

SLIDING WINDOW

..... *Strategy*



What is subarray : It is contiguous part of array

1 2 3 4

(1 2) (1 2 3) (1) (2) (3) (4)

(2, 3) (2 3 4)

(3, 4)

What is subsequence

It is a seq. of array/string which can be obtained by deleting zero or more elements but order should same.

1	2	3	4
---	---	---	---

1

1 2

1 2 3

1 2 3 4

2

2 3

2 3 4

3

3 4

4

Subarray



1 2 3

1 4

1 2 4

1 3 4

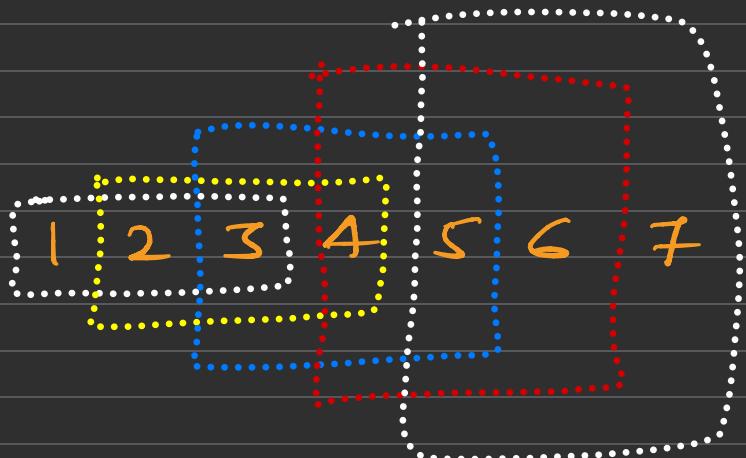
3 4

Sub-seq

Types of Sliding Window (Not an algorithm) Array / String

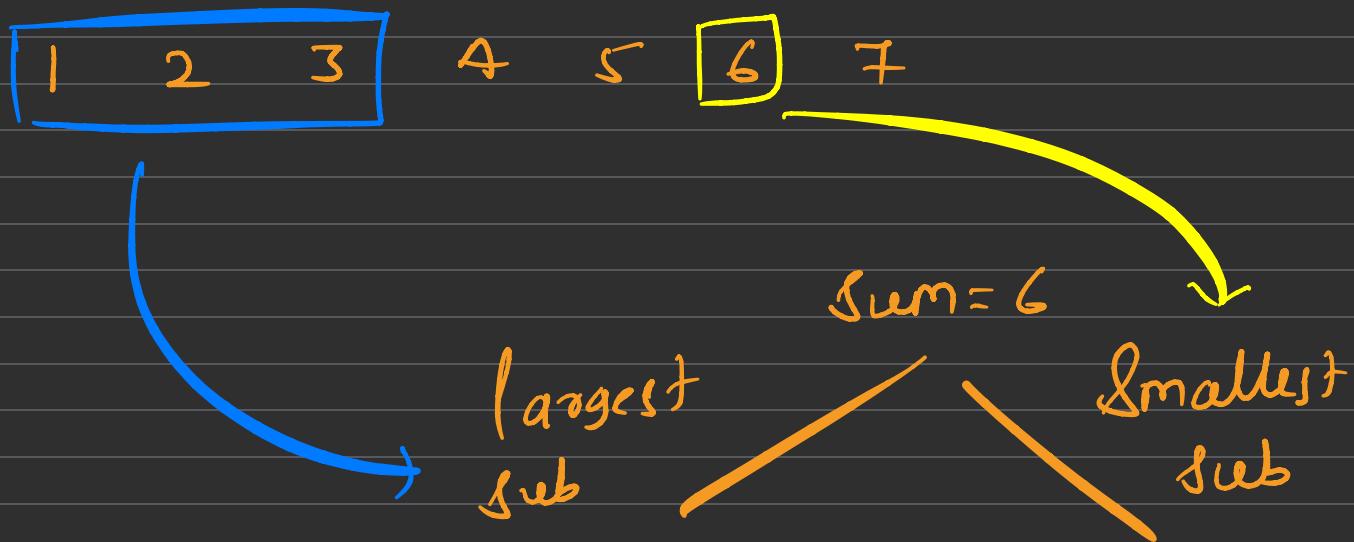
✓ Fixed Window Size

Variable Window Size



$K = 3$ (Window Size)

