

Polytope of Matrices with Invariant Row- and Column- Sums

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Motivated by Aaronson's [1] description of the role of transition matrices in restricting hidden variable theories (HVTs) of quantum mechanics, we address the following general question:

(i) Given $\mathbf{r}, \mathbf{c} \in \mathbb{R}_+^{n \times 1}$, with $\mathbf{1}^T \mathbf{r} = \mathbf{1}^T \mathbf{c} = \gamma$, how does one construct $\mathbf{S} \in \mathbb{R}_+^{n \times n}$ with row and column sums satisfying $\mathbf{S} \mathbf{1} = \mathbf{r}, \mathbf{S}^T \mathbf{1} = \mathbf{c}$? Here, $\mathbf{1}$ is the vector of all ones and the "+" subscript indicates nonnegative matrix elements. The case $\gamma = 1$ corresponds to requiring that \mathbf{S} be a transition matrix.

$\mathcal{S}_{(\mathbf{r}, \mathbf{c})}$, the set of all solutions to (i), is convex. We explicitly construct $\mathcal{V}_{(\mathbf{r}, \mathbf{c})}$, the set of vertices of a polytope $\mathcal{P}_{(\mathbf{r}, \mathbf{c})} \subseteq \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$.

1 Constructing $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$

Without loss of generality, we assume that the elements of \mathbf{r} and \mathbf{c} are arranged in increasing order: $r_i \leq r_{i+1}$ for $i \in [n-1]$, likewise for \mathbf{c} ($[k]$ denotes the set of positive integers not exceeding k). If one solves for this case, then the rows and columns of \mathbf{P} can always be permuted back to their original order without changing the values of the respective sums.

We define $m_i := \min(r_i, c_i)$ and \mathbf{e}_i as the column vector with 1 in the i^{th} position and zeros in the other $(n-1)$ positions.

The polytope $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$ is then defined as the convex hull of the set

$$\mathcal{V}_{(\mathbf{r}, \mathbf{c})} = \mathbf{P}^* \bigcup_{\substack{i, j \in [n], \\ i < j}} \mathbf{P}_{(i, j)} \quad (\diamond)$$

of, in general, $1 + \frac{n(n-1)}{2}$ distinct matrices where $\mathbf{P}_{(i, j)} = \mathbf{P}^* + m_i \cdot (\mathbf{e}_i \mathbf{e}_j^T + \mathbf{e}_j \mathbf{e}_i^T - \mathbf{e}_i \mathbf{e}_i^T - \mathbf{e}_j \mathbf{e}_j^T)$ and \mathbf{P}^* has entries $p_{i, j}^*$ equal to zero except for the following elements:

$p_{i, i}^* = m_i$, $p_{n, i}^* = c_i - m_i$, $p_{i, n}^* = r_i - m_i$ for $i \in [n-1]$ and $p_{n, n}^* = r_n + c_n - \gamma + \sum_{i=1}^{n-1} m_i$.

For example, in the case $n = 3$ with ordering $r_1 \leq c_1 \leq c_2 \leq r_2 \leq r_3 \leq c_3$, we have $m_1 = r_1, m_2 = c_2$ and $m_3 = r_3$. So every element in the polytope defined by the four matrices

$$\mathbf{P}^* = \begin{bmatrix} r_1 & 0 & 0 \\ 0 & c_2 & r_2 - c_2 \\ c_1 - r_1 & 0 & r_3 + c_3 - \gamma + (r_1 + c_2) \end{bmatrix}, \mathbf{P}_{(1, 2)} = \begin{bmatrix} 0 & r_1 & 0 \\ r_1 & c_2 - r_1 & r_2 - c_2 \\ c_1 - r_1 & 0 & r_3 + c_3 - \gamma + (r_1 + c_2) \end{bmatrix},$$

$$\mathbf{P}_{(1, 3)} = \begin{bmatrix} 0 & 0 & r_1 \\ 0 & c_2 & r_2 - c_2 \\ c_1 & 0 & r_3 + c_3 - \gamma + c_2 \end{bmatrix} \text{ and } \mathbf{P}_{(2, 3)} = \begin{bmatrix} r_1 & 0 & 0 \\ 0 & 0 & r_2 \\ c_1 - r_1 & c_2 & r_3 + c_3 - \gamma + r_1 \end{bmatrix}$$

will be a solution to (i). This is generalized for any n and ordering in the following.

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Lemma 1.1. (a) $\mathcal{P}_{(\mathbf{r}, \mathbf{c})} \subseteq \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$. (b) Upto permutation, $\mathcal{P}_{(\mathbf{r}, \mathbf{c})} = \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$ for $n = 2$.

