

ASSIGNMENT 4

Title : Thread Synchronization using Counting Semaphore

Problem Statement : Thread synchronisation using counting semaphores. Application to demonstrate : producer - consumer problem with counting semaphores and mutex.

Theory :-

Semaphore - a variable which is non-negative and shared between threads.

Binary Semaphore - This is also known as mutex lock. It can have only two values 0 and 1. Its value is initialized to 1. It is used to implement the solution of critical section problem with multiple processes.

Counting Semaphore - Its value can range over an unrestricted domain. It is used to control access to a resource that has multiple instances.

Producer - Consumer Problem - Given a buffer of fixed size, a producer can produce an item and place it in the buffer. A consumer can pick items and consume them. We need to ensure that when a producer is placing an item in the buffer, then at the same time consumer should not consume any item. In this problem, buffer is the critical section.

Algorithm :

Consumer Function :-

```
void consumer() {
```

```
// consumes items and finally pops from buffer and processes
```

```
int itemc;
```

```
while (true) {
```

```
    while (count == 0); // buffer empty
```

```
    itemc = Buffer(Out);
```

```
    (Out) = (Out + 1) mod n;
```

```
    count -= 1;
```

```
    process_item(itemc);
```

```
    // I1 load Rc, count
```

```
    // I2 dec Rc
```

```
    // I3 store count, Rc
```

```
}
```

```
}
```

```
int count = 0; // global variable shared by both
```

```
// buffer common to both, shared
```

Producer Function

```
void Producer() {
```

```
// produces items and puts to buffer
```

```
int itemp;
```

```
while (true) {
```

```
    while (count == n); // buffer full
```

```
    Buffer[In] = itemp;
```

```
    In = (In + 1) mod n;
```

```
    count += 1;
```

```
    // I1 load Rp, count
```

```
    // I2 inc Rp
```

```
    // I3 store count, Rc
```

```
}
```

```
}
```

Execution :

- Producer arrives executes I_1, I_2
- Before executing I_3 , consumer arrives, so producer preempts
- Consumer execution begins executes I_1, I_2
- Before executing I_3 , producer resumes and executes I_3
- After producer terminates consumer I_3 executes

Conclusion : I have successfully implemented producer - consumer problem with counting semaphores and mutex.