

# Database Concepts

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# Objective

- Need for a Database
- File – base systems
- Define DBMS
- Features of DBMS
- Use of DBMS
- RDBMS Basics



# Data vs. information

## What is data?

- Facts, statistics used for reference or analysis.
- Numbers, characters, symbols, images etc., which can be processed by a computer.
- Data must be interpreted, by a human or machine, to derive meaning
- Latin 'datum' meaning "that which is given"

## What is information?

- Information is used to reveal the meaning of data.
- Knowledge derived from study, experience (by the senses), or instruction.
- Communication of intelligence.
- Information is interpreted data

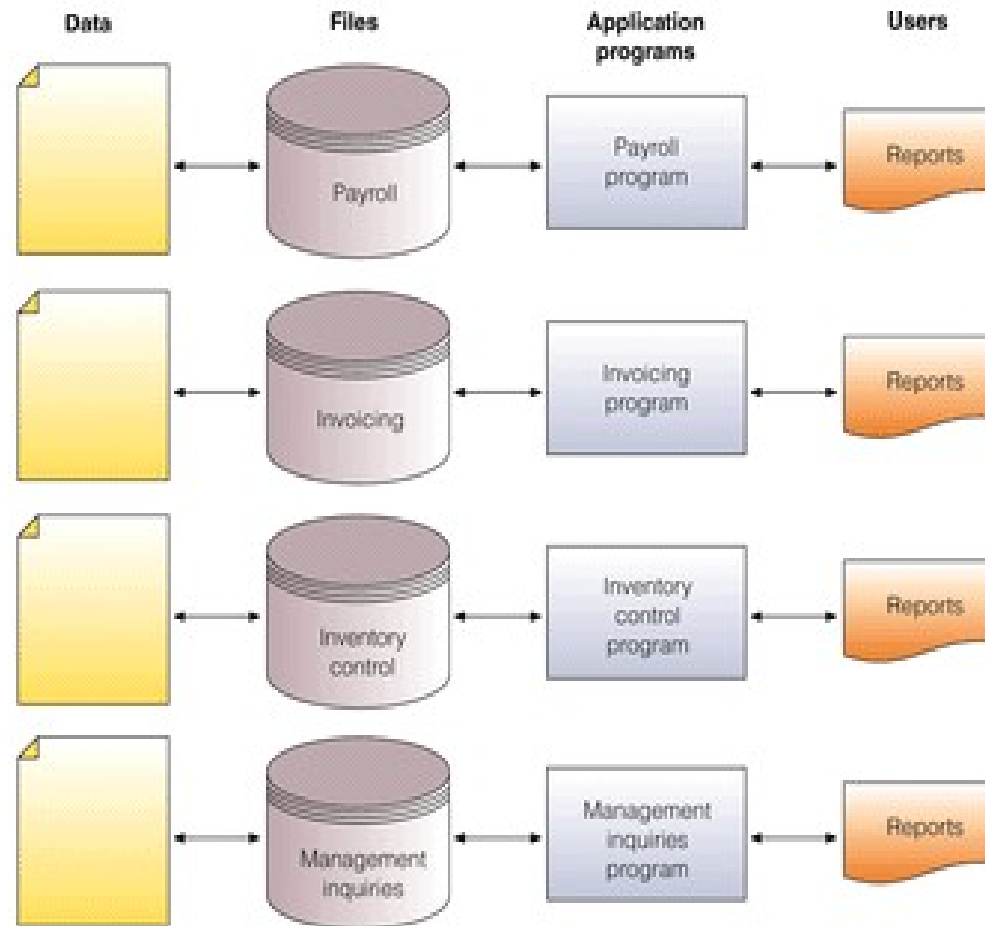


# Why do we need a database?

- Keep records of our:
  - Clients
  - Staff
  - Volunteers
- To keep a record of activities and interventions
- Keep sales records
- Develop reports
- Perform research
- Longitudinal tracking

# File Based Approach

Every program stores and maintains data in separate files





# Limitations – File Based approach


- Separation and isolation of data
- Duplication of data
- Data dependence
- Incompatibility of files
- Fixed queries / proliferation of application programs

# What is DBMS?

- A very large, integrated collection of data.
- Models real-world **enterprise**.
  - Entities (e.g., students, courses)
  - Relationships (e.g., John is taking CS662)
- A **Database Management System** (DBMS) is a software package designed to store and manage databases.



