

EE61012: Convex Optimization for Control and Signal Processing

Instructor: Prof. Ashish R. Hota

- Class Hours: D Slot. Monday: 10am - 10:55pm, Wednesday: 8am - 9:55am, Thursday: 10am - 10:55am
- Venue: NR 313
- Grading Scheme: 50 % Endsem, 30 % Midsem, 20 % Tutorial, Class Tests
- Preferred Mode of Contact: Send email to ahota@ee.iitkgp.ac.in with subject containing [EE61012]. Do not forget to write your name and roll no.
- Any email with a blank subject and without name and roll no. will be ignored.

Weekly Plan

Week 1: 4th - 11th January

- Formal Definition of an Optimization Problem
- Constraints, Feasible solutions, Optimal solution, Optimal value
- Infeasible and unbounded optimization problems
- Local vs. global optimal solutions
- Compact Sets, Continuous Functions, Weierstrass Theorem on existence of global optima
- Gradient, Hessian, Optimality conditions for unconstrained problems

Week 2: 15th - 18th January

- Convex Sets
- Examples
- Operations that preserve convexity of sets
- Convex functions: Definition, Level set Characterization, First order characterization

Weekly Plan

Week 3: 22nd - 25th January

- Convex functions: Second order characterization
- Operations that preserve convexity of functions
- Examples
- Formulate and solve simple convex optimization problems (such as constrained least squares problem) using suitable solvers

Week 4: 29th January - 1st February

- Examples of Convex Optimization Problem Classes
- Equivalent Forms
- Separating Hyperplane Theorems, Theorems of the Alternative, LP Duality

Weekly Plan

Week 5 & 6: 5th - 14th February

- Convex Theorems of the Alternative, Constraint Qualification
- Lagrangian Duality: weak and strong versions
- Saddle Point Formulations
- KKT Optimality Conditions
- Examples
- Properties of Convex Optimization Problems: Global Optimality, Strong Duality, Necessary Conditions being Sufficient
- Regression Problems and applications
- Practice Problems

Mid-semester Examination

Week 7: 26th - 29th February

- Classification via Support Vector Machines
- ML Estimation
- Hypothesis Testing and Optimal Detection

Weekly Plan

Week 8 & 9: 4th - 14th March

- First order algorithms, Accelerated Methods
- Stochastic Gradient Descent
- Distributed Optimization

Week 10 & 11: 18th - 28th March

- Linear Matrix Inequality
- Conic Duality
- Semidefinite Programming
- Applications of SDP in Control: Stability, State Feedback Synthesis, Robust Synthesis

Week 12 & 13: 1st - 11th April

- Constrained Optimal Control, Model Predictive Control
- Applications in System Identification
- Robust Optimization via Duality

References

Primary Reference:

- Convex Optimization by Stephen Boyd and L. Vandenberghe, Cambridge University Press. Available online at: <https://web.stanford.edu/~boyd/cvxbook/>
- Algorithms for Convex Optimization by Nisheeth K. Vishnoi, Cambridge University Press. Available online at: <https://convex-optimization.github.io>

Advanced References on Theory

- Lectures on Modern Convex Optimization, Aharon Ben-Tal and Arkadi Nemirovski, SIAM. Available online at: <https://pubs.siam.org/doi/book/10.1137/1.9780898718829>
- Convex Analysis and Optimization, Bertsekas, Athena Scientific. More information at: <http://www.athenasc.com/convexity.html>
- Convex Analysis and Minimization Algorithms, Jean-Baptiste Hiriart-Urruty, Claude Lemarechal, Springer. Available online at: <https://link.springer.com/book/10.1007/978-3-662-02796-7>

Advanced References on Algorithms

- Optimization for Modern Data Analysis, Benjamin Recht and Stephen J. Wright, Available online at: https://people.eecs.berkeley.edu/~brecht/opt4ml_book/
- Numerical Optimization by Jorge Nocedal, Stephen J. Wright, Springer. Available online at: <https://link.springer.com/book/10.1007/978-0-387-40065-5>
- Introductory Lectures on Convex Optimization A Basic Course, by Yurii Nesterov. Available online at: <https://link.springer.com/book/10.1007/978-1-4419-8853-9>
- First-order Methods in Optimization, by Amir Beck, SIAM. For more information: <https://pubs.siam.org/doi/10.1137/1.9781611974997>

