

Dijkstra's Algorithm

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Dijkstra's algorithm is used to identify single source shortest path for a graph with non-negative edge costs.

Input and parameters

$G(V, E)$ is a graph. A source vertex s is given along with the cost function associated with E . For each vertex $u \in V$, $d[u]$ denotes the upper bound of the distance from s . For each vertex $u \in V$, let $\pi[u]$ denote the shortest parent of u in the shortest path $s \rightarrow u$.

Let, S be the set that holds the finished vertices or the vertices for which the shortest distance has been already computed.

Q be the set of vertices that are not in S . This is maintained as a priority queue. The minimum distance node is kept at the head of the priority queue Q .

Algorithm

Algorithm 1: Dijkstra's

Input: $(G, s, w(*))$

Output: Shortest distance from s to all the points

begin

```
1  Initialize( $G, s$ )
2   $S \leftarrow \phi$ 
3   $Q \leftarrow V$ 
4  while ( $Q \neq \phi$ ) do
5       $u \leftarrow \text{Extract-Min}(Q)$ 
6       $S \leftarrow S \cup \{u\}$ 
7      for each  $v \in V$  do
8          Relax( $u, v, w(*)$ )
9  return  $d$ 
```

Algorithm 2: Initialize(G, s)

Input: (G, s)
begin
1 **for** *each* $u \in V$ **do**
2 $d[u] = \infty$
3 $\pi[u] = NIL$
4 $d[s] = 0$

Algorithm 3: Relax($u, v, w(*)$)

Input: ($u, v, w(*)$)
begin
1 **if** $d[v] > d[u] + w(u \rightarrow v)$ **then**
2 $d[v] = d[u] + w(u \rightarrow v)$
3 $\pi[v] = u$
