vector<int>& v) {
           return v.empty();
           }),
          result.end());
         return result;
```

```
private:
    int _n;
    // root node: -1 * component size
    // otherwise: parent
    std::vector<int> parent_or_size;
};
```

2.3 undo disjoint set

```
1 struct DisjointSet {
    // save() is like recursive
    // undo() is like return
    int n, fa[MXN], sz[MXN];
    vector<pair<int*,int>> h;
    vector<int> sp;
    void init(int tn) {
      for (int i=0; i<n; i++) sz[fa[i]=i]=1;</pre>
      sp.clear(); h.clear();
    void assign(int *k, int v) {
      h.PB({k, *k});
      *k=v;
    void save() { sp.PB(SZ(h)); }
    void undo() {
      assert(!sp.empty());
      int last=sp.back(); sp.pop_back();
      while (SZ(h)!=last) {
        auto x=h.back(); h.pop_back();
        *x.F=x.S;
    int f(int x) {
      while (fa[x]!=x) x=fa[x];
    void uni(int x, int y) {
      x=f(x); y=f(y);
      if (x==y) return ;
      if (sz[x]<sz[y]) swap(x, y);</pre>
      assign(&sz[x], sz[x]+sz[y]);
      assign(&fa[y], x);
36 }djs;
```

2.4 fenwick

```
namespace internal {
template <class T>
using to_unsigned_t = typename to_unsigned<T
>::type;
} // namespace internal

template <class T>
struct FenwickTree {
using U = internal::to_unsigned_t<T>;
}
```

```
public:
 FenwickTree() : n(0) {}
 explicit FenwickTree(int n) : n(n), data( 23
 void add(int p, T x) {
   assert(0 \le p \&\& p < n);
   p++;
   while (p <= n) {
     data[p - 1] += U(x);
    p += p \& -p;
 T sum(int 1, int r) {
   assert(0 <= 1 && 1 <= r && r <= n);
   return sum(r) - sum(1);
private:
int n;
 std::vector<U> data:
 U sum(int r) {
  U s = 0;
   while (r > 0) {
    s += data[r - 1];
    r -= r & -r;
   return s;
```

3 Flow

3.1 min cost flow

```
namespace internal {
  template <class T>
  struct simple queue {
    std::vector<T> payload;
    int pos = 0:
    void reserve(int n) { payload.reserve(n);
    int size() const { return int(payload.size
         ()) - pos; }
    bool empty() const { return pos == int(
         payload.size()); }
    void push(const T& t) { payload.push_back(
         t); }
    T& front() { return payload[pos]; }
    void clear() {
      payload.clear();
      pos = 0;
                                                 73
    void pop() { pos++; }
  template <class E>
20 struct csr {
```

```
32
      for (auto e : edges) {
         elist[counter[e.first]++] = e.second;
34
35
36
  };
  } // namespace internal
  template <class Cap, class Cost>
  struct MinCostFlowGraph {
   public:
    MinCostFlowGraph() {}
    explicit MinCostFlowGraph(int n) : n(n)
                                                    97
                                                    98
    int add edge(int from, int to, Cap cap,
                                                    99
         Cost cost) {
                                                   100
      assert(0 <= from && from < n);</pre>
                                                   101
      assert(0 <= to && to < _n);
                                                   102
      assert(0 <= cap);</pre>
      assert(0 <= cost);</pre>
                                                   103
      int m = int(_edges.size());
      edges.push back({from, to, cap, 0, cost 104
           });
      return m;
53
54
                                                   107
                                                   108
     struct edge {
      int from, to;
                                                   109
      Cap cap, flow;
                                                   110
      Cost cost;
                                                   111
                                                   112
    edge get edge(int i) {
                                                   113
      int m = int( edges.size());
                                                   114
      assert(0 <= i && i < m);
                                                   115
      return edges[i];
    std::vector<edge> edges() { return _edges;
                                                  118
    std::pair<Cap, Cost> flow(int s, int t) {
      return flow(s, t, std::numeric limits<</pre>
                                                   122
           Cap>::max());
                                                   123
                                                   124
    std::pair<Cap, Cost> flow(int s, int t,
                                                   125
          Cap flow limit) {
                                                   126
      return slope(s, t, flow limit).back();
                                                   127
    std::vector<std::pair<Cap, Cost>> slope(
          int s, int t) {
       return slope(s, t, std::numeric_limits<</pre>
           Cap>::max()):
```

std::vector<int> start;

for (auto e : edges) {

auto counter = start;

start[e.first + 1]++;

explicit csr(int n, const std::vector<std

: start(n + 1), elist(edges.size()) {

81

83

::pair<int, E>>& edges)

for (int i = 1; i <= n; i++) {</pre>

start[i] += start[i - 1];

std::vector<E> elist;

```
std::vector<std::pair<Cap, Cost>> slope(
     int s, int t, Cap flow limit) {
  assert(0 <= s && s < n):
  assert(0 <= t && t < _n);
  assert(s != t);
  int m = int( edges.size());
  std::vector<int> edge idx(m);
  auto g = [&]() {
    std::vector<int> degree(_n), redge_idx
    std::vector<std::pair<int, _edge>>
         elist:
    elist.reserve(2 * m);
    for (int i = 0; i < m; i++) {</pre>
      auto e = _edges[i];
      edge_idx[i] = degree[e.from]++;
      redge_idx[i] = degree[e.to]++;
      elist.push_back({e.from, {e.to, -1,
           e.cap - e.flow, e.cost}});
      elist.push_back({e.to, {e.from, -1,
          e.flow, -e.cost}});
    auto g = internal::csr< edge>( n,
         elist);
    for (int i = 0; i < m; i++) {</pre>
      auto e = _edges[i];
      edge_idx[i] += _g.start[e.from];
      redge idx[i] += g.start[e.to];
      _g.elist[edge_idx[i]].rev =
           redge_idx[i];
      g.elist[redge idx[i]].rev =
           edge_idx[i];
    return g;
  }();
  auto result = slope(g, s, t, flow_limit)
  for (int i = 0; i < m; i++) {</pre>
    auto e = g.elist[edge idx[i]];
    edges[i].flow = edges[i].cap - e.cap
  return result;
std::vector<edge> edges;
// inside edge
struct edge {
  int to, rev;
  Cap cap;
  Cost cost;
std::vector<std::pair<Cap, Cost>> slope(
     internal::csr< edge>& g, int s, int t,
      Cap flow limit) {
  // variants (C = maxcost):
```

 $// -(n-1)C \leftarrow dual[s] \leftarrow dual[i] \leftarrow dual$ 183

```
[t] = 0
       // reduced cost (= e.cost + dual[e.from] 184
              - dual[e.to]) >= 0 for all edge
133
       // dual dist[i] = (dual[i], dist[i])
135
       std::vector<std::pair<Cost, Cost>>
             dual dist( n):
       std::vector<int> prev_e(_n);
136
137
       std::vector<bool> vis( n);
138
       struct Q {
                                                     189
139
          Cost key;
         int to;
140
141
          bool operator<(Q r) const { return key</pre>
               > r.kev; }
142
143
       std::vector<int> que min;
144
       std::vector<Q> que;
                                                     195
       auto dual_ref = [&]() {
                                                     196
146
         for (int i = 0; i < _n; i++) {</pre>
                                                     197
147
            dual dist[i].second = std::
                                                     198
                 numeric limits<Cost>::max();
                                                     199
148
                                                     200
149
          std::fill(vis.begin(), vis.end(),
              false);
         que_min.clear();
150
                                                     203
151
         que.clear();
152
153
         // que[0..heap_r) was heapified
154
          size t heap r = 0;
155
156
         dual_dist[s].second = 0;
157
         que min.push back(s);
          while (!que_min.empty() || !que.empty
158
               ()) {
            int v;
159
           if (!que_min.empty()) {
160
                                                     209
161
              v = que min.back();
162
              que_min.pop_back();
                                                     210
163
            } else {
                                                     211
164
              while (heap r < que.size()) {</pre>
                                                     212
165
                heap_r++;
                std::push_heap(que.begin(), que. 214
166
                     begin() + heap r);
              v = que.front().to;
              std::pop heap(que.begin(), que.end 217
                   ());
              que.pop back();
171
              heap_r--;
172
            if (vis[v]) continue;
173
174
           vis[v] = true;
                                                     221
175
           if (v == t) break;
176
           // dist[v] = shortest(s, v) + dual[s]
                 1 - dual[v]
            // dist[v] >= 0 (all reduced cost
                                                     224
                 are positive)
                                                     225
            // dist[v] <= (n-1)C
                                                     226
            Cost dual v = dual dist[v].first.
                 dist_v = dual_dist[v].second;
            for (int i = g.start[v]; i < g.start 229</pre>
                 [v + 1]; i++) {
                                                     230
              auto e = g.elist[i];
181
                                                     231
              if (!e.cap) continue:
                                                     232
```

