

180 Practical - 6

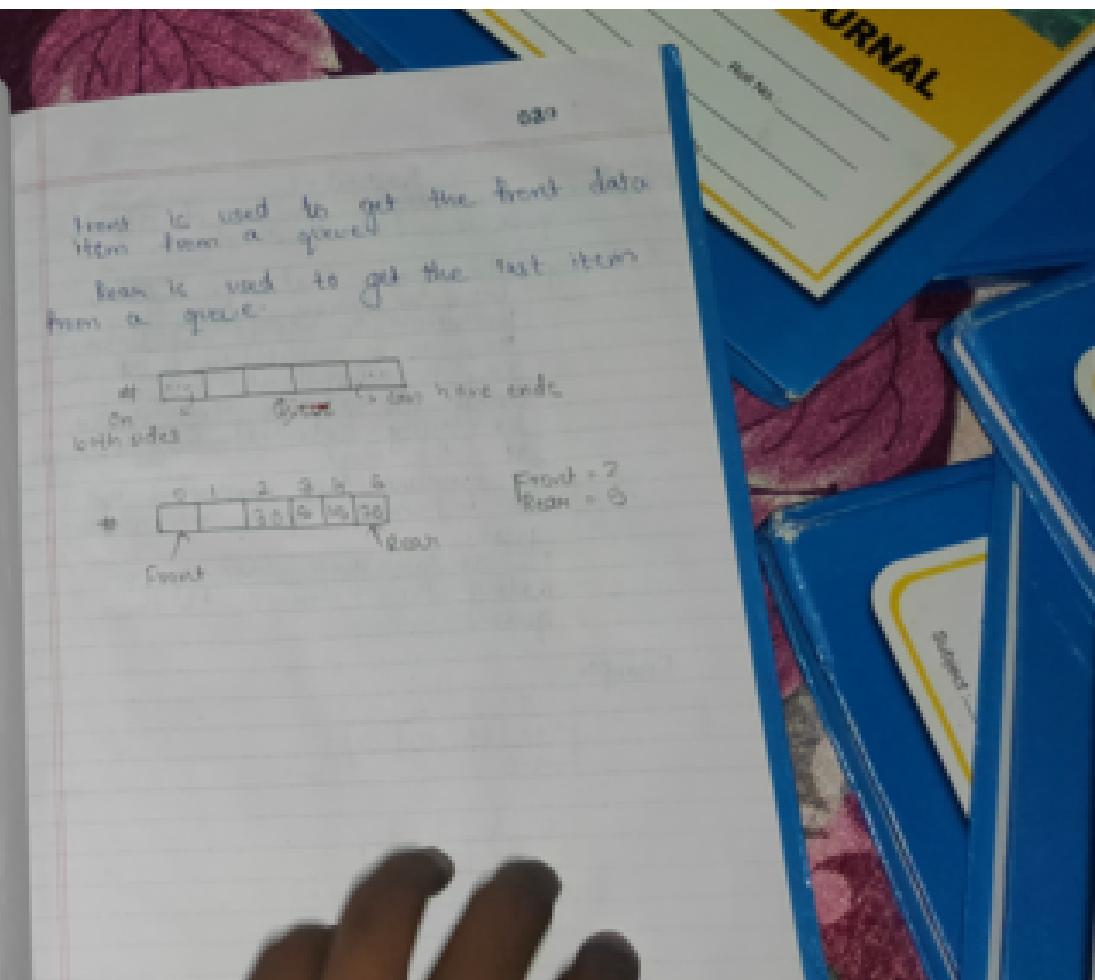
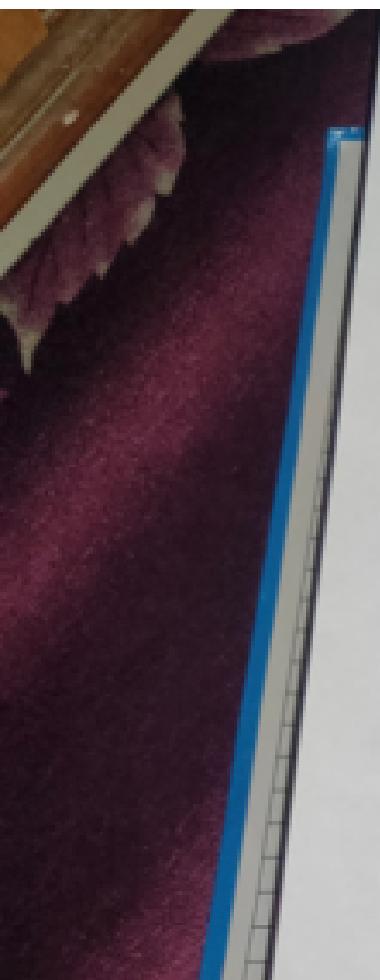
- * Aim To demonstrate Queue add and delete.
- * Theory: Queue is a linear data structure where the first element is inserted from one end called REAR and deleted from the other end called as FRONT.
Front points to the beginning of the queue and rear points to the end of the queue.
Queue follows the FIFO (First In First Out) structure.
According to its FIFO structure elements inserted first will also be removed first.
In a queue one end is always used to insert data (enqueue) and the other is used to delete data (dequeue) because queue is open at both of the ends.
Enqueue() can be termed as add()
in queue i.e. adding a element in queue.
Dequeue() can be termed as
delete or remove i.e. deleting or removing of element.

