

4) Phase/State driven code.

- It uses nested if-then statements, case-statements or a finite state automation to break up the processing of a function into discrete code segments.
- This allows temporary suspension of code segment without loss of critical data, procedure task;

```
begin
```

```
  case flag of
```

```
    1: begin
```

```
      perform_part_1;
```

```
      flag := 2
```

```
    end
```

```
    2: begin
```

```
      perform_part_2;
```

```
      flag := 3
```

```
    end end
```

```
    3: begin
```

```
      perform_part_3;
```

```
      flag := 1
```

```
    end
```

```
end.
```

At the end of each state a flag is set and the process is terminated. Upon restarting, the process resumes where it left off.

5. Co-routines:-

- In this scheme, two or more tasks are coded in the state-driven fashion, just and after each phase is completed, a call is made to a central dispatcher. The dispatcher holds the program counter for a list of tasks that are executed in round robin fashion. i.e. It selects next task to execute. This task then executes until its next phase is completed and the central dispatcher is called again.
- If there is only 1 task, then it becomes cyclic. & The communication between task is achieved by global variable.

```
void task_a(void)
```

```
{
```

```
    switch (state-a)
```

```
{
```

```
    case 1: phase-a1();
```

```
    state-a++;  
    break; /* to dispatcher */
```

```
    case 2: phase-a2();
```

```
    state-a++;  
    break;
```

```
    case 3: phase-a3();
```

```
    state-a++;  
    break;
```

```
}
```

```
void task_b(void)
```

```
{
```

```
    switch (state-b)
```

```
{
```

```
    case 1: phase-b1();
```

```
    break;
```

```
    case 2: phase-b2();
```

```
    break;
```

```
    case 3: phase-b3();
```

```
    break;
```

```
}
```

State-a & state-b are global variable managed by dispatcher.

2. Interrupt-Driven System:-

- Interrupt is concept is:

- Interrupt only system

- H/W interrupt (I/O) Real Time Clock or External device send interrupt to Interrupt controller (IC).

- IC send interrupt signal to CPU depending on order of arrival & priority

- While context switching,
↓
is save

- Context is a snapshot of task being processed

- While context switching, sufficient info is saved so that resuming interrupted task is possible.

- Context switching ko bola hai Interrupt qor qor

3. Preemptive Priority System:-

- higher priority task $\frac{a}{b}$ $\frac{a}{b}$ lower priority task preempt $\frac{a}{b}$

- higher priority task $\frac{a}{b}$ Resource $\frac{a}{b}$ lower priority task's resource starvation $\frac{a}{b}$ $\frac{a}{b}$.

- Rate monotonic System $\frac{a}{b}$ higher the execution frequency, higher the priority.

4) Hybrid Scheduling System:

- ① combination of preemptive & round robin scheduling
 → same priority handled with round robin fashion.
- ② combinations of interrupt & fixed rates or not fixed rate. (sporadically).
- ③ combination of foreground & background systems.

The Task Control Block Model

Rate Monotonic Algorithms:-

	Burst Time	Arrival Time	Period.
P_1	2	0	10
P_2	1	0	5
P_3	5	0	30
P_4	2	0	15

Here $N = 4$

$$\begin{aligned} \text{Least Upper bound of utilization for schedulability} &= N(2^{1/N} - 1) \\ &= 4(2^{1/4} - 1) \\ &= 0.7565. \end{aligned}$$

$$\text{Now; Utilization of } P_1 = \frac{T_1}{P_1} = \frac{2}{10} =$$

$$\text{Utilization of } P_2 = \frac{T_2}{P_2} = \frac{1}{5}$$

$$\text{" " } P_3 = \frac{T_3}{P_3} = \frac{5}{30}$$

$$\text{" " } P_4 = \frac{T_4}{P_4} = \frac{2}{15}$$

$$\therefore \text{Total Utilization} = \frac{2}{10} + \frac{1}{5} + \frac{5}{30} + \frac{2}{15} = 0.7$$

Since $0.7 < 0.7568$, Thus these processes can be scheduled by rate monotonic scheduling.

Highest LCM of Periods = 2, 10, 5, 30, 15

$$5 \begin{array}{r} 5, 5, 15, 15 \\ 3 \end{array}$$

$$3 \begin{array}{r} 1, 1, 3, 3 \\ 1 \end{array}$$

$$1, 1, 1, 1$$

Here, highest common divisor is 5.

P_1, P_2, P_3, P_4

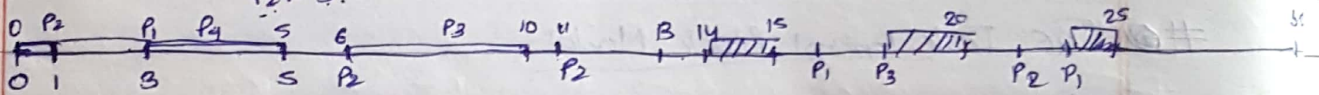
P_1, P_2, P_3

P_2, P_4

P_1, P_2

P_2

P_1, P_2, P_3



P_2