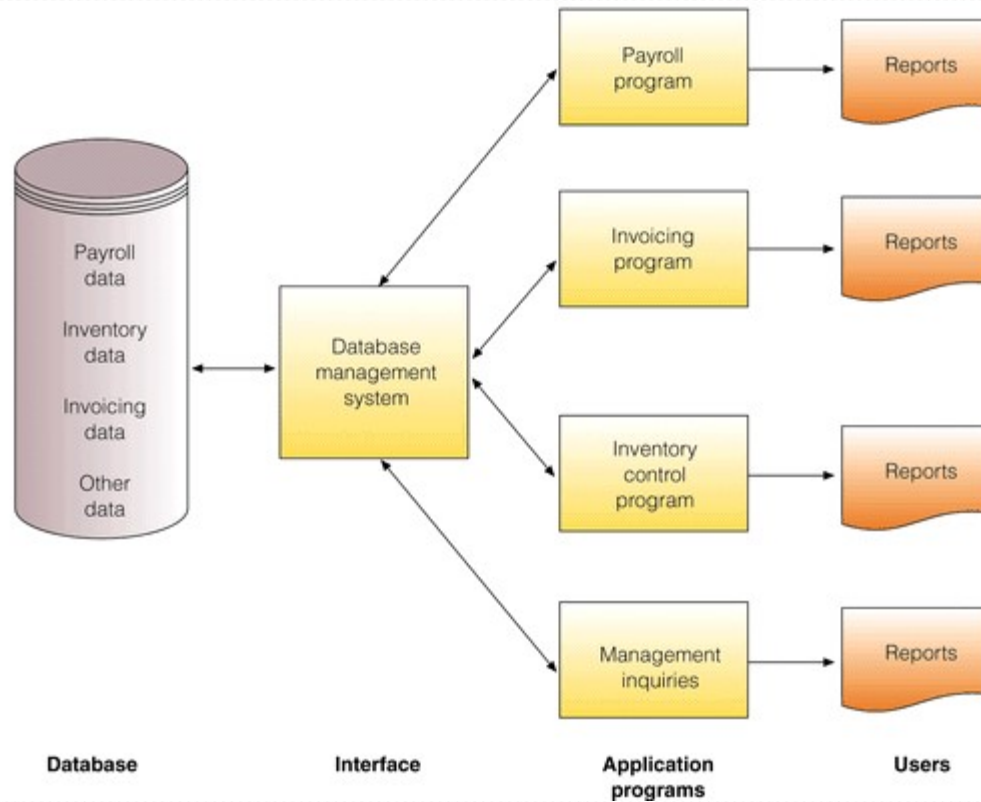
# Purpose of a DBMS

- Is to transform 

```
graph LR; Data[Data] --> Information[Information]; Information --> Knowledge[Knowledge]; Knowledge --> Action[Action];
```
- Data independence and efficient access.
- Data integrity and security.
- Uniform data administration.
- Concurrent access, recovery from crashes.
- Replication control
- Reduced application development time.



# The Database Approach to Manage Data



# Advantages of DBMS

- Control of data redundancy
- Data consistency
- Provides multiple views of the same data
- Sharing of data
- Improved data integrity
- Improved security
- Enforcement of standards
- Improved data accessibility and responsiveness
- Improved productivity & maintenance
- Improved backing and recovery services



# Limitations – DBMS

- Relatively high cost for purchase, operation & maintaining DBMS
- Size (Substantial space in main memory & large disk space.
- Complexity
- Higher impact of a failure

# Database Users

- Database administrators(DBAs):
- Database designers
- End Users
  - Casual
  - Parametric (or naïve)
  - Sophisticated
  - Stand-alone





# Relational Database model

- Developed by E.F. Codd, C.J. Date (70s)
- Table = Entity = Relation
- Table row = tuple = instance
- Table column = attribute
- Table linkage by values
- Entity-Relationship Model



# The Relational Model

- Each attribute has a unique name within an entity
- All entries in the column are examples of it
- Each row is unique
- Ordering of rows and columns is unimportant
- Each position (tuple) is limited to a single entry.



# RDBMS Definitions

- Entity: Object, Concept or event (subject)
- Attribute: a Characteristic of an entity
- Row or Record: the specific characteristics of one entity
- Table: a collection of records
- Database: a collection of tables

# RDBMS Basics

- Consider some information the University maintains:
  - Name
  - Major
  - Tuition Paid
  - Address
  - Courses Taken
  - Tuition Owed
  - SSN
  - Grades Received
  - Grants/Scholarships
- *HOW* is this information stored?

