

# **SOLUTIONS MANUAL**

# **DIGITAL DESIGN**

**WITH AN INTRODUCTION TO THE VERILOG HDL**  
**Fifth Edition**

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**CHAPTER 1**

- 1.1** Base-10: 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32  
 Octal: 20 21 22 23 24 25 26 27 30 31 32 33 34 35 36 37 40  
 Hex: 10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F 20  
 Base-12: 14 15 16 17 18 19 1A 1B 20 21 22 23 24 25 26 27 28
- 1.2** (a) 32,768 (b) 67,108,864 (c) 6,871,947,674
- 1.3**  $(4310)_5 = 4 * 5^3 + 3 * 5^2 + 1 * 5^1 = 580_{10}$   
 $(198)_{12} = 1 * 12^2 + 9 * 12^1 + 8 * 12^0 = 260_{10}$   
 $(435)_8 = 4 * 8^2 + 3 * 8^1 + 5 * 8^0 = 285_{10}$   
 $(345)_6 = 3 * 6^2 + 4 * 6^1 + 5 * 6^0 = 137_{10}$
- 1.4** 16-bit binary: 1111\_1111\_1111\_1111  
 Decimal equivalent:  $2^{16} - 1 = 65,535_{10}$   
 Hexadecimal equivalent: FFFF<sub>16</sub>
- 1.5** Let b = base  
 (a)  $14/2 = (b + 4)/2 = 5$ , so b = 6  
 (b)  $54/4 = (5*b + 4)/4 = b + 3$ , so  $5 * b = 52 - 4$ , and b = 8  
 (c)  $(2 * b + 4) + (b + 7) = 4b$ , so b = 11
- 1.6**  $(x - 3)(x - 6) = x^2 - (6 + 3)x + 6*3 = x^2 - 9x + 18$   
 Therefore:  $6 + 3 = b + 1m$ , so b = 8  
 Also,  $6*3 = (18)_{10} = (22)_8$
- 1.7**  $64CD_{16} = 0110\_0100\_1100\_1101_2 = 110\_010\_011\_001\_101 = (62315)_8$
- 1.8** (a) Results of repeated division by 2 (quotients are followed by remainders):  
 $431_{10} = 215(1); 107(1); 53(1); 26(1); 13(0); 6(1) 3(0) 1(1)$   
 Answer: 1111\_1010<sub>2</sub> = FA<sub>16</sub>  
 (b) Results of repeated division by 16:  
 $431_{10} = 26(15); 1(10)$  (Faster)  
 Answer: FA = 1111\_1010
- 1.9** (a)  $10110.0101_2 = 16 + 4 + 2 + .25 + .0625 = 22.3125$   
 (b)  $16.5_{16} = 16 + 6 + 5*(.0615) = 22.3125$   
 (c)  $26.24_8 = 2 * 8 + 6 + 2/8 + 4/64 = 22.3125$   
 (d)  $DADA.B_{16} = 14*16^3 + 10*16^2 + 14*16 + 10 + 11/16 = 60,138.6875$

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(e)  $1010.1101_2 = 8 + 2 + .5 + .25 + .0625 = 10.8125$

1.10 (a)  $1.10010_2 = 0001.1001_2 = 1.9_{16} = 1 + 9/16 = 1.563_{10}$

(b)  $110.010_2 = 0110.0100_2 = 6.4_{16} = 6 + 4/16 = 6.25_{10}$

Reason:  $110.010_2$  is the same as  $1.10010_2$  shifted to the left by two places.

1.11

$$\begin{array}{r} \underline{1011.11} \\ 101 \overline{) 111011.0000} \\ \underline{101} \phantom{.0000} \\ 01001 \phantom{.0000} \\ \underline{101} \phantom{.0000} \\ 1001 \phantom{.0000} \\ \underline{101} \phantom{.0000} \\ 1000 \phantom{.0000} \\ \underline{101} \phantom{.0000} \\ 0110 \phantom{.0000} \end{array}$$

The quotient is carried to two decimal places, giving  $1011.11$

Checking:  $111011_2 / 101_2 = 59_{10} / 5_{10} \approx 1011.11_2 = 58.75_{10}$

1.12 (a)  $10000$  and  $110111$

$$\begin{array}{r} 1011 \\ +101 \\ \hline 10000 = 16_{10} \end{array}$$

$$\begin{array}{r} 1011 \\ \times 101 \\ \hline 1011 \\ 1011 \\ \hline 110111 = 55_{10} \end{array}$$

(b)  $62_h$  and  $958_h$

$$\begin{array}{r} 2E_h \quad 0010\_1110 \\ +34_h \quad 0011\_0100 \\ \hline 62_h \quad 0110\_0010 = 98_{10} \end{array}$$

$$\begin{array}{r} 2E_h \\ \times 34_h \\ \hline B^38 \\ \hline 8^2A \\ \hline 9\ 5\ 8_h = 2392_{10} \end{array}$$

1.13 (a) Convert  $27.315$  to binary:

	Integer Quotient		Remainder	Coefficient
$27/2 =$	13	+	$\frac{1}{2}$	$a_0 = 1$
$13/2$	6	+	$\frac{1}{2}$	$a_1 = 1$
$6/2$	3	+	0	$a_2 = 0$
$3/2$	1	+	$\frac{1}{2}$	$a_3 = 1$
$\frac{1}{2}$	0	+	$\frac{1}{2}$	$a_4 = 1$

$$27_{10} = 11011_2$$

	Integer	Fraction	Coefficient
.315 x 2 =	0	+ .630	$a_1 = 0$
.630 x 2 =	1	+ .26	$a_2 = 1$
.26 x 2 =	0	+ .52	$a_3 = 0$
.52 x 2 =	1	+ .04	$a_4 = 1$

$$.315_{10} \approx .0101_2 = .25 + .0625 = .3125$$

$$27.315 \approx 11011.0101_2$$

$$(b) \ 2/3 \approx .666666667$$

	Integer	Fraction	Coefficient
.6666_6666_67 x 2 =	1	+ .3333_3333_34	$a_1 = 1$
.333333334 x 2 =	0	+ .6666666668	$a_2 = 0$
.6666666668 x 2 =	1	+ .3333333336	$a_3 = 1$
.3333333336 x 2 =	0	+ .6666666672	$a_4 = 0$
.6666666672 x 2 =	1	+ .3333333344	$a_5 = 1$
.3333333344 x 2 =	0	+ .6666666688	$a_6 = 0$
.6666666688 x 2 =	1	+ .3333333376	$a_7 = 1$
.3333333376 x 2 =	0	+ .6666666752	$a_8 = 0$

$$.666666667_{10} \approx .10101010_2 = .5 + .125 + .0313 + .0078 = .6641_{10}$$

$$.10101010_2 = .1010_{10} = .AA_{16} = 10/16 + 10/256 = .6641_{10} \text{ (Same as (b)).}$$