```
// |-dual[e.to] + dual[v]| <= (n
           -1)C
      // \cos t <= C - -(n-1)C + 0 = nC
      Cost cost = e.cost - dual dist[e.
           to].first + dual v;
      if (dual dist[e.to].second -
          dist v > cost) {
        Cost dist to = dist v + cost;
        dual dist[e.to].second = dist to
        prev e[e.to] = e.rev;
        if (dist_to == dist_v) {
          que_min.push_back(e.to);
       } else {
          que.push_back(Q{dist_to, e.to
              });
  if (!vis[t]) {
   return false;
  for (int v = 0; v < n; v++) {
   if (!vis[v]) continue;
   // dual[v] = dual[v] - dist[t] +
        dist[v]
               = dual[v] - (shortest(s,
        t) + dual[s] - dual[t]) +
               (shortest(s, v) + dual[s]
         - dual[v]) = - shortest(s,
               t) + dual[t] + shortest(s
        , v) = shortest(s, v) -
               shortest(s, t) >= 0 - (n
    dual dist[v].first -= dual dist[t].
        second - dual dist[v].second;
  return true:
Cap flow = 0;
Cost cost = 0, prev_cost_per_flow = -1;
std::vector<std::pair<Cap, Cost>> result
      = {{Cap(0), Cost(0)}};
while (flow < flow_limit) {</pre>
 if (!dual ref()) break;
  Cap c = flow limit - flow;
  for (int v = t; v != s; v = g.elist[
      prev_e[v]].to) {
    c = std::min(c, g.elist[g.elist[
        prev e[v]].rev].cap);
  for (int v = t; v != s; v = g.elist[
      prev e[v]].to) {
    auto& e = g.elist[prev_e[v]];
    e.cap += c:
    g.elist[e.rev].cap -= c;
  Cost d = -dual dist[s].first;
  flow += c;
  cost += c * d;
  if (prev cost per flow == d) {
   result.pop back();
  result.push_back({flow, cost});
```

```
235
236
       return result:
237
    }
238 };
                                                   51
                                                   52
   3.2 max flow
 namespace internal {
   template <class T>
                                                   59
   struct simple queue {
                                                   60
     std::vector<T> payload;
     int pos = 0;
     void reserve(int n) { payload.reserve(n);
     int size() const { return int(payload.size 64
          ()) - pos; }
     bool empty() const { return pos == int(
          payload.size()); }
     void push(const T& t) { payload.push back(
          t); }
     T& front() { return payload[pos]; }
                                                   68
     void clear() {
                                                   69
                                                   70
       payload.clear();
       pos = 0;
     void pop() { pos++; }
 17 };
 19 } // namespace internal
   template <class Cap>
   struct MaxFlowGraph {
22
    public:
23
     MaxFlowGraph() : MaxFlowGraph(0) {}
                                                   80
24
     explicit MaxFlowGraph(int n) : n(n), g(n)
                                                  81
     int add edge(int from, int to, Cap cap) {
                                                   83
       assert(0 <= from && from < _n);
                                                   84
       assert(0 <= to && to < n);
                                                   85
       assert(0 <= cap):
                                                   86
30
       int m = int(pos.size());
31
       pos.push back({from, int(g[from].size())
            });
       int from_id = int(g[from].size());
       int to id = int(g[to].size());
34
       if (from == to) to_id++;
35
       g[from].push back( edge{to, to id, cap})
       g[to].push_back(_edge{from, from_id, 0})
       return m;
 38
 39
     struct edge {
                                                   98
       int from, to;
                                                   99
       Cap cap, flow;
                                                  100
43
44
                                                  101
     edge get edge(int i) {
       int m = int(pos.size());
```

prev cost per flow = d;

234

```
assert(0 <= i && i < m);
  auto _e = g[pos[i].first][pos[i].second
  auto _re = g[_e.to][_e.rev];
  return edge{pos[i].first, _e.to, _e.cap
      + re.cap, re.cap};
std::vector<edge> edges() {
  int m = int(pos.size());
  std::vector<edge> result;
  for (int i = 0; i < m; i++) {</pre>
    result.push back(get edge(i));
 return result;
void change_edge(int i, Cap new_cap, Cap
     new flow) {
  int m = int(pos.size());
  assert(0 <= i && i < m);
  assert(0 <= new_flow && new_flow <=
       new cap);
  auto& _e = g[pos[i].first][pos[i].second
  auto& _re = g[_e.to][_e.rev];
  e.cap = new cap - new flow;
 _re.cap = new_flow;
Cap flow(int s, int t) { return flow(s, t,
     std::numeric limits<Cap>::max()); }
Cap flow(int s, int t, Cap flow_limit) {
  assert(0 <= s && s < n);
  assert(0 <= t && t < n);
  assert(s != t);
  std::vector<int> level(_n), iter(_n);
 internal::simple_queue<int> que;
  auto bfs = [&]() {
   std::fill(level.begin(), level.end(),
    level[s] = 0;
    que.clear();
    aue.push(s):
    while (!que.empty()) {
     int v = que.front();
      que.pop();
     for (auto e : g[v]) {
   if (e.cap == 0 || level[e.to] >=
             0) continue;
        level[e.to] = level[v] + 1;
        if (e.to == t) return;
        que.push(e.to);
  auto dfs = [&](auto self, int v, Cap up)
    if (v == s) return up;
    Cap res = 0:
    int level_v = level[v];
    for (int& i = iter[v]; i < int(g[v]).
         size()); i++) {
       _edge& e = g[v][i];
      if (level v <= level[e.to] || g[e.to</pre>
           ][e.rev].cap == 0) continue;
```

```
- res, g[e.to][e.rev].cap));
           if (d <= 0) continue;</pre>
105
           g[v][i].cap += d;
           g[e.to][e.rev].cap -= d;
106
           res += d:
           if (res == up) return res;
109
110
         level[v] = n;
111
         return res;
112
113
114
       Cap flow = 0;
       while (flow < flow limit) {</pre>
116
         bfs();
         if (level[t] == -1) break;
         std::fill(iter.begin(), iter.end(), 0)
118
119
          Cap f = dfs(dfs, t, flow limit - flow)
         if (!f) break;
120
         flow += f;
121
122
123
       return flow:
124
125
     std::vector<bool> min cut(int s) {
       std::vector<bool> visited( n);
127
       internal::simple_queue<int> que;
129
       que.push(s);
       while (!que.empty()) {
130
         int p = que.front();
132
         que.pop();
133
         visited[p] = true;
134
         for (auto e : g[p]) {
           if (e.cap && !visited[e.to]) {
              visited[e.to] = true;
              que.push(e.to);
139
       return visited;
    private:
     int n;
     struct edge {
       int to, rev;
       Cap cap;
     };
     std::vector<std::pair<int, int>> pos;
     std::vector<std::vector< edge>> g;
```

4 Graph

4.1 tree centroid

```
1 // 这份代码默认节点编号从 1 开始·即 i @ [1,
```

```
Cap d = self(self, e.to, std::min(up 2 int size[MAXN], // 这个节点的「大小」(所有 29|)
                                       子树上节点数 + 该节点)
                                      weight[MAXN], // 这个节点的「重量」,即
                                          所有子树「大小」的最大值
                                      centroid[2]; // 用于记录树的重心(存的
                                          是节点编号)
                                   void GetCentroid(int cur, int fa) { // cur
                                       表示当前节点 (current)
                                     size[cur] = 1:
                                     weight[cur] = 0;
                                     for (int i = head[cur]; i != -1; i = e[i].
                                      if (e[i].to != fa) { // e[i].to 表示这
                                          条有向边所通向的节点。
                                        GetCentroid(e[i].to, cur);
                                        size[cur] += size[e[i].to];
                                        weight[cur] = max(weight[cur], size[e[
                                            i].to]);
                                    weight[cur] = max(weight[cur], n - size[
                                    if (weight[cur] <= n / 2) { // 依照树的重
                                         心的定义统计
                                      centroid[centroid[0] != 0] = cur;
```

4.2 tarjan

```
i int dfn[N], low[N], dfncnt, s[N], in stack[N
      ], tp;
2| int scc[N], sc; // 结点 i 所在 SCC 的编号
                  // 强连通 i 的大小
 int sz[N];
 void tarjan(int u) {
   low[u] = dfn[u] = ++dfncnt, s[++tp] = u,
        in_stack[u] = 1;
   for (int i = h[u]; i; i = e[i].nex) {
     const int &v = e[i].t;
     if (!dfn[v]) {
       tarian(v):
       low[u] = min(low[u], low[v]);
     } else if (in_stack[v]) {
       low[u] = min(low[u], dfn[v]);
   if (dfn[u] == low[u]) {
     ++sc;
     while (s[tp] != u) {
       scc[s[tp]] = sc;
       sz[sc]++;
       in stack[s[tp]] = 0;
       --tp;
     scc[s[tp]] = sc;
     sz[sc]++;
     in_stack[s[tp]] = 0;
     --tp;
```

4.3 heavy light decomposition

 $1 \mid const int kMax = 1e5 + 5;$

