Introduction to Quantum Information and Communication

Take Home Mid-Sem

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Question 5

(a)

To Prove:

$$\sum_{z\in\left\{ 0,1\right\} ^{n}}\left(-1\right) ^{\left(x\oplus y\right) \cdot z}=2^{n}\delta(x,y)$$

Proof:

Case 1: x = y

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z}$$

$$\sum_{z \in \{0,1\}^n} (-1)^{0 \cdot z}$$

$$\sum_{z \in \{0,1\}^n} (-1)^0$$

$$\sum_{z \in \{0,1\}^n} 1$$

$$2^n$$

$$2^n \times 1$$

$$2^n \delta(x,y)$$

Case 2: $x \neq y$

Let k be the number of digits different between x and y, and let the corresponding indices be $\alpha = \{\alpha_1, \alpha_2, \alpha_3, ..., \alpha_k\}$

$$\forall i \in \{1,2,...,k\} \ x_{\alpha_i} \neq y_{\alpha_i}$$

$$\forall i \notin \alpha \ x_i = y_i$$

.

$$\sum_{z \in \{0,1\}^n} (-1)^{(x \oplus y) \cdot z} \\ \sum_{z \in \{0,1\}^n} (-1)^{\bigoplus_{i=1}^n (x_i \oplus y_i) z_i}$$

$$\sum_{z \in \left\{0,1\right\}^n} \left(-1\right)^{\bigoplus_{i=1}^k z_{\alpha_i}}$$

$$\sum_{z \in \left\{0,1\right\}^n} \left(-1\right)^{z_{\alpha_1} \oplus z_{\alpha_2} \oplus \ldots \oplus z_{\alpha_k}}$$

Now, since z is looping through all possible bitstrings of length n, the parity of any subset of its bits will be odd half the times and even half the times.

$$-1+1-1+1...-1+1$$

$$0$$

$$2^{n}\times 0$$

$$2^{n}\delta(x,y)$$

Now, from both the cases we get

$$\sum_{z\in\left\{ 0,1\right\} ^{n}}\left(-1\right) ^{\left(x\oplus y\right) \cdot z}=2^{n}\delta(x,y)$$

Hence, proven

(b)

Given:

$$\begin{split} f: \left\{0,1\right\}^n &\mapsto \left\{0,1\right\}^n \\ U_f\Big(\left|x\right\rangle_Q \otimes \left|y\right\rangle_R\Big) &\coloneqq \left|x\right\rangle_Q \otimes \left|y \oplus f(x)\right\rangle_R \\ \\ V_f\Big(\left|x\right\rangle_Q \otimes \left|y\right\rangle_R\Big) &\coloneqq (-1)^{y \cdot f(x)} |x\right\rangle_Q \otimes \left|y\right\rangle_R \end{split}$$

To Prove:

$$V_f \Big(\left| x \right\rangle_Q \otimes \left| y \right\rangle_R \Big) = \Big(\mathbb{I}_Q \otimes H^{\otimes n} \Big) U_f \Big(\mathbb{I}_Q \otimes H^{\otimes n} \Big) \Big(\left| x \right\rangle_Q \otimes \left| y \right\rangle_R \Big)$$

 $\left(\mathbb{I}_{Q}\otimes H^{\otimes n}\right)U_{f}\left(\mathbb{I}_{Q}\otimes H^{\otimes n}\right)\left(\left|x\right\rangle_{Q}\otimes\left|y\right\rangle_{R}\right)$

Proof: We will be using the identity $H^{\otimes n}|x\rangle=\frac{1}{\sqrt{2^n}}\sum_{z\in\{0,1\}^n}{(-1)}^{x\cdot z}|z\rangle$

$$\begin{split} & \left(\mathbb{I}_{Q} \otimes H^{\otimes n}\right) U_{f} \bigg(\mathbb{I}_{Q} |x\rangle_{Q} \otimes H^{\otimes n} |y\rangle_{R} \bigg) \\ & \left(\mathbb{I}_{Q} \otimes H^{\otimes n}\right) U_{f} \Bigg(|x\rangle_{Q} \otimes \frac{1}{\sqrt{2^{n}}} \sum_{z \in \{0,1\}^{n}} \left(-1\right)^{y \cdot z} |z\rangle_{R} \bigg) \end{split}$$

$$\frac{1}{\sqrt{2^{n}}} \sum_{z \in I_{0,1}\backslash^{n}} \left(-1\right)^{y \cdot z} \left(\mathbb{I}_{Q} \otimes H^{\otimes n}\right) U_{f} \left(\left|x\right\rangle_{Q} \otimes \left|z\right\rangle_{R}\right)$$

$$\frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} \left(-1\right)^{y \cdot z} \left(\mathbb{I}_Q \otimes H^{\otimes n}\right) \left(\left|x\right\rangle_Q \otimes \left|z \oplus f(x)\right\rangle_R\right)$$

