

Name - Swati kesarwani

Class - B.Sc. III (pm) ; sem-6

Sub- Maths

Paper - Discrete

mathematics

Graph Theory

A graph is denoted by G consisting two non-empty set.

V and E i.e. $G = (V, E)$

where V Represent vertex (Node) and E Represented edge (side) vertex is represented as \bullet (Vertex, Node) and Edge is represented as Edge

There are many types of graph -

Simple graph -

- 1). No direction
- 2). No Loop
- 3). No Multiple edge between two vertex.
- 4). It is undirected graph



(Simple graph
Undirected
graph)

Multigraph

- 1). No direction
- 2). No loop
- 3). Multiple Edges exist between vertex
i.e. Multiple edges exist between vertex



(Multigraph)

Pseudo graph

- 1). No direction
- 2). Loop exists
- 3). Multiple edges possible



Directed graph

- 1). Direction contain
- 2). No Multiple Edge
- 3). No Loop contain



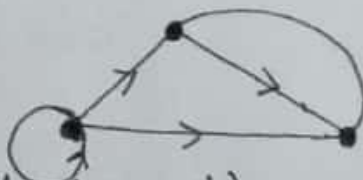
Directed multigraph -

- (1) More than one edge between two nodes
- (2) Direction are available



Mixed graph -

- 1). Direction may be possible
- 2). Multiple edges (May be or not)



Ques. In which Multiple Loop are allowed

- ① Simple graph
- ② Directed graph
- ③ Pseudograph (✓)
- ④ N.G.T

Ques. In a directed graph

- ① graph is fixed
- ② graph is not fixed
- ③ Direction are fixed (✓)
- ④ N.O.T.

Ques. In which of the following graph Loops are not allowed

- ① Pseudograph
- ② Simple graph (✓)
- ③ Mixed graph
- ④ Directed Multigraph

Ques. find the ^{total} number of distinct simple graphs with up to three vertices is

- ① 9
- ② 7 ✓
- ③ 15
- ④ 16

Soln-



Null graph -



(No edges between Nodes)

Trivial graph -



only one Node

Connected graph



If some graph are connected, is called connected graph.

Regular graph (complete graph)

A simple graph G is said to be complete. If every vertex in G has connected with every other vertex.

$$\text{No. of vertex} = 3 = n$$

$$\begin{aligned}\text{No. of edge} &= \frac{n(n-1)}{2} \\ &= \frac{3 \times 2}{2} \\ &= 3\end{aligned}$$



Note - A complete graph (regular graph) has n vertex. Then No. of edges are $\frac{n(n-1)}{2}$



$$\text{Node} = 3$$



$$\text{Node} = 4$$

$$\begin{aligned}\text{No. of edge} &= \frac{n(n-1)}{2} \\ &= \frac{3 \times (3-1)}{2} \\ &= 3\end{aligned}$$

$$\begin{aligned}\text{No. of edge} &= \frac{n(n-1)}{2} = \frac{4 \times (4-1)}{2} \\ &= 2 \times 3 = 6\end{aligned}$$



(complete graph)

$$\begin{aligned}n &= 5 \\ \text{No. of edge} &= \frac{n(n-1)}{2} \\ &= \frac{5 \times 4}{2} \\ &= 10\end{aligned}$$

Ques. Let G be a complete graph of n vertices. Then Total No. of edges are

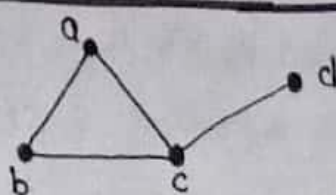
- (a) n (b) $\frac{n(n-1)}{2}$ (c) $n^2 - 1$ (d) $2n$

Degree of vertex (For Undirected graph)

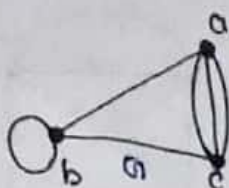
The degree of a vertex V in a graph G is written as $d(V)$ is equal to No. of edges which are incident on V with self loop counted twice.

* Loop is formed at any point (vertex) then degree of loop is 2.