Proof. (a) The convexity of $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$ makes it sufficient to show that matrices in $\mathcal{V}_{(\mathbf{r}, \mathbf{c})}$ are solutions to (i). The construction (\diamond) makes it clear that any $\mathbf{P}_{(i,j)}$ has the same row and columns sums as \mathbf{P}^* and has nonnegative off-diagonal terms. So it suffices to verify that \mathbf{P}^* satisfies (i) and that $p_{i+1,i+1}^* \geq p_{i,i}^*$ for $i \in [n-1]$; the latter inequality ensures that the diagonal elements of every $\mathbf{P}_{(i,j)}$ are nonnegative. We verify these properties of \mathbf{P}^* in order ($n > 1$ is assumed):

- *Verifying row sums:* Denote $\mathbf{P}_{i,\cdot}^*$ to be the i^{th} row of \mathbf{P}^* . For $i \in [n-1]$, $\mathbf{P}_{i,\cdot}^* \mathbf{1} = m_i + r_i - m_i = r_i$. $\mathbf{P}_{n,\cdot}^* \mathbf{1} = \sum_{i=1}^{n-1} (c_i - m_i) + p_{n,n}^* = r_n + (\sum_{i=1}^n c_i) - \gamma = r_n$.
- *Verifying column sums:* Analogous to verifying row sums.
- *Verifying Nonnegativity of off-diagonal terms:* Follows from the definition of \mathbf{P}^* in (\diamond).
- *Verifying $p_{i+1,i+1}^* \geq p_{i,i}^*$ for $i \in [n-2]$ with $n > 2$:* Without loss of generality, let $m_{i+1} = r_{i+1}$. We have $p_{i+1,i+1}^* = m_{i+1} = r_{i+1} \geq r_i \geq \min(r_i, c_i) = m_i = p_{i,i}^*$ where the first inequality is assumed in (\diamond).
- *Verifying $p_{n,n}^* \geq p_{n-1,n-1}^* = m_{n-1}$ for $n \geq 2$:*
 $n = 2$: Requiring $p_{2,2}^* \geq m_1$ is equivalent to requiring that $r_2 + c_2 \geq \gamma$, which is satisfied since $\gamma = \frac{r_2+c_2}{2} + \frac{r_1+c_1}{2}$ and the orderings $r_1 \leq r_2$ and $c_1 \leq c_2$ are assumed in (\diamond).
 $n > 2$: Requiring $p_{n,n}^* \geq m_{n-1}$ is equivalent to requiring $r_n + c_n + \sum_{i=1}^{n-2} m_i \geq \gamma$. By the inequality of arithmetic and geometric means, the left hand side is minimized when each of the n terms is equal to, say, M . The orderings $r_i \leq r_{i+1}$ and $c_i \leq c_{i+1}$ then imply that $r_i = c_i = M$ for $i \in [n]$. Since $\sum_{i=1}^n r_i = \gamma$, it follows that $nM = \gamma$.

(b) $\mathcal{V}_{(\mathbf{r}, \mathbf{c})}$ contains two elements \mathbf{P}^* and $\mathbf{P}_{(1,2)}$. For any $\mathbf{S} \in \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$, construct $\mathbf{S}' = \mathbf{P}_R \mathbf{S} \mathbf{P}_C$ where \mathbf{P}_R and \mathbf{P}_C are permutation matrices which are chosen to ensure that the row and column sums of \mathbf{S}' are in ascending order, as required in (\diamond). Assuming $m_1 > 0$, $\mathbf{S}' = \phi \mathbf{P}^* + (1 - \phi) \mathbf{P}_{(1,2)}$ where $0 \leq \phi = \frac{s'_{11}}{m_1} \leq 1$ and s'_{11} is the $(1, 1)$ entry of \mathbf{S}' . The case $m_1 = 0$ is trivial. □

2 Regarding $\mathcal{S}_{(\mathbf{r}, \mathbf{c})} \setminus \mathcal{P}_{(\mathbf{r}, \mathbf{c})}$

- For $n > 2$, $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$ does not, in general, contain *all* solutions to (i) since its construction shows that any $\mathbf{P} \in \mathcal{P}_{(\mathbf{r}, \mathbf{c})}$ has a symmetric principal minor of order $n-1$ (this does not matter for the $n \leq 2$). But, can one extend the above procedure to construct *any* matrix $\mathbf{S} \in \mathcal{S}_{(\mathbf{r}, \mathbf{c})}$?
- Define the *diameter* of a bounded convex set of matrices as the maximum of $\|\mathbf{M}_i - \mathbf{M}_j\|_\infty$, for any two matrices \mathbf{M}_i and \mathbf{M}_j in the set. Does $\mathcal{S}_{(\mathbf{r}, \mathbf{c})}$ have a larger diameter than $\mathcal{P}_{(\mathbf{r}, \mathbf{c})}$?
- \mathbf{P}^* has at most $2n-1$ nonzero entries and any solution in $\mathcal{S}_{(\mathbf{r}, \mathbf{c})}$ must in general have at least as many. Does $\mathcal{S}_{(\mathbf{r}, \mathbf{c})} \setminus \mathcal{P}_{(\mathbf{r}, \mathbf{c})}$ contain matrices of equal sparsity?

References

- [1] S. Aaronson, “Quantum computing and hidden variables”, *Physical Review A*, vol. 71, no. 3, Mar. 2005. DOI: 10.1103/physreva.71.032325. [Online]. Available: <https://doi.org/10.1103/physreva.71.032325>.