Merkle Trees (1979)

- hash is collision-resistant MHT root of Merkle hash tree = h(cld) dada Hems $d_1 \xrightarrow{h} b_2 \xrightarrow{h} i \xrightarrow{j} h \rightarrow MHT?=MHT$ $d_1 + d_1 \Rightarrow h(d_1) + h(d_1)$ $d_1 + d_2 \Rightarrow h(d_1) + h(d_1)$ $h(a,b) \Rightarrow i + i$ $h(a,b) \Rightarrow i + i$ $h(a,b) \Rightarrow i + i$ $h(a,b) \Rightarrow i + i$ $h(h(a,h(d_1)),j) \stackrel{?}{=} MHT$

Property: Gilven MHT, one can prove that data item of is in the tree in a logarithmu'c # of steps and using a log number of nodes in the tree

Audit proof for a nodelle = sillings of nodes on the path from N to root

Client: have MHT de Attacker: Says de is in the tree, a', j' is the audit proof

Assume that

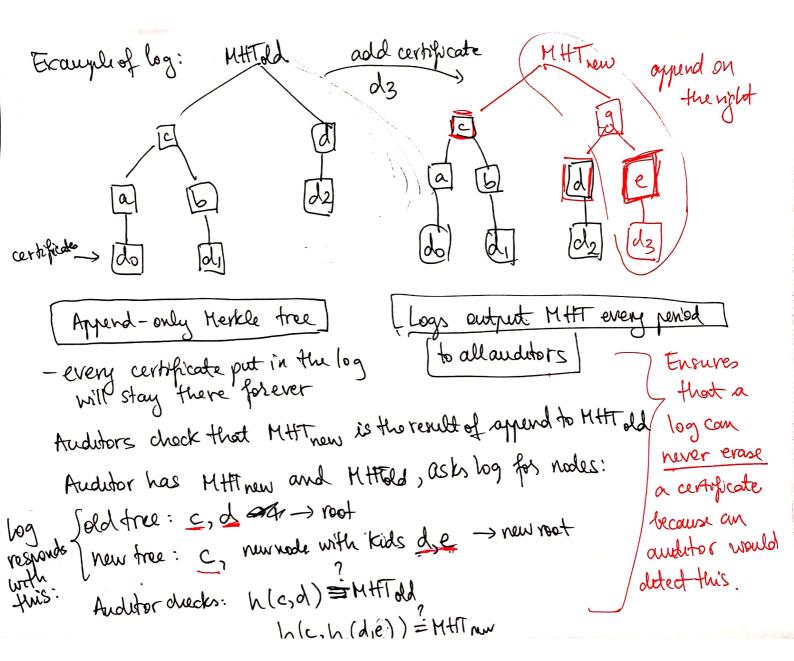
Example use case: secure storage Want Integrity Client MHT > MHT latest HHT @MHT Ask for siblings of path of > 1 Venfy a, h(Fz) hlah(h(F2) 30 h(F3)) = HHT Stramman: 2. Compute MHT F2; MACy (fleid; F2) $F_2 \xrightarrow{h} h(F_2)$ Chent fileid 2 Chard $h(a,h(h(F_2'),h(F_3))) \rightarrow MHT'$ check T2', MACK (fledd2', F2)

MACR (fledd2', F2)

Freshness

All id Replay attacks attack

MAC Cloud can respond with old Alternatively, fle versions because they had a cowed MAC



Certificate Transparency (CT)

Problem: -2011 diginator signed fake certs for \$0, gogle.com

CA3 can get compromised

- ledger - hosted centrally (eg. Google) + high-throughput, scalable, no proof of work

Lecentralized security

Thisight: transparency: CT enables anyone to detect if there is a fake certificates

them; it does not prevent creation of fake certificates

3 parties: -logs: store certificates, all corts in the world = Google

- monetors

-auditors

Logs: lach log is a Merkle hash tree over curtificates

