

## Lecture 24: 1 July, 2021

Madhavan Mukund

<https://www.cmi.ac.in/~madhavan>

Data Mining and Machine Learning  
April–July 2021

# Information retrieval on the Internet

- Traditional IR
  - Books published after editing, review — trustworthy content
- IR for Internet
  - Internet documents are self-published, unverified
  - Economic incentive to boost rankings through fraudulent means
  - Ranking algorithms should try not to be fooled
- Easy to add invisible content in HTML to misdirect search
  - Merging text and background colour, overlay text with images, unreadable font size
- Self published documents may omit useful search terms
  - IBM webpage did not mention the word “computer”

# Exploiting hypertext

- Hypertext links refer from one document to another
  - `<a href="https://www.cmi.ac.in"> CMI webpage </a>`
  - Target location : `https://www.cmi.ac.in`
  - Anchor text : `CMI webpage`
- Use anchor text to index document at target location
  - Reliable indicator of what target document is about
- Hyperlinks also connect internet documents as a directed graph
  - Reason about the World Wide Web (WWW) as a gigantic graph
  - Use techniques from **social network analysis**

# Social network analysis — prestige

- Consider the film industry
  - When is an actor a star? When is a director famous?
  - Stars are sought out by famous directors
  - Famous directors get stars to work in their films
  - Recursive definition
- Network (graph) of actors and directors, matrix  $M$

$$\begin{array}{c} \text{Directors} \\ j \\ \vdots \\ 1 \end{array} \quad \begin{array}{c} \text{Actors} \\ i \end{array} \left[ \begin{array}{c} \dots \\ \dots \\ 1 \end{array} \right]$$

$M[i,j] = 1$  if Actor  $i$  works in a film directed by Director  $j$

# Social network analysis — prestige

- Each actor  $i$  has star value  $S[i]$
- Each director  $j$  has fame  $F[j]$
- Actors derive star value from the famous directors they work with

$$S[i] = \sum_j M[i,j] \cdot F[j], \text{ or } S = M \cdot F$$

- Directors derive fame from the stars who work with them

$$F[j] = \sum_i M[i,j] \cdot S[i], \text{ or } F = M^T \cdot S$$

- Substituting  $F$  from second equation,  $S = M \cdot M^T \cdot S$
- Substituting  $S$  from first equation,  $F = M^T \cdot M \cdot F$
- Solve for  $S$ ,  $F$  to compute star ratings, fame

# Prestige for webpages

- Each document  $i$  has prestige  $P[i]$
- Prestigious (reliable) documents confer prestige on documents they link to
  - $P[i]$  is shared equally among all outgoing links
- A document derives prestige from documents that link to it
  - $P[i]$  is sum of prestige transferred by incoming links
- Structure of the internet, adjacency matrix  $A$

$$\begin{array}{c} \text{Webpages} \\ j \\ \vdots \\ \text{Webpages } i \left[ \begin{array}{c} \dots \\ 1 \end{array} \right] \end{array}$$

$A[i, j] = 1$  if webpage  $i$  has a link to webpage  $j$

## Prestige for webpages ...

- Suppose  $A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix}$
- Each document initially has prestige 1,  $P = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$
- If a webpage points to  $n$  other pages, each of them gets  $1/n$  of  $P[i]$
- Prestige transfer matrix,  $A^* = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix}$
- One step:  $P^\top \cdot A^* = \begin{bmatrix} 1 & 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} 1.5 & 1 & 0.5 \end{bmatrix}$

# Page rank

- Stable solution:  $P^\top \cdot A^* = P^\top$
- $P[i]$  is Page rank of webpage  $i$ 
  - Larry Page, co-founder of Google with Sergey Brin
- How do we compute  $P^\top$ ?
- $A^*$  is a stochastic matrix — each row sums to 1

$$\forall i \sum_j A^*[i,j] = 1$$

- Interpret  $A^*[i,j]$  as probability of moving from document  $i$  to document  $j$  — random web surfer
- Use theory of Markov chains