```
2 int n, m, r, kMod, op, x, y, z, cnt = 0;
3 int a[kMax] = {}, dep[kMax] = {}, fa[kMax] =
        \{\}, siz[kMax] = \{\},
      hson[kMax] = {};
  int nid[kMax] = {}, na[kMax] = {}, ltop[kMax
  int seg[kMax << 2] = {}, tag[kMax << 2] =</pre>
  vector<int> g[kMax];
  void dfs1(int u, int f) {
    dep[u] = dep[f] + 1;
    fa[u] = f;
                                                   67
    siz[u] = 1;
                                                   68
    for (auto v : g[u]) {
      if (v == f) continue;
      dfs1(v, u);
      siz[u] += siz[v];
      if (siz[v] > siz[hson[u]]) hson[u] = v;
19
   void dfs2(int u, int cur ltop) {
    nid[u] = ++cnt;
    na[cnt] = a[u];
                                                   75
    ltop[u] = cur ltop;
    if (hson[u]) dfs2(hson[u], cur ltop);
    for (auto v : g[u]) {
  if (v == fa[u] || v == hson[u]) continue
      dfs2(v, v);
29
30 }
                                                   81
                                                   82
  void build(int u, int l, int r) {
                                                   83
    if (1 == r) {
      seg[u] = na[1] % kMod;
34
35
      return:
    int mid = (1 + r) >> 1:
38
    build(u << 1, 1, mid);
    build(u \ll 1 | 1, mid + 1, r);
    seg[u] = (seg[u << 1] + seg[u << 1 | 1]) %
  void push_down(int u, int 1, int r) {
43
    if (tag[u])
                                                   93
      int mid = (1 + r) >> 1;
      seg[u << 1] = (seg[u << 1] + tag[u] * (
           mid - 1 + 1)) % kMod;
      seg[u << 1 | 1] = (seg[u << 1 | 1] + tag
           [u] * (r - mid)) % kMod;
      tag[u << 1] = (tag[u << 1] + tag[u]) %
      tag[u << 1 | 1] = (tag[u << 1 | 1] + tag 100
           [u]) % kMod;
      tag[u] = 0;
```

```
51 }
52 }
53
54 int query(int u, int l, int r, int ql, int
    if (ql <= 1 && r <= qr) return seg[u];</pre>
    push down(u, 1, r);
    int mid = (1 + r) >> 1, res = 0;
    if (ql <= mid) res = (res + query(u << 1,</pre>
         1, mid, ql, qr)) % kMod;
    if (qr > mid) res = (res + query(u << 1 |
         1, mid + 1, r, ql, qr)) % kMod;
    return res;
61 }
63 void update(int u, int l, int r, int ql, int
         qr, int k) {
    if (q1 <= 1 && r <= qr) {</pre>
       seg[u] = (seg[u] + k * (r - 1 + 1)) %
       tag[u] = (tag[u] + k) % kMod;
      return:
    push_down(u, 1, r);
    int mid = (1 + r) >> 1;
    if (ql <= mid) update(u << 1, 1, mid, ql,</pre>
          qr, k);
    if (qr > mid) update(u << 1 | 1, mid + 1,</pre>
          r, ql, qr, k);
    seg[u] = (seg[u << 1] + seg[u << 1 | 1]) %
           kMod;
74
   int query_path(int u, int v) {
    int res = 0:
     while (ltop[u] != ltop[v]) {
       if (dep[ltop[u]] < dep[ltop[v]]) swap(u,</pre>
             v);
       res = (res + query(1, 1, n, nid[ltop[u
            ]], nid[u])) % kMod;
       u = fa[ltop[u]];
     if (dep[u] > dep[v]) swap(u, v);
     res = (res + query(1, 1, n, nid[u], nid[v])
         ])) % kMod;
    return res;
   void update path(int u, int v, int k) {
    k %= kMod;
     while (ltop[u] != ltop[v]) {
       if (dep[ltop[u]] < dep[ltop[v]]) swap(u,</pre>
       update(1, 1, n, nid[ltop[u]], nid[u], k)
      u = fa[ltop[u]];
    if (dep[u] > dep[v]) swap(u, v);
    update(1, 1, n, nid[u], nid[v], k);
  int query subtree(int u) { return query(1,
       1, n, nid[u], nid[u] + siz[u] - 1); }
void update subtree(int u, int k) {
```

19 void sieve(void){

primecnt = 0:

phi[1] = mu[1] = 1;

if(!iscom[i]) {

mu[i] = -1;

phi[i] = i-1;

memset(iscom,0,sizeof(iscom));

for(int i=2;i<MAXPRIME;++i) {</pre>

prime[primecnt++] = i;

int k = i * prime[i];

if(k>=MAXPRIME) break;

iscom[k] = prime[j];

if(i%prime[j]==0) {

mu[k] = -mu[i];

mu[k] = 0;

break;

} else {

for(int j=0;j<primecnt;++j) {</pre>

phi[k] = phi[i] * prime[j];

phi[k] = phi[i] * (prime[j]-1);

```
update(1, 1, n, nid[u], nid[u] + siz[u] -
          1, k);
103 }
104
   int32 t main() {
105
     cin.tie(nullptr)->sync with stdio(false);
     cin >> n >> m >> r >> kMod;
     for (int i = 1; i <= n; i++) cin >> a[i];
     for (int i = 1; i < n; i++) {</pre>
       cin >> x >> y;
111
       g[x].push_back(y);
112
       g[y].push_back(x);
113
114
     dfs1(r, 0);
115
     dfs2(r, r);
     build(1, 1, n);
     while (m--) {
       cin >> op >> x;
120
121
       if (op == 1) {
122
          cin >> y >> z;
          update_path(x, y, z);
123
       } else if (op == 2) {
124
125
          cin >> y;
126
          cout << query_path(x, y) << '\n';</pre>
127
       } else if (op == 3) {
128
          cin >> z;
129
          update_subtree(x, z);
130
       } else {
131
          cout << query_subtree(x) << '\n';</pre>
132
133
134
     return 0;
```

5 Number Theory

5.1 basic

```
1 template<tvpename T>
2 void gcd(const T &a,const T &b,T &d,T &x,T &
    if(!b) d=a, x=1, y=0;
    else gcd(b,a%b,d,y,x), y-=x*(a/b);
 long long int phi[N+1];
  void phiTable(){
    for(int i=1;i<=N;i++)phi[i]=i;</pre>
    for(int i=1;i<=N;i++)for(x=i*2;x<=N;x+=i)</pre>
         phi[x]-=phi[i];
void all_divdown(const LL &n) {// all n/x
   for(LL a=1;a<=n;a=n/(n/(a+1))){</pre>
      // dosomething;
16 const int MAXPRIME = 1000000;
int iscom[MAXPRIME], prime[MAXPRIME],
       primecnt:
int phi[MAXPRIME], mu[MAXPRIME];
```

```
bool g test(const LL &g, const LL &p, const
    vector<LL> &v) {
  for(int i=0;i<v.size();++i)</pre>
   if(modexp(g,(p-1)/v[i],p)==1)
      return false;
  return true;
LL primitive_root(const LL &p) {
 if(p==2) return 1;
  vector<LL> v;
  Factor(p-1,v);
  v.erase(unique(v.begin(), v.end()), v.end
  for(LL g=2;g<p;++g)</pre>
   if(g test(g,p,v))
      return g;
  puts("primitive root NOT FOUND");
  return -1:
int Legendre(const LL &a, const LL &p) {
    return modexp(a\%p,(p-1)/2,p); }
LL inv(const LL &a, const LL &n) {
 LL d,x,y;
  gcd(a,n,d,x,y);
 return d==1 ? (x+n)%n : -1;
int inv[maxN];
LL invtable(int n,LL P){
 inv[1]=1;
  for(int i=2;i<n;++i)</pre>
    inv[i]=(P-(P/i))*inv[P%i]%P;
LL log mod(const LL &a, const LL &b, const
    LL &p) {
  // a ^ x = b \pmod{p}
 int m=sart(p+.5), e=1;
 LL v=inv(modexp(a,m,p), p);
```

```
map<LL,int> x;
     x[1]=0;
     for(int i=1:i<m:++i) {</pre>
       e = LLmul(e,a,p);
85
       if(!x.count(e)) x[e] = i;
     for(int i=0;i<m;++i) {</pre>
       if(x.count(b)) return i*m + x[b];
       b = LLmul(b,v,p);
 90
     return -1;
92 }
 94 LL Tonelli Shanks(const LL &n, const LL &p)
     // x^2 = n \pmod{p}
     if(n==0) return 0:
     if(Legendre(n,p)!=1) while(1) { puts("SQRT
           ROOT does not exist"); }
     int S = 0;
     LL 0 = p-1;
     while( !(Q&1) ) { Q>>=1; ++S; }
     if(S==1) return modexp(n\%p,(p+1)/4,p);
102
     LL z = 2:
     for(;Legendre(z,p)!=-1;++z)
     LL c = modexp(z,Q,p);
     LL R = modexp(n\%p,(Q+1)/2,p), t = modexp(n
          %p,Q,p);
     int M = S;
     while(1) {
107
       if(t==1) return R;
108
       LL b = modexp(c,1L << (M-i-1),p);
109
       R = LLmul(R,b,p);
       t = LLmul( LLmul(b,b,p), t, p);
       c = LLmul(b,b,p);
112
113
       M = i;
     return -1;
   template<typename T>
   T Euler(T n){
     T ans=n;
120
     for(T i=2;i*i<=n;++i){</pre>
121
122
       if(n%i==0){
         ans=ans/i*(i-1);
123
124
         while(n%i==0)n/=i;
125
126
127
     if(n>1)ans=ans/n*(n-1);
     return ans;
128
129
   //Chinese remainder theorem
   template<typename T>
   T pow mod(T n, T k, T m){
134
     T ans=1:
     for(n=(n)=m?n\%m:n);k;k>>=1){
       if(k&1)ans=ans*n%m;
136
137
       n=n*n%m;
138
     return ans;
140 }
141 template<typename T>
142 T crt(vector<T> &m.vector<T> &a){
    T M=1,tM,ans=0;
```