$\deg(a) = 2$
 $\deg(b) = 2$
 $\deg(c) = 3$
 $\deg(d) = 1$



In G
 $\deg(a) = 4$
 $\deg(b) = 4$
 $\deg(c) = 4$



Isolated Vertex - A vertex in a graph G having No edge incident onto is an Isolated Vertex. It's degree is 0.

Ques. The No. of Vertices in odd degree of a graph is
 (a) always even (b) always odd (c) even as well as odd (d) N.O.T

$\deg = 3$



$\deg V = 3$

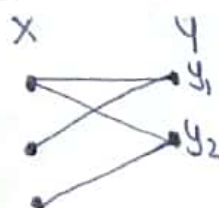
No. of Nodes (Vertex) = 4 (even)

No. of vertices even



degree $V = 5$

BIPARTITE Graph -



It arrange the set and the edge between Node same set is not allowed and the edge between Node of different set can be allowed.

Planar graph - In Planar graph edges does not intersect.



Planar graph



Planar graph

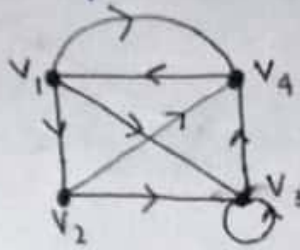
These are Planar graph because edges does not intersect in these graph.

degree for directed graph - In directed graph there are two types of degree

1). In-degree - No. of edges incident into vertex (v_i) $\{d^-(v_i)\}$

2). Out-degree - No. of edges incident out of a vertex (v_i) $\{d^+(v_i)\}$

Ques. Find in-degree and out-degree of the following graph.



In-degree - $d^-(v_1) = 1$
 $d^-(v_2) = 1$
 $d^-(v_3) = 3$
 $d^-(v_4) = 3$

Out-degree - $d^+(v_1) = 3$
 $d^+(v_2) = 2$
 $d^+(v_3) = 1 + 1$ (due to loop)
 $\quad \quad \quad = 2$
 $d^+(v_4) = 1$

Note - degree of loop in directed graph is divided in one is in-degree and other one is out-degree

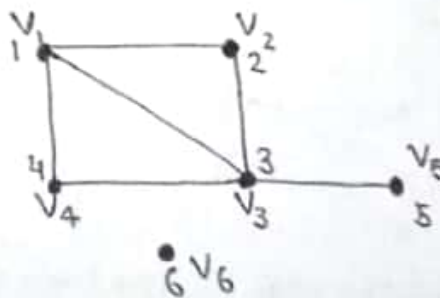
Handshaking Theorem / Sum of degree Theorem -

Let $G = (V, E)$ be a undirected graph with e edges.

Then $\sum_{i=1}^n \deg(v_i) = 2e$

$[d(v_1) + d(v_2) + d(v_3) + \dots + d(v_n) = 2e]$

Example - $\deg(v_1) = 3$
 $\deg(v_2) = 2$
 $\deg(v_3) = 4$
 $\deg(v_4) = 2$
 $\deg(v_5) = 1$
 $\deg(v_6) = 0$