Program

```
#include <iostream.h>
#include <conio.h>
class Queue
{
public:
    int front;
    int rear;
    int queue[10];
    Queue()
    {
        front = -1;
        rear = -1;
    }
    void enqueue()
    {
        if (front == -1)
        {
            cout << "Queue is empty";
            return;
        }
        else
        {
            cout << "Enter element to be enqueued";
            cin >> queue[rear];
            rear++;
        }
    }
    void dequeue()
    {
        if (front == -1)
        {
            cout << "Queue is empty";
            return;
        }
        else
        {
            cout << "Element dequeued";
            cout << queue[front];
            front++;
        }
    }
};
```

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```
int main()
{
    Queue q;
    q.enqueue();
    q.enqueue();
    q.enqueue();
    q.enqueue();
    q.enqueue();
    q.dequeue();
    q.dequeue();
    q.dequeue();
    q.dequeue();
    q.dequeue();
}
```



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- * Practical - 6
How to demonstrate the use of queue in data structures.
- * Theory : The queue that we implement using an array suffer from the limitation. In what implementation there is a possibility that the queue is reported as full even though it's actually not. It can be empty state but there might be empty slots at the beginning of the queue.
To overcome this limitation we can implement queue as circular queue.
In circular queue we go on adding the element to the queue.

Example:

0	1	2	3
→ AA	BB	CC	DD

Front = 0

Rear = 3

0	1	2	3
	BB	CC	DD

Front = 1

Rear = 3

Practical

```
pushFront(20)  
pushFront(10)  
pop()  
push(1)  
push(2)  
delFront()  
delFront()  
push(3)  
push(4)  
push(5)  
delFront()  
delFront()  
push(6)  
push(7)  
push(8)  
push(9)  
delFront()  
delFront()  
delFront()  
delFront()  
delFront()
```

10

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```
def remove():
    global head
    if head == None:
        print("List is empty")
        exit(0)
    if head.next == None:
        print("List is empty")
        exit(0)
    else:
        print("Element removed")
        head = head.next
        print(head.data)
        print(head.next)
        print(head.next.next)
        print(head.next.next.next)
        print(head.next.next.next.next)
        print(head.next.next.next.next.next)
        print(head.next.next.next.next.next.next)
```

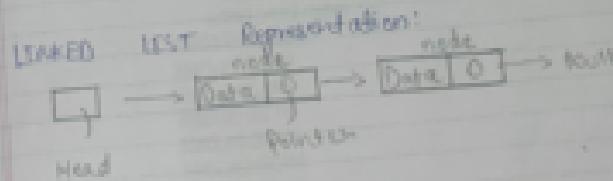
Output:

```
----- REPORT: C:\Users\Kapil\Downloads\Python\140.py -----
Primary Thread
data added: 44
data added: 33
data added: 44
data added: 44
data added: 44
data added: 33
data added: 33
data removed: 44
33
```

041

Practical-4

- * Aim:- To demonstrate the use of linked list in data structure.
- * Theory:- A linked data is a sequence of data structures. linked list is a sequence of links which contains data. Each link contains a connection to another link. Each link of a linked list can store a data called an element.
- * NEXT:- Each link of a linked list contains a link to the next link called NEXT.
- * FIRST:- A linked list contains the connection link to the first link called first.