<b>1.14</b>	<b>(a)</b>	0001_0000 1s comp: 1110_1111 2s comp: 1111_0000	<b>(b)</b>	0000_0000 1s comp: 1111_1111 2s comp: 0000_0000	<b>(c)</b>	1101_1010 1s comp: 0010_0101 2s comp: 0010_0110
	<b>(d)</b>	1010_1010 1s comp: 0101_0101 2s comp: 0101_0110	<b>(e)</b>	1000_0101 1s comp: 0111_1010 2s comp: 0111_1011	<b>(f)</b>	1111_1111 1s comp: 0000_0000 2s comp: 0000_0001

<b>1.15</b>	<b>(a)</b>	25,478,036 9s comp: 74,521,963 10s comp: 74,521,964	<b>(b)</b>	63,325,600 9s comp: 36,674,399 10s comp: 36,674,400
	<b>(c)</b>	25,000,000 9s comp: 74,999,999 10s comp: 75,000,000	<b>(d)</b>	00000000 9s comp: 99999999 10s comp: 100000000

<b>1.16</b>	C3DF	C3DF: 1100_0011_1101_1111
	15s comp: 3C20	1s comp: 0011_1100_0010_0000
	16s comp: 3C21	2s comp: 0011_1100_0010_0001 = 3C21

**1.17** **(a)**  $2,579 \rightarrow 02,579 \rightarrow 97,420$  (9s comp)  $\rightarrow 97,421$  (10s comp)  
 $4637 - 2,579 = 2,579 + 97,421 = 2058_{10}$

**(b)**  $1800 \rightarrow 01800 \rightarrow 98199$  (9s comp)  $\rightarrow 98200$  (10 comp)  
 $125 - 1800 = 00125 + 98200 = 98325$  (negative)  
 Magnitude: 1675  
 Result:  $125 - 1800 = 1675$

(c)  $4,361 \rightarrow 04361 \rightarrow 95638$  (9s comp)  $\rightarrow 95639$  (10s comp)  
 $2043 - 4361 = 02043 + 95639 = 97682$  (Negative)  
 Magnitude: 2318  
 Result:  $2043 - 6152 = -2318$

(d)  $745 \rightarrow 00745 \rightarrow 99254$  (9s comp)  $\rightarrow 99255$  (10s comp)  
 $1631 - 745 = 01631 + 99255 = 0886$  (Positive)  
 Result:  $1631 - 745 = 886$

**1.18** Note: Consider sign extension with 2s complement arithmetic.

<p>(a) <math>0\_10010</math>            1s comp: <math>1\_01101</math>            2s comp: <math>1\_01110</math>  <math>0\_10011</math>            Diff: <math>0\_00001</math> (Positive)            Check: <math>19 - 18 = +1</math></p>	<p>(b) <math>0\_100110</math>            1s comp: <math>1\_011001</math> with sign extension            2s comp: <math>1\_011010</math>  <math>0\_100010</math>  <math>1\_111100</math> sign bit indicates that the result is negative  <math>0\_000011</math> 1s complement  <math>0\_000100</math> 2s complement  <math>000100</math> magnitude            Result: -4            Check: <math>34 - 38 = -4</math></p>
<p>(c) <math>0\_110101</math>            1s comp: <math>1\_001010</math>            2s comp: <math>1\_001011</math>  <math>0\_001001</math>            Diff: <math>1\_010100</math> (negative)  <math>0\_101011</math> (1s comp)  <math>0\_101100</math> (2s complement)  <math>101100</math> (magnitude)  <math>-44_{10}</math> (result)</p>	<p>(d) <math>0\_010101</math>            1s comp: <math>1\_101010</math> with sign extension            2s comp: <math>1\_101011</math>  <math>0\_101000</math>  <math>0\_010011</math> sign bit indicates that the result is positive            Result: <math>19_{10}</math>            Check: <math>40 - 21 = 19_{10}</math></p>

**1.19**  $+9286 \rightarrow 009286$ ;  $+801 \rightarrow 000801$ ;  $-9286 \rightarrow 990714$ ;  $-801 \rightarrow 999199$

(a)  $(+9286) + (801) = 009286 + 000801 = 010087$

(b)  $(+9286) + (-801) = 009286 + 999199 = 008485$

(c)  $(-9286) + (+801) = 990714 + 000801 = 991515$

(d)  $(-9286) + (-801) = 990714 + 999199 = 989913$

**1.20**  $+49 \rightarrow 0\_110001$  (Needs leading zero extension to indicate + value);  
 $+29 \rightarrow 0\_011101$  (Leading 0 indicates + value)  
 $-49 \rightarrow 1\_001110 + 0\_000001 \rightarrow 1\_001111$   
 $-29 \rightarrow 1\_100011$  (sign extension indicates negative value)

(a)  $(+29) + (-49) = 0\_011101 + 1\_001111 = 1\_101100$  (1 indicates negative value.)  
 Magnitude =  $0\_010011 + 0\_000001 = 0\_010100 = 20$ ; Result  $(+29) + (-49) = -20$

(b)  $(-29) + (+49) = 1\_100011 + 0\_110001 = 0\_010100$  (0 indicates positive value)  
 $(-29) + (+49) = +20$

- (c) Must increase word size by 1 (sign extension) to accomodate overflow of values:  
 $(-29) + (-49) = 11\_100011 + 11\_001111 = 10\_110010$  (1 indicates negative result)  
 Magnitude:  $01\_001110 = 78_{10}$   
 Result:  $(-29) + (-49) = -78_{10}$

**1.21**  $+9742 \rightarrow 009742 \rightarrow 990257$  (9's comp)  $\rightarrow 990258$  (10s) comp  
 $+641 \rightarrow 000641 \rightarrow 999358$  (9's comp)  $\rightarrow 999359$  (10s) comp

(a)  $(+9742) + (+641) \rightarrow 010383$

(b)  $(+9742) + (-641) \rightarrow 009742 + 999359 = 009102$   
 Result:  $(+9742) + (-641) = 9102$

(c)  $(-9742) + (+641) = 990258 + 000641 = 990899$  (negative)  
 Magnitude: 009101  
 Result:  $(-9742) + (+641) = -9101$

(d)  $(-9742) + (-641) = 990258 + 999359 = 989617$  (Negative)  
 Magnitude: 10383  
 Result:  $(-9742) + (-641) = -10383$

**1.22** 6,514  
 BCD: 0110\_0101\_0001\_0100  
 ASCII: 0\_011\_0110\_0\_011\_0101\_1\_011\_0001\_1\_011\_0100  
 ASCII: 0011\_0110\_0011\_0101\_1011\_0001\_1011\_0100