$$\begin{split} \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} \Big(\mathbb{I}_Q | x \rangle_Q \otimes H^{\otimes n} | z \oplus f(x) \rangle_R \Big) \\ \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} (-1)^{y \cdot z} | x \rangle_Q \otimes \Bigg(\frac{1}{\sqrt{2^n}} \sum_{w \in \{0,1\}^n} (-1)^{(z \oplus f(x)) \cdot w} | w \rangle_R \Bigg) \\ \frac{1}{2^n} \sum_{z,w \in \{0,1\}^n} (-1)^{(y \cdot z)} (-1)^{(z \oplus f(x)) \cdot w} | x \rangle_Q \otimes | w \rangle_R \\ \frac{1}{2^n} \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} | x \rangle_Q \otimes | w \rangle_R \sum_{z \in \{0,1\}^n} (-1)^{(y \oplus w) \cdot z} \\ \frac{1}{2^n} \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} | x \rangle_Q \otimes | w \rangle_R 2^n \delta(w,y) \\ \sum_{w \in \{0,1\}^n} (-1)^{w \cdot f(x)} | x \rangle_Q \otimes | w \rangle_R \delta(w,y) \\ (-1)^{y \cdot f(x)} | x \rangle_Q \otimes | y \rangle_R \\ V_f \Big(| x \rangle_Q \otimes | y \rangle_R \Big) \end{split}$$

Hence, proven

Question 6

(a)

Before the first Hadamard, the state is

$$|0\rangle_A |\psi\rangle_B |\varphi\rangle_C$$

After the first Hadamard, the state is

$$\begin{split} H_A |0\rangle_A |\psi\rangle_B |\varphi\rangle_C \\ \frac{1}{\sqrt{2}} \Big(|0\rangle_A + |1\rangle_A \Big) |\psi\rangle_B |\varphi\rangle_C \\ \frac{1}{\sqrt{2}} |0\rangle_A |\psi\rangle_B |\varphi\rangle_C + \frac{1}{\sqrt{2}} |1\rangle_A |\psi\rangle_B |\varphi\rangle_C \end{split}$$

After the Controlled-SWAP, the state is

$$\frac{1}{\sqrt{2}}|0\rangle_A|\psi\rangle_B|\varphi\rangle_C + \frac{1}{\sqrt{2}}|1\rangle_A|\varphi\rangle_B|\psi\rangle_C$$

After the second Hadamard, we get the required state

$$\left. |\psi'\rangle_{ABC} = H_A \bigg(\frac{1}{\sqrt{2}} |0\rangle_A |\psi\rangle_B |\varphi\rangle_C + \frac{1}{\sqrt{2}} |1\rangle_A |\varphi\rangle_B |\psi\rangle_C \bigg)$$

$$\begin{split} |\psi'\rangle_{ABC} &= \frac{1}{\sqrt{2}} H_A |0\rangle_A |\psi\rangle_B |\varphi\rangle_C + \frac{1}{\sqrt{2}} H_A |1\rangle_A |\varphi\rangle_B |\psi\rangle_C \\ |\psi'\rangle_{ABC} &= \frac{1}{2} \Big(|0\rangle_A + |1\rangle_A \Big) |\psi\rangle_B |\varphi\rangle_C + \frac{1}{2} \Big(|0\rangle_A - |1\rangle_A \Big) |\varphi\rangle_B |\psi\rangle_C \\ |\psi'\rangle_{ABC} &= \frac{1}{2} |0\rangle_A \Big(|\psi\rangle_B |\varphi\rangle_C + |\varphi\rangle_B |\psi\rangle_C \Big) + \frac{1}{2} |1\rangle_A \Big(|\psi\rangle_B |\varphi\rangle_C - |\varphi\rangle_B |\psi\rangle_C \Big) \end{split}$$

This is the required tripartite state

(b)

$$\begin{split} p_0 &= \frac{1}{2} \Big(\langle \psi |_B \langle \varphi |_C + \langle \varphi |_B \langle \psi |_C \Big) \frac{1}{2} \Big(|\psi \rangle_B |\varphi \rangle_C + |\varphi \rangle_B |\psi \rangle_C \Big) \\ p_0 &= \frac{1}{4} \Big(\langle \psi |_B \langle \varphi |_C |\psi \rangle_B |\varphi \rangle_C + \langle \psi |_B \langle \varphi |_C |\varphi \rangle_B |\psi \rangle_C + \langle \varphi |_B \langle \psi |_C |\psi \rangle_B |\varphi \rangle_C + \langle \varphi |_B \langle \psi |_C |\varphi \rangle_B |\psi \rangle_C \Big) \\ p_0 &= \frac{1}{4} \Big(\langle \psi |\psi \rangle_B \otimes \langle \varphi |\varphi \rangle_C + \langle \psi |\varphi \rangle_B \otimes \langle \varphi |\psi \rangle_C + \langle \varphi |\psi \rangle_B \otimes \langle \psi |\varphi \rangle_C + \langle \varphi |\varphi \rangle_B \otimes \langle \psi |\psi \rangle_C \Big) \\ p_0 &= \frac{1}{4} \Big(1 + |\langle \psi |\varphi \rangle|^2 + |\langle \psi |\varphi \rangle|^2 + 1 \Big) \end{split}$$