10

Empirical observations:
1. Limited time.
2. Delays
3. Disruption
4. Unknown time
5. Delays

卷之三

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2018-2019
2019-2020
2020-2021
2021-2022
2022-2023
2023-2024

```

print(head.data)
head = head.next
print(head.data)
start.addlast(1)
start.addlast(2)
start.addlast(3)
start.addlast(4)
start.addlast(5)
start.addlast(6)
start.addlast(7)
start.addlast(8)
start.addlast(9)
print("Pranay Bawali(-17479")

```

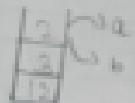
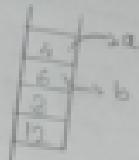
042

Practical - 8.

- To evaluate postfix expression using stack
- Theory: Stack is an (array) and works on LIFO (Last-In, First-Out) principle by operations.
- A postfix expression is a collection of operators and operands in which the operator is placed after the operands.
- Steps to be followed:
1. Read all the symbols one by one from left to right in the given postfix expression.
 2. If the reading symbol is operand then push it on to the stack.
 3. If the reading symbol is operation (+, -, /, *) then perform the pop operation and store the two popped operand in two different variables (operator is stored in global 2) and push result back onto the stack.
 4. Finally, perform a pop operation and display the popped value as final result.

1

Stage: 6-4 = + +



$$b_0 = g_0 - (g_1 + \frac{g_2}{2}) = 60$$

Report

Practical - 9

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- Practical

 - 4. Aim: To sort given random data by using bubble sort.
 - 5. Theory: Definition: A type in which any random data is sorted in ascending or descending order. Bubble sort considers element as an unitary sort.
 - 6. It is a simple sorting algorithm that repeatedly steps through the list, compares adjacent elements and swaps them if they are in wrong order. The pass through the list is repeated until the list is sorted. The algorithm which is a comparison sort is based on the way smaller or larger elements "bubble" to the top of the list.
 - 7. Working: The algorithm is simple. It is like this: we compare one item after another. If found false then swap them. Else move goes on.

at Example

$$\text{example: } \text{fact} \quad \text{push} \rightarrow (2 \leftarrow 3 \leftarrow 1)$$

$(1 \ 2 \ 3 \ 4) \rightarrow (2 \ 3 \ 4 \ 1)$

None algorithm compares the first

Here algorithm compares two segments and swaps since $b > 2$:
 $(2 \ 6 \ 5 \ 3 \ 9) \rightarrow (2 \ 5 \ 6 \ 3 \ 9)(2 \ 3 \ 6 \ 5 \ 9)$

Pass
 $(2, 3, 5, 6, 7) \rightarrow (2, 3, 6, 7)$ Swap 5, 6, 2.
 $(2, 3, 6, 7) \rightarrow (2, 3, 6, 7)$ Now the elements are already in order
 $(2, 3, 6, 7)$ algorithm doesn't swap them.

Second pass:
 $(2, 3, 6, 7) \rightarrow (2, 6, 3, 7)$
 $(2, 6, 3, 7) \rightarrow (2, 3, 6, 7)$ Swap 6, 3.
 $(2, 3, 6, 7) \rightarrow (2, 3, 6, 7)$

Third pass:
 $(2, 3, 6, 7)$ - 24 checks and give the data in sorted order.