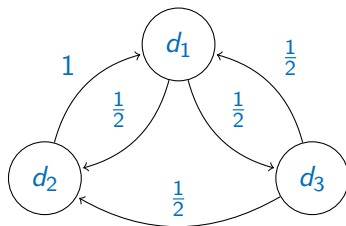


# Markov chains

- Finite set of states, with transition probabilities between states
- For us, states are documents
  - Henceforth, write  $A^*$  as  $A$  for convenience

$$A = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix}$$

Three  
state  
Markov  
chain



- $P[j]$  is probability of being in document  $j$
- Start in document 1, so initially  $P = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$

# Markov chains ...

- After one step:  $P^T A = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix}$

- After second step:  $\begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 \end{bmatrix} = \begin{bmatrix} \frac{3}{4} & \frac{1}{4} & 0 \end{bmatrix}$

- After  $k$  steps,  $P[j]$  is probability of being in state  $j$

- Continuing our example,

$$\begin{bmatrix} \frac{3}{4} & \frac{1}{4} & 0 \end{bmatrix} \rightarrow \begin{bmatrix} \frac{1}{4} & \frac{3}{8} & \frac{3}{8} \end{bmatrix} \rightarrow \begin{bmatrix} \frac{9}{16} & \frac{5}{16} & \frac{1}{8} \end{bmatrix}$$

- Is it the case that  $P[j] > 0$  for all  $j$  continuously, after some point?

- Markov chain  $A$  is **ergodic** if there is some  $t_0$  such that for every  $P$ , for all  $t > t_0$ , for every  $j$ ,  $(P^\top A^t)[j] > 0$ .
  - No matter where we start, after  $t > t_0$  steps, every state has a nonzero probability of being visited in step  $t$
- Properties of ergodic Markov chains
  - There is a stationary distribution  $\pi$  such that  $\pi^\top A = \pi^\top$ 
    - $\pi^\top$  is a **left eigenvector** of  $A$
  - For *any* starting distribution  $P$ ,  $\lim_{t \rightarrow \infty} P^\top A^t = \pi^\top$

# Ergodicity ...

- How can ergodicity fail?
  - Starting from  $i$ , we reach a set of states from which there is no path back to  $i$
  - We have a cycle  $i \rightarrow j \rightarrow k \rightarrow i \rightarrow j \rightarrow k \dots$ , so we can only visit some states periodically
- Sufficient conditions for ergodicity
  - **Irreducibility**: When viewed as a directed graph,  $A$  is strongly connected
    - For all states  $i, j$ , there is a path from  $i$  to  $j$  and a path from  $j$  to  $i$
  - **Aperiodicity**: For any pair of vertices  $i, j$ , the gcd of the lengths of all paths from  $i$  to  $j$  is 1
    - In particular, paths (loops) from  $i$  to  $i$  do not all have lengths that are multiples of some  $k \geq 2$
    - Prevents bad cycles

# Making the web graph ergodic

- No reason why web graph is irreducible and aperiodic
- Web graph has dead ends — terminal documents, no outgoing links
- Solution: Add random jumps between documents — **teleportation**
- Teleportation matrix  $T$ : For all  $i, j$ ,  $T[i, j] = 1/N$ , where  $N$  is the total number of documents
  - The random surfer ignores all the links in the current document and types a new URL
- Let  $\alpha$  be the probability of teleportation:  $M = \alpha T + (1 - \alpha)A$ 
  - Check that  $M$  is stochastic
- By construction,
  - $M$  is strongly connected — direct edge between each pair of documents
  - $M$  is aperiodic — paths of any length exist between  $i$  and  $j$
  - $M$  has no dead ends

# Page Rank

- In the modified web graph, stationary distribution is the Page rank,  $\pi^T M = \pi^T$
- Compute using  $\lim_{t \rightarrow \infty} P^T M^t$
- Use recursive doubling to accelerate computation of  $\lim_{t \rightarrow \infty} P^T M^t$ 
  - Compute  $M, M^2, (M^2)^2 = M^4, \dots, (M^{2^i})^2 = M^{4^i}, \dots$
  - Set a threshold for progress to stop the process
- Some limitations of Page rank
  - Universal property of a webpage, independent of a query
  - Define a topic-sensitive page rank
- Page rank was one the keys to the initial success of Google
  - Constant tweaks to ranking algorithm to keep ahead of search engine optimizers (SEO)