$d(v_1) + d(v_2) + d(v_3) + d(v_4) + d(v_5) + d(v_6) = 12$

$\sum_{i=1}^6 d(v_i) = 12$

$e = 6$
 $2e = 12 \Rightarrow \left[\sum_{i=1}^6 d(v_i) = 2e \right]$

Example -

$$\deg(1) = 2$$

$$\deg(2) = 2$$

$$\deg(3) = 2$$

$$\deg(1) + \deg(2) + \deg(3) = 6$$

$$\sum_{i=1}^3 d(v_i) = 6$$

$$e = 3$$

$$2e = 6$$

$$\left[\sum_{i=1}^3 d(v_i) = 2e \right]$$



Ques. An isolated vertex of a graph has

- ① one degree ② zero degree (✓) ③ two degree ④ NOT

Ques. find the no. of edges in a graph with 10 vertices each of degree 4 is

- ① 10 ② 20 (✓) ③ 40 ④ 80

$$\sum_{i=1}^{10} d(v_i) = 2e$$

$$d(v_1) + d(v_2) + \dots + d(v_{10}) = 2e$$

$$4 \times 10 = 2e$$

$$e = 20$$

Euler's formula Let G be a connected Planar simple graph with edge e and

Vertices v . Let γ be the No. of ^{Region} (bounded and unbounded) of G , then

$$[\gamma = e - v + 2]$$

$$\text{or } [v - e + \gamma = 2]$$

Example -

Here $v = 4$

$$e = 6$$

$$\gamma = 4$$

$$v - e + \gamma = 4 - 6 + 4$$

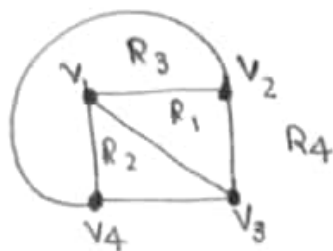
$$v - e + \gamma = 2$$

$$\gamma = e - v + 2$$

$$= 6 - 4 + 2$$

$$= 8 - 4$$

$$\gamma = 4$$



Ques. If There are 20 vertex each of degree 3. Then How many Region of a Planar graph

Soln -

$$\gamma = e - v + 2 \quad \text{if } v = 20,$$

$$2e = 20 \times 3 \Rightarrow e = 30$$

$$\chi = 30 - 20 + 2 = 32 - 20$$

$$\chi = 12$$

Ques. In theory of graph, The Euler formula is

(a) $\chi + e - v = 2$

(b) $v - e + \chi = 3$

(c) $e - \chi + v = 2$

(d) $v - e + \chi = 2$ (✓)

Graph Representation

1). The Adjacency Matrix (for Undirected graph) - Let G is a graph. Then Adjacency matrix is

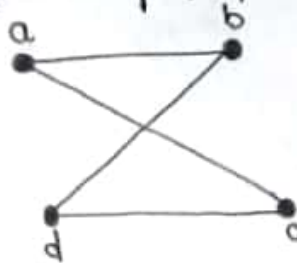
Represented as

$$A = [O_{ij}] = \begin{cases} 1 & \text{If } (i, j) \text{ is an edge of } G \\ 0 & \text{If there is no edge b/w them} \end{cases}$$

Ques. Find the adjacency matrix of following graphs

(i)

Vertex	Adjacency
a	b, c
b	d, a
c	d, a
d	b, c

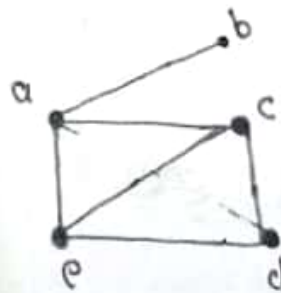


Matrix =

$$\begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

(ii)

Vertex	Adjacency
a	b, c, e
b	a
c	d, e, a
d	e, c
e	c, d, a



Matrix =

$$\begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 \end{bmatrix} \end{matrix}$$

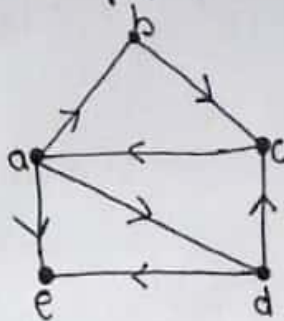
Adjacency Matrix of directed graph

The adjacency Matrix of a directed graph is represented as

$$A = [a_{ij}] = \begin{cases} 1, & \text{if an edge directed from } i^{\text{th}} \text{ vertex to } j^{\text{th}} \text{ vertex} \\ 0, & \text{otherwise} \end{cases}$$

Ques. Find the Adjacency matrix of following graph

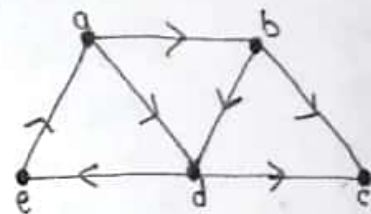
Vertex	Adjacency
a	b, d, e
b	c
c	a
d	e, c
e	



$$A = \begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Ques. Find the adjacency matrix of following graph

Vertex	Adjacency
a	b, d
b	c, d
c	—
d	e, c
e	a



$$A = [a_{ij}] = \begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Incidence Matrix of undirected graph - Let $G = (V, E)$ be an undirected graph
let $1, 2, 3, \dots, n$ have n vertices and
 $e_1, e_2, e_3, \dots, e_m$ edges of G .

