

**PSG College of Technology**  
**Department of Applied Mathematics and Computational Sciences**  
**MSc Theoretical Computer Science(2017-18)**  
**15XT68 -Security in Computing Lab**  
**Problem Sheet No 3**

### **1.Cyclic attack in RSA**

Perform a cyclic attack on a Cipher text C to find the plain text P. Consider a short message and a low value of n.

### **2.Broadcast Attack using CRT**

An RSA public-key is a pair (n, e) where n = p, q is the product of two primes.

$$\text{RSA}_{n,e}(m) = m^e \bmod n$$

Alice, Bob, and Carol use RSA public-keys  $(n_A, 3)$ ,  $(n_B, 3)$ , and  $(n_C, 3)$ , respectively. David wants to send the same message 'm' to three of them. So, David computes

$$y_A = m^3 \bmod n_A$$

$$y_B = m^3 \bmod n_B$$

$$y_C = m^3 \bmod n_C$$

and sends the cipher text Y to the respective users. Show how an eavesdropper Eve can now compute the message 'm' even without knowing any of the private keys of Alice, Bob, and Carol.

### **3. Chosen Cipher text attack**

This is an attack against the textbook version of the RSA algorithm. In this attack, an attacker first chooses a message and encrypts it the victim's public key. Then, the attacker asks the victim to sign (decrypt) for him a specially crafted related message. Due to the following property of RSA

$$E_{PU}(M_1) \times E_{PU}(M_2) = E_{PU}(M_1 \times M_2) \text{ -----}, (1)$$

the attacker can easily recover any message encrypted with the victim's private key, without ever learning this private key. For example, the attacker wants to decrypt the following ciphertext  $C = M^e \bmod N$ , without knowing the private key d. The attacker proceeds as follows. Knowing the victim's public key e, he prepares the following message

$$X = (C \times 2^e) \bmod N,$$

gives it to the victim and asks her to sign it. The victim signs message X with its private key and sends the result Y back to the attacker.

$$Y = X^d \bmod N$$

$$\begin{aligned} X^d &= \left( (C \bmod N) \times (2^e \bmod N) \right)^d \\ &= \left( (M^e \bmod N) \times (2^e \bmod N) \right)^d \\ &= \left( (2 \times M)^e \bmod N \right)^d \\ &= (2 \times M)^{ed} \bmod N \\ &= 2 \times M. \end{aligned}$$

Using Y and equation (1), the attacker can retrieve the encrypted message M as follows:

1. Show by example that equation (1) holds for the RSA encryption algorithm.
2. Demonstrate by example the chosen ciphertext attack against RSA.