Quick Sort

- Quick Sort follows the **Divide and Conquer** approach.
- It selects a **pivot**, partitions the array around the pivot, and recursively sorts the left and right sub-arrays.

Steps of Quick Sort:

- 1. Divide: Choose a pivot and partition the array into elements smaller than and greater than the pivot.
- 2. Conquer (Sort Recursively): Apply Quick Sort recursively on the left and right subarrays.
- 3. **Combine:** Since sorting happens in-place, no explicit merging is needed.

Example

Consider the array: [8, 3, 7, 4, 9, 2, 6, 5]

Step 1: First Partition

Pivot = 6

- Left: [3, 4, 2, 5] (elements < 6)
- Middle: [6]
- Right: [8, 7, 9] (elements > 6)

Step 2: Sorting Left [3, 4, 2, 5]

Pivot = 4

- Left: [3, 2]
- Middle: [4]
- Right: [5]

Step 3: Sorting [3, 2]

Pivot = 3

- Left: [2]
- Middle: [3]
- Right: []

Sorted left partition: [2, 3, 4, 5]



Step 4: Sorting Right [8, 7, 9]

Pivot = 7

• Left: []

• Middle: [7]

• Right: [8, 9]

Step 5: Sorting [8, 9]

Pivot = 8

• Left: []

• Middle: [8]

• Right: [9]

Sorted right partition: [7, 8, 9]

Final Sorted Array:

Merging all partitions: [2, 3, 4, 5, 6, 7, 8, 9]

Time Complexity

• Best Case: O(n log n)

• Average Case: O(n log n)

• Worst Case: O(n²) (If pivot selection is poor)

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Why Use Quick Sort?

- **✓ Faster for most cases** (In-place sorting, unlike Merge Sort).
- No extra space needed (Except for recursive calls).
- **✓** Good for large datasets with efficient pivot selection.
- X Not stable (May not preserve order of equal elements).
- \bowtie Worst case $O(n^2)$ if pivot selection is bad.