**1.23**

0111	1001	0001 ( 791)
0110	0101	1000 (+658)
1101	1110	1001
0110	0110	
0001 0011	0100	
0001 0001		
0001 0100	0100	1001 (1,449)

<b>1.24</b>	(a)	(b)
	6 3 1 1    Decimal	6 4 2 1    Decimal
	0 0 0 0    0	0 0 0 0    0
	0 0 0 1    1	0 0 0 1    1
	0 0 1 0    2	0 0 1 0    2
	0 1 0 0    3	0 0 1 1    3
	0 1 1 0    4 (or 0101)	0 1 0 0    4
	0 1 1 1    5	0 1 0 1    5
	1 0 0 0    6	1 0 0 0    6 (or 0110)
	1 0 1 0    7 (or 1001)	1 0 0 1    7
	1 0 1 1    8	1 0 1 0    8
	1 1 0 0    9	1 0 1 1    9

**1.25**

(a)	6,248 <sub>10</sub>	BCD:	0110_0010_0100_1000
(b)		Excess-3:	1001_0101_0111_1011
(c)	2421:		0110_0010_0100_1110
(d)	6311:		1000_0010_0110_1011

1.26      6,248   9s Comp:   3,751  
             2421 code:   0011\_0111\_0101\_0001  
             1s comp c:   1001\_1101\_1011\_0001 (2421 code alternative #1)  
  
             6,248<sub>2421</sub>   0110\_0010\_0100\_1110 (2421 code alternative #2)  
             1s comp c   1001\_1101\_1011\_0001 Match

**1.27** For a deck with 52 cards, we need 6 bits ( $2^5 = 32 < 52 < 64 = 2^6$ ). Let the msb's select the suit (e.g., diamonds, hearts, clubs, spades are encoded respectively as 00, 01, 10, and 11. The remaining four bits select the "number" of the card. Example: 0001 (ace) through 1011 (9), plus 101 through 1100 (jack, queen, king). This a jack of spades might be coded as 11\_1010. (Note: only 52 out of 64 patterns are used.)

**1.28** G (dot) (space) B o o l e  
11000111\_11101111\_01101000\_01101110\_00100000\_11000100\_11101111\_11100101

**1.29** Steve Jobs

**1.30** 73 F4 E5 76 E5 4A EF 62 73

73: 0\_111\_0011 s  
F4: 1\_111\_0100 t  
E5: 1\_110\_0101 e  
76: 0\_111\_0110 v  
E5: 1\_110\_0101 e  
4A: 0\_100\_1010 j  
EF: 1\_110\_1111 o  
62: 0\_110\_0010 b  
73: 0\_111\_0011 s

**1.31**  $62 + 32 = 94$  printing characters

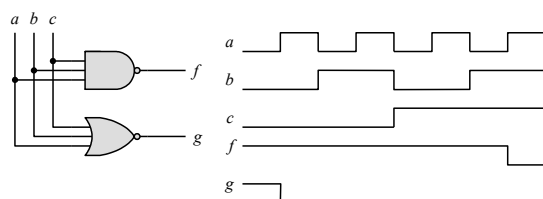
**1.32** bit 6 from the right

**1.33** (a) 897 (b) 564 (c) 871 (d) 2,199

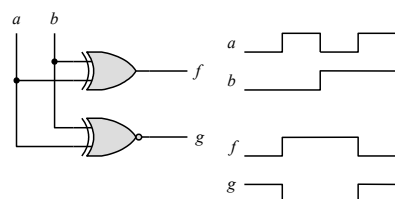
**1.34** ASCII for decimal digits with even parity:

(0): 00110000 (1): 10110001 (2): 10110010 (3): 00110011  
(4): 10110100 (5): 00110101 (6): 00110110 (7): 10110111  
(8): 10111000 (9): 00111001

**1.35** (a)



**1.36**





**CHAPTER 2****2.1 (a)**

$x y z$	$x + y + z$	$(x + y + z)'$	$x'$	$y'$	$z'$	$x' y' z'$	$x y z$	$(xyz)$	$(xyz)'$	$x'$	$y'$	$z'$	$x' + y' + z'$
0 0 0	0	1	1	1	1	1	0 0 0	0	1	1	1	1	1
0 0 1	1	0	1	1	0	0	0 0 1	0	1	1	1	0	1
0 1 0	1	0	1	0	1	0	0 1 0	0	1	1	0	1	1
0 1 1	1	0	1	0	0	0	0 1 1	0	1	1	0	0	1
1 0 0	1	0	0	1	1	0	1 0 0	0	1	0	1	1	1
1 0 1	1	0	0	1	0	0	1 0 1	0	1	0	1	0	1
1 1 0	1	0	0	0	1	0	1 1 0	0	1	0	0	1	1
1 1 1	1	0	0	0	0	0	1 1 1	1	0	0	0	0	0

**(b)**

$x y z$	$x + yz$	$(x + y)$	$(x + z)$	$(x + y)(x + z)$
0 0 0	0	0	0	0
0 0 1	0	0	1	0
0 1 0	0	1	0	0
0 1 1	1	1	1	1
1 0 0	1	1	1	1
1 0 1	1	1	1	1
1 1 0	1	1	1	1
1 1 1	1	1	1	1

**(c)**

$x y z$	$x(y + z)$	$xy$	$xz$	$xy + xz$
0 0 0	0	0	0	0
0 0 1	0	0	0	0
0 1 0	0	0	0	0
0 1 1	0	0	0	0
1 0 0	0	0	0	0
1 0 1	1	0	1	1
1 1 0	1	1	0	1
1 1 1	1	1	1	1

**(e)**

$x y z$	$x$	$y + z$	$x + (y + z)$	$(x + y)$	$(x + y) + z$
0 0 0	0	0	0	0	0
0 0 1	0	1	1	0	1
0 1 0	0	1	1	1	1
0 1 1	0	1	1	1	1
1 0 0	1	0	1	1	1
1 0 1	1	1	1	1	1
1 1 0	1	1	1	1	1
1 1 1	1	1	1	1	1

**(d)**

$x y z$	$yz$	$x(yz)$	$xy$	$(xy)z$
0 0 0	0	0	0	0
0 0 1	0	0	0	0
0 1 0	0	0	0	0
0 1 1	1	0	0	0
1 0 0	0	0	0	0
1 0 1	0	0	0	0
1 1 0	0	0	1	0
1 1 1	1	1	1	1