You are an entity with attributes which vary. Within the University, different areas have different interests in you (i.e., the Registrar, the Bursar, etc.). Nonetheless, you are still part of the University as a whole.

- *HOW* does this relate to a database?



# RDBMS Basics

## *Database Components:*

You are an entity



Record

with attributes



Fields

which vary



Fields can contain characters, numbers, symbols, etc.

Within the University, different areas,  
have different interests in you



Files

(i.e.,. The Registrar, Bursar, etc.)

Nonetheless, you are still part of the

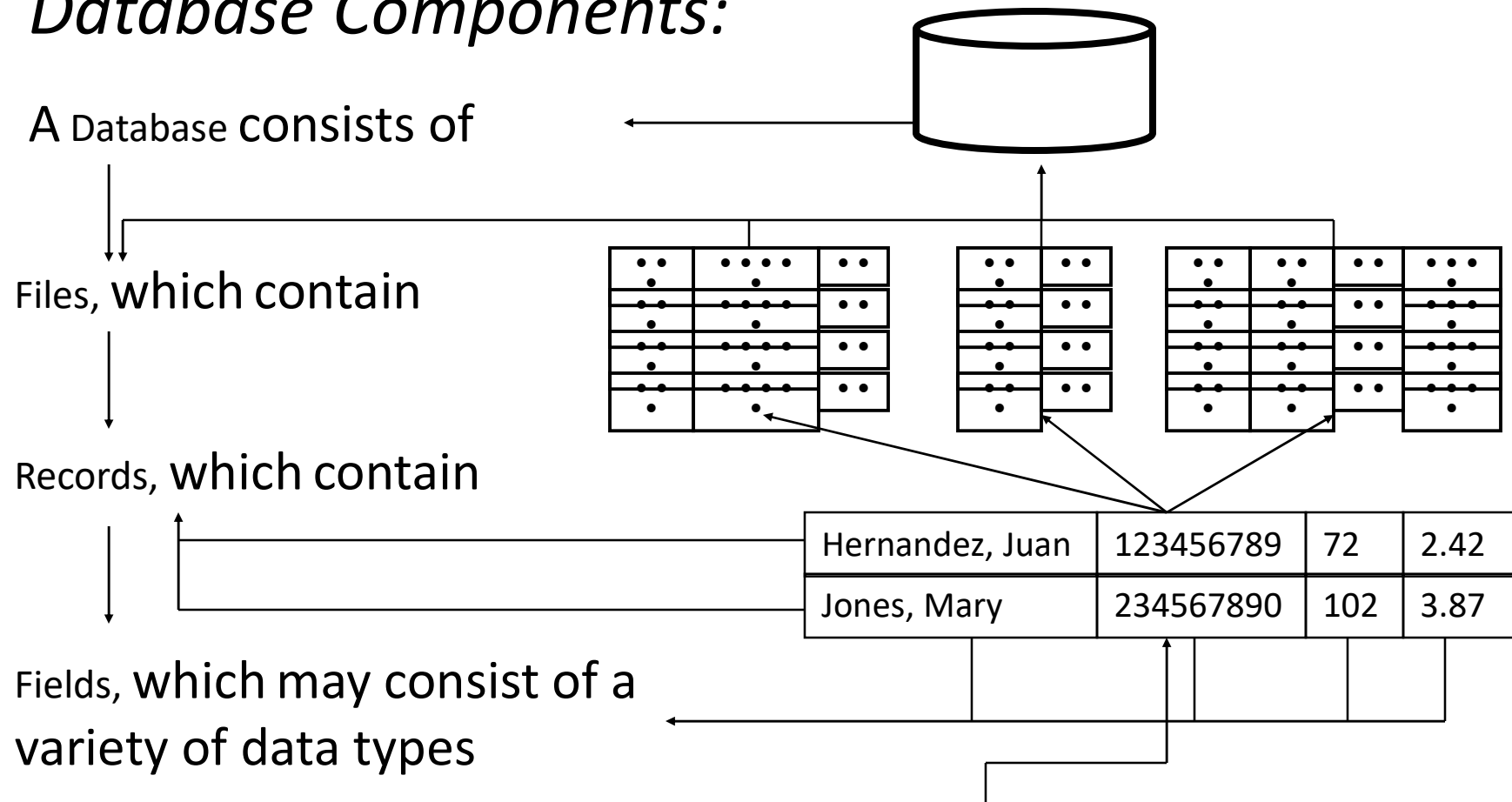
University



Database

# RDBMS Basics

## *Database Components:*



Notice that there should always be a Key (Unique) Field



# RDBMS Basics

## *RDBMS Restrictions/Conventions:*

- Each Relation *MUST* have a unique name

Table Student

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting
456789012	Elam, Mary	123-22 E St.	Accounting
.....	.....	.....	.....