Program :-
#include <iostream>
using namespace std;
int arr[5] = {2, 3, 5, 6, 7};
int n = 5;
int i, j, temp;
int main()
{
 for (i = 0; i < n; i++)
 {
 for (j = i + 1; j < n; j++)
 {
 if (arr[i] > arr[j])
 {
 temp = arr[i];
 arr[i] = arr[j];
 arr[j] = temp;
 }
 }
 }
 cout << "The sorted array is : " << endl;
 for (i = 0; i < n; i++)
 {
 cout << arr[i] << " ";
 }
 return 0;
}

Output :-

DESKTOP: 2023-07-10 10:57:21
1.07
Program Executed
Wall time: 0.016s
{2, 3, 5, 6, 7}
{2, 3, 5, 6, 7}
0.00

Java # Program

```
public class QuickSort {  
    public static void main(String[] args) {  
        int arr[] = {12, 45, 64, 32, 76, 13, 50, 80};  
        quickSort(arr, 0, arr.length - 1);  
        System.out.println("Sorted array:");  
        for (int i : arr) {  
            System.out.print(i + " ");  
        }  
    }  
  
    public static void quickSort(int arr[], int low, int high) {  
        if (low < high) {  
            int pivotIndex = partition(arr, low, high);  
            quickSort(arr, low, pivotIndex - 1);  
            quickSort(arr, pivotIndex + 1, high);  
        }  
    }  
  
    public static int partition(int arr[], int low, int high) {  
        int pivot = arr[high];  
        int i = low - 1;  
        for (int j = low; j < high; j++) {  
            if (arr[j] < pivot) {  
                i++;  
                swap(arr, i, j);  
            }  
        }  
        swap(arr, i + 1, high);  
        return i + 1;  
    }  
  
    public static void swap(int arr[], int i, int j) {  
        int temp = arr[i];  
        arr[i] = arr[j];  
        arr[j] = temp;  
    }  
}
```

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Theoretical - 10.

- a) How to 'partition' i.e. to sort the given data in Quick Sort.
- b) Theory: Quicksort is an efficient sorting algorithm type of a divide & conquer algorithm. It picks an element as pivot and partitions the given array around the pivot point. There are many different variants of quick sort that pick pivot in different ways.
 - i) Always pick first elements as pivot.
 - ii) Always pick last elements as pivot.
 - iii) Pick a random element as pivot.
 - iv) Pick median as pivot.

The key process in quicksort is partition(). Target of partition is given an array and an element X of array as pivot, put X at its current position in sorted array and put all smaller elements (smaller than X) before X, & put all greater elements (greater than X) after X. All this should be done in linear time.

048

• August 1 -

REMARKS: C. fuscus? Immature. Found in dense brushy vegetation.
1-07
Pawnee National
Grassland National
Park, CO, USA, 08, 08, 08, 08, 08
1988

JK

JURNAL

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Merge Sort

Aim:- To sort given data using merge sort.

Theory:- Merge sort uses the divide and conquer technique. In merge sort we divide the list into nearly equal subset such that each of which are then again divided into two subsets. -

* We continue this procedure until there is only one element in the individual subset.
Once we have individual elements in the subset, we consider these subset as sub problems which can be solved.

Since there is only one element in the subset, the subset is already sorted.

* Now while merging the sub problems, we compare the 2 subsets and create the combined list by comparing the element of the subset. After comparison we place the element in their correct position in the original subset.

W.C

90
JOURNAL

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Practical 12

Type : Binary Tree

Binary Search Tree is a tree which satisfies
any node within the tree for the
fact that any particular node can
have either 0 or 1 or 2 children.
Left Node : Block with no left
node has any children.

Internal Node : Node which do have
nodes.

Traversing can be defined as a
process of visiting every node of
the tree exactly once.

Preorder: ① Visit root node

② Traverse the left subtree. The
left subtree will always include
both left and right subtrees.

③ Traverse the right subtree.

The right subtree includes
right, then left and right
subtrees.

Inorder: ① Traverse left. The left subtree
might have right subtree.