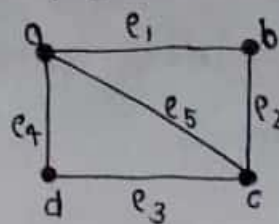
Then incidence Matrix

$B = \{b_{ij}\}$ is $n \times m$ matrix

$$B = [b_{ij}] = \begin{cases} 1 & \text{when edge } e_j \text{ is incident with } i \\ 0 & \text{otherwise i.e. } e_j \text{ is not incident with vertex } i \end{cases}$$

Ques. Find the incidence Matrix of following graph

Vertex	Adjacency
a	e_1, e_4, e_5
b	e_2, e_1
c	e_2, e_3, e_5
d	e_4, e_3

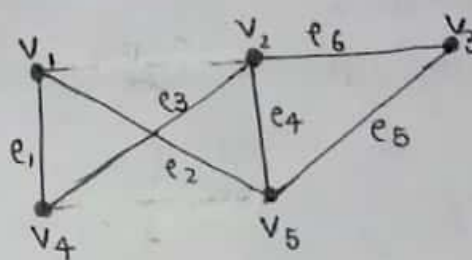


Incidence matrix

	e_1	e_2	e_3	e_4	e_5
a	1	0	0	1	1
b	1	1	0	0	0
c	0	1	1	0	1
d	0	0	1	1	0

Ques. Find the Incidence Matrix.

Vertex	Adjacency
v_1	e_1, e_2
v_2	e_3, e_4, e_6
v_3	e_5, e_6
v_4	e_1, e_3
v_5	e_2, e_4, e_5

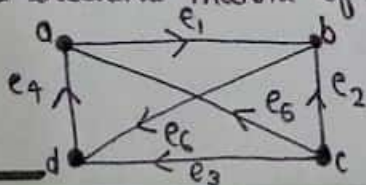


	e_1	e_2	e_3	e_4	e_5	e_6
v_1	1	1	0	0	0	0
v_2	0	0	1	1	0	1
v_3	0	0	0	0	1	1
v_4	1	0	1	0	0	0
v_5	0	1	0	1	1	0

Incidence matrix of directed graph - The incidence matrix of directed graph is $B = \{b_{ij}\}$.

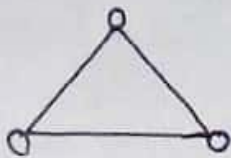
$$\text{where } b_{ij} = \begin{cases} 1 & \text{if edge } e_j \text{ is directed away from vertex } i \\ -1 & \text{" " " " toward " " } \\ 0 & \text{otherwise} \end{cases}$$

Ques. Find the incidence matrix of following graph

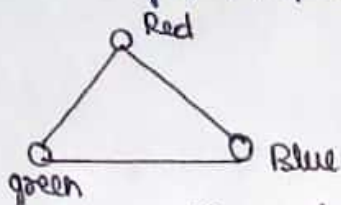


Vertex	adjacency	matrix
a	e_1, e_5, e_4	$ \begin{matrix} & e_1 & e_2 & e_3 & e_4 & e_5 & e_6 \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 1 & 0 & 0 & -1 & -1 & 0 \\ -1 & -1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & -1 & 1 & 0 & -1 \end{bmatrix} \end{matrix} $
b	e_2, e_1, e_6	
c	e_5, e_2, e_3	
d	e_4, e_3, e_6	