```
for(int i=0;i<(int)m.size();++i)M*=m[i];</pre>
145
     for(int i=0;i<(int)a.size();++i){</pre>
146
       tM=M/m[i]:
       ans=(ans+(a[i]*tM%M)*pow_mod(tM,Euler(m[
            i])-1,m[i])%M)%M;
       /*如果m[i]是質數·Euler(m[i])-1=m[i]-2·
            就不用算Euler了*/
149
150
     return ans;
151 }
152
153 //java code
154 / / 求 s g r t (N) 的 連 分 數
public static void Pell(int n){
     BigInteger N,p1,p2,q1,q2,a0,a1,a2,g1,g2,h1
          ,h2,p,q;
     g1=q2=p1=BigInteger.ZERO;
     h1=q1=p2=BigInteger.ONE;
159
     a0=a1=BigInteger.valueOf((int)Math.sqrt
          (1.0*n));
     BigInteger ans=a0.multiply(a0);
160
     if(ans.equals(BigInteger.valueOf(n))){
       System.out.println("No solution!");
162
163
       return ;
164
     while(true){
165
       g2=a1.multiply(h1).substract(g1);
       h2=N.substract(g2.pow(2)).divide(h1);
168
       a2=g2.add(a0).divide(h2);
       p=a1.multiply(p2).add(p1);
169
       q=a1.multiply(q2).add(q1);
170
       if(p.pow(2).substract(N.multiply(q.pow
171
            (2))).compareTo(BigInteger.ONE)==0)
            break:
172
       g1=g2;h1=h2;a1=a2;
173
       p1=p2;p2=p;
174
       q1=q2;q2=q;
175
176
     System.out.println(p+" "+q);
177 }
```

5.2 bit set

```
1 void sub_set(int S){
    int sub=S;
    do{
      //對某集合的子集合的處理
      sub=(sub-1)&S;
    }while(sub!=S);
8 void k sub set(int k,int n){
    int comb=(1<<k)-1,S=1<<n;</pre>
    while(comb<S){</pre>
      //對大小為k的子集合的處理
      int x=comb&-comb,v=comb+x;
12
      comb = ((comb \& y)/x > 1) | y;
13
14
15 }
```

5.3 FFT

```
i template<typename T, typename VT=vector<</pre>
       complex<T>>>
  struct FFT{
   const T pi;
   FFT(const T pi=acos((T)-1)):pi(pi){}
    unsigned bit reverse(unsigned a,int len){
  a=((a&0x55555555U)<<1)|((a&0xAAAAAAAAU)>>1);
  a=((a&0x33333333U)<<2)|((a&0xCCCCCCCU)>>2);
8 a=((a&0x0F0F0F0FU)<<4)|((a&0xF0F0F0F0U)>>4); 32
9 a=((a&0x00FF00FFU)<<8)|((a&0xFF00FF00U)>>8); 33
a = ((a\&0x0000FFFFU) < (16))((a\&0xFFFF0000U))
       >>16);
      return a>>(32-len);
    void fft(bool is_inv,VT &in,VT &out,int N)
      int bitlen= lg(N), num=is inv?-1:1;
      for(int i=0;i<N;++i)out[bit reverse(i,</pre>
           bitlen)]=in[i];
      for(int step=2;step<=N;step<<=1){</pre>
        const int mh=step>>1;
        for(int i=0;i<mh;++i){</pre>
          complex<T> wi=exp(complex<T>(0,i*num
               *pi/mh));
          for(int j=i;j<N;j+=step){</pre>
            int k=j+mh;
            complex<T> u=out[j],t=wi*out[k];
            out[j]=u+t;
            out[k]=u-t;
      if(is inv)for(int i=0;i<N;++i)out[i]/=N;</pre>
30 };
```

5.4 質因數分解

```
1 | LL func(const LL n,const LL mod,const int c)
    return (LLmul(n,n,mod)+c+mod)%mod;
5 LL pollorrho(const LL n, const int c) {//循
       環節長度
    LL a=1, b=1;
    a=func(a,n,c)%n:
    b=func(b,n,c)%n; b=func(b,n,c)%n;
    while(gcd(abs(a-b),n)==1) {
     a=func(a,n,c)%n;
     b=func(b,n,c)%n; b=func(b,n,c)%n;
   return gcd(abs(a-b),n);
void prefactor(LL &n, vector<LL> &v) {
   for(int i=0;i<12;++i) {</pre>
      while(n%prime[i]==0) {
        v.push back(prime[i]);
        n/=prime[i];
```

```
5.5 find real root
```

2 int sign(double x){

 $1 / / an*x^n + ... + a1x + a0 = 0;$

return x < -eps ? -1 : x > eps;

86 }

void smallfactor(LL n, vector<LL> &v) {

for(int i=0;i<primecnt&&prime[i]*prime[i</pre>

v.push_back(isp[(int)n]);

if(n<MAXPRIME) {</pre>

while(isp[(int)n]) {

n/=isp[(int)n];

]<=n;++i) {</pre>

n/=prime[i];

smallfactor(n,v);

if(n!=1) v.push_back(n);

while(n%prime[i]==0) {

v.push back(prime[i]);

void comfactor(const LL &n, vector<LL> &v) {

void Factor(const LL &x, vector<LL> &v) {

if(n==1) { puts("Factor 1"); return; }

void AllFactor(const LL &n, vector<LL> &v) {

for(int i=0;i<tmp.size();++i) {</pre>

len = v.size():

for(int j=0;j<len;++j)</pre>

v.push back(v[j]*now);

if(i==0 || tmp[i]!=tmp[i-1]) {

v.push back(n);

} else {

if(n<1e9) {

return;

return:

LL d;

if(Isprime(n)) {

v.push back(n);

for(int c=3;;++c) {
 d = pollorrho(n,c);

if(d!=n) break;

comfactor(d,v);

prefactor(n,v);

if(n==1) return;

comfactor(n,v);

vector<LL> tmp;

v.push back(1);

now = 1;

now*=tmp[i];

Factor(n,tmp);

v.clear();

int len;

sort(v.begin(),v.end());

comfactor(n/d,v);

```
6 double get(const vector<double>&coef, double
    double e = 1, s = 0:
    for(auto i : coef) s += i*e, e *= x;
    return s;
12 double find(const vector<double>&coef, int n
       , double lo, double hi){
    double sign_lo, sign_hi;
    if( !(sign lo = sign(get(coef,lo))) )
         return lo;
    if( !(sign_hi = sign(get(coef,hi))) )
         return hi:
    if(sign lo * sign hi > 0) return INF;
    for(int stp = 0; stp < 100 && hi - lo >
         eps; ++stp){
      double m = (lo+hi)/2.0;
      int sign_mid = sign(get(coef,m));
      if(!sign_mid) return m;
      if(sign lo*sign mid < 0) hi = m;</pre>
      else lo = m:
23
    return (lo+hi)/2.0;
  vector<double> cal(vector<double>coef, int n
    vector<double>res;
    if(n == 1){
      if(sign(coef[1])) res.pb(-coef[0]/coef
           [1]);
      return res;
32
    vector<double>dcoef(n);
    for(int i = 0; i < n; ++i) dcoef[i] = coef</pre>
         [i+1]*(i+1);
    vector<double>droot = cal(dcoef, n-1);
    droot.insert(droot.begin(), -INF);
    droot.pb(INF);
    for(int i = 0; i+1 < droot.size(); ++i){</pre>
      double tmp = find(coef, n, droot[i],
           droot[i+1]);
      if(tmp < INF) res.pb(tmp);</pre>
    return res;
  int main () {
    vector<double>ve;
    vector<double>ans = cal(ve, n);
    // 視情況把答案 +eps,避免 -0
```

5.6 中國剩餘定理

```
int exgcd(int a, int b, int &x, int &y) {
    int x1 = 1, x2 = 0, x3 = 0, x4 = 1;
    while (b != 0) {
      int c = a / b;
      std::tie(x1, x2, x3, x4, a, b) =
          std::make tuple(x3, x4, x1 - x3 * c,
                x2 - x4 * c, b, a - b * c);
    x = x1, y = x2;
    return a;
11 LL CRT(int k, LL* a, LL* r) {
    LL n = 1, ans = 0:
    for (int i = 1; i <= k; i++) n = n * r[i];</pre>
    for (int i = 1; i <= k; i++) {</pre>
      LL m = n / r[i], b, y;
      exgcd(m, r[i], b, y); // b * m mod r[i]
      ans = (ans + a[i] * m * b % n) % n;
18
19
    return (ans % n + n) % n;
```

5.7 linear sieve

```
const int N = 10000000;
vector<int> lp(N+1);
vector<int> pr;

for (int i=2; i <= N; ++i) {
    if (lp[i] == 0) {
        lp[i] = i;
        pr.push_back(i);
    }
    for (int j = 0; i * pr[j] <= N; ++j) {
        lp[i * pr[j]] = pr[j];
        if (pr[j] == lp[i]) {
            break;
    }
}</pre>
```

5.8 Matrix

```
template<typename T>
struct Matrix{
    using rt = std::vector<T>;
    using mt = std::vector<rt>;
    using mt = std::vector<rt>;
    using mt = std::vector<rt>;
    using mt = std::vector<rt>;
    using matrix = Matrix<T>;
    int r,c;
    mt m;
    Matrix(int r,int c):r(r),c(c),m(r,rt(c)){}
    rt& operator[](int i){return m[i];}
    matrix operator+(const matrix &a){
        matrix rev(r,c);
        for(int i=0;i<r;++i)
        for(int j=0;j<c;++j)
        rev[i][j]=m[i][j]+a.m[i][j];</pre>
```