**2.2 (a)**  $xy + xy' = x(y + y') = x$

**(b)**  $(x + y)(x + y') = x + yy' = x(x + y') + y(x + y') = xx + xy' + xy + yy' = x$

**(c)**  $xyz + x'y + xyz' = xy(z + z') + x'y = xy + x'y = y$

**(d)**  $(A + B)(A' + B')' = (A'B')(AB) = (A'B')(BA) = A'(B'B)A = 0$

**(e)**  $(a + b + c')(a'b' + c) = aa'b' + ac + ba'b' + bc + c'a'b' + c'c = ac + bc + a'b'c'$

**(f)**  $a'bc + abc' + abc + a'bc' = a'b(c + c') + ab(c + c') = a'b + ab = (a' + a)b = b$

**2.3 (a)**  $ABC + A'B + ABC' = AB + A'B = B$

(b)  $x'yz + xz = (x'y + x)z = z(x + x')(x + y) = z(x + y)$

(c)  $(x + y)'(x' + y') = x'y'(x' + y') = x'y'$

(d)  $xy + x(wz + wz') = x(y + wz + wz') = x(w + y)$

(e)  $(BC' + A'D)(AB' + CD') = BC'AB' + BC'CD' + A'DAB' + A'DCD' = 0$

(f)  $(a' + c')(a + b' + c') = a'a + a'b' + a'c' + c'a + c'b' + c'c' = a'b' + a'c' + ac' + b'c' = c' + b'(a' + c')$   
 $= c' + b'c' + a'b' = c' + a'b'$

2.4 (a)  $A'C' + ABC + AC' = C' + ABC = (C + C')(C' + AB) = AB + C'$

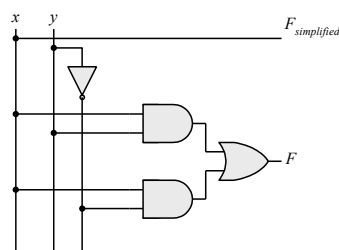
(b)  $(x'y' + z)' + z + xy + wz = (x'y')'z' + z + xy + wz = [(x + y)z' + z] + xy + wz =$   
 $= (z + z')(z + x + y) + xy + wz = z + wz + x + xy + y = z(1 + w) + x(1 + y) + y = x + y + z$

(c)  $A'B(D' + C'D) + B(A + A'CD) = B(A'D' + A'C'D + A + A'CD)$   
 $= B(A'D' + A + A'D(C + C')) = B(A + A'(D' + D)) = B(A + A') = B$

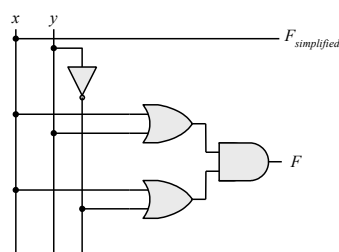
(d)  $(A' + C)(A' + C')(A + B + C'D) = (A' + CC')(A + B + C'D) = A'(A + B + C'D)$   
 $= AA' + A'B + A'C'D = A'(B + C'D)$

(e)  $ABC'D + A'BD + ABCD = AB(C + C')D + A'BD = ABD + A'BD = BD$

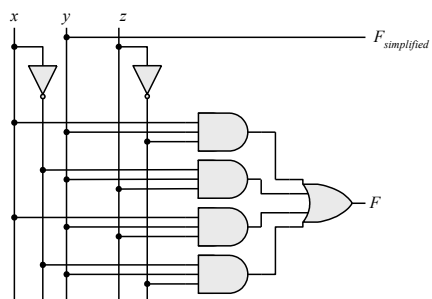
2.5 (a)



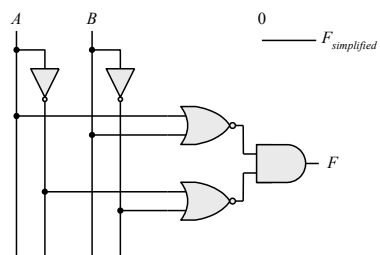
(b)



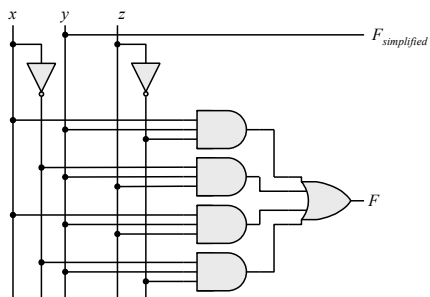
(c)



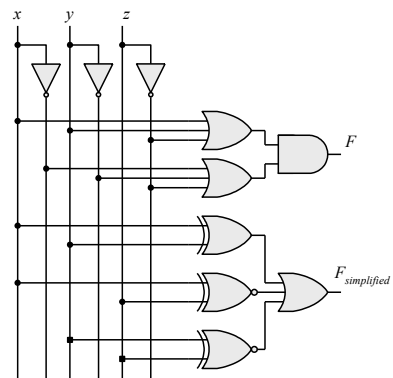
(d)



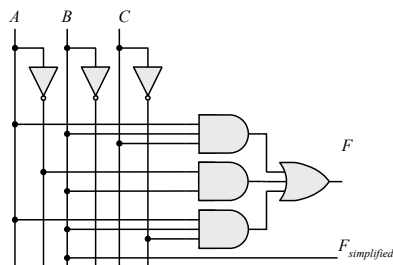
(e)



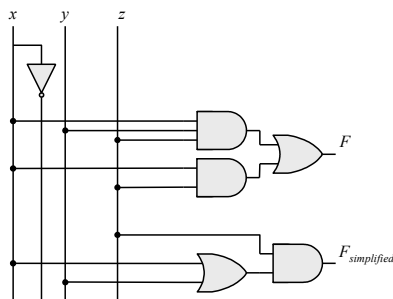
(f)



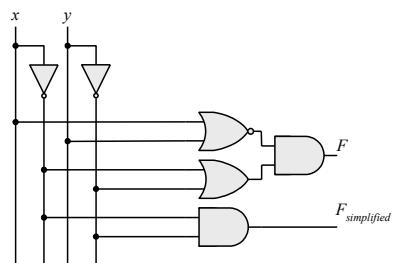
2.6 (a)



(b)

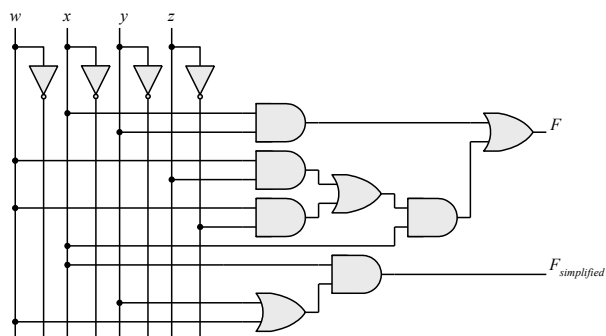


(c)

