

Department of Computer Science and Engineering
I Sem. M.Tech(Computer Science and Information Security)
Advanced Cryptography CSE 5171
Formula Sheet

RSA Digital Signature Scheme

$$S = M^d \bmod n$$

$$\text{Verification } M' = S^e \bmod n$$

Is M congruent to M' ?

ElGamal Digital Signature Scheme

$$\text{first signature } S1 = e1^r \bmod p$$

$$\text{second point } e2 = e1^d \bmod p$$

$$\text{second signature } S2 = (M - d \times S1) \times r^{-1} \bmod (p - 1)$$

$$V1 = e1^M \bmod p$$

$$V2 = e2^{S1} \times S1^{S2} \bmod p$$

Is $V1 = V2$?

Schnorr Digital Signature Scheme

$$\text{Signature } S1 = h(M || e1^r \bmod p)$$

$$\text{Signature } S2 = r + d \times S1 \bmod q$$

$$\text{Verification } V = h(M || e1^{S2} e2^{-S1} \bmod p)$$

Is $S1$ congruent to V ?

Digital Signature Standard (DSS)

$$e1 = e0^{(p-1)/q} \bmod p$$

$$e2 = e1^d \bmod p$$

$$\text{Signature } S1 = (e1^r \bmod p) \bmod q$$

$$\text{Signature } S2 = (h(M) + d \times S1) r^{-1} \bmod q$$

$$\text{Verification } V = [(e1^{h(M)} S2^{-1} e2^{S1} S2^{-1}) \bmod p] \bmod q$$

The ECDSS scheme

$$P(u, v) = r \times e1 (... , ...)$$

$$\text{Second point } e2 = d \times e1(... , ...)$$

$$\text{Signature } S1 = u \bmod q$$

$$\text{Signature } S2 = (h(M) + d \times S1) r^{-1} \bmod q.$$

$$A = h(M) S2^{-1} \bmod q \text{ and } B = S2^{-1} S1 \bmod q$$

$$\text{Third Point } T(x, y) = A \times e1 (... , ...) + B \times e2(... , ...)$$

Is x congruent to $S1 \bmod q$?

Blind Signature

$$\text{Blinded message } B = M \times b^e \bmod n$$

$$\text{Signature on the blind version of the message } S_b = B^d \bmod n$$

$$\text{Signature is } S = S_b b^{-1} \bmod n$$

Chaum and van Antwerpen Undeniable signature

Signature $s = m^k \pmod{p}$

Challenge $c = s^a (g^k)^b \pmod{p}$

Response $r = c^{k^{-1}}$

Check whether $r = m^a g^b \pmod{p}$

Fiat-Shamir Protocol

$v = s^2 \pmod{n}$

$x = r^2 \pmod{n}$

$y = rs^c \pmod{n}$

Is y^2 congruent to xv^c ?

Feige-Fiat-Shamir Protocol

$[s_1, s_2, \dots, s_k]$ - a vector of private keys

$[v_1, v_2, \dots, v_k]$ - a vector of public keys

a vector of challenges (c_1, c_2, \dots, c_k)

$x = r^2 \pmod{n}$

$v_i = (s_i^2)^{-1} \pmod{n}$

$y = (rs_1^{c_1} s_2^{c_2} s_1^{c_3} \dots s_k^{c_k}) \pmod{n}$

Is $y^2 v_1^{c_1} v_2^{c_2} v_1^{c_3} \dots v_k^{c_k}$ congruent to x

Guillou-Quisquater Protocol

$s^e \times v = 1 \pmod{n}$

$x = r^e \pmod{n}$

$y = rs^c \pmod{n}$

Is x congruent to $y^e v^c$