References

Advanced References on Applications in Control

- Linear Matrix Inequalities in System and Control Theory, by Stephen Boyd, Laurent El Ghaoui, E. Feron, and V. Balakrishnan, Society for Industrial and Applied Mathematics (SIAM), 1994. Available online at: <https://web.stanford.edu/~boyd/lmibook/>
- A Course in Robust Control Theory: A Convex Approach, Springer. Available online at: <https://link.springer.com/book/10.1007/978-1-4757-3290-0>
- Predictive Control for Linear and Hybrid Systems, Cambridge University Press. More information at: <http://www.mpc.berkeley.edu/mpc-course-material>

Advanced References on Applications in Signal Processing and Machine Learning

- Convex Optimization in Signal Processing and Communications, Cambridge University Press. More information at: <https://www.cambridge.org/in/academic/subjects/engineering/communications-and-signal-processing/convex-optimization-signal-processing-and-communications?format=HB&isbn=9780521762229>
- Optimization for Machine Learning, by Suvrit Sra, Stephen J. Wright, Sebastian Nowozin, MIT Press. More information at: <https://mitpress.mit.edu/9780262537766/optimization-for-machine-learning/>
- Recent Special Issue of Proceedings of the IEEE: <https://ieeexplore.ieee.org/xpl/tocresult.jsp?isnumber=9241485&punumber=5>

Computing Resources

MATLAB Toolbox

- YALMIP: <https://yalmip.github.io/>
- CVX: <http://cvxr.com/cvx/>

Python Toolbox

- CVXPY: <https://www.cvxpy.org/>
- PYOMO: <http://www.pyomo.org/>

Solvers

- MOSEK: <https://www.mosek.com/>
- Gurobi: <https://www.gurobi.com/>
- IPOPT: <https://github.com/coin-or/Ipopt>
- COIN-OR: <https://github.com/coin-or/>
- For optimal control, Casadi: <https://web.casadi.org/>

Preliminaries

See https://www.stat.cmu.edu/~ryantibs/convexopt/prerequisite_topics.pdf for refresher.

Please also see the Appendices of Boyd's Book and Chapter 2 of ACO Book.

Optimization in Abstract Form

$|X|$ is finite.

An optimization problem can be stated as

$$\min_{x \in X} f(x),$$

$$\min a x^2 \quad (1)$$

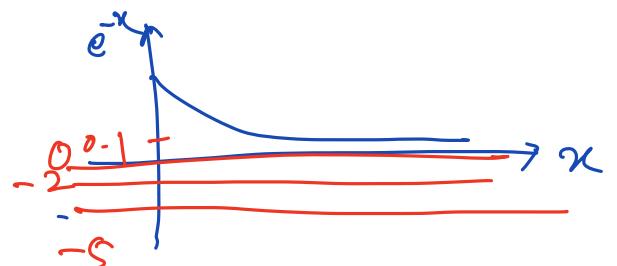
where

- x decision variable, often a vector in \mathbb{R}^n
- X set of feasible solutions, often a subset of \mathbb{R}^n
 - often specified in terms of equality and inequality constraints
$$X := \{x \in \mathbb{R}^n \mid g_i(x) \leq 0, h_j(x) = 0, i \in \{1, 2, \dots, m\}, j \in \{1, 2, \dots, p\}\}.$$
- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ cost function

Goal:

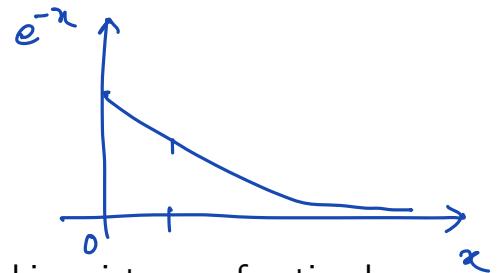
- Find $x^* \in X$ that minimizes the cost function, i.e., $f(x^*) \leq f(x)$ for every $x \in X$.
- Optimal value: $f^* := \inf_{x \in X} f(x)$
- Optimal solution: $x^* \in X$ if $f(x^*) = f^*$.

What is $\inf_{x \in X} f(x)$? : Greatest lower bound.



Examples

- Let $f(x) = e^{-x}$ and $X = [0, \infty)$. Find f^* and x^* . x^* is not defined
- minimum possible value e^{-1}
- What if $X = [0, 1]$? and $x^* = 1$.
- What if $X = [0, 1)$? $\{x \in \mathbb{R} \mid x \leq 1, x \geq 0\}$



Moral of the story: Properties of feasibility set X is critical in existence of optimal solution.

Infimum vs. Minimum

$f^* := \inf_{x \in X} f(x)$ if f^* is the greatest lower bound on the value of the function $f(x)$ over $x \in X$.

- For any $\epsilon > 0$, there exists some $\bar{x} \in X$ such that $f(\bar{x}) < f^* + \epsilon$.

