

Tutorial-7

Priyanka Pande
D-28, 20/6/2006

Q1 Greedy Algo Paradigm:- Greedy is an algorithmic paradigm that builds up a solution piece by piece always choosing the next most promising piece..

Greedy Algos are simply heuristic algorithms used for optimization (either maximize or minimize) problems.

Q2 (i) Activity Selection:-
Time Complexity = $O(n \log n)$ (if input activities may not be sorted).
= $O(n)$ (when input activities are sorted)

(ii) Job Scheduling \Rightarrow Time Complexity = $O(n \log n)$.
Space Complexity = $O(n)$.

(iii) Fractional Knapsack \Rightarrow TC $\Rightarrow O(n \log n)$
SC $\Rightarrow O(1)$

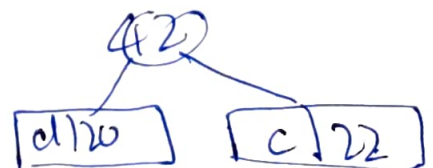
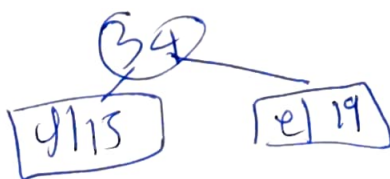
(iv) Huffman's Coding \Rightarrow TC $\Rightarrow O(n \log n)$
SC $\Rightarrow O(n)$.

Q3
 $a = 45$ $c = 22$ $e = 19$
 $b = 23$ $d = 20$ $f = 15$

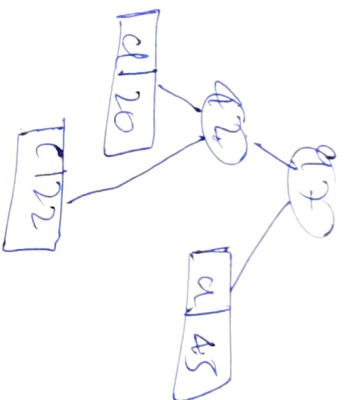
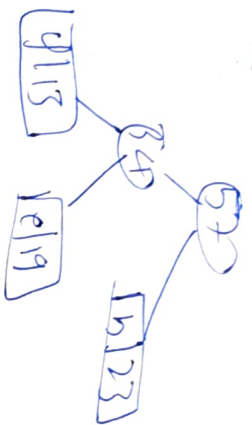
Step 1



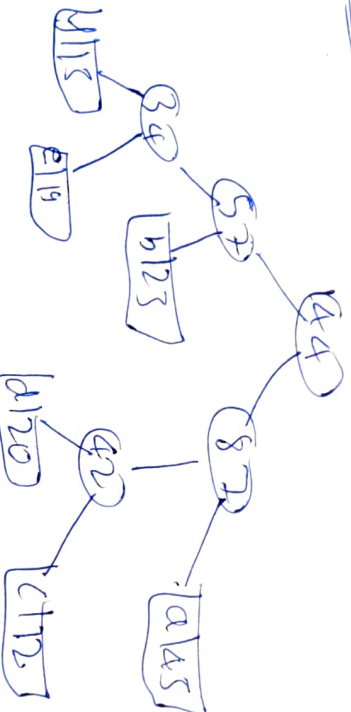
Step 2



Step 5-



Step 4:-



$$\begin{array}{llll}
 a = 11 & c = 101 & e = 001 \\
 b = 01 & d = 100 & f = 000
 \end{array}$$

Q8 We can optimize the approach we do solve the job sequencing problem by using priority queue (max heap)

Algo.

- Sort the job based on deadline.
- Generate from the end calculate the max every 2 consecutive deadlines.
- Which the slot are available & there are job get slot in the max heap.
- Sort the result array based on their deadline.

Q.1 Disadvantages of Greedy Approach.

→ It is not suitable for problems where a solution is required for every subproblem. The greedy strategy can be wrong in most cases even when it gives an optimal solution.

Ex: i) Dijkstra's Algorithm does not find path with -ve weights.

Q.2 We can use greedy approach to solve the problem by always going to the nearest possible city. We select any of the city as the 1st one.

(i) We can build a position of cities in a way that the greedy strategy gives the worst possible solution.

Q.3

Start Time	1	2	0	6	9	10
End Time	3	5	7	8	11	12

void PrintMaxActivities (Activity Arr[], int n)

{

Sort (arr, arr+n, activity Compare).

cout << "Following activities are selected";

int i=0;

cout << " (" << arr[i].start << ", " << arr[i].

for (int j=1; j<n; j++) {

if (arr[j].start >= arr[i].finish) {

cout << ", " << arr[j].start << ", " << arr[j].

finish $c^{\prime\prime}, j^{\prime\prime}$;
 $c^{\prime} = j^{\prime}$;

}
}

Print maxLen () {

 int arr[] = { (1,3), (2,5), (6,7), (6,8),
 (11,12) };

 int n = size of arr (size of arr[a]);

 Print max Activity (arr, n);

 return 0;

}

2016906