* Variable Window Size



How to approach the problem

1. Generate all the subarrays/substrings and pick the optimal one
2. Then optimise by moving the window slowly until certain conditions.
3. For variable window size, use the two pointer approach.

Given an array of integers, find all the subarrays

1	8	9	3	6	7	3	4	6
0	1	2	3	4	5	6	7	8

```
for ( i=0 : i<n ; i++ )  
{  
    for ( j=i : j<n ; j++ )  
    {  
        for ( k=i : k≤j : k++ )  
        {  
            cout << a(k) )  
        }  
        cout << endl  
    }  
}
```

Given a string, generate all the substrings.

String = abcdfh

Sub - strings

a

abc

ab c d

abc df

abc dfh

b

b c

b c d

b c d f

b c d f h

c

c d

c d f

c d f h

d

df

dfh

f

f h

h

1	2	3	4
---	---	---	---

1
1 2 3 4

2
2 3 4

3
3 4

4

```
for ( i=0 : i<n ; i++ )
{
    for ( j=i : j<n ; j++ )
    {
        for ( k=j : k<=j ; k++ )
        {
            cout << a(k)
        }
        cout << endl
    }
}
```

1	2	3	4
0	1	2	3

N=4

.

```
for ( i=0 : i<n ; i++ )
{
    for ( j=i : j<n ; j++ )
}
```

1 2 3 4

3 4

4

.

- given an array, print sum of all subarray
- given an array, find min from all subarray
(print min. of each sub-array)
- $= // =$ (max)

Max Sum Subarray of size K



Difficulty: Easy

Accuracy: 49.6%

Submissions: 194K+

Points: 2

Given an array of integers **arr[]** and a number **k**. Return the **maximum sum** of a subarray of size **k**.

Note: A subarray is a contiguous part of any given array.

Examples:

Input: arr[] = [100, 200, 300, 400] , k = 2

Output: 700

Explanation: $\text{arr}_3 + \text{arr}_4 = 700$, which is maximum.

Input: arr[] = [100, 200, 300, 400] , k = 4

Output: 1000

Explanation: $\text{arr}_1 + \text{arr}_2 + \text{arr}_3 + \text{arr}_4 = 1000$, which is maximum.

Input: arr[] = [100, 200, 300, 400] , k = 1

Output: 400

Explanation: $\text{arr}_4 = 400$, which is maximum.

fixed window

$N=8$

7	2	1	1	4	3	9	1
0	1	2	3	4	5	6	7

$K=3$

$\text{maxSum} = 0$

$O(N-K)$



```
for (i=0 ; i <= n-K ; i++)  
    }  
    sum = 0
```

$O(K)$



```
for (j=i ; j < i+K ; j++)  
    }  
    sum = sum + a(j)
```

$O(N-K) \times O(K)$

$$\boxed{O(N \times K)} - K^2 =$$

\uparrow if ($\text{sum} > \text{maxSum}$)