```
return rev;
16
    matrix operator-(const matrix &a){
       matrix rev(r,c);
      for(int i=0;i<r;++i)</pre>
         for(int j=0;j<c;++j)</pre>
           rev[i][j]=m[i][j]-a.m[i][j];
22
      return rev:
    matrix operator*(const matrix &a){
      matrix rev(r,a.c);
      matrix tmp(a.c,a.r);
      for(int i=0;i<a.r;++i)</pre>
         for(int j=0;j<a.c;++j)</pre>
           tmp[j][i]=a.m[i][j];
       for(int i=0;i<r;++i)</pre>
         for(int j=0;j<a.c;++j)</pre>
           for(int k=0;k<c;++k)</pre>
             rev.m[i][j]+=m[i][k]*tmp[j][k];
      return rev;
    bool inverse(){
      Matrix t(r,r+c);
      for(int y=0;y<r;y++){</pre>
         t.m[y][c+y] = 1;
         for(int x=0;x<c;++x)</pre>
           t.m[y][x]=m[y][x];
      if(!t.gas())
         return false;
       for(int y=0;y<r;y++)</pre>
         for(int x=0;x<c;++x)</pre>
           m[y][x]=t.m[y][c+x]/t.m[y][y];
       return true;
    T gas(){
      vector<T> lazy(r,1);
      bool sign=false;
      for(int i=0;i<r;++i){</pre>
         if( m[i][i]==0 ){
           int j=i+1;
           while(j<r&&!m[j][i])j++;</pre>
           if(j==r)continue;
           m[i].swap(m[j]);
           sign=!sign;
         for(int j=0;j<r;++j){</pre>
           if(i==j)continue;
           lazy[j]=lazy[j]*m[i][i];
           T mx=m[j][i];
           for(int k=0;k<c;++k)</pre>
             m[j][k]=m[j][k]*m[i][i]-m[i][k]*mx
       T det=sign?-1:1;
       for(int i=0;i<r;++i){</pre>
         det = det*m[i][i];
         det = det/lazy[i];
         for(auto &j:m[i])j/=lazy[i];
      return det;
```

5.9 數位統計

```
ı | 11 d[65], dp[65][2];//up區間是不是完整
2 11 dfs(int p,bool is8,bool up){
    if(!p)return 1; // 回傳0是不是答案
    if(!up&&~dp[p][is8])return dp[p][is8];
    int mx = up?d[p]:9;//可以用的有那些
    for(int i=0;i<=mx;++i){</pre>
      if( is8&&i==7 )continue;
      ans += dfs(p-1,i==8,up&&i==mx);
    if(!up)dp[p][is8]=ans;
    return ans;
14 11 f(11 N){
   int k=0;
15
    while(N){ // 把數字先分解到陣列
      d[++k] = N%10;
      N/=10:
    return dfs(k,false,true);
```

6 String

6.1 manacher

6.2 reverseBWT

```
const int MAXN = 305, MAXC = 'Z';
int ranks[MAXN], tots[MAXC], first[MAXC];
void rankBWT(const string &bw){
    memset(ranks,0,sizeof(int)*bw.size());
    for(size_t i=0;i<bw.size();++i)
    ranks[i] = tots[int(bw[i])]++;
}
void firstCol(){
    memset(first,0,sizeof(first));
    int totc = 0;
    for(int c='A';c<='Z';++c){
        if(!tots[c]) continue;
    first[c] = totc;
    totc += tots[c];</pre>
```

```
17 }
string reverseBwt(string bw.int begin){
    rankBWT(bw), firstCol();
    int i = begin; //原字串最後一個元素的位置
20
    string res;
21
22
    do{
23
      char c = bw[i];
      res = c + res;
      i = first[int(c)] + ranks[i];
25
    }while( i != begin );
    return res;
```

6.3 suffix array lcp

```
1 #define radix_sort(x,y){\
    for(i=0;i<A;++i)c[i]=0;\</pre>
    for(i=0;i<n;++i)c[x[y[i]]]++;\</pre>
    for(i=1;i<A;++i)c[i]+=c[i-1];\</pre>
    for(i=n-1;~i;--i)sa[--c[x[y[i]]]]=y[i];\
  #define AC(r,a,b)\
    r[a]!=r[b]||a+k>=n||r[a+k]!=r[b+k]
  void suffix_array(const char *s,int n,int *
       sa,int *rank,int *tmp,int *c){
     int A='z'+1,i,k,id=0;
    for(i=0;i<n;++i)rank[tmp[i]=i]=s[i];</pre>
    radix sort(rank,tmp);
    for(k=1;id<n-1;k<<=1){</pre>
      for(id=0,i=n-k;i<n;++i)tmp[id++]=i;</pre>
      for(i=0;i<n;++i)</pre>
         if(sa[i]>=k)tmp[id++]=sa[i]-k;
       radix sort(rank,tmp);
      swap(rank,tmp);
      for(rank[sa[0]]=id=0,i=1;i<n;++i)</pre>
         rank[sa[i]]=id+=AC(tmp,sa[i-1],sa[i]);
21
      A=id+1;
22
23 }
24 //h: 高度數組 sa:後綴數組 rank:排名
void suffix_array_lcp(const char *s,int len,
       int *h,int *sa,int *rank){
     for(int i=0;i<len;++i)rank[sa[i]]=i;</pre>
     for(int i=0,k=0;i<len;++i){</pre>
      if(rank[i]==0)continue;
      if(k)--k;
29
      while(s[i+k]==s[sa[rank[i]-1]+k])++k;
30
      h[rank[i]]=k;
31
32
    h[0]=0;// h[k]=lcp(sa[k],sa[k-1]);
```

6.4 KMP

```
1 /*產生fail function*/
2 void kmp_fail(char *s,int len,int *fail){
3   int id=-1;
4  fail[0]=-1;
```

```
for(int i=1;i<len;++i){</pre>
      while(~id&&s[id+1]!=s[i])id=fail[id];
      if(s[id+1]==s[i])++id;
      fail[i]=id;
111 /*以字串B匹配字串A·傳回匹配成功的數量(用B的
       fail)*/
int kmp_match(char *A,int lenA,char *B,int
       lenB.int *fail){
    int id=-1, ans=0;
    for(int i=0;i<lenA;++i){</pre>
      while(~id&&B[id+1]!=A[i])id=fail[id];
15
      if(B[id+1]==A[i])++id;
16
17
      if(id==lenB-1){/*匹配成功*/
        ++ans, id=fail[id];
18
19
20
    return ans;
```

6.5 hash

```
1 | #define MAXN 1000000
 2 #define mod 1073676287
 3 /*mod 必須要是質數*/
 4 typedef long long T;
 5 char s[MAXN+5];
 6 T h[MAXN+5]; /*hash 陣 列 */
 7 T h base[MAXN+5];/*h base[n]=(prime^n)%mod*/
 8 void hash init(int len,T prime){
    h base[0]=1;
     for(int i=1:i<=len:++i){</pre>
      h[i]=(h[i-1]*prime+s[i-1])%mod;
11
       h_base[i]=(h_base[i-1]*prime)%mod;
13
14 }
15 T get hash(int 1,int r){/*閉區間寫法,設編號
        為0 ~ Len-1*/
    return (h[r+1]-(h[1]*h_base[r-l+1])%mod+
         mod)%mod;
```

6.6 AC 自動機

```
1 template < char L='a', char R='z'>
2 class ac automaton{
    struct joe{
      int next[R-L+1], fail, efl, ed, cnt dp, vis;
      joe():ed(0),cnt_dp(0),vis(0){
         for(int i=0;i<=R-L;++i)next[i]=0;</pre>
    };
  public:
    std::vector<joe> S;
    std::vector<int> q;
11
12
    int qs,qe,vt;
13
    ac automaton():S(1),qs(0),qe(0),vt(0){}
    void clear(){
```

for(auto v:rG[u]) if(v=dfn[v]){

idom[v] = semi[best[v]]==pa[i]

tree[id[idom[i]]].push_back(id[i]);

? pa[i] : best[v];

tree[semi[i]].push_back(i);

for(auto v:tree[pa[i]]){

for(int i=2; i<=dfnCnt; ++i){</pre>

idom[i] = idom[idom[i]];

find(v, pa[i]);

tree[pa[i]].clear();

if(idom[i] != semi[i])

橋連涌分量

int u,v;

semi[i]=min(semi[i],semi[best[v]]);

find(v,i);

