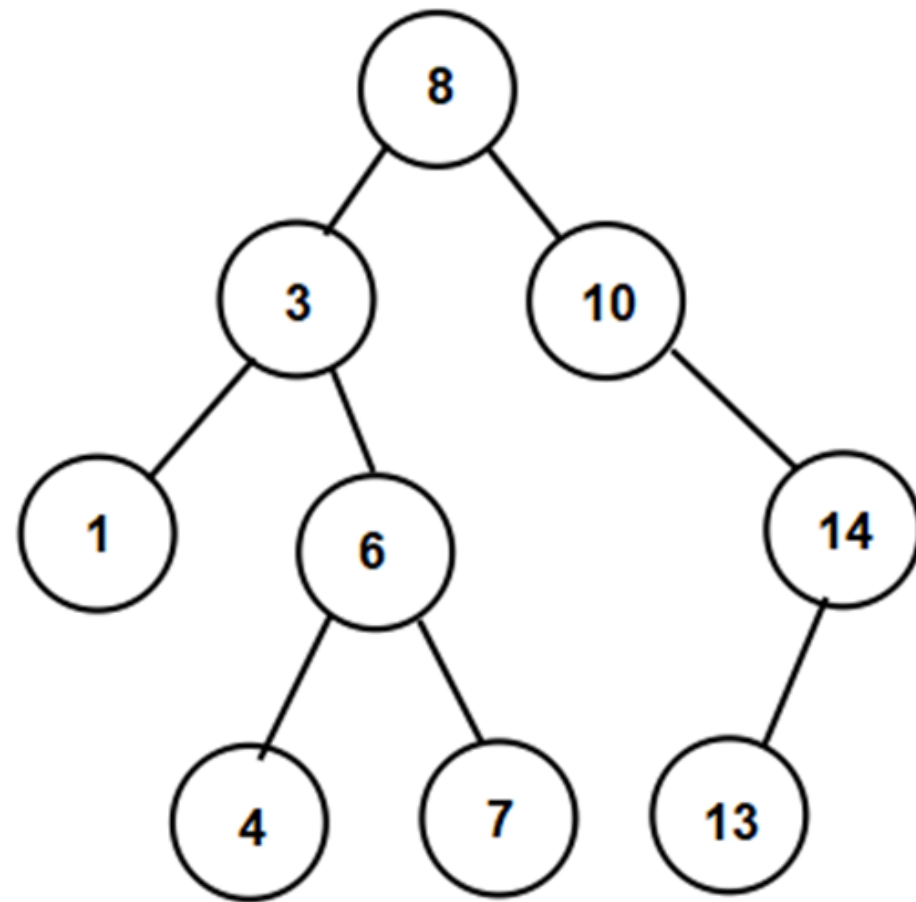


## BST - Height



//0. if left or right sub tree is absent

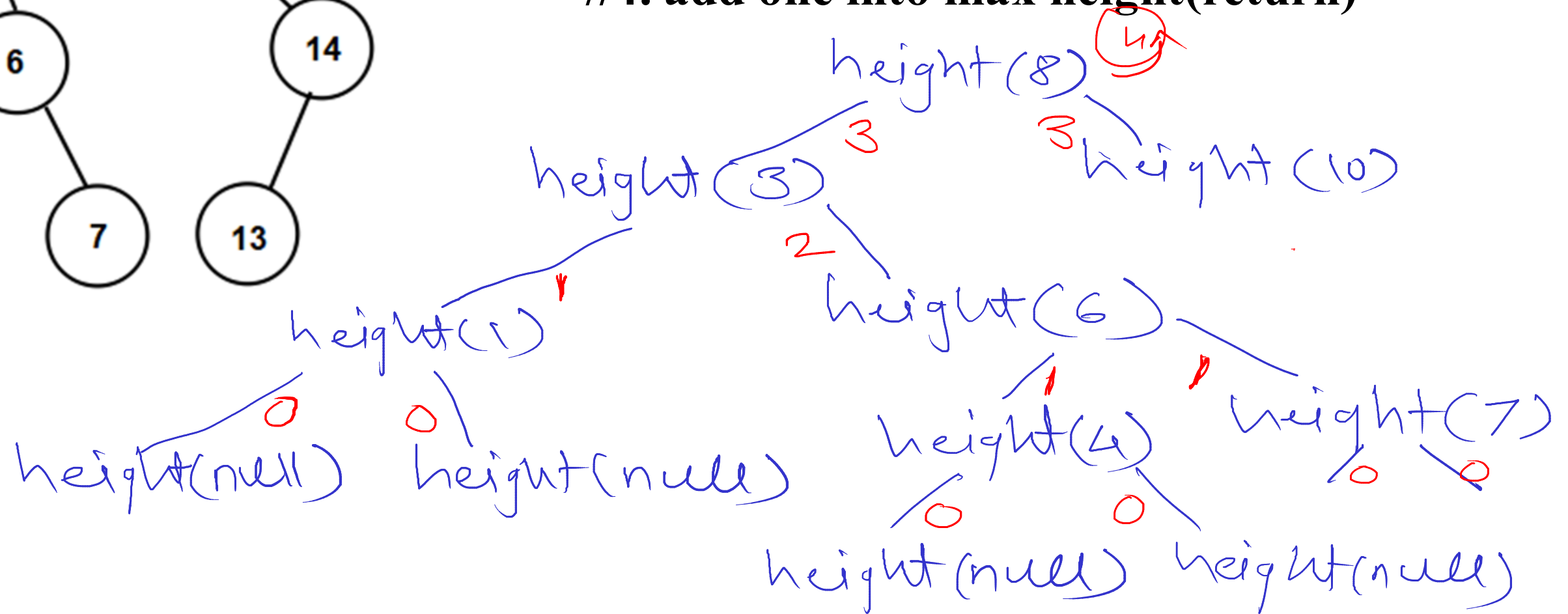
//then return 0