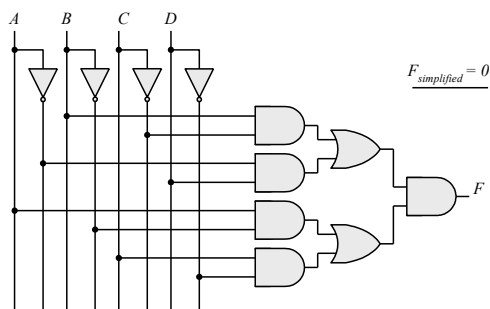


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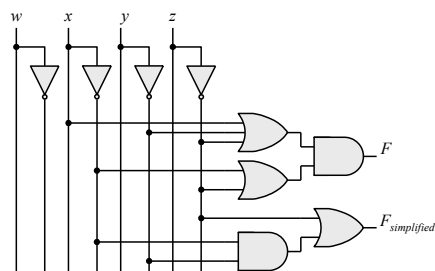
(d)



(e)

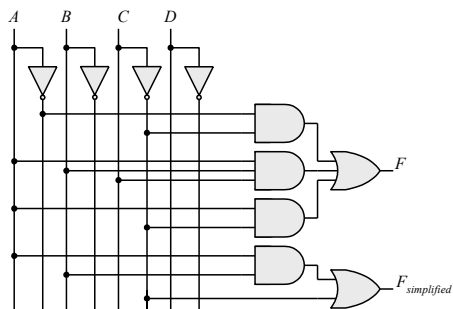


(f)



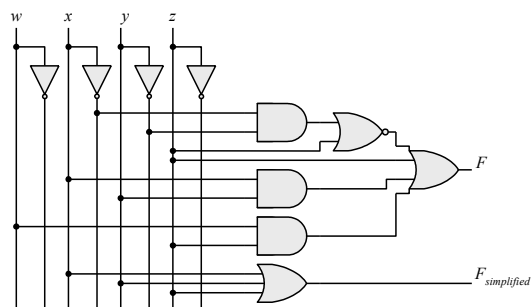
2.7

(a)

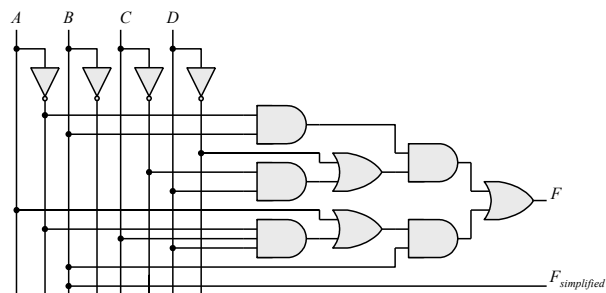


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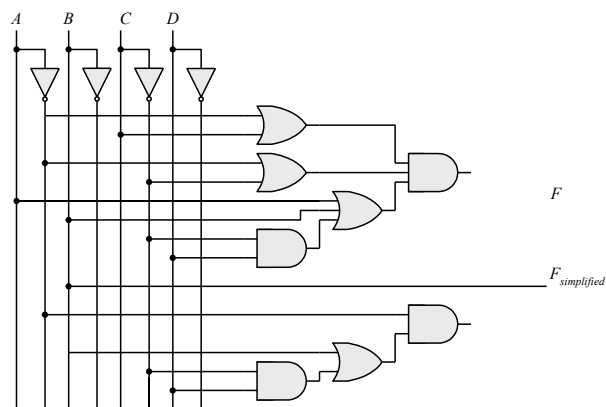
(b)



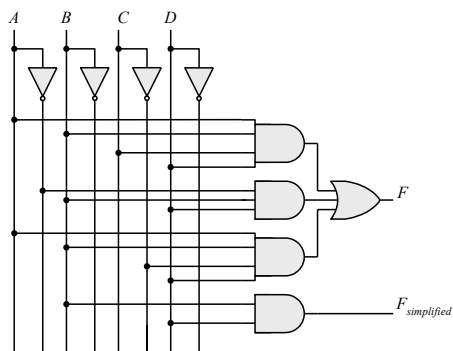
(c)



(d)



(e)



2.8  $F' = (wx + yz)' = (wx)'(yz)' = (w' + x')(y' + z')$

$$FF' = wx(w' + x')(y' + z') + yz(w' + x')(y' + z') = 0$$

$$F + F' = wx + yz + (wx + yz)' = A + A' = 1 \text{ with } A = wx + yz$$

2.9 (a)  $F' = (xy' + x'y)' = (xy')'(x'y)' = (x' + y)(x + y') = xy + x'y'$

(b)  $F' = [(a + c)(a + b')(a' + b + c')] = (a + c)' + (a + b')' + (a' + b + c')'$   
 $= a'c' + a'b + ab'c$

(c)  $F' = [z + z'(v'w + xy)]' = z'[z'(v'w + xy)]' = z'[z'v'w + xyz]'$   
 $= z'[(z'v'w)'(xyz)'] = z'[(z + v + w)'(x' + y' + z)]$   
 $= z'z + z'v + z'w' + z'x' + z'y' + z'z = z'(v + w' + x' + y')$

2.10 (a)  $F_1 + F_2 = \sum m_{1i} + \sum m_{2i} = \sum (m_{1i} + m_{2i})$

(b)  $F_1 F_2 = \sum m_i \sum m_j$  where  $m_i m_j = 0$  if  $i \neq j$  and  $m_i m_j = 1$  if  $i = j$

2.11 (a)  $F(x, y, z) = \sum(1, 4, 5, 6, 7)$

(b)  $F(a, b, c) = \sum(0, 2, 3, 7)$

$$F = xy + xy' + y'z$$

$$F = bc + a'c'$$

x y z	F	a b c	F
0 0 0	0	0 0 0	1
0 0 1	1	0 0 1	0
0 1 0	0	0 1 0	1
0 1 1	0	0 1 1	1
1 0 0	1	1 0 0	0
1 0 1	1	1 0 1	0
1 1 0	1	1 1 0	0
1 1 1	1	1 1 1	1

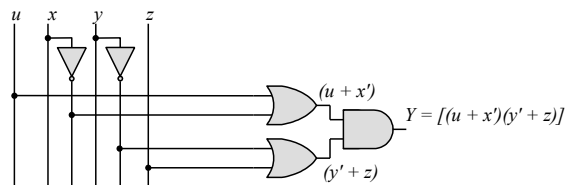
2.12  $A = 1011\_0001$

$$B = 1010\_1100$$

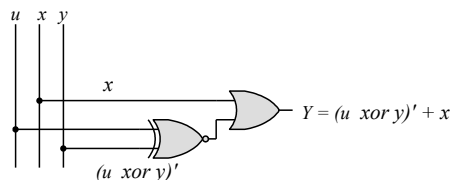
- (a)  $A \text{ AND } B = 1010\_0000$   
 (b)  $A \text{ OR } B = 1011\_1101$   
 (c)  $A \text{ XOR } B = 0001\_1101$   
 (d)  $\text{NOT } A = 0100\_1110$   
 (e)  $\text{NOT } B = 0101\_0011$

