

Introduction to Quantum Information and Communication

Take Home Mid-Sem

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Question 5

(a)

To Prove:

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z} = 2^n \delta(x, y)$$

Proof:

Case 1: $x = y$

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z}$$

$$\sum_{z \in \{0,1\}^n} (-1)^{0 \cdot z}$$

$$\sum_{z \in \{0,1\}^n} (-1)^0$$

$$\sum_{z \in \{0,1\}^n} 1$$

$$2^n$$

$$2^n \times 1$$

$$2^n \delta(x, y)$$

Case 2: $x \neq y$

Let k be the number of digits different between x and y , and let the corresponding indices be

$$\alpha = \{\alpha_1, \alpha_2, \alpha_3, \dots, \alpha_k\}$$

$$\forall i \in \{1, 2, \dots, k\} \ x_{\alpha_i} \neq y_{\alpha_i}$$

$$\forall i \notin \alpha \ x_i = y_i$$

.

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z}$$

$$\sum_{z \in \{0,1\}^n} (-1)^{\oplus_{i=1}^n (x_i \oplus y_i) z_i}$$

$$\sum_{z \in \{0,1\}^n} (-1)^{\bigoplus_{i=1}^k z_{\alpha_i}}$$

$$\sum_{z \in \{0,1\}^n} (-1)^{z_{\alpha_1} \oplus z_{\alpha_2} \oplus \dots \oplus z_{\alpha_k}}$$

Now, since z is looping through all possible bitstrings of length n , the parity of any subset of its bits will be odd half the times and even half the times.

$$-1 + 1 - 1 + 1 \dots - 1 + 1$$

$$0$$

$$2^n \times 0$$

$$2^n \delta(x, y)$$

Now, from both the cases we get

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z} = 2^n \delta(x, y)$$

Hence, proven

(b)

Given:

$$f : \{0, 1\}^n \mapsto \{0, 1\}^n$$

$$U_f(|x\rangle_Q \otimes |y\rangle_R) := |x\rangle_Q \otimes |y \oplus f(x)\rangle_R$$

$$V_f(|x\rangle_Q \otimes |y\rangle_R) := (-1)^{y \cdot f(x)} |x\rangle_Q \otimes |y\rangle_R$$

To Prove:

$$V_f(|x\rangle_Q \otimes |y\rangle_R) = (\mathbb{I}_Q \otimes H^{\otimes n}) U_f(\mathbb{I}_Q \otimes H^{\otimes n}) (|x\rangle_Q \otimes |y\rangle_R)$$

Proof: We will be using the identity $H^{\otimes n}|x\rangle = \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{x \cdot z} |z\rangle$

$$(\mathbb{I}_Q \otimes H^{\otimes n}) U_f(\mathbb{I}_Q \otimes H^{\otimes n}) (|x\rangle_Q \otimes |y\rangle_R)$$

$$(\mathbb{I}_Q \otimes H^{\otimes n}) U_f(\mathbb{I}_Q |x\rangle_Q \otimes H^{\otimes n} |y\rangle_R)$$

$$(\mathbb{I}_Q \otimes H^{\otimes n}) U_f \left(|x\rangle_Q \otimes \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} |z\rangle_R \right)$$

$$\frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} (\mathbb{I}_Q \otimes H^{\otimes n}) U_f (|x\rangle_Q \otimes |z\rangle_R)$$

$$\frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} (\mathbb{I}_Q \otimes H^{\otimes n}) (|x\rangle_Q \otimes |z \oplus f(x)\rangle_R)$$

$$\begin{aligned}
& \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} \left(\mathbb{I}_Q |x\rangle_Q \otimes H^{\otimes n} |z \oplus f(x)\rangle_R \right) \\
& \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} |x\rangle_Q \otimes \left(\frac{1}{\sqrt{2^n}} \sum_{w \in \{0,1\}^n} (-1)^{(z \oplus f(x)) \cdot w} |w\rangle_R \right) \\
& \frac{1}{2^n} \sum_{z, w \in \{0,1\}^n} (-1)^{(y \cdot z)} (-1)^{(z \oplus f(x)) \cdot w} |x\rangle_Q \otimes |w\rangle_R \\
& \frac{1}{2^n} \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} |x\rangle_Q \otimes |w\rangle_R \sum_{z \in \{0,1\}^n} (-1)^{(y \oplus w) \cdot z} \\
& \frac{1}{2^n} \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} |x\rangle_Q \otimes |w\rangle_R 2^n \delta(w, y) \\
& \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} |x\rangle_Q \otimes |w\rangle_R \delta(w, y) \\
& (-1)^{y \cdot f(x)} |x\rangle_Q \otimes |y\rangle_R \\
& V_f \left(|x\rangle_Q \otimes |y\rangle_R \right)
\end{aligned}$$

Hence, proven

Question 7

Given:

$$f : \{0, 1\}^n \mapsto \{0, 1\}^n$$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n, d\}$$

$$U_f \left(|x\rangle_Q \otimes |y\rangle_R \right) := |x\rangle_Q \otimes |y \oplus f(x)\rangle_R$$

(a)

To Prove: f is one-to-one when $d = 0^n$ and two-to-one otherwise

Proof:

Case 1: $d = 0^n$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n, 0^n\}$$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n\}$$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x \oplus y = 0^n$$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x = y$$

Thus, f is one-one in this case

Case 2: $d \neq 0^n$

To prove that f is two-one, we need to show that $\forall z \in \text{range}(f)$, we have exactly two elements x, y such that $f(x) = f(y) = z$

$$\forall x, y \in \{0, 1\}^n \quad f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n, d\}$$

(i) $x \oplus y = 0^n$

$x = y$, thus $f(x) = f(y)$

(ii) $x \oplus y = d$ with $d \neq 0^n$

$$y = d \oplus x$$

Since $d \neq 0^n$, we get $y \neq x$, and $f(x) = f(y)$

Clearly, two distinct values x and y give the same output. Now, we need to prove that no more than two distinct inputs give the same output.

Consider distinct $a, b, c \in \{0, 1\}^n$ such that $f(a) = f(b) = f(c)$

Since a, b, c are distinct, their xor cannot be 0^d , thus we have

$$a \oplus b = b \oplus c = d$$

$$a = d \oplus b, c = d \oplus b$$

$$a = c$$

This is a contradiction. Thus, there only exist exactly two input values for each output value.

Thus, f is a two-one function in this case

Hence, proven