Table Balance

Student	Owed	Department
103456678	1,502.36	Marketing
123456789	COBA219	Finance
456789012	COBA232	Accounting
.....	.....	.....

Table Student

StudentID	.....	Depart
987654321	.....	Finance
876543210	.....	INFOSYS
765432109	.....	Accounting
.....	.....	.....

NOT Allowed

# RDBMS Basics

## *RDBMS Restrictions /Conventions :*

- Each Relation *MUST* have a unique name
- All Columns (Tuples) *MUST* have Unique names

Table Student

StudentID	Name	Address	Address	Major
123456789	Saenz, Lupe	123 Mesa	Arlington	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	New York	INFOSYS
345678901	Adams, John	54B Hague	Dallas	Accounting
456789012	Elam, Mary	123-22 E St.	Ft. Worth	INFOSYS

One of the names *MUST* be changed

\* NOTE: The same field names *CAN* be used in different Relations



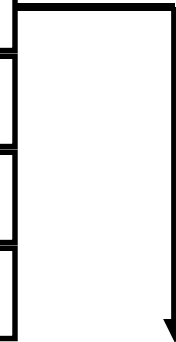
# RDBMS Basics

## *RDBMS Restrictions /Conventions :*

- Each Relation *MUST* have a unique name
- All Columns (Tuples) *MUST* have Unique names
- All Column Elements *MUST* be of the same data type

Table Student

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	24.34
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting
456789012	Elam, Mary	123-22 E St.	INFOSYS



Unless this is stored as Character String “24.34” (and *NOT* as the real number 24.34), it *MUST* be changed

# RDBMS Basics

## *RDBMS Restrictions /Conventions :*

- Each Relation *MUST* have a unique name
- All Columns (Tuples) *MUST* have Unique names
- All Column Elements *MUST* be of the same data type
  - Notice that this means each record requires the *SAME* number of Bytes of Storage

Table Student

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance

A Character String Requiring 9 Bytes of Storage  
A Character String Requiring 20 Bytes of Storage  
A Character String Requiring 25 Bytes of Storage  
A Character String Requiring 10 Bytes of Storage

A Total 64  
Bytes of  
Storage



# RDBMS Basics

## *RDBMS Restrictions /Conventions :*

- Each Relation *MUST* have a unique name
- All Columns (Tuples) *MUST* have Unique names
- All Column Elements *MUST* be of the same data type
- The order of Rows is *NOT* important

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting

Is the Same as

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finace
345678901	Adams, John	54B Hague	Accounting
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS

# RDBMS Basics

## *RDBMS Restrictions /Conventions :*

- Each Relation *MUST* have a unique name
- All Columns (Tuples) *MUST* have Unique names
- All Column Elements *MUST* be of the same data type
- The order of Rows is *NOT* important
- The Number of Bytes/Record Must be the same

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting

Each contains  
64 Bytes



# RDBMS Basics

## *RDBMS Keys:*

- Purpose
  - Define entity relationships
  - Determination
    - Knowing the value of a key field means you also know (Determine) the values of the other fields
    - E.g., knowing StudentID means you know StudentName, StudentAddress, etc.

StudentID       $\longrightarrow$       StudentName, StudentAddress  
(StudentID Determines StudentName and StudentAddress)

# RDBMS Basics

## *RDBMS Keys:*

- Purpose
  - Define entity relationships
  - Determination
  - Functional Dependence
    - An attribute is functionally dependent on another if can be determined by that attribute

StudentID  $\longrightarrow$  StudentAddress

(StudentID Determines Student Address)

- NOTE:

StudentAddress  $\not\longrightarrow$  StudentID

(Two Students *MAY* live at the same address)



# RDBMS Basics

## *RDBMS Keys:*

- Each Relation *MUST* have a unique identifier or *PRIMARY KEY*

Table Student

StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting
456789012	Elam, Mary	123-22 E St.	INFOSYS
•••••	•••••	•••••	•••••

No two students can have the same StudentID

# RDBMS Basics

## *RDBMS Keys:*

- Each Relation *MUST* have a unique identifier or *PRIMARY KEY*
- A *COMPOSITE KEY* is a combination of keys (Multi-key attributes) used to produce uniqueness

