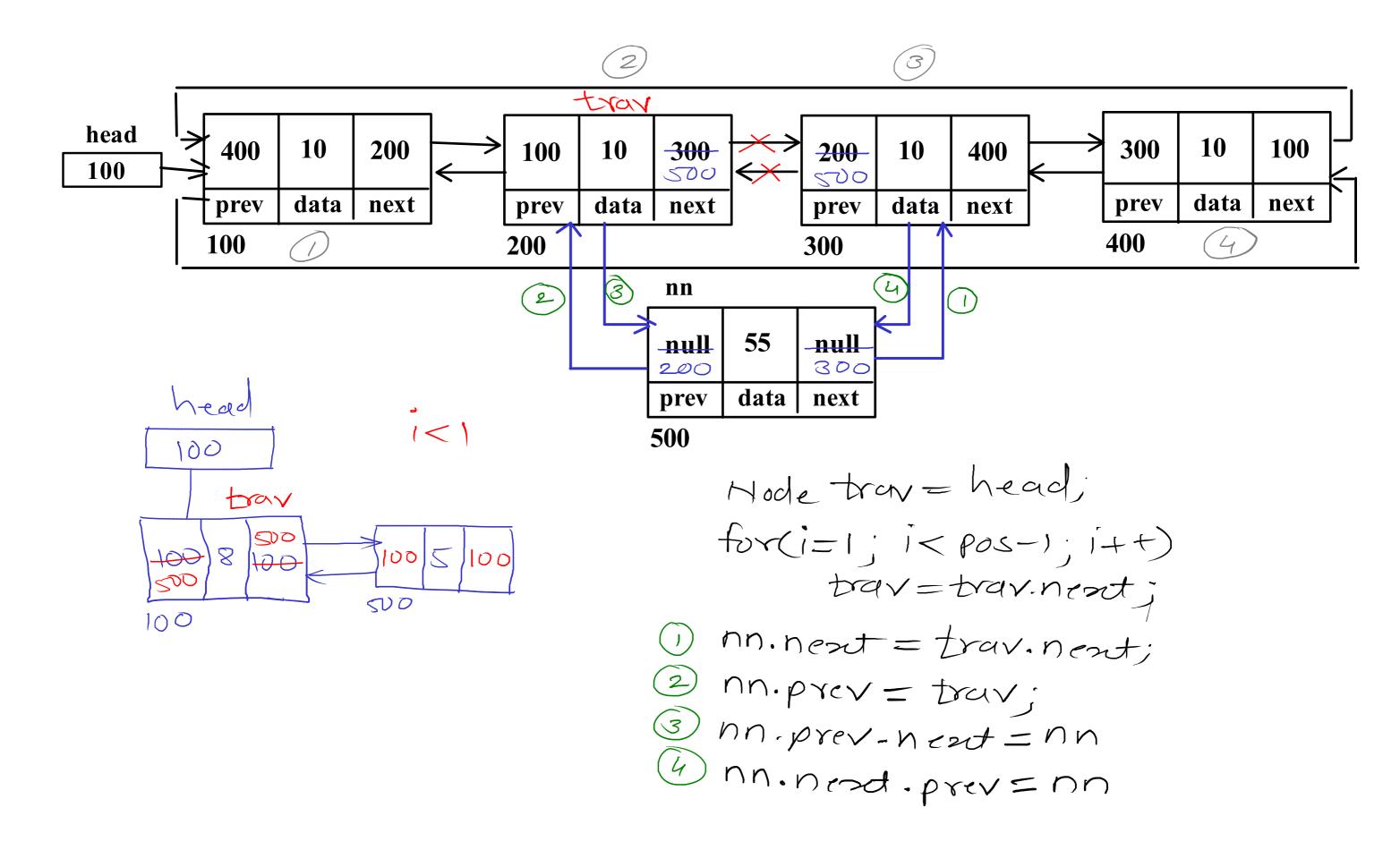
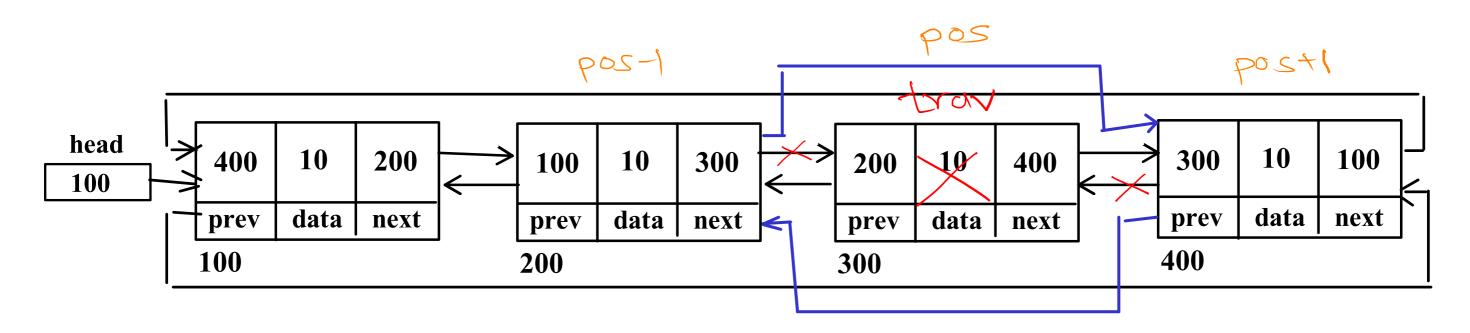
Time Complexity Analysis of Linked List