2.13

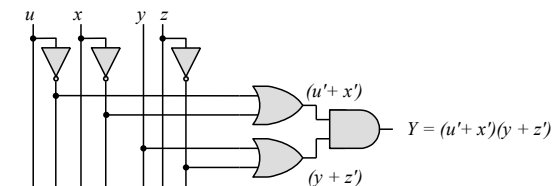
(a)



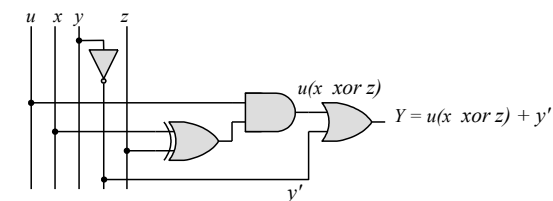
(b)



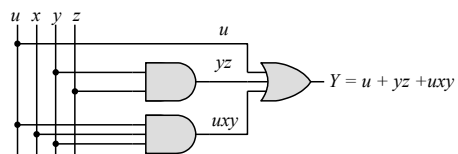
(c)



(d)

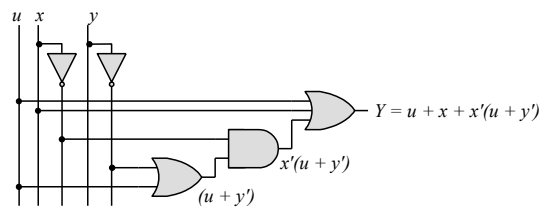


(e)

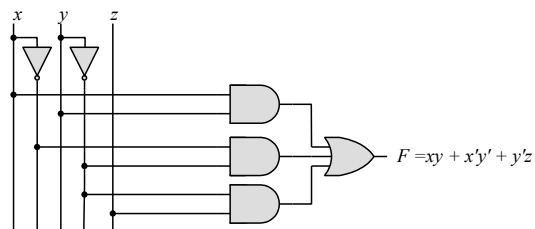


(f)

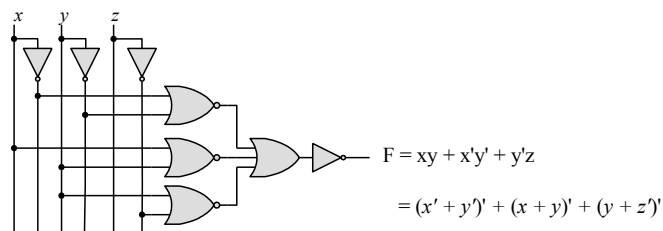




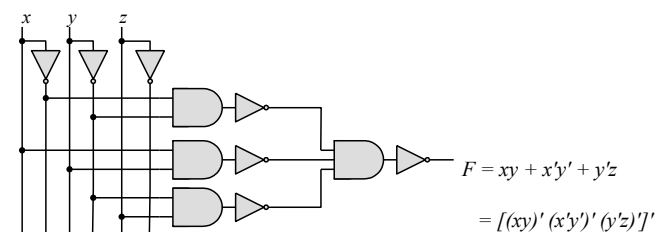
2.14 (a)



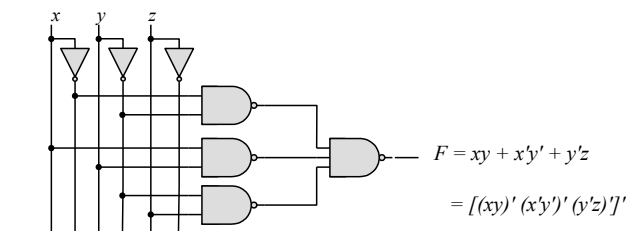
(b)



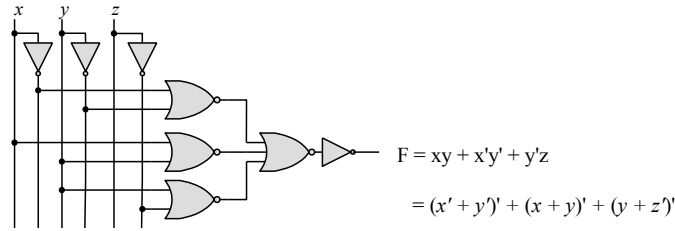
(c)



(d)



(e)



2.15 (a)  $T_1 = A'B'C' + A'B'C + A'BC' = A'B'(C' + C) + A'C'(B' + B) = A'B' + A'C' = A'(B' + C')$

(b)  $T_2 = T_1' = A'BC + AB'C' + AB'C + ABC' + ABC$   
 $= BC(A' + A) + AB'(C' + C) + AB(C' + C)$   
 $= BC + AB' + AB = BC + A(B' + B) = A + BC$

$$\Sigma(3, 5, 6, 7) = \Pi(0, 1, 2, 4)$$

$$T_1 = A'B'C' + A'B'C + A'BC'$$

$$\begin{array}{c} \swarrow \quad \searrow \\ A'B' \quad A'C' \end{array}$$

$$T_1 = A'B' A'C' = A'(B' + C')$$

$$T_2 = A'BC + AB'C' + AB'C + ABC' + ABC$$

$$\begin{array}{c} \swarrow \quad \searrow \quad \swarrow \quad \searrow \\ AC' \quad AC \end{array}$$

$$BC$$

$$T_2 = AC' + BC + AC = A + BC$$

2.16 (a)  $F(A, B, C) = A'B'C' + A'B'C + A'BC' + A'BC + AB'C' + AB'C + ABC' + ABC$   
 $= A'(B'C' + B'C + BC' + BC) + A((B'C' + B'C + BC' + BC)$   
 $= (A' + A)(B'C' + B'C + BC' + BC) = B'C' + B'C + BC' + BC$   
 $= B'(C' + C) + B(C' + C) = B' + B = 1$

(b)  $F(x_1, x_2, x_3, \dots, x_n) = \Sigma m_i$  has  $2^{n-1}/2$  minterms with  $x_1$  and  $2^{n-1}/2$  minterms with  $x_1'$ , which can be factored and removed as in (a). The remaining  $2^{n-1}$  product terms will have  $2^{n-1}/2$  minterms with  $x_2$  and  $2^{n-1}/2$  minterms with  $x_2'$ , which can be factored to remove  $x_2$  and  $x_2'$ . Continue this process until the last term is left and  $x_n + x_n' = 1$ . Alternatively, by induction,  $F$  can be written as  $F = x_n G + x_n' G$  with  $G = 1$ . So  $F = (x_n + x_n')G = 1$ .

