

# Polytope of Matrices with Invariant Row- and Column- Sums

Pramod Mathai \*

May 31, 2019

Motivated by Aaronson's [1] description of the role of transition matrices in restricting hidden variable theories (HVTs) of quantum mechanics, we address the following general question:

(i) Given  $\mathbf{r}, \mathbf{c} \in \mathbb{R}_+^{n \times 1}$ , with  $\mathbf{1}^T \mathbf{r} = \mathbf{1}^T \mathbf{c} = \gamma$ , how does one construct  $\mathbf{S} \in \mathbb{R}_+^{n \times n}$  with row and column sums satisfying  $\mathbf{S} \mathbf{1} = \mathbf{r}, \mathbf{S}^T \mathbf{1} = \mathbf{c}$ ? Here,  $\mathbf{1}$  is the vector of all ones and the "+" subscript indicates nonnegative matrix elements. The case  $\gamma = 1$  corresponds to requiring that  $\mathbf{S}$  be a transition matrix.

$\mathcal{S}_{(\mathbf{r}, \mathbf{c})}$ , the set of all solutions to (i), is convex. We explicitly construct  $\mathcal{V}_{(\mathbf{r}, \mathbf{c})}$ , the set of vertices of a polytope  $\mathcal{P}_{(\mathbf{r}, \mathbf{c})} \subseteq \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$ .

## 1 Constructing $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$

Without loss of generality, we assume that the elements of  $\mathbf{r}$  and  $\mathbf{c}$  are arranged in increasing order:  $r_i \leq r_{i+1}$  for  $i \in [n-1]$ , likewise for  $\mathbf{c}$  ( $[k]$  denotes the set of positive integers not exceeding  $k$ ). If one solves for this case, then the rows and columns of  $\mathbf{P}$  can always be permuted back to their original order without changing the values of the respective sums.

We define  $m_i := \min(r_i, c_i)$  and  $\mathbf{e}_i$  as the column vector with 1 in the  $i^{th}$  position and zeros in the other  $(n-1)$  positions.

The polytope  $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$  is then defined as the convex hull of the set

$$\mathcal{V}_{(\mathbf{r}, \mathbf{c})} = \mathbf{P}^* \bigcup_{\substack{i, j \in [n], \\ i < j}} \mathbf{P}_{(i, j)} \quad (\diamond)$$

of, in general,  $1 + \frac{n(n-1)}{2}$  distinct matrices where  $\mathbf{P}_{(i, j)} = \mathbf{P}^* + m_1 \cdot (\mathbf{e}_i \mathbf{e}_j^T + \mathbf{e}_j \mathbf{e}_i^T - \mathbf{e}_i \mathbf{e}_i^T - \mathbf{e}_j \mathbf{e}_j^T)$  and  $\mathbf{P}^*$  has entries  $p_{i, j}^*$  equal to zero except for the following elements:

$p_{i, i}^* = m_i$ ,  $p_{n, i}^* = c_i - m_i$ ,  $p_{i, n}^* = r_i - m_i$  for  $i \in [n-1]$  and  $p_{n, n}^* = r_n + c_n - \gamma + \sum_{i=1}^{n-1} m_i$ .

For example, in the case  $n = 3$  with ordering  $r_1 \leq c_1 \leq c_2 \leq r_2 \leq r_3 \leq c_3$ , every element in the polytope defined by the four matrices

$$\mathbf{P}^* = \begin{bmatrix} r_1 & 0 & 0 \\ 0 & c_2 & r_2 - c_2 \\ c_1 - r_1 & 0 & r_3 + c_3 - \gamma + (r_1 + c_2) \end{bmatrix}, \mathbf{P}_{(1, 2)} = \begin{bmatrix} 0 & r_1 & 0 \\ r_1 & c_2 - r_1 & r_2 - c_2 \\ c_1 - r_1 & 0 & r_3 + c_3 - \gamma + (r_1 + c_2) \end{bmatrix},$$

$$\mathbf{P}_{(1, 3)} = \begin{bmatrix} 0 & 0 & r_1 \\ 0 & c_2 & r_2 - c_2 \\ c_1 & 0 & r_3 + c_3 - \gamma + c_2 \end{bmatrix} \text{ and } \mathbf{P}_{(2, 3)} = \begin{bmatrix} r_1 & 0 & 0 \\ 0 & c_2 - r_1 & r_2 - c_2 + r_1 \\ c_1 - r_1 & r_1 & r_3 + c_3 - \gamma + c_2 \end{bmatrix}$$

will be a solution to (i). This is generalized for any  $n$  and ordering in the following.

---

\*pramod.m@gmail.com

**Lemma 1.1.** (a)  $\mathcal{P}_{(\mathbf{r},\mathbf{c})} \subseteq \mathcal{S}_{(\mathbf{r},\mathbf{c})}$  for  $n > 2$ . (b) Up to permutation,  $\mathcal{P}_{(\mathbf{r},\mathbf{c})} = \mathcal{S}_{(\mathbf{r},\mathbf{c})}$  for  $n = 2$ .