	SI	LLL	SCLL		DLLL		DCLL
1. Add First	O (1)	O (1)	O(n)	O(1)	O(1)	O(1)	O (1)
2. Add Last	O(n)	O (1)	O(n)	O(1)	O(n)	O(1)	O (1)
3. Add Position	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)
4. Del First	O (1)	O (1)	O(n)	O(1)	O(1)	O (1)	O (1)
5. Del Last	O(n)	O(n)	O(n)	O(n)	O(n)	O(1)	O (1)
6. Del Position	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)
7. Display	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)	O(n)





Node trav= head;
for(i=1; i < pos; i++)

trav = trav. next;

trav. prev. next = trav. next;

trav. next. prev = trav. prev;

Linked List Applications

- dynamic data structure grow / shrink at runtime
- due to this dynamic nature, it is used to implement other data structures
 - 1. Stack
 - 2. Queue
 - 3. Hash Table (Seperate chaining)
 - 4. Graph (Adjacency list)
- Operating system job queue, ready queue, waiting queues (Doubly circular linked list)

Deque (Double Ended Queue)

	Stack	Queue	Push front —> Pop front ←	→ Pop rear ← Push rear	
1.	Add First	1. Add First	front	rear	
	Del First	Del Last	1. Push front Add First	1. Push rear Add Last	
2.	Add Last Del Last	2. Add Last Del First	2. Pop front Del First	2. Pop rear Del Last	