② Visit root node

Postorder:

```
class Node
    int data
    Node left
    Node right

class Tree
    Node root

    void insert(int data)
    {
        if (root == null)
            root = new Node(data);
        else
        {
            Node current = root;
            while (true)
            {
                if (data < current.data)
                    if (current.left == null)
                        current.left = new Node(data);
                    else
                        current = current.left;
                else if (data > current.data)
                    if (current.right == null)
                        current.right = new Node(data);
                    else
                        current = current.right;
                else
                    break;
            }
        }
    }

    void printInorder(Node node)
    {
        if (node == null)
            return;
        printInorder(node.left);
        System.out.print(node.data + " ");
        printInorder(node.right);
    }

    void printPreorder(Node node)
    {
        if (node == null)
            return;
        System.out.print(node.data + " ");
        printPreorder(node.left);
        printPreorder(node.right);
    }

    void printPostorder(Node node)
    {
        if (node == null)
            return;
        printPostorder(node.left);
        printPostorder(node.right);
        System.out.print(node.data + " ");
    }
}
```

COLLEGE
INTER JOURNAL

10

④ Travel right after the right turn
right back left and right side

Reproduction: Measures the left subtrex.
The left subtrex when right
hand left and right subtrex
Measures right subtrex the right
subtrex might have left
and right subtrex
at first the root note

With the next reader

180 Practical No: 13

Aim: To simulate how to sort given random data by using selection sort.

Theory:- Selection sort is a simple sorting algorithm. This sorting algorithm is an in-place comparison-based algorithm in which the list is divided into two part the sorted part at the left end and the unsorted part at the right end initially, the unsorted part is empty and the sorted part is the entire list. The entire smallest element is taken from the unsorted array and placed with the leftmost element and the element becomes a part of sorted array. This process continues moving unsorted array boundary by one element to the right. This algorithm is not suitable for huge data set as its coverage and worst case complexity are $O(n^2)$, where n is the number of items.

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(a) Input:-

```
print("Please Enter the list (100)\n",end="")\narr=[23,31,39,18,17]\nprint()\n\narr[0]=min(arr[0])\nfor i in range(1,5):\n    if arr[i]<arr[0]:\n        arr[0]=arr[i]\n    else:\n        arr[1]=arr[i]\nprint()
```

(b) Output:-

```
Python 3.7.4 (tags/v3.7.4:e09359112e, Jul 9 2019, 19:04:11)\n[Clang 6.0 (clang-600.0.57)] on darwin\nType "help", "copyright", "credits" or "license" for more information.\n>>>\n>>> arr=[23,31,39,18,17]\n>>> arr[0]=min(arr[0])\n>>> arr\n[17, 23, 31, 18, 39]\n>>>
```

JOURNAL

150 Practical 2

- 1. How to search a number from the list using linear search.
- 2. Trace the process of identifying or finding a particular record if it's called sorted. There are two types of search:
 - a) Linear search
 - b) Binary searchThe linear search is further classified as:
 - a) Sorted & Unsorted.Here we will look on the UNSORTED linear search.
Linear Search, also known as sequential search, is a process that checks every element in the list sequentially until the desired element is found. When the elements to be searched are not specifically arranged in ascending or descending order, it is called unsorted linear search.
 - ↳ The data is stored in random manner.
 - ↳ One needs to specify the element to be searched in the sorted list.
 - ↳ Check the condition that whether the entered number matches if it matches then display the location else increment 1 as add 1 stored from location again.

```
Program  
-----  
#include<iostream.h>  
#include<conio.h>  
#include<math.h>  
#include<string.h>  
#include<limits.h>  
#include<stdio.h>  
#include<math.h>  
#include<conio.h>
```

10

Output:

```
System: Dell OptiPlex 7040, RAM: 8 GB (8000 MB), Processor: Intel(R) Core(TM) i5-10210U CPU @ 1.60 GHz, OS: Windows 10 Pro  
Type "exit", "copyright", "credits" or "license()" for more information.  
100  
----  
SEARCHING C:\Users\Bhupesh\Downloads\Search\Programs\Python\Python27\T.py  
----  
Process [1]:  
Enter the number to find:  
Number not found  
Number not found  
Number found  
100
```

Q3 Practical 2:

- a) How To search a number from the list using bubble sort method.
- b) Theory explanation and Pseudo code for different kinds or types of data structures.
Bubble sort - To basically sort the inputted data in ascending or descending manner.
Quicksort - To make elements tend to display the same.
In calculating that too in BINARY SEARCH to search the data is arranged in ascending to descending or descending to ascending that is all what it want by searching through sorted list is all of arranged data.