Graph coloring



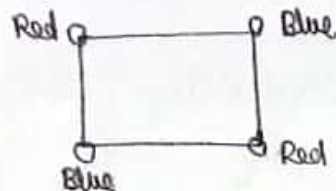
coloring of graph constitute vertex (Node) of graph. In coloring of graph it is necessary the two adjacent vertex are not in the same color. The least no. of color needed for coloring the graph is called chromatic Number. It is also known as K-chromatic graph, Represented as $\chi(G)$ or $X(G)$



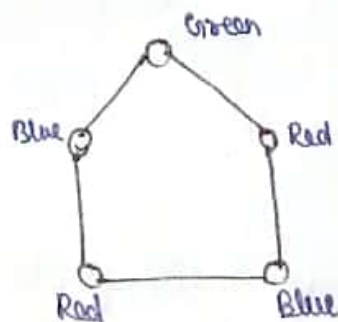
chai

Ques find the chromatic no. of the following graphs

- (i) Then Minimum color is 2
 $[\chi(G) = 2]$



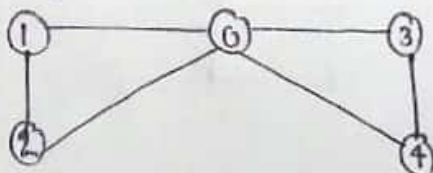
- (ii) Minimum color is 3
 $[\chi(G) = 3]$



Trail - A walk in which no edge is Repeated (only vertex can be repeated).

Euler circuit - A Trail which start and ends at same vertex. The circuit is called Euler circuit.

A graph which has Euler circuit is called Euler graph.



2 → 1 → 0 → 3 → 4 → 0 → 2

This is a Euler circuit

* For Euler graph Every Vertex (Node) has even degree.

$$\deg(1) = 2$$

$$\deg(0) = 4$$

$$\deg(2) = 2$$

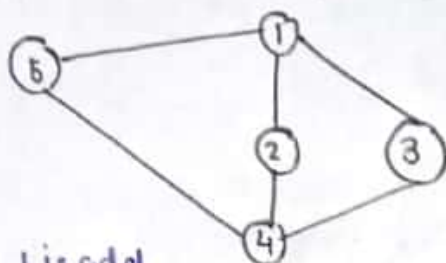
$$\deg(3) = 2$$

$$\deg(4) = 2$$

Each node has even degree. So graph is Euler graph.

Note - Any Node which has 0 degree is Euler circuit.

Ques.

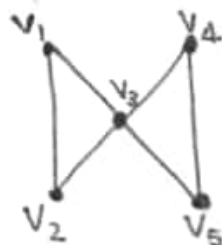


$$\deg(5) = 2$$

$$\deg(1) = 3$$

Since degree of 1 is odd
so it is not a

Ques. Identify the graph for Eulerian circuit and examine the graph is Euler graph or not.



Euler circuit

$$V_1 \rightarrow V_3 \rightarrow V_4 \rightarrow V_5 \rightarrow V_3 \rightarrow V_2 \rightarrow V_1$$

$$\deg(V_1) = 2$$

$$\deg(V_5) = 2$$

$$\deg(V_2) = 2$$

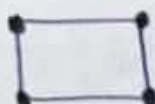
$$\deg(V_4) = 2$$

$$\deg(V_3) = 4$$

deg. of all vertex are even so graph is Euler graph.

Hamiltonian graph - A graph is said to be Hamiltonian graph if it contains (Node can't be repeated) Hamiltonian circuit. In Hamiltonian circuit graph passes through every vertex exactly once.

(i) In Hamiltonian graph $n \geq 3$ (No. of vertices)



(ii) Now we find $\frac{n}{2}$

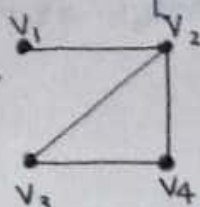
(iii) Then graph is called Hamiltonian graph [when degree of vertex $\geq \frac{n}{2}$]

Ques. Examine the following graphs is Hamiltonian or not.

(i)

$$n = 4 > 3$$

$$\frac{n}{2} = \frac{4}{2} = 2$$



i.e. for Hamiltonian graph degree of each vertex ≥ 2 .

$$\text{Degree } V_1 = 1 < 2$$

So it is not Hamiltonian graph.