Types: 1. Input restricted Deque

2. Output restricted Deque

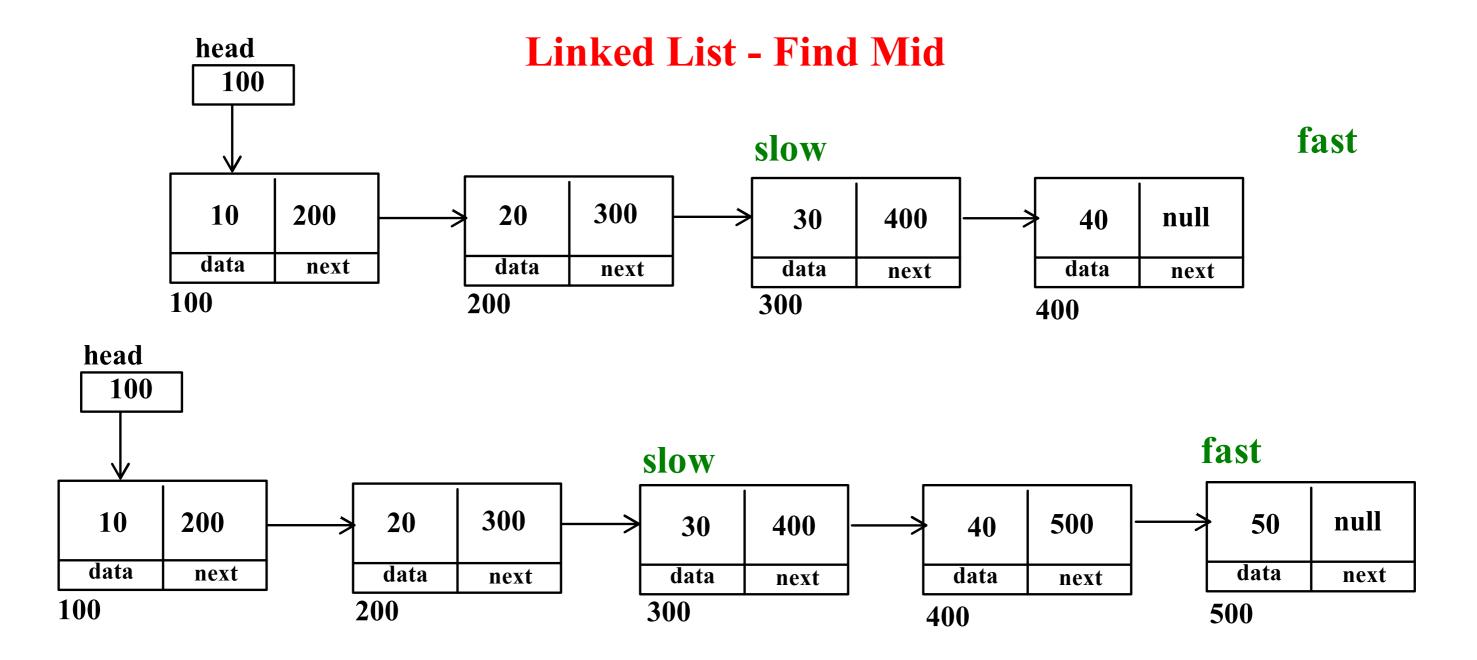
Array Vs Linked List

Array

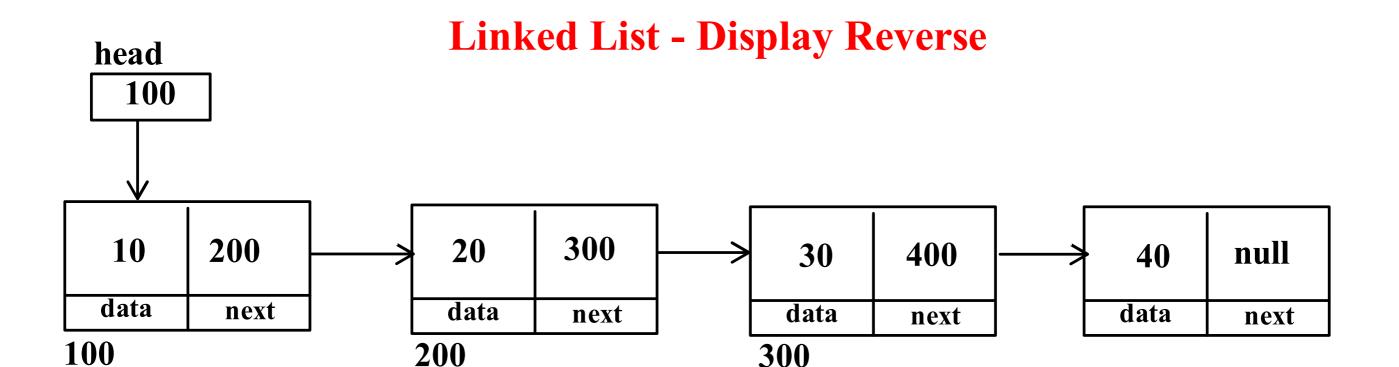
- 1. Array space in memory is contiguous
- 2. Array can not grow or shrink at runtime
- 3. Random access of elements is allowed
- 4. Insert or Delete, needs shifting of array elements
- 5. Array needs less space