Since $p_0 + p_1 = 1$,

$$p_1 = \frac{1}{2} - \frac{1}{2} |\langle \psi | \varphi \rangle|^2$$

Question 7

Given:

$$\begin{split} f:\left\{0,1\right\}^n &\mapsto \left\{0,1\right\}^n \\ \forall x,y \in \left\{0,1\right\}^n & f(x) = f(y) \leftrightarrow x \oplus y \in \left\{0^n,d\right\} \\ & U_f\Big(\left|x\right\rangle_Q \otimes \left|y\right\rangle_R\Big) \coloneqq \left|x\right\rangle_Q \otimes \left|y \oplus f(x)\right\rangle_R \end{split}$$

(a)

To Prove: f is one-to-one when $d = 0^n$ and two-to-one otherwise

Proof:

Case 1: $d = 0^n$

$$\forall x, y \in \{0, 1\}^n \ f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n, 0^n\}$$
$$\forall x, y \in \{0, 1\}^n \ f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n\}$$
$$\forall x, y \in \{0, 1\}^n \ f(x) = f(y) \leftrightarrow x \oplus y = 0^n$$
$$\forall x, y \in \{0, 1\}^n \ f(x) = f(y) \leftrightarrow x = y$$

Thus, f is one-one in this case

Case 2: $d \neq 0^n$

To prove that f is two-one, we need to show that $\forall z \in \text{range}(f)$, we have exactly two elements x, y such that f(x) = f(y) = z

$$\forall x, y \in \{0, 1\}^n \ f(x) = f(y) \leftrightarrow x \oplus y \in \{0^n, d\}$$

(i) $x \oplus y = 0^n$

x = y, thus f(x) = f(y)

(ii) $x \oplus y = d$ with $d \neq 0^n$

$$y = d \oplus x$$

Since $d \neq 0^n$, we get $y \neq x$, and f(x) = f(y)

Clearly, two distinct values x and y give the same output. Now, we need to prove that no more than two distinct inputs give the same output.

Consider distinct $a, b, c \in \{0, 1\}^n$ such that f(a) = f(b) = f(c)

Since a, b, c are distinct, their xor cannot be 0^d , thus we have

$$a \oplus b = b \oplus c = d$$

$$a = d \oplus b, c = d \oplus b$$

$$a = c$$

This is a contradiction. Thus, there only exist exactly two input values for each output value.

Thus, f is a two-one function in this case

Hence, proven

(b)

To Find: $\left|\psi'\right\rangle_{QR}$

Solution:

Initially, the state is

$$\left|0^{n}\right\rangle_{Q}\otimes\left|O^{n}\right\rangle_{R}$$

After the first Hadamard, the state is

$$H^{\otimes n}|0^n\rangle_Q\otimes|0^n\rangle_R$$

$$\frac{1}{\sqrt{2^n}} \sum_{x \in \left\{0,1\right\}^n} \left| x \right\rangle_Q \otimes \left| 0^n \right\rangle_R$$

After the oracle, the state is

$$U_f \frac{1}{\sqrt{2^n}} \sum_{x \in \left\{0,1\right\}^n} \left| x \right\rangle_Q \otimes \left| 0^n \right\rangle_R$$

$$\begin{split} &\frac{1}{\sqrt{2^n}} \sum_{x \in \left\{0,1\right\}^n} U_f \big| x \big\rangle_Q \otimes \big| 0^n \big\rangle_R \\ &\frac{1}{\sqrt{2^n}} \sum_{x \in \left\{0,1\right\}^n} \big| x \big\rangle_Q \otimes \big| 0^n \oplus f(x) \big\rangle_R \\ &\frac{1}{\sqrt{2^n}} \sum_{x \in \left\{0,1\right\}^n} \big| x \big\rangle_Q \otimes \big| f(x) \big\rangle_R \end{split}$$

After the second Hadamard, the required state is

$$\begin{split} \left|\psi'\right\rangle_{QR} &= H^{\otimes n} \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} \left|x\right\rangle_Q \otimes \left|f(x)\right\rangle_R \\ \left|\psi'\right\rangle_{QR} &= \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} H^{\otimes n} |x\rangle_Q \otimes \left|f(x)\right\rangle_R \\ \left|\psi'\right\rangle_{QR} &= \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} \frac{1}{\sqrt{2^n}} \sum_{z \in \{0,1\}^n} \left(-1\right)^{x \cdot z} |z\rangle_Q \otimes \left|f(x)\right\rangle_R \\ \left|\psi'\right\rangle_{QR} &= \frac{1}{2^n} \sum_{x,z \in \{0,1\}^n} \left(-1\right)^{x \cdot z} |z\rangle_Q \otimes \left|f(x)\right\rangle_R \end{split}$$