There are two possibilities:

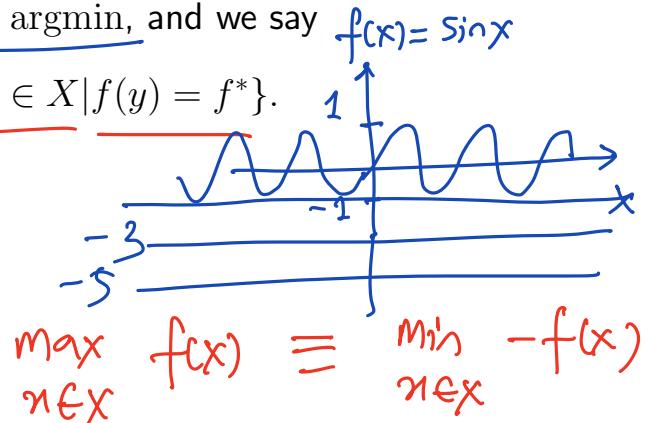
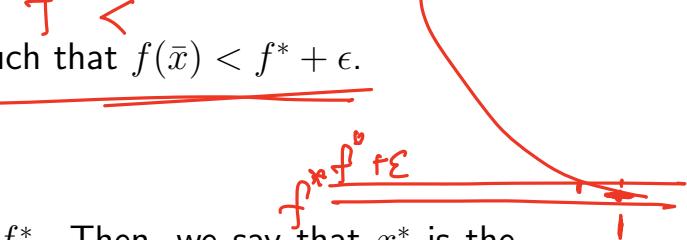
- There exists $x^* \in X$ for which $f(x^*) = f^*$. Then, we say that x^* is the optimal solution and $f^* := \min_{x \in X} f(x)$ is the optimal value.
- $f(x) \neq f^*$ for any $x \in X$. We then say that the infimum is not attained for this problem.
- If $|X|$ is finite, then infimum is always attained.
- The set of optimal solutions is denoted by argmin , and we say

$$\min_{x \in [0,1]} f(x) = 0$$

$$X = [0, 10\pi] \quad \text{inf}_{x \in X} f(x) = \sin x$$

$$-1 + \epsilon = -1$$

$$\text{argmin}_{x \in X} f(x) = \left\{ \frac{3\pi}{2}, \frac{7\pi}{2}, \dots \right\}$$



$$g(x) = -f(x)$$

Example

$$\min_{x \in X} f(x) \quad , \quad f(x) = \begin{cases} 1 & , \text{ when } x = 0 \\ x & , \text{ when } x > 0 \end{cases}$$

$x \in [0, 1]$

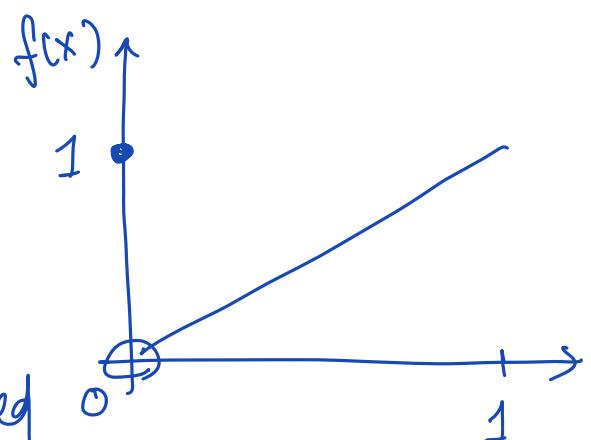
Moral of the story:

$$\inf_{x \in X} f(x) = 0$$

optimal solution is not defined

optimal value is not attained

$$\arg \min_{x \in X} f(x) = \emptyset$$



Infeasible optimization problem

- The problem is infeasible when \underline{X} is an empty set.

- In this case, $f^* := +\infty$.

- Example:

$$\underline{X} = \left\{ \underline{x} \in \mathbb{R}^2 \mid \begin{array}{l} x_1 \geq 0, x_2 \geq 1, \\ x_1 + x_2 = -1 \end{array} \right\}$$

$$X_1 = \left\{ x_1 \geq 0 \right\}$$

$$X_2 = \left\{ x_2 \geq 1 \right\}$$

$$X_3 = \left\{ x_1 + x_2 = -1 \right\}$$

$$X = X_1 \cap X_2 \cap X_3 = \emptyset$$

Unbounded optimization problem

- The problem is unbounded when $f^* = -\infty$.