Linked List

- 1. Linked list space in memory is not contiguous
- 2. Linked list can grow or shrink at runtime
- 3. Random access of elements is not allowed(sequential)
- 4. Insert or Delete, do not need shifting of nodes
- 5. Linked lists need more space



```
public Node findMid(){
    Node fast = head;
    Node slow = head;
    while(fast != null && fast.next != null){
        fast = fast.next.next;
        slow = slow.next;
    }
    return slow;
}
```



```
void forwardDisplay(Node trav){
       if(trav == null)
               return;
       sysout(trav.data);
       forwardDisplay(trav.next);
 10, 20, 30, 40

Forward Display (100)

Forward Display (200)

Forward Display (300)

Sforward Display (400)

Sforward Display (400)
```

```
void backwardDisplay(Node trav){
        if(trav == null)
                return;
        backwardDisplay(trav.next);
        sysout(trav.data);
  > backward Display (100)

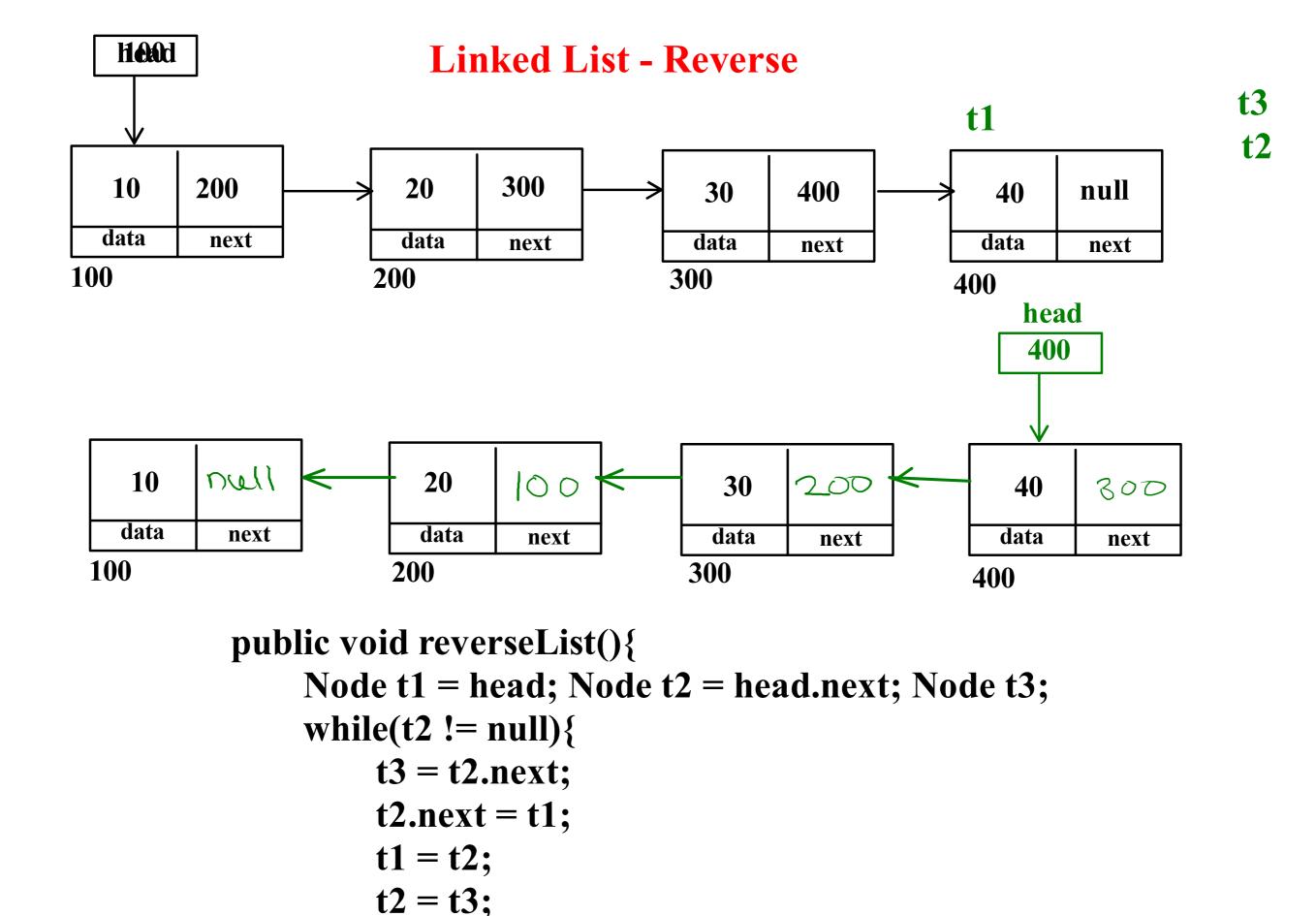
> backward Display (200)

> backward Display (300)

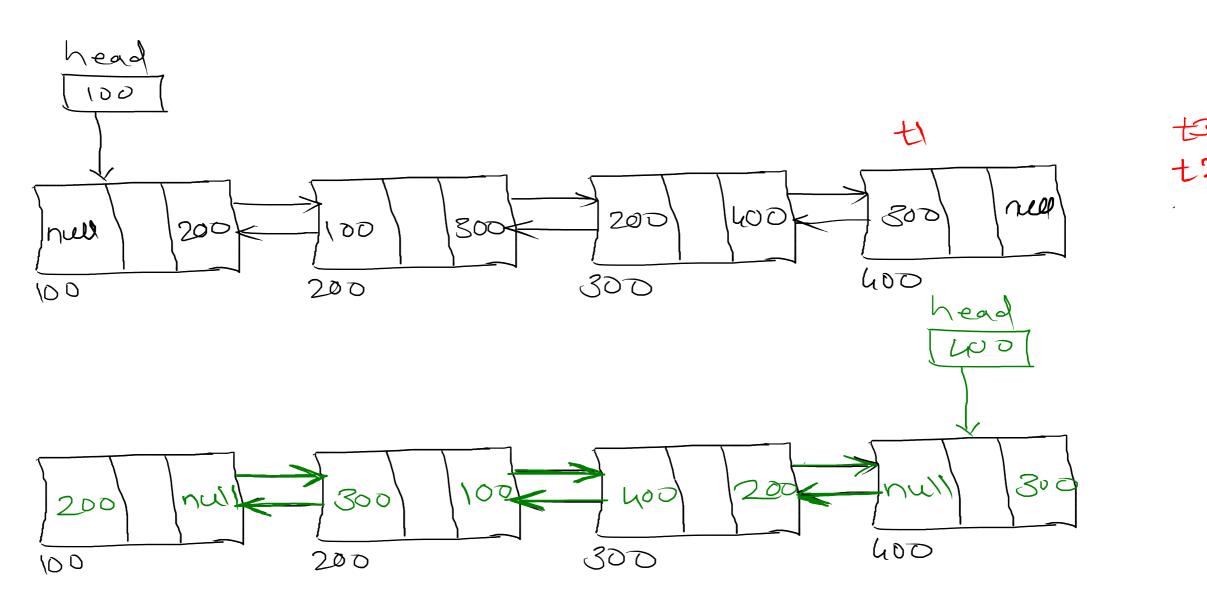
> backward Display (400)

> backward Display (400)

> backward Display (null)
```