StudentName, StudentAddress  $\longrightarrow$  StudentMajor  
(StudentName AND StudentAddress Determines StudentMajor)

- NOTE: If Student Major is functionally dependent upon StudentName AND StudentAddress , BUT not on either StudentName or StudentName it is FULLY FUNCTIONAL DEPENDENT on the Concatenated key

(Attributes are fully functionally dependent on *PRIMARY KEYS*)

# RDBMS Basics

## *RDBMS Keys:*

- Each Relation *MUST* have a unique identifier or *PRIMARY KEY*
- A *COMPOSITE KEY* is a combination of keys (Multi-key attributes) used to produce uniqueness
- A *SUPER KEY* is either a *PRIMARY* or *COMPOSITE KEY* that uniquely identifies an entity

StudentID  $\longrightarrow$  StudentAddress

AND

StudentName, StudentAddress  $\longrightarrow$  StudentMajor

Are BOTH Superkeys



# RDBMS Basics

## RDBMS Keys:

- Each Relation *MUST* have a unique identifier or *PRIMARY KEY*
- A *COMPOSITE KEY* is a combination of keys (Multi-key attributes) used to produce uniqueness
- A *SUPER KEY* is either a *PRIMARY* or *COMPOSITE KEY* that uniquely identifies an entity
- A *CANDIDATE KEY* is any key or group of keys that could become a *SUPER KEY*
  - StudentID is a candidate key
  - StudentName, StudentAddress is a candidate key
  - StudentID, StudentAddress is *NOT* a candidate key

(StudentID by itself is a *CANDIDATE KEY*)

# RDBMS Basics

## *RDBMS Keys:*

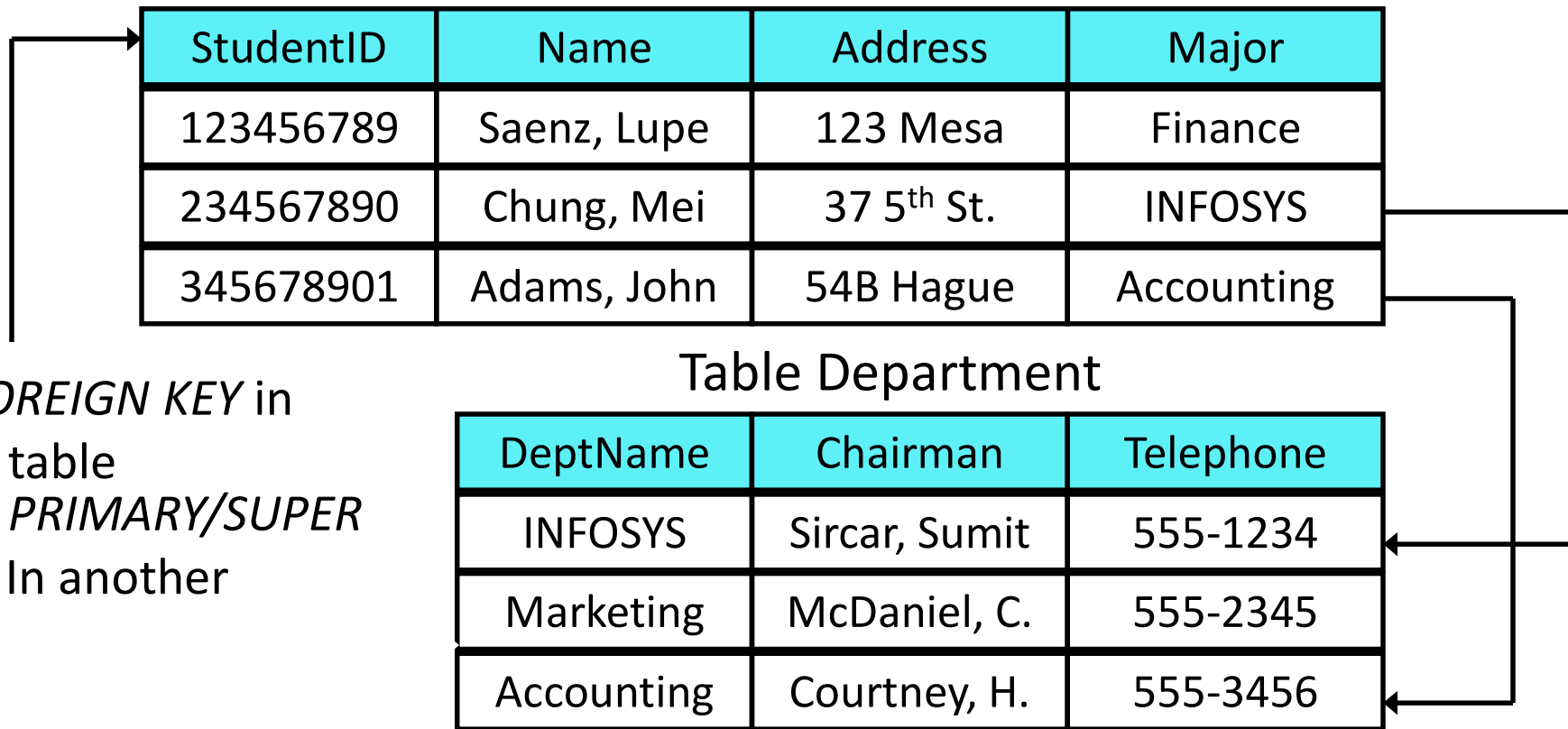
- Each Relation *MUST* have a unique identifier or *PRIMARY KEY*
- A *COMPOSITE KEY* is a combination of keys (Multi-key attributes) used to produce uniqueness
- A *SUPER KEY* is either a *PRIMARY* or *COMPOSITE KEY* that uniquely identifies an entity
- A *CANDIDATE KEY* is any key or group of keys that could become a *SUPER KEY*
- A *SECONDARY KEY* is any field, or combination of fields, which does NOT yield a unique value
  - Used for retrieval/Narrowing purposes only
  - StudentName, StudentZip may yield several records