```
q.clear();
                                                       p=S[p].next[id];
      S.resize(1);
                                                       if(S[p].ed)ans+=S[p].ed;
      for(int i=0;i<=R-L;++i)S[0].next[i]=0;</pre>
                                                       for(t=S[p].efl;~t;t=S[t].efl){
      S[0].cnt_dp=S[0].vis=qs=qe=vt=0;
                                                        ans+=S[t].ed;/*因為都走efL邊所以保證
                                              77
                                                             匹配成功*/
    void insert(const char *s){
      int o=0:
      for(int i=0,id;s[i];++i){
                                                     return ans;
        id=s[i]-L;
        if(!S[o].next[id]){
                                                   /*枚舉(s的子字串nA)的所有相異字串各恰一次
         S.push back(joe());
                                                        並傳回次數O(N*M^(1/3))*/
          S[o].next[id]=S.size()-1;
                                                   int match_2(const char *s){
                                                    int ans=0,id,p=0,t;
        o=S[o].next[id];
                                                     ++vt;
                                                     /*把戳記vt+=1,只要vt沒溢位,所有S[p].
      ++S[o].ed;
                                                         vis==vt就會變成false
                                                     這種利用vt的方法可以0(1)歸零vis陣列*/
    void build_fail(){
     S[0].fail=S[0].efl=-1;
                                                     for(int i=0;s[i];++i){
      q.clear();
                                                       id=s[i]-L;
      q.push_back(0);
                                                       while(!S[p].next[id]&&p)p=S[p].fail;
      ++qe;
                                                       if(!S[p].next[id])continue;
      while(qs!=qe){
                                                       p=S[p].next[id];
       int pa=q[qs++],id,t;
                                                       if(S[p].ed&&S[p].vis!=vt){
        for(int i=0;i<=R-L;++i){</pre>
                                                        S[p].vis=vt;
         t=S[pa].next[i];
                                                        ans+=S[p].ed;
         if(!t)continue;
          id=S[pa].fail;
                                                       for(t=S[p].efl;~t&&S[t].vis!=vt;t=S[t
          while(~id&&!S[id].next[i])id=S[id].
                                                           ].efl){
              fail;
                                                        S[t].vis=vt;
          S[t].fail=~id?S[id].next[i]:0;
                                                        ans+=S[t].ed;/*因為都走efl邊所以保證
         S[t].efl=S[S[t].fail].ed?S[t].fail:S
                                                             匹配成功*/
              [S[t].fail].efl;
          q.push_back(t);
                                              101
          ++qe;
                                                     return ans;
                                                   /*把AC自動機變成真的自動機*/
                                                   void evolution(){
    /*DP出每個前綴在字串s出現的次數並傳回所有
                                                     for(qs=1;qs!=qe;){
         字串被s匹配成功的次數O(N+M)*/
                                                       int p=q[qs++];
    int match 0(const char *s){
                                                       for(int i=0;i<=R-L;++i)</pre>
      int ans=0,id,p=0,i;
                                              109
                                                        if(S[p].next[i]==0)S[p].next[i]=S[S[
      for(i=0;s[i];++i){
                                                             p].fail].next[i];
        id=s[i]-L;
        while(!S[p].next[id]&&p)p=S[p].fail;
                                              111
        if(!S[p].next[id])continue;
                                              112 };
        p=S[p].next[id];
        ++S[p].cnt_dp;/*匹配成功則它所有後綴都
             可以被匹配(DP計算)*/
                                                       minimal string rotation
      for(i=qe-1;i>=0;--i){
        ans+=S[q[i]].cnt_dp*S[q[i]].ed;
        if(~S[q[i]].fail)S[S[q[i]].fail].
                                               i| int min_string_rotation(const string &s){
            cnt_dp+=S[q[i]].cnt_dp;
                                                   int n=s.size(),i=0,j=1,k=0;
                                                   while(i<n&&j<n&&k<n){</pre>
      return ans;
                                                    int t=s[(i+k)%n]-s[(j+k)%n];
66
                                                     ++k;
   /*多串匹配走efL邊並傳回所有字串被s匹配成功
                                                     if(t){
         的 次 數 O(N*M^1.5)*/
                                                      if(t>0)i+=k;
    int match_1(const char *s)const{
                                                       else j+=k;
                                                       if(i==j)++j;
      int ans=0,id,p=0,t;
      for(int i=0;s[i];++i){
                                                       k=0;
        id=s[i]-L;
        while(!S[p].next[id]&&p)p=S[p].fail;
```

if(!S[p].next[id])continue;

14 } **6.8** Z void z alg(char *s,int len,int *z){ int 1=0, r=0; z[0]=len;for(int i=1;i<len;++i){</pre> z[i]=i>r?0:(i-1+z[i-1]< z[1]?z[i-1]:r-i51 while(i+z[i]<len&&s[i+z[i]]==s[z[i]])++z **if**(i+z[i]-1>r)r=i+z[i]-1,l=i; 54 55 56 57 | }dom; Tarjan dominator tree 1 #define N 1005 2 struct edge{ 1 struct dominator tree{ static const int MAXN=5005; int n;// 1-base vector<int> G[MAXN], rG[MAXN]; int pa[MAXN], dfn[MAXN], id[MAXN], dfnCnt; int semi[MAXN], idom[MAXN], best[MAXN]; vector<int> tree[MAXN]; // tree here void init(int _n){

19

20

21

23

24

26

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28

29

32

33

34

35

36

return min(i,j);//最小循環表示法起始位置

```
bool is bridge;
                                                 edge(int u=0,int v=0):u(u),v(v),is_bridge
                                                       (0){}
                                                vector<edge> E;
                                               vector<int> G[N];// 1-base
                                               int low[N], vis[N], Time;
                                             int bcc id[N], bridge cnt, bcc cnt; // 1-base
                                             11 int st[N],top;//BCC用
  for(int i=1; i<=n; ++i)</pre>
                                             12 void add_edge(int u,int v){
    G[i].clear(), rG[i].clear();
                                                 G[u].push_back(E.size());
                                                 E.emplace back(u,v);
void add_edge(int u, int v){
                                                 G[v].push_back(E.size());
                                             15
  G[u].push_back(v);
                                                 E.emplace_back(v,u);
                                             17 }
  rG[v].push_back(u);
                                             | void dfs(int u,int re=-1){//u當前點,re為u連
void dfs(int u){
                                                     接前一個點的邊
  id[dfn[u]=++dfnCnt]=u;
                                                  int v;
  for(auto v:G[u]) if(!dfn[v])
                                                  low[u]=vis[u]=++Time;
    dfs(v),pa[dfn[v]]=dfn[u];
                                                  st[top++]=u;
                                                  for(int e:G[u]){
int find(int y,int x){
                                                   v=E[e].v;
  if(y <= x) return y;</pre>
                                                    if(!vis[v]){
                                             24
  int tmp = find(pa[y],x);
                                             25
                                                      dfs(v,e^1);//e^1反向邊
  if(semi[best[y]] > semi[best[pa[y]]])
                                             26
                                                      low[u]=min(low[u],low[v]);
    best[y] = best[pa[y]];
                                             27
                                                      if(vis[u]<low[v]){</pre>
  return pa[y] = tmp;
                                                        E[e].is_bridge=E[e^1].is_bridge=1;
                                                        ++bridge cnt;
void tarjan(int root){
  dfnCnt = 0;
                                             31
                                                    }else if(vis[v]<vis[u]&&e!=re)</pre>
  for(int i=1; i<=n; ++i){</pre>
                                             32
                                                      low[u]=min(low[u], vis[v]);
    dfn[i] = idom[i] = 0;
                                             33
    tree[i].clear();
                                                  if(vis[u]==low[u]){//處理BCC
    best[i] = semi[i] = i;
                                                    ++bcc_cnt;// 1-base
                                                    do bcc_id[v=st[--top]]=bcc_cnt;//每個點
  dfs(root);
                                                         所在的BCC
  for(int i=dfnCnt; i>1; --i){
                                                    while(v!=u);
    int u = id[i];
```

```
40 void bcc init(int n){
   Time=bcc_cnt=bridge_cnt=top=0;
   E.clear();
    for(int i=1;i<=n;++i){</pre>
      G[i].clear();
      vis[i]=bcc id[i]=0;
```

雙連誦分量 & 割點

```
1 #define N 1005
vector<int> G[N];// 1-base
3 vector (int > bcc[N]; // 存每塊雙連通分量的點
4 int low[N], vis[N], Time;
5 int bcc_id[N],bcc_cnt;// 1-base
6|bool is cut[N];//是否為割點
7 int st[N],top;
s|void dfs(int u,int pa=-1){//u當前點,pa父親
   int t, child=0;
   low[u]=vis[u]=++Time;
    st[top++]=u;
    for(int v:G[u]){
     if(!vis[v]){
        dfs(v,u),++child;
        low[u]=min(low[u],low[v]);
        if(vis[u]<=low[v]){</pre>
          is_cut[u]=1;
          bcc[++bcc cnt].clear();
            bcc_id[t=st[--top]]=bcc_cnt;
            bcc[bcc_cnt].push_back(t);
          }while(t!=v);
          bcc id[u]=bcc cnt;
          bcc[bcc_cnt].push_back(u);
      }else if(vis[v]<vis[u]&&v!=pa)//反向邊
        low[u] = min(low[u], vis[v]);
    }//u是dfs樹的根要特判
   if(pa==-1&&child<2)is cut[u]=0;</pre>
  void bcc_init(int n){
   Time=bcc cnt=top=0;
   for(int i=1;i<=n;++i){</pre>
     G[i].clear();
     is_cut[i]=vis[i]=bcc_id[i]=0;
37 }
```

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```
| #include < bits / stdc++.h>
2 using namespace std:
3 #define MAXN 8001
4 #define MAXN2 MAXN*4
5 #define n(X) ((X)+2*N)
6 vector<int> v[MAXN2], rv[MAXN2], vis_t;
```

```
void addedge(int s,int e){
    v[s].push back(e);
    rv[e].push_back(s);
12 int scc[MAXN2];
  bool vis[MAXN2]={false};
  void dfs(vector<int> *uv.int n.int k=-1){
    for(int i=0;i<uv[n].size();++i)</pre>
      if(!vis[uv[n][i]])
        dfs(uv,uv[n][i],k);
    if(uv==v)vis_t.push_back(n);
    scc[n]=k;
  void solve(){
    for(int i=1;i<=N;++i){</pre>
      if(!vis[i])dfs(v,i);
      if(!vis[n(i)])dfs(v,n(i));
    memset(vis,0,sizeof(vis));
    int c=0:
    for(int i=vis_t.size()-1;i>=0;--i)
      if(!vis[vis t[i]])
        dfs(rv,vis t[i],c++);
  int main(){
    int a,b;
    scanf("%d%d",&N,&M);
    for(int i=1;i<=N;++i){</pre>
      // (A or B)&(!A & !B) A^B
      a=i*2-1;
      b=i*2;
      addedge(n(a),b);
      addedge(n(b),a);
      addedge(a,n(b));
      addedge(b,n(a));
    while(M--){
      scanf("%d%d",&a,&b);
      a = a>0?a*2-1:-a*2;
      b = b>0?b*2-1:-b*2;
      // A or B
      addedge(n(a),b);
       addedge(n(b),a);
    solve();
    bool check=true;
    for(int i=1;i<=2*N;++i)</pre>
      if(scc[i]==scc[n(i)])
        check=false;
    if(check){
      printf("%d\n",N);
      for(int i=1;i<=2*N;i+=2){</pre>
        if(scc[i]>scc[i+2*N]) putchar('+');
        else putchar('-');
```

puts("");

return 0:

