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Lab-9 First Order Logic Forward Chaining

Considering Statement

As per law, it is a crime for an American to sell weapons to hostile nations. Country A, an enemy of America, has some missiles, and all missiles were sold ^{to it} by Robert, who is American.

Prove Robert is a criminal.

Deduced Statements:

$$\text{American}(x) \wedge \text{Weapon}(y) \wedge \text{Sold}(z, y, x) \Rightarrow \text{Criminal}(x) \quad \wedge \text{Hostile}(z)$$

Enemy(America, America)

~~Has (miss)~~

$$\forall x \text{ Missile}(x) \rightarrow \text{to}$$

$$\exists x \text{ Missile}(x) \wedge \text{Own}(A, x)$$

$$\forall x \text{ Missile}(x) \wedge \text{Own}(A, x) \Rightarrow \text{Sold}(A, x, \text{Robert})$$

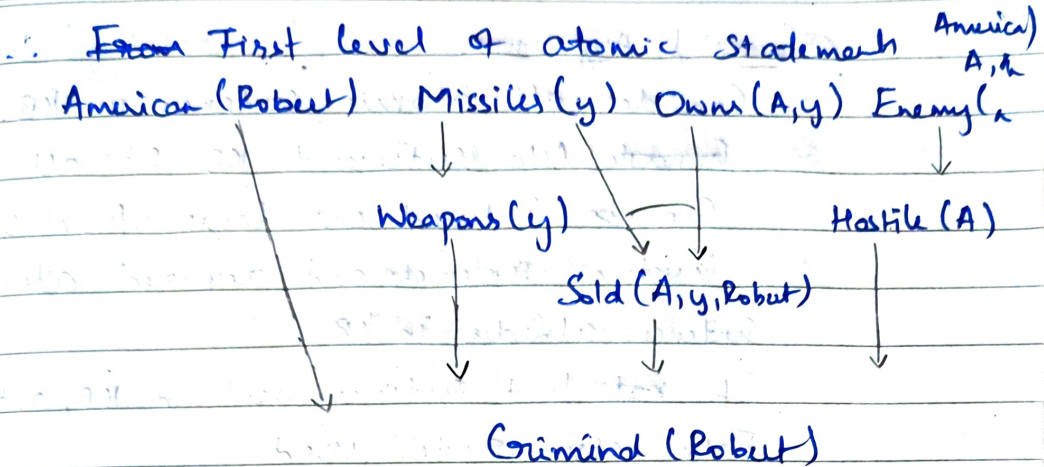
American(Robert)

We can deduce that missiles are weapons

$$\forall x \text{ Missile}(x) \rightarrow \text{Weapon}(x)$$

We can also deduce that an enemy of America is hostile

$$\forall x \text{ Enemy}(x, \text{America}) \Rightarrow \text{Hostile}(x)$$



Algorithm:

~~forward-chain~~ (KB, α)

Assuming KB represents all current knowledge

α represents atomic query to be validated

while new != {}

for rule in KB do (note)

$(p_1, p_2, \dots, p_n \rightarrow q) \Leftrightarrow$ Standardize-variable

for θ in new st

$(\theta, p_1, \dots, p_n) = (\theta, p_1, p_2, \dots, p_n)$

for p'_1, p'_2, \dots in KB

$q' = \text{subst}(\theta, q)$

if q' does not unify some sentence in KB or new

new += q'

$\phi = \text{Unify}(q', \alpha)$

if not ϕ return False

else return ϕ

~~new~~ += KB += new

return False

General Algorithm:

- Consider all atomic sentences in KB
- ~~Generate~~ Add sentences to KB s.t all can be inferred one layer from current level of sentences and required atomic sentences already in KB.
- If not find sentence not in KB, return False else rules used.

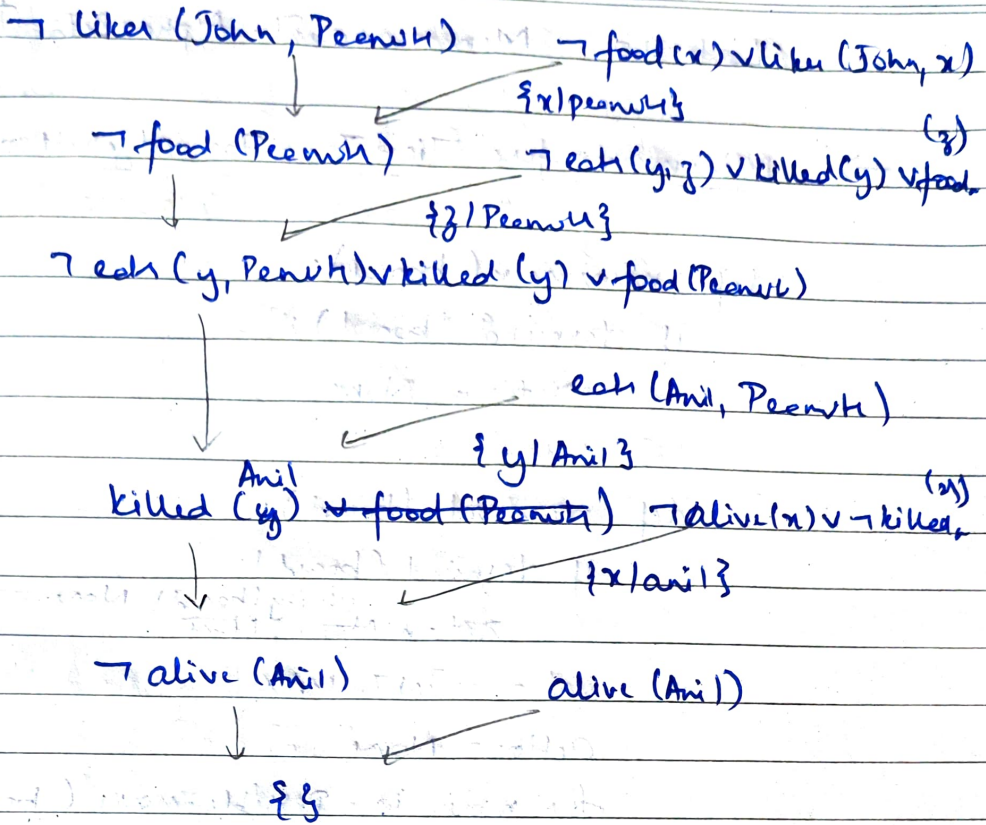
→ Resolution using FOL

Consider the following sentences:

- 1) $\text{food}(x) \rightarrow \text{likes}(\text{John}, x)$
- 2) $\text{food}(\text{Apple})$
- 3) $\text{food}(\text{Vegetables})$
- 4) $\neg \text{eats}(y, z) \wedge \neg \text{killed}(zy) \rightarrow \text{food}(z)$
- 5) $\text{eats}(\text{Anil}, \text{Peanuts})$
- 6) $\text{alive}(\text{Anil})$
- 6) $\text{killed}(x) \vee \text{alive}(x)$
- 7) $\neg \text{killed}(x) \vee \neg \text{alive}(x)$

To prove 4: ~~$\text{likes}(\text{John}, \text{Peanuts})$~~

Adding negation ($\neg \text{likes}(\text{John}, \text{Peanuts})$)
to KB and resolving



Results in empty set, therefore contradiction
 $\therefore \text{likes}(\text{John}, \text{Peanut})$ is true.