# RDBMS Basics

## *RDBMS Keys:*

- In Order to relate two (or more) tables *FOREIGN KEYS* must be used

Table Student



StudentID	Name	Address	Major
123456789	Saenz, Lupe	123 Mesa	Finance
234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
345678901	Adams, John	54B Hague	Accounting

The diagram illustrates a foreign key relationship. A line starts from the 'StudentID' column header in the 'Table Student' and extends downwards. Another line starts from the 'DeptName' column header in the 'Table Department' and extends upwards. These two lines meet at a point, indicating that the 'StudentID' in the 'Table Student' is a foreign key that references the 'DeptName' in the 'Table Department'.

A *FOREIGN KEY* in one table is a *PRIMARY/SUPER KEY* In another

Table Department

DeptName	Chairman	Telephone
INFOSYS	Sircar, Sumit	555-1234
Marketing	McDaniel, C.	555-2345
Accounting	Courtney, H.	555-3456



# RDBMS Basics

## *Database Integrity:*

- Maintaining wholeness and Unity
- Entity Integrity:
  - All Entries MUST be Unique
  - No NULL values in primary key fields

Table Student

					StudentID	Name	Address	Major
Illegal Entries					123456789	Saenz, Lupe	123 Mesa	Finance
					234567890	Chung, Mei	37 5 <sup>th</sup> St.	INFOSYS
					123456789	Adams, John	54B Hague	Accounting
					456789012	Elam, Mary	123-22 E St.	INFOSYS
						Bush, G.W.	555 Austin	Marketing

# RDBMS Basics

## *Database Integrity:*

- Maintaining wholeness and Unity
- Entity Integrity
- Relational Integrity:
  - Foreign Key MUST have a valid entry in the corresponding table (or be NULL)

Table Student

StudentID	Major
123456789	Scamming
234567890	INFOSYS
345678901	

Table Department

DeptName	Chairman	Telephone
INFOSYS	Sircar, Sumit	555-1234
Marketing	McDaniel, C.	555-2345
Accounting	Courtney, H.	555-3456

???



NOT Allowed

# RDBMS Basics

## Database Integrity:

- Maintaining wholeness and Unity
- Entity Integrity
- Relational Integrity:
  - Foreign Key MUST have a valid entry in the corresponding table (or be NULL)
  - Primary Key entry CAN NOT be deleted if a foreign key refers to it

Table Student

StudentID	Major
123456789	Accounting
234567890	INFOSYS
345678901	Finance

Table Department

DeptName	Chairman	Telephone
INFOSYS	Sircar, Sumit	555-1234
Marketing	McDaniel, C.	555-2345

INFOSYS can NOT be deleted





# Summary

- Data - one of the most valuable resources a firm possesses.
- DBMS - a group of programs used as an interface between a database and application programs.
- The database approach to data management provides significant advantages over the traditional file-based approach.