

Program : 02

Lexical Analyzer Using C

Aim:

To develop a lexical analyzer to identify identifiers, constants, comments, operators etc. using C program.

Algorithm:

Step 1: Start the Program

Step 2: Declare all the variables and file Pointers.

Step 3: Display the input program

Step 4: Separate the keyword in the program and display it.

Step 5: Display the header files of the input program

Step 6: Separate the operators of the input program and display it.

Step 7: Print the Punctuation marks.

Step 8: Print the constants that are present in input program.

Step 9: Print the identifiers of the input program

Step 10: Stop the program.

Calculator Using LEX

Aim:

To write an algorithm for an calculator and execute program using lexical analyzer tool

Algorithm:

Step 1: Take the input string

Step 2: Call the function yylexc(). The control will given to rule section.

Step 3: Define the rules.

dig [0-9] + / ([0-9] *) " ." ([0-9] +)

add "+"

sub "-"

mul "*"

div "/"

pow "^n"

Step 4: check the rules as

{dig} {dig1} ; }

{add} {OP = 1 ; }

{sub} {OP = 2 ; }

{mul} {OP = 3 ; }

{div} {OP = 4 ; }

{pow} {OP = 5 ; }

{ln} {printf("In The Answer : %f \n", a) ; }

switch (OP)

{

Case 1: $a \neq a+b;$

break;

Case 2: $a = a-b;$

break;

Case 3: $a = a \times b;$

break;

Case 4: $a = a/b;$

break.

Cases : for ($i=a; b>1; b--$)

$a = a \times i;$

break;

3 OP = 0; 33

Step 5: Call yywrap() to finish up the session

Step 6: stop.

Result:

Program run successfully and output verified



Counting Vowels And Consonants

Aim:

To write an algorithm for counting vowels and execute the program using lexical analyzer tool.

Algorithm:

Step 1: Start

Step 2: Take the input string

Step 3: Call the function yyleng(). The control will given to rule section

Step 4: Check the rule if input is [aeiouAEIOU] then increment count as vt++

Else

Step 5: Increment count as ct++

Step 6: Output the no of vowels and consonants

Step 7: Stop.

Result:

Program run successfully and output verified.

Counting Number of Words, Lines and Character, etc

Aim:

To write an algorithm for counting number of words, lines and characters and execute the program using lexical analyzer tool.

Algorithm :

Step 1: Start

Step 2: Take a string as input

Step 3: Call yylex(), Step 4: check for rule section

Step 5: if \n then lc++

Else check [] then sc++

Else if check \t then tc++

Else if characters then ch++

Step 6: Call yywrap for wrapping the lex section

Step 7: Print the output

Step 8: STOP

Result:

Program run successfully and output verified

Experiment - 6

Intermediate Code Generation

Aim:

To Implement Intermediate code generation for simple expression using Program in C

Algorithm:

Step 1: Start the Program

Step 2: Open the input file in read mode

Step 3: Open the output file in write mode

Step 4: In Input file scan for operator, argument 1, argument 2 and result

Step 5: If the operator is '+', move arg1 to R0, move R0 to result ^{Add arg2 and R0}

Step 6: If the operator is '-', move arg1 to R0, Subtract arg2 and R0, Mov R0 to result

Step 7: If the operator is '*' move arg1 to R0, multiply arg2 and R0, Move R0 to result

Step 8: If the operator is '/', move arg1 to R0, Divide arg2 and R0, mov R0 to result

Step 9: If the operator is '=', move arg1 to R0 mov R0 to result

Step 10: Close both the files

Step-11: Stop the program.

NFA to DFA Conversion

Aim:

To write an algorithm for NFA to DFA converter
and execute the program in C.

Algorithm :

Step 1: start the program

Step 2: Input the required array i.e, set of alphabets, set of states,
initial state, final state,

Step 3: Initially $Q' = \emptyset$ and input

Step 4: Input the no. of states, final state and rules as
input

Step 5: Add q_0 of NFA to Q' . Then find the transitions from
this start state.

Step 6: In Q' , find the possible set of states for each input
symbol. If this set of states is not in Q' , then add
it to Q' .

Step 7: Solving according to DFA

Step 8: In DFA, the final state will be all the states which
contain final states of NFA.

Step 9: Print the output state as table

Step 10: STOP

Experiment - 8

Constant Propagation

Aim:

To write an algorithm for Constant Propagation and execute the program in C.

Algorithm:

Step 1: Start the Program

Step 2: Input the maximum no. of expressions and the inputs

Step 3: Set arr[i].Flag = 0

Step 4: Construct a control flow graph (CFG)

Step 5: Associate transfer functions with the edges of CFG

Step 6: At every node, maintain the values of the program's variable at that point. Then initialize those to 1

Step 7: Iterate until the values of the variables stabilize

Step 8: Stop

Experiment - 9

YACC Specification to Recognize a Valid Arithmetic Expression

Aim:

To write a Yacc Program to valid arithmetic expression using Yacc.

Algorithm:

Step 1: Start the program

Step 2: Enter the expression (Reading)

Step 3: Checking the validating of the given expression according to the rule using Yacc.

Step 4: Using expression rule, Print the result of the given values

Step 5: Stop the program

Experiment - 10

Back-end of a Compiler

Aim:

To implement the back end of the compiler
using C

Algorithm :

Step 1: Start the program

Step 2: Get / Read the three address code

Step 3: Check the operation in st 1 and copy to variable

Step 4: Print MOV and variable R

Step 5: Print operation second variable R

Step 6: Print MOV R, result variable

Step 7: STOP

Experiment - 11

First and Follow

Aim:

To Write a Program to find first and follow of any given grammar.

Algorithm

Step 1: Start the Program

Step 2: Calculating First, $\alpha \rightarrow t\beta$

Step 3: If α is a terminal, then $\text{FIRST}(\alpha) = \{\alpha\}$

Step 4: If α is a non-terminal and $\alpha \rightarrow E_j$ is a production, then $\text{FIRST}(\alpha) = \{E_j\}$

Step 5: If α is a non-terminal and $\alpha \rightarrow r_1, r_2, r_3, \dots, r_n$ and any $\text{FIRST}(r_i)$ contain t then t is in $\text{FIRST}(\alpha)$

Step 6: Calculating Follow

Step 7: If α is a start symbol, then $\text{FOLLOW}(\alpha) = \$$

Step 8: If α is a non-terminal and has a production $\alpha \rightarrow AB$ then $\text{FIRST}(B)$ is in $\text{Follow}(\alpha)$ except E_j

Step 9: If α is a non-terminal and has a production $\alpha \rightarrow ABB$, where $B \neq E_j$, then $\text{Follow}(A)$ is in $\text{Follow}(\alpha)$

Step 10: Stop the program.

Experiment - 12

Shift Reduce Parser

Aim:

To construct a shift reduce parser for a given language in C.

Algorithm :

Step 1: Start the program

Step 2: Assume the grammar be,

$$E \rightarrow E + E$$

$$E \rightarrow E * E$$

$$E \rightarrow (E)$$

$$E \rightarrow id$$

Step 3: Read input variables

Step 4: ~~Stack~~ Set top element of stack as \$

Step 5: Push an element from input string to stack

Step 6: If the pushed element is operator, no action

Step 7: Else, reduce the element in the stack to the given grammar

Step 8: Repeat steps 5 to 7 until input string is \$

Step 9: If the input string is \$ and stack is reduced to only E, then Accept, Else, Reject

Step 10: STOP.