*Proof.* (a) To show that every matrix in  $\mathcal{P}_{(\mathbf{r},\mathbf{c})}$  is a solution of (i), the convexity of  $\mathcal{P}_{(\mathbf{r},\mathbf{c})}$  makes it sufficient to show that matrices in  $\mathcal{V}_{(\mathbf{r},\mathbf{c})}$  are solutions to (i). The construction ( $\diamond$ ) makes it clear that any  $\mathbf{P}_{(i,j)}$  has the same row and column sums as  $\mathbf{P}^*$ . So it suffices to show that  $\mathbf{P}^*$  satisfies (i) and that  $p_{i,i}^* \geq m_1$  for  $i \in [n]$ , the additional condition ensuring that the diagonal elements of  $\mathbf{P}_{(i,j)}$  are nonnegative. We verify these properties in order:

- *Row sums:* Denote  $\mathbf{P}_{i,\cdot}^*$  to be the  $i^{\text{th}}$  row of  $\mathbf{P}^*$ . For  $i \in [n-1]$ ,  $\mathbf{P}_{i,\cdot}^* \mathbf{1} = m_i + r_i - m_i = r_i$ .  
 $\mathbf{P}_{n,\cdot}^* \mathbf{1} = \sum_{i=1}^{n-1} (c_i - m_i) + p_{n,n}^* = r_n + (\sum_{i=1}^n c_i) - \gamma = r_n$ .
- *Column sums:* Verification is analogous to *Row sums*.
- *Nonnegativity of off-diagonal terms; diagonal terms  $\geq m_1$ :* Readily follows from ( $\diamond$ ) for all elements except  $p_{n,n}^*$ .  
 Requiring  $p_{n,n}^* \geq m_1$  is equivalent to requiring  $r_n + c_n + \sum_{i=2}^{n-1} m_i \geq \gamma$ . The inequality of arithmetic and geometric means implies that the left hand side is minimized when each of the  $n$  terms is equal, say, to  $M$ . The orderings  $r_i \leq r_{i+1}$  and  $c_i \leq c_{i+1}$  then imply that  $r_i = c_i = M$  for  $2 \leq i \leq n$ ,  $M \geq m_1$  and thus  $(n-1)M + m_1 = \gamma$ .  $p_{n,n}^* \geq m_1$  follows.

(b) *Case  $n = 2$ :*  $\mathcal{V}_{(\mathbf{r},\mathbf{c})}$  contains two elements  $\mathbf{P}^*$  and  $\mathbf{P}_{(1,2)}$ . For any  $\mathbf{S} \in \mathcal{S}_{(\mathbf{r},\mathbf{c})}$ , construct  $\mathbf{S}' = \mathbf{P}_R \mathbf{S} \mathbf{P}_C$  where  $\mathbf{P}_R$  and  $\mathbf{P}_C$  are permutation matrices which are chosen to ensure that the row and column sums of  $\mathbf{S}'$  are in ascending order, as required in ( $\diamond$ ). But  $\mathbf{S}' = \phi \mathbf{P}^* + (1 - \phi) \mathbf{P}_{(1,2)}$  where  $0 \leq \phi = \frac{s'_{11}}{m_1} \leq 1$  and  $s'_{11}$  is the  $(1,1)$  entry of  $\mathbf{S}'$ , assuming  $m_1 > 0$ . The case  $m_1 = 0$  is trivial. □

## 2 Regarding $\mathcal{S}_{(\mathbf{r},\mathbf{c})} \setminus \mathcal{P}_{(\mathbf{r},\mathbf{c})}$

- For  $n > 2$ ,  $\mathcal{P}_{(\mathbf{r},\mathbf{c})}$  does not, in general, contain *all* solutions to (i) since its construction shows that any  $\mathbf{P} \in \mathcal{P}_{(\mathbf{r},\mathbf{c})}$  has a symmetric principal minor of order  $n-1$  (this does not matter for the  $n \leq 2$ ). But, can one extend the above construction procedure iteratively to construct *any* matrix  $\mathbf{S} \in \mathcal{S}_{(\mathbf{r},\mathbf{c})}$ ?
- Define the *diameter* of a bounded convex set of matrices as the maximum of  $\|\mathbf{M}_i - \mathbf{M}_j\|_\infty$ , for any two matrices  $\mathbf{M}_i$  and  $\mathbf{M}_j$  in the set.  $\mathcal{S}_{(\mathbf{r},\mathbf{c})}$  is bounded and  $\mathcal{P}_{(\mathbf{r},\mathbf{c})}$  has diameter  $m_1$ . Does  $\mathcal{S}_{(\mathbf{r},\mathbf{c})}$  have a larger diameter?
- $\mathbf{P}^*$  has at most  $2n-1$  nonzero entries and any solution in  $\mathcal{S}_{(\mathbf{r},\mathbf{c})}$  must in general have at least as many. Is  $\mathbf{P}^*$  uniquely sparsest in  $\mathcal{S}_{(\mathbf{r},\mathbf{c})}$ ?

## References

- [1] S. Aaronson, “Quantum computing and hidden variables”, *Physical Review A*, vol. 71, no. 3, Mar. 2005. DOI: 10.1103/physreva.71.032325. [Online]. Available: <https://doi.org/10.1103/physreva.71.032325>.