head.next = null; head = t1;



Mode ti = head; Node ts;

Node tz = head. next;

While (tz != null) &

t3 = tz. next

t1.prev=tz

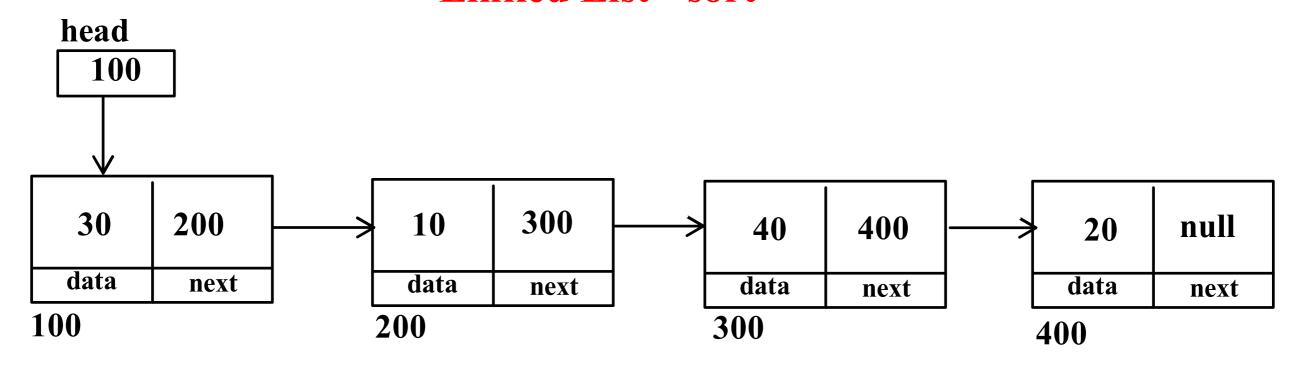
tz.next=t1

ti=tz; tz=ts;

head. next= null; t1.prev= null;

head=t1;

Linked List - sort



for (Mode i = head; i-next != null; i=i-next)

for (Mode j=i-next; j!= null; j=j-next)

[f(i.data > j.data) {
 int temp=i.data;
 i.data = j.data;
 j.data = temp;
}

ADD Node into BST

```
//1. create node with given data
//2. if BStree is empty
     // add newnode into root
//3. if BSTree is not empty
     //3.1 create one tray pointer and start from root
     //3.2 if data is less than current node data
          //3.2.1 if current node do not have left child
               // add newnode into left of current node
          //3.2.2 if current node have left child
               // go on left child
     //3.3 if data is greater than current node data
          //3.3.1 if current node do not have right child
               // add newnode into right of current node
          //3.3.2 if current node has right child
               // go on right child
     //3.4 repeat step 3.2 and 3.3 till node is not added into BST
```