BINARY SEARCH

- b) The user is supposed to enter data in sorted manner.
- b) User has to give an element for comparing through sorted list.
- b) If element is found display with an update as value is sorted from large to small.
- b) If sorted order list of elements user can be check the condition that whether the entered number lies from starting point to the last element.

Program:

Q32

```
using System;
class Program
{
    static void Main(string[] args)
    {
        int[] arr = { 10, 20, 30, 40, 50 };
        int num;
        Console.WriteLine("Enter the number to search");
        num = int.Parse(Console.ReadLine());
        if (BinarySearch(arr, num))
            Console.WriteLine("Number found in list");
        else
            Console.WriteLine("Number not found");
    }

    static bool BinarySearch(int[] arr, int num)
    {
        int start = 0;
        int end = arr.Length - 1;
        while (start <= end)
        {
            int mid = (start + end) / 2;
            if (arr[mid] == num)
                return true;
            else if (arr[mid] < num)
                start = mid + 1;
            else
                end = mid - 1;
        }
        return false;
    }
}
```

Output:

```
----- OUTPUT: c:\Users\kamal\Desktop\Semester 2\Q32.cs -----
Enter 100
Enter the number to search
Number not found
100 |
```

023

is gg data or element not found quote the
same.

L&O Thailand 2

2. Now To switch a random item from the given array
by using lottery method.

4. Theory: A linear search also known as a left-internal search, is an algorithm used by computer science to locate a specified value(key) written on array. For the search to be binary, the array must be sorted in either descending or ascending order.

At each step of the algorithm a comparison is made and the process moves in one of two directions, specifically, the key value is compared to the middle element of the array.

If the key value is less than or greater than this middle element, the algorithm knows which half of the array to continue searching in because the array is sorted.

This process is repeated on smaller segments until the value is reached.

Because each step in the algorithm divides the search space in half a binary search will complete

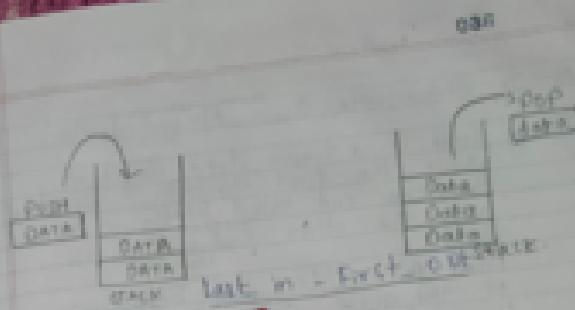
Output

RECEIVED: C:\Users\Bengt\Documents\Scanning\Scans\001-001
2000-01-01
FEBRUARY 1943
Lower number to be used from March 29
Number Found at Location 7
001
RECEIVED: C:\Users\Bengt\Documents\Scanning\Scans\001-001
2000-01-01
FEBRUARY 1943
Lower number to be used from March 29
Number not in range 1100
002
RECEIVED: C:\Users\Bengt\Documents\Scanning\Scans\001-001
2000-01-01
FEBRUARY 1943
Lower number to be used from March 29
Number not in range 1100
003

10.2 Stack(11)

- Aim To demonstrate the use of Stack.
- Theory In computer science, a stack is an abstract data type that stores all a collection of elements with two principal operations, push, which adds an element to the collection and pop, which removes the most recently added element that was just inserted.
The order may be LIFO (last in first out) or FILO (first in last out).
These basic operations are pushed in the stack.
- Push: Adds an item in the . If the Stack is full said to be over flow Condition.
- Pop: Removes an item from the stack. The last one pushed in the reversed order in which they are pushed.
If the stack is empty, then it is said to be an underflow condition.
- peek or top: Returns top element of stack.
- IsEmpty: Returns true if Stack is empty else false.

```
print("Pranay 17412")
class stack:
    global tos
    def __init__(self):
        self.tos=[0,0,0,0,0,0]
        self.tos=-1
    def push(self,data):
        n=len(self)
        if self.tos==n-1:
            print("stack is full")
        else:
            self.tos=self.tos+1
            self.tos[n]=data
```



Unit 10 - Page