- Example:

$$\left\{ \begin{array}{l} \text{unbounded} \\ f(x) = \log x \\ X = [0, 5] \end{array} \right.$$

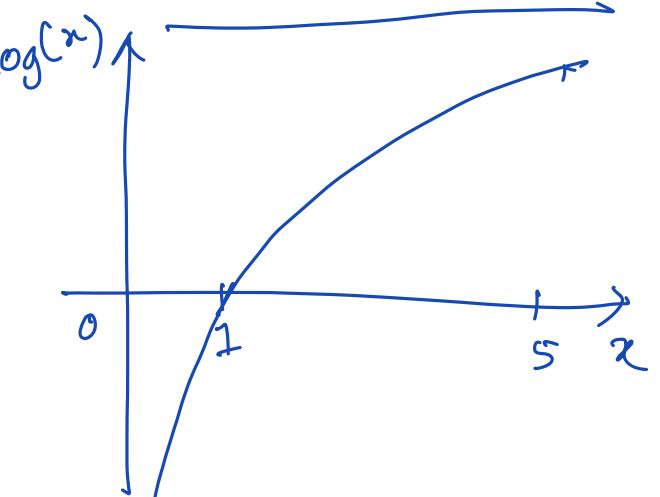
$$\left\{ \begin{array}{l} f(x) = \log x \\ X_1 = [1, 5] \end{array} \right.$$

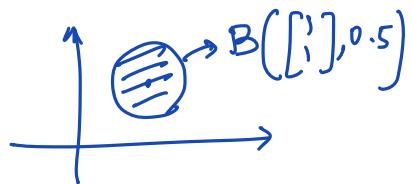
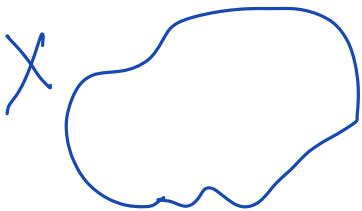
$\min_{x \in X_1} f(x)$ is not unbounded.

$f^* = 0$, and optimal solution $x^* = 1$

$X_2 = [0, 5]$: also unbounded

over the feasibility set X .





Basic Topology of Sets

Let $B(x_0, r) := \{x \in \mathbb{R}^n \mid \|x - x_0\|_2 \leq r\}$ denote the ball around point $x_0 \in \mathbb{R}^n$ with radius $r > 0$.

- Interior of the set X , denoted $\text{int}(X) = \{x_0 \in X \mid \exists \text{ a radius } r \text{ for which } B(x_0, r) \subset X\}$
- Set X is called an **open** set if $X = \text{int}(X)$.
- Set X is called **closed** if and only if its complement is open.
- Intersection of arbitrary number of closed sets is closed.

Examples of Open and Closed Sets':

Example of a set which is neither open nor closed $\boxed{x = \{1\}} \Rightarrow x \text{ is not open.}$

$$x = (0, 1]$$

$$\text{int}(x) = (0, 1)$$

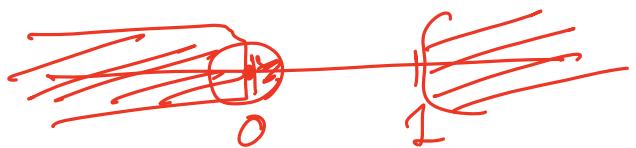
x is not open

$$x^c = (-\infty, 0] \cup (1, \infty)$$

$$\text{int}(x^c) = (-\infty, 0) \cup (1, \infty)$$

$$x^c \neq \text{int}(x^c)$$

x^c is not open \Rightarrow x is not closed.



Bounded and Compact Set

- A set X is bounded if there exists $B \in (0, \infty)$ such that for any $x_1, x_2 \in X$,
 $\|x_1 - x_2\|_2 \leq B$.

Is $[0, 10]$ bounded? Yes

Is \mathbb{R}^2 bounded? No.

- A set X is compact if it is closed and bounded.

Global and Local Optimum

Definition 1 (Global Optimum). A feasible solution $x^* \in X$ is a global optimum if $f(x^*) \leq f(x)$ for all $x \in X$. In this case, $f^* = f(x^*)$. The set of global optima is denoted by

$$\arg\min_{x \in X} f(x) := \{z \in X | f(z) = f^*\}.$$

Definition 2 (Local Optimum). A feasible solution $x^* \in X$ is a local optimum if $f(x^*) \leq f(x)$ for all $x \in B(x^*, r)$ for some $r > 0$.

Existence of Optimal Solution:

Theorem 1: Weierstrass Theorem

If the cost function f is continuous and the feasible region X is compact (closed and bounded), then (at least one global) optimal solution x^* exists.