- 2.17 (a)  $F = (b + cd)(c + bd)$   $bc + bd + cd + bcd = \Sigma(3, 5, 6, 7, 11, 14, 15)$   
 $F' = \Sigma(0, 1, 2, 4, 8, 9, 10, 12, 13)$   
 $F = \Pi(0, 1, 2, 4, 8, 9, 10, 12, 13)$

a	b	c	d	F
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	0
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

- (b)  $(cd + b'c + bd')(b + d) = bcd + bd' + cd + b'cd = cd + bd'$   
 $= \Sigma(3, 4, 7, 11, 12, 14, 15)$   
 $= \Pi(0, 1, 2, 5, 6, 8, 9, 10, 13)$

a	b	c	d	F
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	1
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	1

- (c)  $(c' + d)(b + c') = bc' + c' + bd + c'd = (c' + bd)$   
 $= \Sigma(0, 1, 4, 5, 7, 8, 12, 13, 15)$   
 $F = \Pi(2, 3, 6, 9, 10, 11, 14)$

(d)  $bd' + acd' + ab'c + a'c' = \Sigma(0, 1, 4, 5, 10, 11, 14)$

$$F' = \Sigma(2, 3, 6, 7, 8, 9, 12, 13, 15)$$

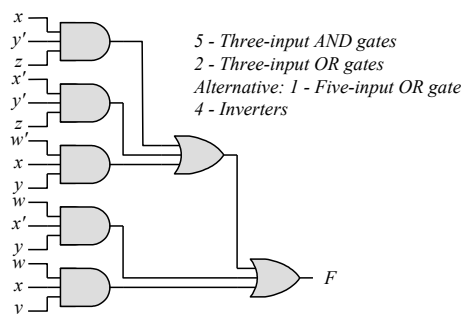
$$F = \Pi(02, 3, 6, 7, 8, 12, 13, 15)$$

a	b	c	d	F
0	0	0	0	1
0	0	0	1	1
0	0	1	0	0
0	0	1	1	0
0	1	0	0	1
0	1	0	1	1
0	1	1	0	0
0	1	1	1	0
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	0

2.18 (a)

wx y z	F	$F = xy'z + x'y'z + w'xy + wx'y + wxy$ $F = \Sigma(1, 5, 6, 7, 9, 10, 11, 13, 14, 15)$
00 0 0	0	
00 0 1	1	
00 1 0	0	
00 1 1	0	
01 0 0	0	
01 0 1	1	
01 1 0	1	
01 1 1	1	
10 0 0	0	
10 0 1	1	
10 1 0	1	
10 1 1	1	
11 0 0	0	
11 0 1	1	
11 1 0	1	
11 1 1	1	

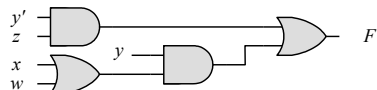
(b)



(c)  $F = xy'z + x'y'z + w'xy + wx'y + wxy = y'z + xy + wy = y'z + y(w + x)$

(d)  $F = y'z + yw + yx = \Sigma(1, 5, 9, 13, 10, 11, 13, 15, 6, 7, 14, 15)$   
 $= \Sigma(1, 5, 6, 7, 9, 10, 11, 13, 14, 15)$

(e)



1 - Inverter, 2 - Two-input AND gates, 2 - Two-input OR gates

2.19  $F = B'D + A'D + BD$

$ABCD$	$ABCD$	$ABCD$
$\neg B'D$	$A'\neg D$	$\neg B'D$
0001 = 1	0001 = 1	0101 = 5
0011 = 3	0011 = 3	0111 = 7
1001 = 9	0101 = 5	1101 = 13
1011 = 11	0111 = 7	1111 = 15

$$F = \Sigma(1, 3, 5, 7, 9, 11, 13, 15) = \Pi(0, 2, 4, 6, 8, 10, 12, 14)$$

2.20 (a)  $F(A, B, C, D) = \Sigma(2, 4, 7, 10, 12, 14)$   
 $F'(A, B, C, D) = \Sigma(0, 1, 3, 5, 6, 8, 9, 11, 13, 15)$

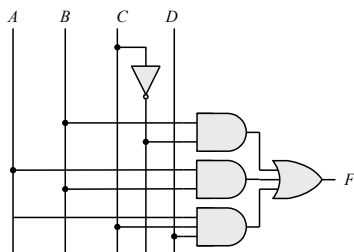
(b)  $F(x, y, z) = \Pi(3, 5, 7)$   
 $F' = \Sigma(3, 5, 7)$

2.21 (a)  $F(x, y, z) = \Sigma(1, 3, 5) = \Pi(0, 2, 4, 6, 7)$   
 (b)  $F(A, B, C, D) = \Pi(3, 5, 8, 11) = \Sigma(0, 1, 2, 4, 6, 7, 9, 10, 12, 13, 14, 15)$

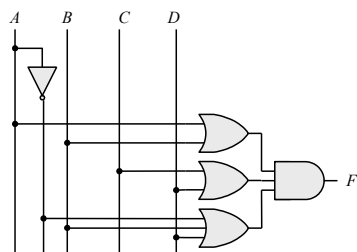
2.22 (a)  $(u + xw)(x + u'v) = ux + uu'v + xw + xwu'v = ux + xw + xwu'v$   
 $= ux + xw = x(u + w)$   
 $= ux + xw$  (SOP form)  
 $= x(u + w)$  (POS form)

(b)  $x' + x(x + y')(y + z') = x' + x(xy + xz' + y'y + y'z')$   
 $= x' + xy + xz' + xy'z' = x' + xy + xz'$  (SOP form)  
 $= (x' + y + z')$  (POS form)

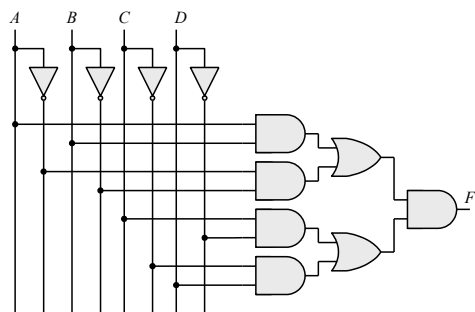
2.23 (a)  $B'C + AB + ACD$



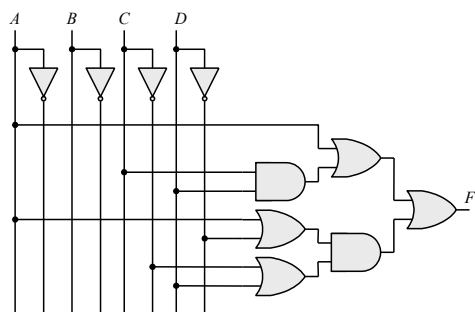
(b)  $(A + B)(C + D)(A' + B + D)$



(c)  $(AB + A'B')(CD' + C'D)$



(d)  $A + CD + (A + D')(C' + D)$



2.24  $x \oplus y = x'y + xy'$  and  $(x \oplus y)' = (x + y')(x' + y)$

Dual of  $x'y + xy' = (x' + y)(x + y') = (x \oplus y)'$

2.25 (a)  $x|y = xy' \neq y|x = x'y$  Not commutative  
 $(x|y)|z = xy'z' \neq x|(y|z) = x(yz')' = xy' + xz$  Not associative