//1. find height of left subtree

//2. find height of right subtree

//3. find max height

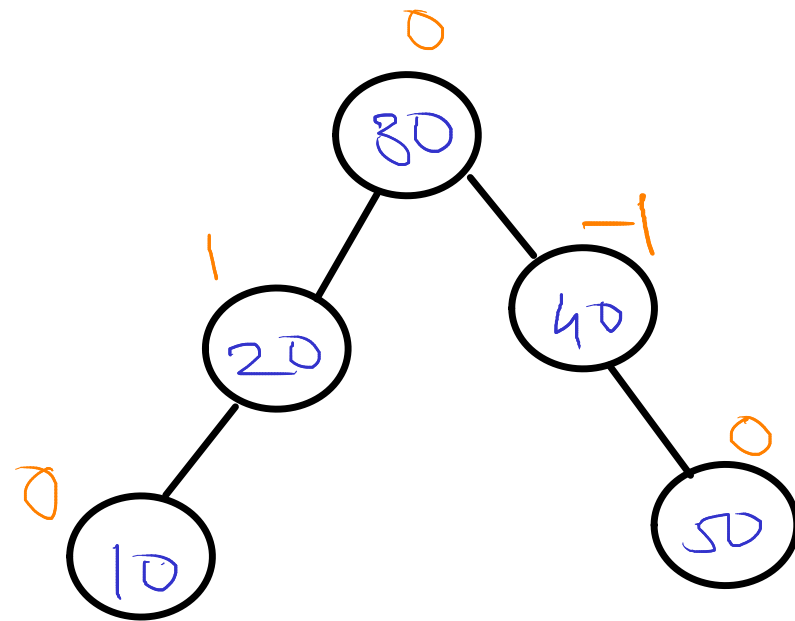
//4. add one into max height(return)



