

Sol 1

```
int LinearSearch (int arr, int n, int key)
{
    for i >= 0 to n-1
        if arr[i] == key
            return i
    } return -1.
```

Sol 2

Iterative Insertion Sort

```
void insertion-sort (int arr[], int n)
{
    int j, temp, g;
    for j < 1 to n
        temp ← arr[j]
        j ← j - 1
        while (j >= 0 AND arr[j] > temp)
            arr[j+1] ← arr[j]
        arr[j+1] ← temp
```

Recursive Insertion Sort

```
void insertion-sort (int arr[], int n)
{
    if (n ≤ 1) return
    insertion-sort (arr, n-1)
    last ← arr[n-1]
    g ← n-2
    while (j >= 0 && arr[j] > last)
        arr[j+1] = arr[j]
    arr[j+1] = last.
```

Sol 3

Selection Sort:

Time Complexity: Best case: $O(n^2)$, worst case = $O(n^2)$
Space Complexity: $O(1)$

Insertion Sort
TC: best case: $O(n)$, worst case = $O(n^2)$
SC: $O(1)$

Merge Sort:

$T(n) \Rightarrow \text{Best} \Rightarrow O(n \log n)$

Worst $\Rightarrow O(n \log n)$

$S.C \Rightarrow O(n)$

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Quick Sort:

$T(n) \Rightarrow \text{Best } O(n \log n)$

Worst $O(n^2)$

$S.C \Rightarrow O(n)$

Solⁿ 4

Sorting	Inplace	Stable
Selection	✓	✓
Insertion	✓	✓
Merge	✓	
Quick		✓
Heap	✓	
Bubble	✓	

Solⁿ 5

Quick Sort is the fastest general purpose sort, O_n most practical situation quick sort is the molecule of choice as stability is important and space is bubble sort, merge sort, might be best.

Sol 6 the worst $T.C$ of quick sort is $O(n^2)$. the worst case occurs when the picked first is always an extreme (smallest or largest) element. the ~~problem~~ happens when input array is sorted or reverse sorted array and often it or least element is picked as pivot.

Sol 7 Recurrence relation of:

a) Merge Sort $\Rightarrow T(n) = 2T(n/2) + n$

b) Quick Sort $\Rightarrow T(n) = 2T(n/2) + n$

Merge Sort is more efficient & uniform than quick sort in case of large array as 3 or 4 data sets.