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(b)  $(x \oplus y) = xy' + x'y = y \oplus x = yx' + y'x$  Commutative

$(x \oplus y) \oplus z = \sum(1, 2, 4, 7) = x \oplus (y \oplus z)$  Associative

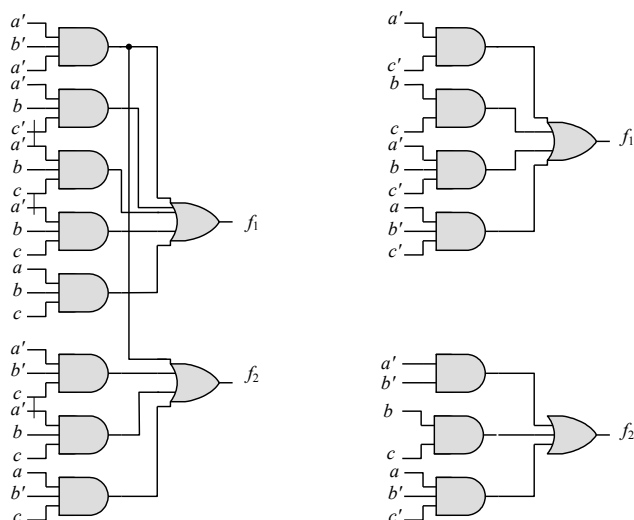
2.26

Gate		NAND (Positive logic)		NOR (Negative logic)	
x	y	x	y	x	y
L	L	H	0 0	0	1 1
L	H	H	0 1	0	1 0
H	L	H	1 0	0	0 1
H	H	L	1 1	1	0 0

Gate		NOR (Positive logic)		NAND (Negative logic)	
x	y	x	y	x	y
L	L	H	0 0	0	1 1
L	H	L	0 1	1	1 0
H	L	L	1 0	1	0 1
H	H	L	1 1	1	0 0

2.27  $f_1 = a'b'c' + a'bc' + a'bc + ab'c' + abc = a'c' + bc + a'bc' + ab'c'$

$f_2 = a'b'c' + a'b'c + a'bc + ab'c' + abc = a'b' + bc + ab'c'$





2.28 (a)  $y = a(bcd)'e = a(b' + c' + d')e$

$$y = a(b' + c' + d')e = ab'e + ac'e + ad'e$$

$$= \Sigma(17, 19, 21, 23, 25, 27, 29)$$

a bcde	y	a bcde	y
0 0000	0	1 0000	0
0 0001	0	1 0001	1
0 0010	0	1 0010	0
0 0011	0	1 0011	1
0 0100	0	1 0100	0
0 0101	0	1 0101	1
0 0110	0	1 0110	0
0 0111	0	1 0111	1
0 1000	0	1 1000	0
0 1001	0	1 1001	1
0 1010	0	1 1010	0
0 1011	0	1 1011	1
0 1100	0	1 1100	0
0 1101	0	1 1101	1
0 1110	0	1 1110	0
0 1111	0	1 1111	0

(b)  $y_1 = a \oplus (c + d + e) = a'(c + d + e) + a(c'd'e') = a'c + a'd + a'e + ac'd'e'$

$$y_2 = b'(c + d + e)f = b'cf + b'df + b'ef$$

$$y_1 = a(c + d + e) = a'(c + d + e) + a(c'd'e') = a'c + a'd + a'e + ac'd'e'$$

$$y_2 = b'(c + d + e)f = b'cf + b'df + b'ef$$

$a'-c---$	$a'--d--$	$a'---e-$	$a-c'd'e'-$
001000 = 8	000100 = 8	000010 = 2	100000 = 32
001001 = 9	000101 = 9	000011 = 3	100001 = 33
001010 = 10	000110 = 10	000110 = 6	110000 = 34
001011 = 11	000111 = 11	000111 = 7	110001 = 35

001100 = 12	001100 = 12	001010 = 10
001101 = 13	001101 = 13	001011 = 11
001110 = 14	001110 = 14	001110 = 14
001111 = 15	001111 = 15	001111 = 15

011000 = 24	010100 = 20	010010 = 18	001001 = 9	001001 = 9	000011 = 3
011001 = 25	010101 = 21	010011 = 19	001011 = 11	001011 = 11	000111 = 7
011010 = 26	010110 = 22	010110 = 22	001101 = 13	001101 = 13	001011 = 11
011011 = 27	010111 = 23	010111 = 23	001111 = 15	001111 = 15	001111 = 15
			101001 = 41	101001 = 41	100011 = 35
			101011 = 43	101011 = 43	100111 = 39
			101101 = 45	101101 = 45	101011 = 51
			101111 = 47	101111 = 47	101111 = 55

$$y_1 = \Sigma (2, 3, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 18, 19, 22, 23, 24, 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35)$$

$$y_2 = \Sigma (3, 7, 9, 13, 15, 35, 39, 41, 43, 45, 47, 51, 55)$$

<i>ab cdef</i>	<i>y<sub>1</sub> y<sub>2</sub></i>	<i>ab cdef</i>	<i>y<sub>1</sub> y<sub>2</sub></i>	<i>ab cdef</i>	<i>y<sub>1</sub> y<sub>2</sub></i>	<i>ab cdef</i>	<i>y<sub>1</sub> y<sub>2</sub></i>
00 0000	0 0	01 0000	0 0	10 0000	1 0	11 0000	0 0
00 0001	0 0	01 0001	0 0	10 0001	1 0	11 0001	0 0
00 0010	1 0	01 0010	1 0	10 0010	1 0	11 0010	0 0
00 0011	1 1	01 0011	1 0	10 0011	1 1	11 0011	0 1
00 0100	0 0	01 0100	0 0	10 0100	0 0	11 0100	0 0
00 0101	0 0	01 0101	0 0	10 0101	0 0	11 0101	0 0
00 0110	1 0	01 0110	1 0	10 0110	0 0	11 0110	0 0
00 0111	1 1	01 0111	1 0	10 0111	0 1	11 0111	0 1
00 1000	1 0	01 1000	1 0	10 1000	0 0	11 1000	0 0
00 1001	1 1	01 1001	1 0	10 1001	0 1	11 1001	0 0
00 1010	1 0	01 1010	1 0	10 1010	0 0	11 1010	0 0
00 1011	1 0	01 1011	1 0	10 1011	0 1	11 1011	0 0
00 1100	1 0	01 1100	1 0	10 1100	0 0	11 1100	0 0
00 1101	1 1	01 1101	1 0	10 1101	0 1	11 1101	0 0
00 1110	1 0	01 1110	1 0	10 1110	0 0	11 1110	0 0
00 1111	1 1	01 1111	1 0	10 1111	0 1	11 1111	0 0