**Height of node = Max(Height(left sub tree), Height(right sub tree)) + 1**

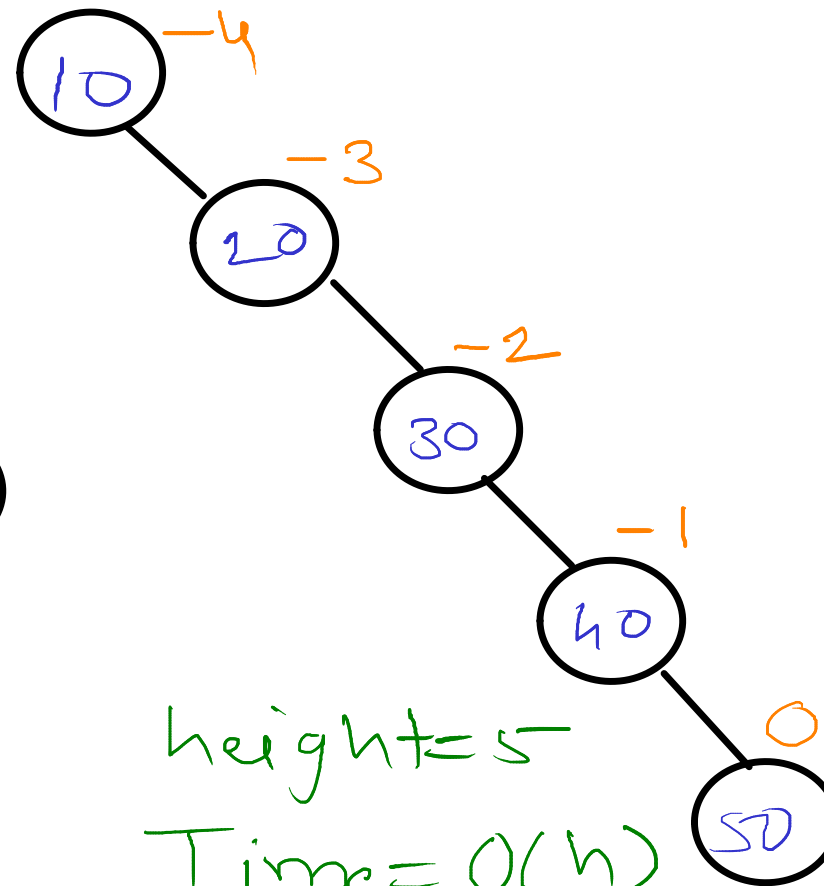
## Skewed BST

Keys : 30, 40, 20, 50, 10



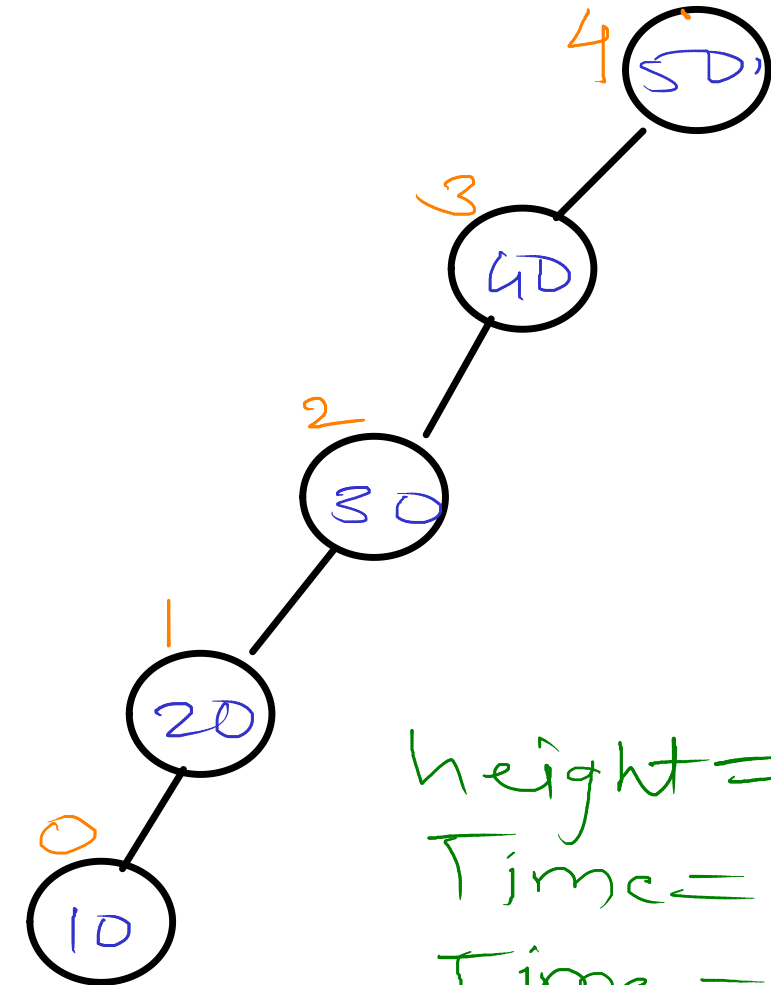
height = 3  
Time =  $O(h)$   
Time =  $O(\log n)$

Keys : 10, 20, 30, 40, 50



height = 5  
Time =  $O(h)$   
Time =  $O(n)$

Key : 50, 40, 30, 20, 10



height = 5  
Time =  $O(h)$   
Time =  $O(n)$

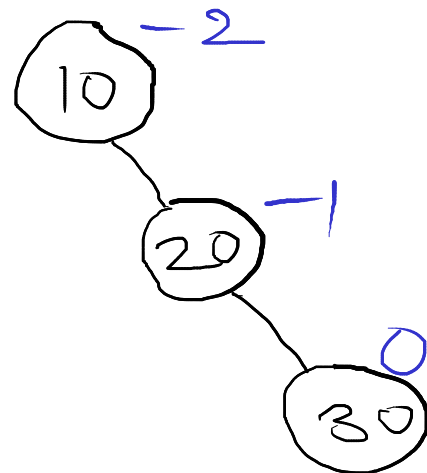
- tree which grows in only one direction is called as skewed BST.
- if tree grows in left direction, it is called as left skewed BST.
- if tree grows in right direction, it is called as right skewed BST.