Example:

$$B(-2, \varepsilon) = [-2 - \varepsilon, -2 + \varepsilon]$$



When X is not bounded, then the above theorem still holds when an α -sublevel set of f , defined as

$$S_\alpha(f) := \{x \in X | f(x) \leq \alpha\},$$

(for at least one $\alpha \in \mathbb{R}$)

is non-empty and bounded.

If $f(x) = x^2$, $\alpha = 36$,

$$S_\alpha(f) = \{x \in \mathbb{R} | f(x) \leq 36\}$$

$$= [-6, 6]$$

The story so far

Given an optimization problem,
first determine

$$\min_{x \in X} f(x)$$

- i) decision variable $x \in ?$
- ii) feasibility set X
- iii) cost function $f: X \rightarrow \mathbb{R}$

To check whether a globally optimal solution exists, check whether

- i) f is continuous
- ii) X is closed
- iii) X is bounded or any sub-level set of X is bounded.

convince yourself that

$$\min_{x \in X} f(x)$$

$$\equiv \boxed{\min_{x \in S_\alpha(f)} f(x)}$$

for any $\alpha \in \mathbb{R}$.
at which $S_\alpha(f) \neq \emptyset$

Today's lecture:

How do we verify if $x^* \in X$ is indeed an optimal solution?

Derivative of a function $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, denoted $Df(x)$
is such that $f(x + \Delta x) \simeq f(x) + Df(x) \cdot \Delta x$
 $\Rightarrow Df(x) \in \mathbb{R}^{m \times n}$

Chain rule: Let $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$, $g: \mathbb{R}^m \rightarrow \mathbb{R}^p$, $h(x) = g(f(x)) \in \mathbb{R}^p$

$$\underbrace{Dh(x_0)}_{\in \mathbb{R}^{p \times n}} = \underbrace{Dg(y_0)}_{\mathbb{R}^{p \times m}} \underbrace{Df(x_0)}_{\mathbb{R}^{m \times n}} \quad \text{where } \underbrace{y_0 = f(x_0)}$$

Gradient ($\nabla f(x)$)

For a function $f : \mathbb{R}^n \rightarrow \mathbb{R}$, its gradient is defined as:

$$\nabla f(y) = \begin{bmatrix} \frac{\partial f}{\partial x_1} \Big|_y \\ \frac{\partial f}{\partial x_2} \Big|_y \\ \vdots \\ \frac{\partial f}{\partial x_n} \Big|_y \end{bmatrix} = (\nabla f(y))^T$$

Compute gradient of

- $f(x) = x^T a$, $\nabla f(x) = a$

- $f(x) = x^T A x$

- $f(x) = \|Ax - b\|_2^2$

$$\Rightarrow \nabla f(x) = 2 A^T A x - 2 A^T b.$$

$$f(x) = [x_1 \ x_2 \ \dots \ x_n]^T \begin{bmatrix} \sum_{i=1}^n a_{1i} x_i \\ \sum_{i=1}^n a_{2i} x_i \\ \vdots \\ \sum_{i=1}^n a_{ni} x_i \end{bmatrix} = \sum_{j=1}^n \sum_{i=1}^n x_i a_{ij} x_j = \sum_{j=1}^n \left[a_{jj} x_j^2 + \sum_{i \neq j} x_i a_{ij} x_j \right]$$

$$\frac{\partial f}{\partial x_k} = \frac{\partial}{\partial x_k} \left[a_{kk} x_k^2 + \sum_{i \neq k} x_i a_{ik} x_k + \sum_{j \neq k} \left[a_{jj} x_j^2 + \sum_{i \neq j} x_i a_{ij} x_j \right] \right]$$

$$= 2a_{kk} x_k + \sum_{i \neq k} x_i a_{ik} + \sum_{j \neq k} a_{jj} x_j$$

$$= \sum_{i=1}^n x_i a_{ik} + \sum_{j=1}^n a_{kj} x_j = [Ax]_k + [A^T x]_k$$

$$\nabla f(x) = (A + A^T)x$$

$$\nabla f(x) : \mathbb{R}^n \rightarrow \mathbb{R}^n$$

Hessian ($H(x)$)

