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1(a)

In this task I created a list of total inputs in next_row and sum in another list name Sum. After testing every element to get the required sum in a nested loop we finally can find it by a if condition and to keep track of where the pointers were we use two pointers. If we can find the sum we simply write 'impossible'. And in this code the time complexity is $O(n^2)$.

1(b)

In this task I started by setting two pointers left and right at the first and last index respectively of a sorted list. Then in the initialization I set as the left pointer is less than total number of input and right pointer is greater than 0 to check the sum. Then if the sum of elements at left and right pointers equal to sum we find the solution and record the position in check1 and 2 and break the loop. Next I adjusted if the sum is less than left then the pointer will move one step right and if more than right right will move one step left. Then I find the solution.

if not found then IMPOSSIBLE. In this code
time complexity is $O(n)$.

2(a)

In this task firstly I read the length of Alice and bobs
and also both the list of inputs, arr1 and arr2
one from 2 lists. Then I created a empty list to
store all elements. I converted each elements to
int in the list lst. Using a sort function I sorted
in ascending order, and thus I get the output
and hence the time complexity is $O(n \log n)$.

2(b)

In this task I read the length of Alice and bobs
list and then split to get arr1 and arr2
two lists containing 2 sets of inputs. After that I
created a empty list to store all outputs and created
two pointers for both arr1 and arr2. Then I used
a while loop which will run until both the pointers
reach the end of their respective list. Then I used

a comparison by. Conditional statement where they were being appended to. list and increasing their pointer value as well in each iteration.

After the loop ends I added the remaining elements of both list to the merged list and get the final output here Time complexity is $O(n)$.

(3)

by using greedy algorithm and In this task, by read line I take the total number input and used in for loop as a range and for every iteration a group of numbers were being appended to an empty list. Then I used bubble sort to sort them according to their end time.

Then I set a start time and end time as a 0th index of list and set a counter and and from list, empty list. Then I appended the start and end time. Then I run another loop from 1 to len(list) and get a of out condition was satisfied it will be append to list 1 and out output is list 1.

(4)

In this task I again used greedy algorithm. Firstly in my code I read the file to find how many ^{work} ~~time~~ are there and how many people. Each work is sorted by Merge Sort based on their end time. ~~the~~ Then I iterate the sorted tasks, for each task the code will find a earliest ~~time~~ available person who can complete it without overlapping with other tasks assigned to them. If a suitable person is found, then the task was assigned to him and the count was increased by 1 each time and. finally we get the output.

Brainstorming

To get $O(n \log n)$ for task 4 I will use the following steps:

① Sort by end time: I will start by sorting the tasks based on their end times. This operation will take $O(n \log n)$ time.

② Greedy Task Selection: I iterate through the sorted tasks at each step I select the task with the earliest end time that doesn't overlap with other tasks (previous). This can be done efficiently since the tasks are sorted by end time.

③ Counting: Add a counter to keep track of the selected tasks

④ Output: Finally I get the desired output

So from this steps it can be ensured that I will always assign the next activity to the person who becomes available first and by sorting the activities by their end times, I can ensure that I am always looking at the earliest possible end time for each location. And so we get a $O(n \log n)$.