Design & Analysis of Algorithum Pranau Umakant Prijon 1001965075 [2.1-1) given array 2 (1,4,5,10,16,17,213 BST constraint: less children must be & parent right duit must be > parent Tre of height 2: (3 layers, first layer is soot) 4 17 1 5 16 21

Tree of height = 3:

Tree of height 4 Tree of hight 5: Tree of height 6:

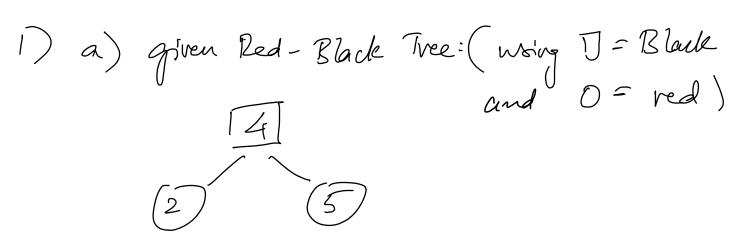
1-4-5-10-16-17-21

12.1-2) RST Property: All left children are lessen than the parent rode in value and all right children are greater than the parent

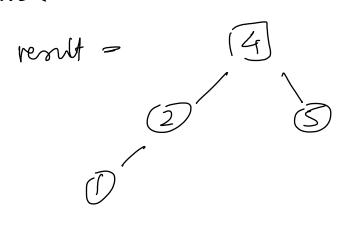
node in value. This means that an in-order DFG transversal terrough the BFG tree will return the costed order of elements in ascending order. Min-tleap Property: Every node has value that is lesser than or agual to that of its parent. The root node is the minimum value.

RST is already RoArd in-order, so it is possible to retrieve a Goded sequence in O(n) time. Moneury a min heap must heapity (O(logn)) every time the min " element is popped. This happens in time, hence O(nlogn) is the minimum The for this operation woing a min-heap. There is no relative ordening maintained between the noder at a particular level in a nin-heap, unlike à RST.

HW - Ch 13

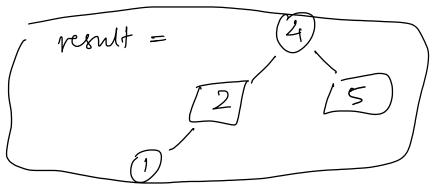


to insert 1, we can ty inserting (1) as a red node.



this violates property that red node cannot have a red child.

to repolve this, lets convert (1 to a red node, Give "4" is not roost, this is possible.



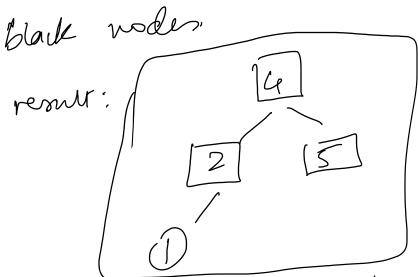
final check: No consecutive red nodes, black-height is maintained from noot to all NIL modes.

This is Now a valid Ped-Black Free.

b) Now, since 4 is the soot, it will have to remain a black usde.

So, to add red made (T), we need to have its parent mades be black. Hence,

(5) and 2 need to become [5] and [2]



this is now a valid red-black tree.

C) given tree:

incoming red node (5):-

res ut: this violates property that a red node counst have another red usde. Changing parent of 6 to black:

this tree catisfies the Red- black properties!