For a function $f : \mathbb{R}^n \rightarrow \mathbb{R}$, its Hessian is defined as:

$$H(x) = \nabla^2 f(x)$$

Compute Hessian of

- $f(x) = x^T a \Rightarrow Hf(x) = \nabla a = 0_{n \times n}$

- $f(x) = x^T Ax$

- $f(x) = \|Ax - b\|_2^2$

$$\nabla^2 f(x) = (A + A^T)x$$

$$H(x) = \nabla^2 [(A + A^T)x]$$

$$= A + A^T$$

$$\left[\begin{array}{c} \nabla^2 \frac{\partial f}{\partial x_1} \\ \nabla^2 \frac{\partial f}{\partial x_2} \\ \vdots \\ \nabla^2 \frac{\partial f}{\partial x_n} \end{array} \right] = \left[\begin{array}{cccc} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_1 \partial x_n} \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2^2} & \cdots & \frac{\partial^2 f}{\partial x_2 \partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f}{\partial x_n \partial x_1} & \frac{\partial^2 f}{\partial x_n \partial x_2} & \cdots & \frac{\partial^2 f}{\partial x_n^2} \end{array} \right] \in \mathbb{R}^{n \times n}$$

$$[H(x)]_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j}$$

$$\nabla^2 f(x) = 2A^T A x - 2A^T b$$

$$H(x) = 2A^T A$$

Directional Derivative and Descent Direction

Consider a function $f: \mathbb{R}^n \rightarrow \mathbb{R}$

Let $d \in \mathbb{R}^n$ be the direction of interest.

Definition: directional derivative of f at point $x_0 \in \mathbb{R}^n$ along $d \in \mathbb{R}^n$

is

$$\lim_{\varepsilon \rightarrow 0} \frac{f(x_0 + \varepsilon d) - f(x_0)}{\varepsilon} = \underline{\phi'(0)}$$

Define $\underline{\phi(t)} := f(x + td)$, $t \in \mathbb{R}$,

Compute $\phi'(0)$:

$$\begin{aligned} \frac{d\phi}{dt} &= \underline{\phi'(t)} = Df(y) \frac{d}{dt}(x+td) \quad , \quad y = x+td \\ &= \nabla f(x+td)^T \frac{d}{dt}(x+td) \\ &= \nabla f(x+td)^T \underline{d} \end{aligned}$$

$$\underline{\underline{\phi'(0) = \nabla f(x)^T d}}$$

'dash' will only be used for derivatives when the function is from \mathbb{R} to \mathbb{R} .

$$\begin{aligned} g(t) &= x+td \\ g &: \mathbb{R} \rightarrow \mathbb{R}^n \end{aligned}$$

$$\begin{aligned} Dg(t) &\in \mathbb{R}^{n \times 1} \\ &= d \end{aligned}$$

feasibility space
 $\Rightarrow X = \mathbb{R}^n$.

Necessary Condition of Optimality for Unconstrained Problems

Theorem 2

If x^* is a local optimum for the problem $\min_{x \in \mathbb{R}^n} f(x)$, then $\nabla f(x^*) = 0$.

Proof by contradiction:

Suppose x^* is a local optimum, yet $\nabla f(x^*) \neq 0$.

Recall: directional derivative along d : $\underline{\nabla f(x^*)^T d}$

If $d = -\nabla f(x^*)$, then directional derivative
 $\rightarrow \nabla f(x^*)^T \nabla f(x^*)$

$$\begin{aligned} f(x^* + td) &\approx f(x^*) + t \nabla f(x^*)^T d + (\text{Hot}) = -\|\nabla f(x^*)\|_2^2 \\ &= f(x^*) - t \|\nabla f(x^*)\|_2^2 + (\text{Hot}) \end{aligned}$$

There always exists \bar{t} sufficiently small such that

$$-\bar{t} \|\nabla f(x^*)\|_2^2 + (\text{Hot}) < 0$$

$$\Rightarrow f(x^* + \bar{t}d) < f(x^*)$$

$\Rightarrow x^*$ is not a local optimum \Rightarrow contradiction!

Sufficient Condition of Optimality for Unconstrained Problems

Let f be twice continuously differentiable over \mathbb{R}^n .

Theorem 3

If for $x^* \in \mathbb{R}^n$, we have $\nabla f(x^*) = 0$ and the Hessian of the cost function f at x^* is a positive definite matrix, then x^* is a local optimum for the problem $\min_{x \in \mathbb{R}^n} f(x)$.

Recall: Taylor Series expansion

$$f(x) = f(x_0) + \nabla f(x_0)^T (x - x_0) + \frac{1}{2} (x - x_0)^T \nabla^2 f(x_0) (x - x_0) + \dots$$

If at x^* , $\nabla f(x^*) = 0$, then

$$f(x^* + \varepsilon d) = f(x^*) + \underline{\underline{(\varepsilon d)^T H(x^*) \varepsilon d}} + \text{not tve } + d.$$

Therefore $f(x^* + \varepsilon d) > f(x^*)$.