}else puts("0");

other

上下最大正方形

```
void solve(int n,int a[],int b[]){// 1-base
   int ans=0:
   deque<int>da,db;
    for(int l=1,r=1;r<=n;++r){</pre>
     while(da.size()&&a[da.back()]>=a[r]){
       da.pop_back();
     da.push back(r);
     while(db.size()&&b[db.back()]>=b[r]){
       db.pop back();
     db.push_back(r);
     for(int d=a[da.front()]+b[db.front()];r-
          l+1>d;++l){
       if(da.front()==1)da.pop_front();
       if(db.front()==1)db.pop front();
       if(da.size()&&db.size()){
          d=a[da.front()]+b[db.front()];
     ans=max(ans,r-l+1);
   printf("%d\n",ans);
```

8.2 WhatDay

```
i int whatday(int y,int m,int d){
   if(m<=2)m+=12,--y;
   if(y<1752||y==1752&&m<9||y==1752&&m==9&&d
     return (d+2*m+3*(m+1)/5+y+y/4+5)%7;
   return (d+2*m+3*(m+1)/5+y+y/4-y/100+y/400)
```

最大矩形

```
1 LL max rectangle(vector<int> s){
   stack<pair<int,int >> st;
    st.push(make pair(-1,0));
    s.push back(0);
    LL ans=0;
    for(size t i=0;i<s.size();++i){</pre>
      int h=s[i];
      pair<int,int > now=make pair(h,i);
      while(h<st.top().first){</pre>
        now=st.top();
        st.pop();
11
        ans=max(ans,(LL)(i-now.second)*now.
12
             first):
13
      if(h>st.top().first){
14
```

```
st.push(make_pair(h,now.second));
16
17
    }
18
   return ans;
19 }
```

zformula

9.1 formula

9.1.1 Pick 公式

給定頂點坐標均是整點的簡單多邊形,面積 = 內部格點數 + 邊上格點數/2-1

9.1.2 圖論

- 對於平面圖 $\cdot F = E V + C + 1 \cdot \mathbf{C}$ 是連通分量數對於平面圖 $\cdot E < 3V 6$ 對於連通圖 $\mathbf{G} \cdot$ 最大獨立點集的大小設為 $\mathbf{I}(\mathbf{G}) \cdot$ 最大 匹配大小設為 M(G) · 最小點覆蓋設為 Cv(G) · 最小 邊覆蓋設為 Ce(G)。對於任意連通圖:
 - (a) I(G) + Cv(G) = |V|(b) M(G) + Ce(G) = |V|
- 4. 對於連通二分圖:
 - (a) I(G) = Cv(G)(b) M(G) = Ce(G)
- 5. 最大權閉合圖:
 - $\begin{array}{ll} \text{(a)} & C(u,v) = \infty, (u,v) \in E \\ \text{(b)} & C(S,v) = W_v, W_v > 0 \\ \text{(c)} & C(v,T) = -W_v, W_v < 0 \\ \text{(d)} & \operatorname{ans} = \sum_{W_v > 0} W_v flow(S,T) \end{array}$
- 6. 最大密度子圖:
 - (a) $\Re \max\left(\frac{W_e+W_v}{|V'|}\right), e \in E', v \in V'$ (b) $U = \sum_{v \in V} 2W_v + \sum_{e \in E} W_e$
 - (c) $C(u,v) = W_{(u,v)}, (u,v) \in E$. 雙向邊
 - (d) $C(S, v) = U, v \in V$
 - (e) $D_u = \sum_{(u,v) \in E} W_{(u,v)}$
 - (f) $C(v,T) = U + 2g D_v 2W_v, v \in V$
 - (g) 二分搜 g: $l = 0, r = U, eps = 1/n^2$ $if((U \times |V| - flow(S, T))/2 > 0) l = mid$ else r = mid
 - (h) ans= $min_cut(S, T)$
 - (i) |E| = 0 要特殊判斷
- 7. 弦圖:

 - (a) 點數大於 3 的環都要有一條弦 (b) 完美消除序列從後往前依次給每個點染色·給
 - 每個點來上可以染的最小顏色 最大團大小=色數 最大獨立集: 完美消除序列從前往往 最小團覆蓋: 最大獨立集的點和他

 - 成 區間圖是弦圖 區間圖的完美消除序列: 將區間按造又端點由
 - 小到大排序 (h) 區間圖染色: 用線段樹做

9.1.3 dinic 特殊圖複雜度

1. 單位流: $O\left(min\left(V^{3/2},E^{1/2}\right)E\right)$ 2. 二分圖: $O\left(V^{1/2}E\right)$

9.1.4 0-1 分數規劃

 $x_i = \{0,1\} \cdot x_i$ 可能會有其他限制 · 求 $max\left(rac{\sum B_i x_i}{\sum C_i x_i}
ight)$

- 1. $D(i,g) = B_i g \times C_i$
- 2. $f(g) = \sum D(i, g)x_i$
- 3. f(g) = 0 時 g 為最佳解 f(g) < 0 沒有意義
- 4. 因為 f(g) 單調可以二分搜 g
- 5. 或用 Dinkelbach 通常比較快

```
i| binary_search(){
   while(r-1>eps){
    g=(1+r)/2;
    for(i:所有元素)D[i]=B[i]-g*C[i];//D(i,g)
    找出一組合法x[i]使f(g)最大;
    if(f(g)>0) l=g;
    else r=g;
   Ans = r;
Dinkelbach(){
   g=任意狀態(通常設為0);
   do{
    Ans=g;
    for(i:所有元素)D[i]=B[i]-g*C[i];//D(i,g)
    找出一組合法x[i]使f(g)最大;
    p=0,q=0;
    for(i:所有元素)
      if(x[i])p+=B[i],q+=C[i];
    g=p/q;//更新解·注意q=0的情況
   }while(abs(Ans-g)>EPS);
   return Ans;
```

9.1.5 學長公式

- 1. $\sum_{d|n} \phi(n) = n$
- 2. $g(n) = \sum_{d \mid n} f(d) => f(n) = \sum_{d \mid n} \mu(d) \times g(n/d)$
- 3. Harmonic series $H_n = \ln(n) + \gamma + 1/(2n) 1/(12n^2) + 1/(120n^4)$
- 4. $\gamma = 0.57721566490153286060651209008240243104215$
- 5. 格雷碼 = $n \oplus (n >> 1)$
- 6. $SG(A+B) = SG(A) \oplus SG(B)$
- 7. 選轉矩陣 $M(\theta) = \begin{pmatrix} cos\theta & -sin\theta \\ sin\theta & cos\theta \end{pmatrix}$

9.1.6 基本數論

- 1. $\sum_{d|n} \mu(n) = [n == 1]$
- 2. $g(m) = \sum_{d|m} f(d) \Leftrightarrow f(m) = \sum_{d|m} \mu(d) \times g(m/d)$
- 3. $\sum_{i=1}^{n} \sum_{j=1}^{m} \underline{\Delta} = \sum_{i=1}^{m} \mu(d) \left| \frac{n}{d} \right| \left| \frac{m}{d} \right|$
- 4. $\sum_{i=1}^{n} \sum_{j=1}^{n} lcm(i,j) = n \sum_{d|n} d \times \phi(d)$