\uparrow } $\text{maxSum} = \text{sum}$

return maxSum

7	2	1	1	4	3	9	1
0	1	2	3	4	5	6	7

$$N = 8 \quad K = 3$$

| 2 3 | 4 5 6 $K=3$

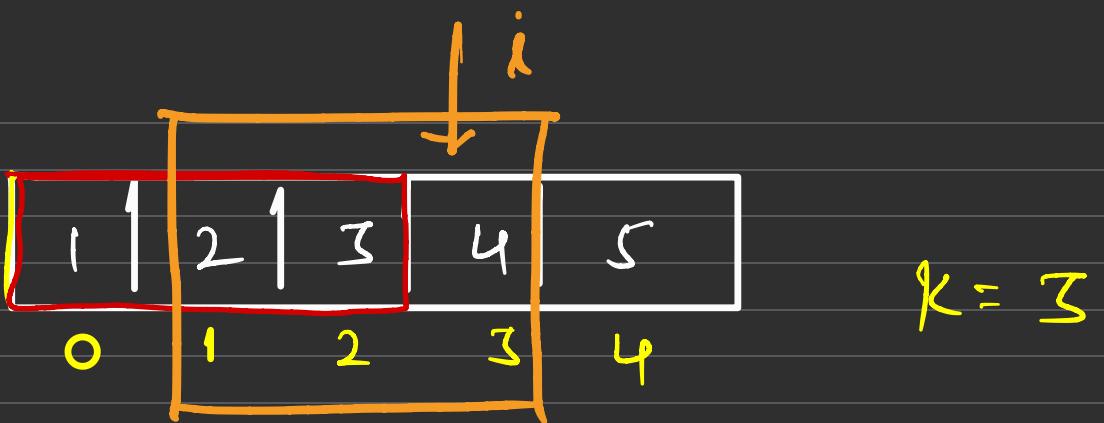
$$1 + \underline{2 + 3} = 6$$

$$\underline{2 + 3} + \underline{4} = 9$$

$$\underline{3 + 4} + 5 = 12$$

$$\underline{4 + 5} + 6 = 15$$

$$\text{Sum} = 9 + 5 - 2 = \underline{\underline{12}}$$



$$k = 3$$

$$i = 3 - k$$

$$= 3 - 3 = 0$$

$$\text{sum} = 6$$

$$\text{sum} = \text{sum} - a(i-k) + a(i)$$

- Steps for fixed Window

- ① first perform operation on first k elements. ($0, 1, \dots, k-1^{\text{th}}$ index) ($i < k$)
- ② start from k ($i = k$) & slide forward
but remove out of windows element
 $\{ q[i-k] \}$ & add current element
in window

int maxSubarraySum (arr[], n, K)
} sum = 0

for (i=0 ; i < K ; i++) — Step 1
 sum += arr[i]

maxSum = sum

for (j=K ; j < n ; j++) — Step 2
} sum = sum - arr[j-K] — Removed
 out of window.

sum = sum + arr[i]

maxSum = max (sum, maxSum) — Add current element

}

Given an array **arr[]** and an integer **K**, the task is to calculate the sum of all subarrays of size **K**.

Given an array of integer, task is to find the subarray with minimum sum of size k

given an array, find min No. from every sub.
of size K.

7	1	1	1	2	3	7
0	1	2	3	4	5	6

K=2

for (i=0 ; i < N-K ; i++)

 num = a(i)

 for (j=i ; j < i+K ; j++)

 if (a(j) < num)
 num = a(j)

print \n / num

7	1	1	1	.	2	3	7
0	1	2	3	4	5	6	

$k = 3$

map

$\{7, 1\}$
$\{3, 1\}$
$\{2, 1\}$

1.4 perform ppn first k

② $i = k$

$i = 3486$

map
 $\{ \text{key, value} \}$



```
void printMin (arr[], N, K)
```

```
{ map<int, int> mp
```

```
    for (i=0; i < K; i++)
```

```
        mp(a(i))++
```

— Step ①

begin of
(min) {
key first
value second

```
cout << (*mp.begin()).first
```

```
for (i = K; i < n; i++)
```

```
} int out = a(i-K)
```

```
mp(out)--
```

```
if (mp(out) == 0)
```

```
    mp.erase(out)
```

a(i-K)

x = a(i-K)

int cur = a(i)

mp[{cur}]++

cout << (*mp.begin()).first

}

}