Least-squares problem / Linear Regression

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2$$

$$\text{Hence } f(x) = \|Ax - b\|_2^2$$

$$\nabla f(x) = 2A^T A x - 2A^T b.$$

$$H(x) = 2A^T A.$$

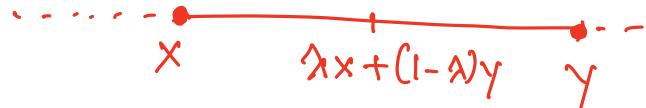
If x^* is an optimal solution, then $\nabla f(x^*) = 0$
 $\Rightarrow \boxed{A^T A x^* = A^T b}$

If $H(x)$ is positive definite, $x^* = (A^T A)^{-1} A^T b$ is the unique solution which is a local optimum.

Note: $A^T A$ is always positive semi-definite, but may not always be positive definite.

Convex Sets

Definition 3. Given a collection of points x_1, x_2, \dots, x_k , the combination $\lambda_1 x_1 + \lambda_2 x_2 + \dots + \lambda_k x_k$ is called **Convex** if $\lambda_i \geq 0$ and $\sum_{i=1}^k \lambda_i = 1$. A set X is convex if all convex combinations of its elements are in the set.



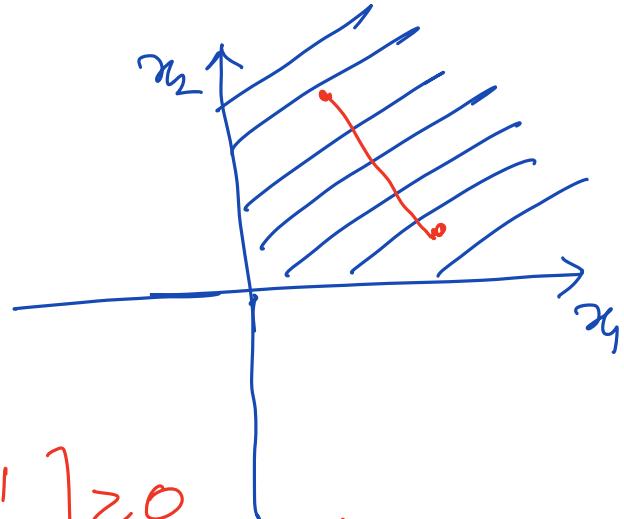
Equivalently, X is a convex set if

- for every $x, y \in X$, $\lambda x + (1 - \lambda)y \in X$ for any $\lambda \in [0, 1]$.
- it contains all convex combinations of any two of its elements.

Examples :

i) $X_1 = \{x \in \mathbb{R}^2 \mid x_1 \geq 0, x_2 \geq 0\}$

$y_1, z_1 \in X \Rightarrow y_1 \geq 0, z_1 \geq 0$
 $y_2, z_2 \geq 0$



Let $\lambda \in [0, 1]$.

$$p = \lambda y_1 + (1 - \lambda)z_1 = \begin{bmatrix} \lambda y_1 + (1 - \lambda)z_1 \\ \lambda y_2 + (1 - \lambda)z_2 \end{bmatrix} \geq 0 \Rightarrow p \in X$$

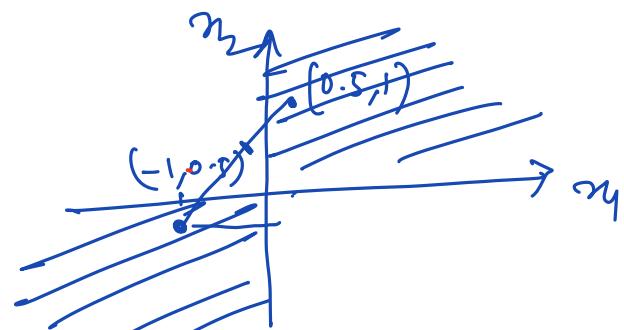
ii) $X_2 = \{x \in \mathbb{R}^2 \mid x_1, x_2 \geq 0\}$

$$y_1 = \begin{bmatrix} -1 \\ -0.5 \end{bmatrix}, z_1 = \begin{bmatrix} 0.5 \\ 1 \end{bmatrix}$$

$$\lambda = 1, p = \lambda y_1 + (1 - \lambda)z_1,$$

$$= \frac{y_1 + z_1}{2} = \begin{bmatrix} -0.25 \\ +0.25 \end{bmatrix} \notin X_2$$

$\Rightarrow X_2$ is not a convex set.