9.1.7 排組公式

- 1. k 卡特蘭 $\frac{C_n^{kn}}{n(k-1)+1} \cdot C_m^n = \frac{n!}{m!(n-m)!}$
- 2. $H(n,m) \cong x_1 + x_2 ... + x_n = k, num = C_k^{n+k-1}$
- 3. Stirling number of 2^{nd} ,n 人分 k 組方法數目
 - (a) S(0,0) = S(n,n) = 1
 - (b) S(n,0) = 0
 - (c) S(n,k) = kS(n-1,k) + S(n-1,k-1)
- 4. Bell number.n 人分任意多組方法數目
 - (a) $B_0 = 1$
 - (b) $B_n = \sum_{i=0}^n S(n,i)$
 - (c) $B_{n+1} = \sum_{k=0}^{n} C_k^n B_k$
 - (d) $B_{p+n} \equiv B_n + B_{n+1} mod p$, p is prime
 - (e) $B_{p^m+n} \equiv mB_n + B_{n+1} mod p$, p is prime
 - $\begin{array}{ll} \text{(f)} & \text{From } B_0: 1, 1, 2, 5, 15, 52, \\ & 203, 877, 4140, 21147, 115975 \end{array}$
- 5. Derangement, 錯排, 沒有人在自己位置上
 - (a) $D_n = n!(1 \frac{1}{1!} + \frac{1}{2!} \frac{1}{3!} \dots + (-1)^n \frac{1}{n!})$
 - (b) $D_n = (n-1)(D_{n-1} + D_{n-2}), D_0 = 1, D_1 = 0$
 - (c) From $D_0: 1, 0, 1, 2, 9, 44, 265, 1854, 14833, 133496$
- 6. Binomial Equality
 - (a) $\sum_{k} {r \choose m+k} {s \choose n-k} = {r+s \choose m+n}$
 - (b) $\sum_{k} {l \choose m+k} {s \choose n+k} = {l+s \choose l-m+n}$
 - (c) $\sum_{k} {l \choose m+k} {s+k \choose n} (-1)^k = (-1)^{l+m} {s-m \choose n-l}$
 - (d) $\sum_{k \le l} {l \choose m} {s \choose k-n} (-1)^k = (-1)^{l+m} {s-m-1 \choose l-n-m}$
 - (e) $\sum_{0 \le k \le l} {l-k \choose m} {q+k \choose n} = {l+q+1 \choose m+n+1}$
 - (f) $\binom{r}{k} = (-1)^k \binom{k-r-1}{k}$
 - (g) $\binom{r}{m}\binom{m}{k} = \binom{r}{k}\binom{r-k}{m-k}$
 - (h) $\sum_{k \le n} {r+k \choose k} = {r+n+1 \choose n}$
 - (i) $\sum_{0 \le k \le n} {k \choose m} = {n+1 \choose m+1}$
 - (j) $\sum_{k \le m} {m+r \choose k} x^k y^k$ $\sum_{k \le m} {-r \choose k} (-x)^k (x+y)^{m-k}$

9.1.8 冪次, 冪次和

- 1. $a^{b}\%P = a^{b\%\varphi(p) + \varphi(p)}, b > \varphi(p)$
- 2. $1^3 + 2^3 + 3^3 + \ldots + n^3 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4}$
- 3. $1^4 + 2^4 + 3^4 + \ldots + n^4 = \frac{n^5}{5} + \frac{n^4}{2} + \frac{n^3}{3} \frac{n}{30}$
- 4. $1^5 + 2^5 + 3^5 + \ldots + n^5 = \frac{n^6}{6} + \frac{n^5}{2} + \frac{5n^4}{12} \frac{n^2}{12}$
- 5. $0^k + 1^k + 2^k + \dots + n^k = P(k), P(k) = \frac{(n+1)^{k+1} \sum_{i=0}^{k-1} C_i^{k+1} P(i)}{k+1}, P(0) = n+1$
- 6. $\sum_{k=0}^{m-1} k^n = \frac{1}{n+1} \sum_{k=0}^n C_k^{n+1} B_k m^{n+1-k}$
- 7. $\sum_{j=0}^{m} C_j^{m+1} B_j = 0, B_0 = 1$
- 8. 除了 $B_1 = -1/2$,剩下的奇數項都是 0
- 9. $B_2=1/6, B_4=-1/30, B_6=1/42, B_8=-1/30, B_{10}=5/66, B_{12}=-691/2730, B_{14}=7/6, B_{16}=-3617/510, B_{18}=43867/798, B_{20}=-174611/330,$

9.1.9 Burnside's lemma

- 1. $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$
- 2. $X^g = t^{c(g)}$
- 3. G 表示有幾種轉法 \cdot X^g 表示在那種轉法下 \cdot 有幾種 是會保持對稱的 \cdot t 是顏色數 \cdot c(g) 是循環節不動的 面數 \cdot
- 4. 正立方體塗三顏色·轉 0 有 3^6 個元素不變·轉 90 有 6 種·每種有 3^3 不變·180 有 $3 \times 3^4 \cdot 120$ (角) 有 $8 \times 3^2 \cdot 180$ (邊) 有 $6 \times 3^3 \cdot 2$ 全部 $\frac{1}{24} \left(3^6 + 6 \times 3^3 + 3 \times 3^4 + 8 \times 3^2 + 6 \times 3^3 \right) = 57$

9.1.10 Count on a tree

- 1. Rooted tree: $s_{n+1} = \frac{1}{n} \sum_{i=1}^{n} (i \times a_i \times \sum_{j=1}^{\lfloor n/i \rfloor} a_{n+1-i \times j})$
- 2. Unrooted tree:
 - (a) Odd: $a_n \sum_{i=1}^{n/2} a_i a_{n-i}$
 - (b) Even: $Odd + \frac{1}{2}a_{n/2}(a_{n/2} + 1)$
- 3. Spanning Tree
 - (a) 完全圖 $n^n 2$
 - (b) 一般圖 (Kirchhoff's theorem) $M[i][i] = degree(V_i), M[i][j] = -1,$ if have E(i, j), 0 if no edge. delete any one row and col in A, ans = det(A)

ACM ICPC	2	Data Structure	3			nd real root	8		7.4 tnfshb017	2 sat	11
		2.1 segment tree	3		5.6 🗜	□國剩餘定理	8				
Toom Defended		2.2 DisjointSet	3		5.7 lii	near sieve	8	8	other		11
Team Reference -		2.3 undo disjoint set	4			Matrix	8		8.1 上下最大		
. ~		2.4 fenwick	4		5.9 數	效位統計	9		8.2 WhatDay		11
Angry Crow	2	Flow	4	6	String		0		8.3 最大矩形		11
<i>S</i> - <i>J</i>	3		4	U	0		0				
Tologo Elight		3.1 min cost flow	4			nanacher	9	9	zformula		11
Takes Flight!		3.2 max flow	5			everseBWT			9.1 formula .		11
\mathcal{C}						uffix array lcp	9		9.1.1 Pic	ck 公式	11
	4	Graph			6.4 K	XMP	9			論	
		4.1 tree centroid	6		6.5 ha	ash	9			nic 特殊圖複雜度	
Contents		4.2 tarjan	6		C 自動機						
		4.3 heavy light decomposition	6			minimal string rotation				1 分數規劃	12
		,g							9.1.5 學	長公式	12
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1 Computational Geometry 1		5.1 basic	7	7	Tarjan		10		9.1.7 排	組公式	12
1.1 Geometry 1		5.2 bit set	7		•	ominator tree	10		9.1.8 冪	次, 冪次和	12
1.2 SmallestCircle 3		5.3 FFT	8		7.2 橇	喬連通分量	10		9.1.9 B u	ırnside's lemma	12
1.3 最近點對 3		5.4 質因數分解	8		7.3 雙	雙連通分量 & 割點	11		9.1.10 Co	ount on a tree	12

ACM ICPC Judge Test Angry Crow Takes Flight!

C++ Resource Test

```
#include <bits/stdc++.h>
using namespace std;

namespace system_test {

const size_t KB = 1024;
const size_t MB = KB * 1024;
const size_t GB = MB * 1024;
```

```
chrono::duration<double> diff = end -
10 size t block size, bound;
                                                          begin;
  void stack size dfs(size t depth = 1) {
                                                     return diff.count():
   if (depth >= bound)
                                                   void runtime error 1() {
    int8_t ptr[block_size]; // 若無法編譯將
                                                     // Segmentation fault
         block size 改成常數
                                                     int *ptr = nullptr;
    memset(ptr, 'a', block_size);
                                                     *(ptr + 7122) = 7122;
    cout << depth << endl;</pre>
                                                 42 }
    stack_size_dfs(depth + 1);
                                                   void runtime_error_2() {
                                                     // Segmentation fault
  void stack_size_and_runtime_error(size_t
                                                     int *ptr = (int *)memset;
       block size, size t bound = 1024) {
                                                     *ptr = 7122;
    system test::block size = block size;
                                                 48 }
    system_test::bound = bound;
    stack size dfs();
                                                   void runtime_error_3() {
                                                     // munmap_chunk(): invalid pointer
                                                     int *ptr = (int *)memset;
  double speed(int iter num) {
                                                     delete ptr;
    const int block_size = 1024;
    volatile int A[block size];
    auto begin = chrono::high resolution clock
                                                   void runtime_error_4() {
         ::now();
                                                     // free(): invalid pointer
    while (iter_num--)
                                                     int *ptr = new int[7122];
      for (int j = 0; j < block_size; ++j)</pre>
                                                     ptr += 1;
                                                     delete[] ptr;
    auto end = chrono::high resolution clock::
```

```
63 void runtime error 5() {
    // maybe illegal instruction
    int a = 7122, b = 0;
    cout << (a / b) << endl;</pre>
  void runtime error 6() {
    // floating point exception
    volatile int a = 7122, b = 0;
    cout << (a / b) << endl;
73 }
  void runtime error 7() {
    // call to abort.
    assert(false);
78 }
  } // namespace system test
82 #include <sys/resource.h>
void print_stack_limit() { // only work in
       Linux
    struct rlimit 1;
    getrlimit(RLIMIT STACK, &1);
    cout << "stack_size = " << l.rlim_cur << "</pre>
          byte" << endl;</pre>
87 }
```