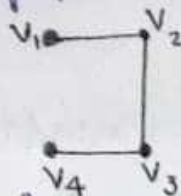
(ii)

$$n = 4$$

$$\frac{n}{2} = \frac{4}{2} = 2$$

$$\text{deg}(V_1) = 1 < 2$$

So graph is not Hamiltonian graph.



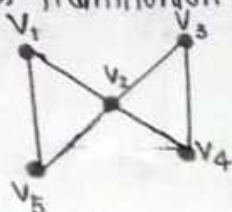
Ques. Examine the graph is Hamiltonian or not.

$$n = 5$$

$$\frac{n}{2} = \frac{5}{2} = 2.5$$

i.e. 2

for Hamiltonian graph degree of each vertex ≥ 2



$$\text{d}(V_1) = 2$$

$$\text{d}(V_4) = 2$$

$$\text{d}(V_2) = 4$$

$$\text{d}(V_5) = 2$$

$$\text{d}(V_3) = 2$$

degree of each vertex is ≥ 2 . So graph is Hamiltonian graph.

Ques. Examine the following graph is Hamiltonian or not.

$$(i) \quad n = 5 \geq 3$$

$$\frac{n}{2} = 2.5 \text{ i.e. } 2$$

i.e. for Hamiltonian graph degree of each vertex ≥ 2

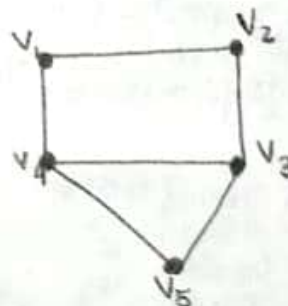
$$\text{deg}(V_1) = 2$$

$$\text{deg}(V_2) = 3$$

$$\text{deg}(V_3) = 3$$

$$\text{deg}(V_4) = 4$$

$$\text{deg}(V_5) = 2$$

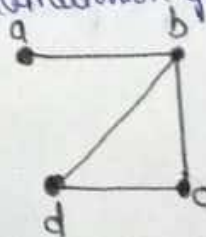


\Rightarrow degree of each vertex is ≥ 2 . So graph is Hamiltonian graph.

$$(ii) \quad n = 4 \geq 3$$

$$\frac{n}{2} = 2$$

for Hamiltonian graph degree of each vertex ≥ 2



$$\deg(a) = 1 < 2$$

so it is not a Hamiltonian graph.

Isomorphic graph Two graph $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ are said to be isomorphic graph, when it contain

(a) Same No. of vertices.

(b) Same No. of edges.

(c) Equal No. of Vertices of same degree.

If these properties are satisfied then graph are isomorphic graph.

Ques. Examine the following graphs is isomorphic or not.

(i) No. of vertex in $G_1 = 4$

↳ No. of vertex in $G_2 = 4$

No. of vertex in $G_1 =$ No. of vertex in G_2

↳ No. of edge in $G_1 = 4$

" " " in $G_2 = 4$

No. of edge in $G_1 =$ No. of edge in G_2

↳ $\deg U_1 = 2$ $\deg V_1 = 2$

$\deg U_2 = 2$ $\deg V_2 = 2$

$\deg U_3 = 2$ $\deg V_3 = 2$

$\deg U_4 = 2$ $\deg V_4 = 2$

degree are same in each graph of. Hence all the properties of isomorphic graph so it is called isomorphic graph.

(ii) Examine the following graph is isomorphic or not.

↳ No. of vertices in $G_1 = 5$

No. of vertices in $G_2 = 5$

} \Rightarrow No. of vertices in G_1 and G_2 are same

↳ No. of edges in $G_1 = 4$

↳ No. of edges in $G_2 = 4$

No. of edges in $G_1 =$ No. of edges in G_2

↳ $\deg(V_1) = 1 = \deg(U_1)$

$\deg(V_2) = 2 = \deg(U_2)$

$\deg(V_3) \neq \deg(U_3)$

So the given graph is not a Isomorphic graph.