Basic Examples of Convex Sets

Sets Defined by Linear Inequalities:

- Hyperplane: $H = \{x \in \mathbb{R}^n | a^T x = b\}$ for some $a \in \mathbb{R}^n, b \in \mathbb{R}$.

$$\text{let } x_1, x_2 \in H \Rightarrow a^T x_1 = b \quad \text{let } \lambda \in [0, 1) \\ a^T x_2 = b. \quad p = \lambda x_1 + (1-\lambda) x_2$$

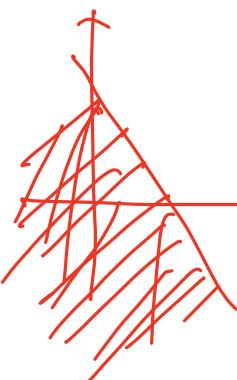
- Halfspaces: $\{x \in \mathbb{R}^n | a^T x \leq b\}$ for some $a \in \mathbb{R}^n, b \in \mathbb{R}$.

we need to check whether $a^T p = b$.

also a convex set $a^T p = a^T (\lambda x_1 + (1-\lambda) x_2)$

$$= \lambda a^T x_1 + (1-\lambda) a^T x_2 = \lambda b + (1-\lambda)b \\ = b$$

$\Rightarrow p \in H \Rightarrow H \text{ is a convex set.}$

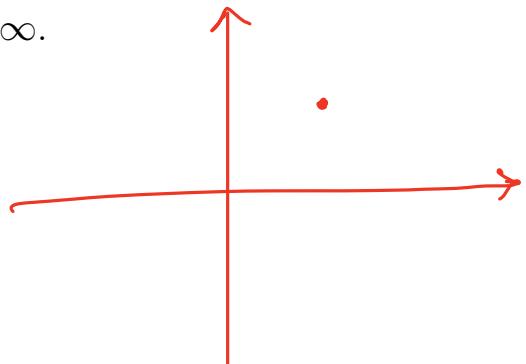


Sets Defined by Norms

Consider the Ball $B_p(c, R) := \{x \in \mathbb{R}^n \mid \|\underline{x} - c\|_p \leq R\}$ where

$$\|\underline{x}\|_p := \begin{cases} \left(\sum_{i \in [n]} |x_i|^p\right)^{\frac{1}{p}}, & 1 \leq p < \infty, \\ \max_{i \in [n]} |x_i|, & p = \infty. \end{cases}$$

We define $[n] := \{1, 2, \dots, n\}$.



Proposition 1. $B_p(c, R)$ is a convex set.

Let $x_1, x_2 \in B_p(c, R)$

$$\Rightarrow \|x_1 - c\|_p \leq R, \|x_2 - c\|_p \leq R$$

Let $\lambda \in [0, 1]$, $z := \lambda x_1 + (1-\lambda)x_2$.

Let us check whether z belongs to $B_p(c, R)$.

$$\begin{aligned} \|z - c\|_p &= \|\lambda x_1 + (1-\lambda)x_2 - c\|_p \\ &= \|\lambda x_1 + (1-\lambda)x_2 - \lambda c - (1-\lambda)c\|_p \\ &= \|\lambda(x_1 - c) + (1-\lambda)(x_2 - c)\|_p \\ &\leq \|\lambda(x_1 - c)\|_p + \|(1-\lambda)(x_2 - c)\|_p \quad (\text{triangle inequality}) \\ &= \lambda \|x_1 - c\|_p + (1-\lambda) \|x_2 - c\|_p \quad (\text{the homogeneity}) \\ &\leq \lambda R + (1-\lambda)R = R. \quad (\because x_1, x_2 \in B_p(c, R)) \end{aligned}$$

$\Rightarrow z \in B_p(c, R) \Rightarrow B_p(c, R)$ is a convex set.

Positive Semidefinite Matrices

$$A \succeq B \Rightarrow (A - B) \succeq 0$$

\downarrow
A - B is the semidefinite

Proposition 2. Set of symmetric positive semidefinite matrices, denoted by $S_n^+ := \{X \in S^n \mid X \succeq 0_{n \times n}\}$, is a convex set.

Let $X_1, X_2 \in S_n^+ \Rightarrow X_1 \succeq 0$
 $X_2 \succeq 0$

Let $\lambda \in [0, 1]$

$$Z = \lambda X_1 + (1 - \lambda) X_2$$

Let $v \in \mathbb{R}^n$.

$$\begin{aligned} v^T Z v &= v^T (\lambda X_1 + (1 - \lambda) X_2) v \\ &= \underbrace{\lambda v^T X_1 v}_{\geq 0} + (1 - \lambda) \underbrace{v^T X_2 v}_{\geq 0} \geq 0 \end{aligned}$$

$\Rightarrow Z$ is positive semidefinite -

$\Rightarrow S_n^+$ is a convex set -