## Balanced BST

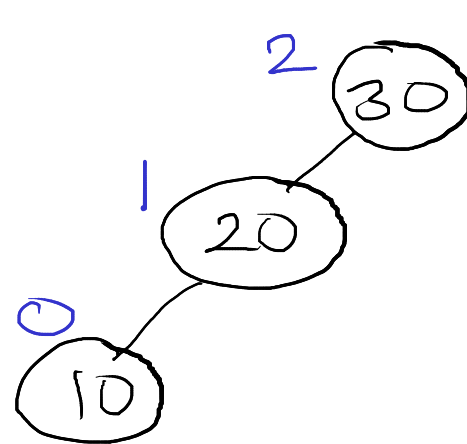
$$\text{Balance Factor} = \text{height}(\text{left sub tree}) - \text{height}(\text{right sub tree})$$

- tree is balanced if balance factor of all the nodes is either -1, 0 or +1
- balance factor = {-1, 0, +1}

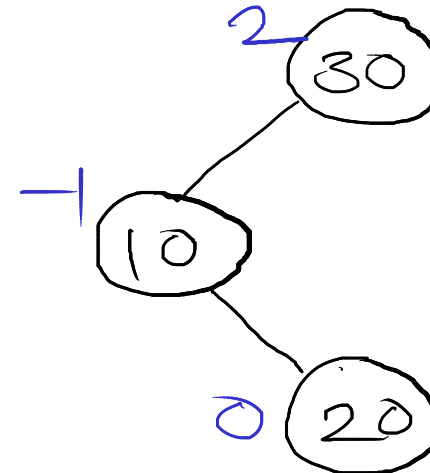
Keys : 10, 20, 30



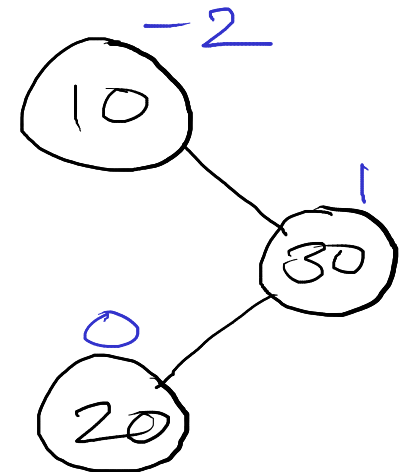
Keys : 30, 20, 10



Keys : 30, 10, 20

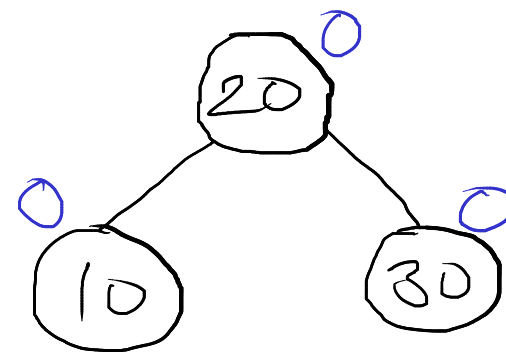


Keys : 10, 30, 20



Keys : 20, 10, 30

Keys : 20, 30, 10

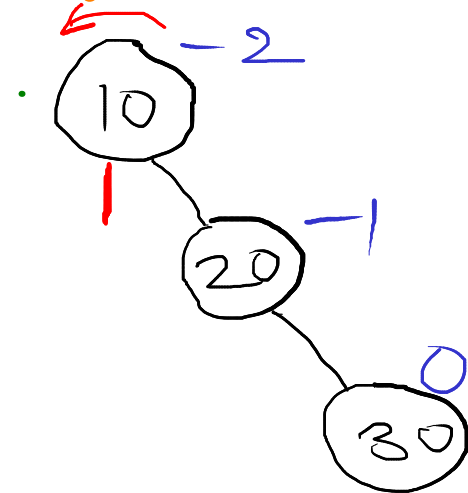


Balanced BST

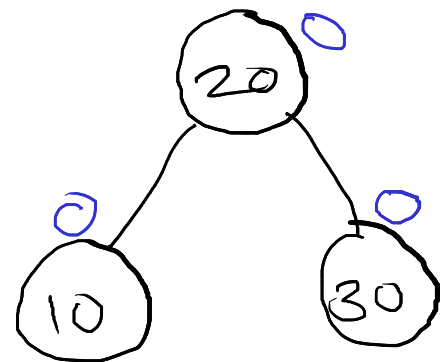
# Rotations

## RR Imbalance

Keys : 10, 20, 30



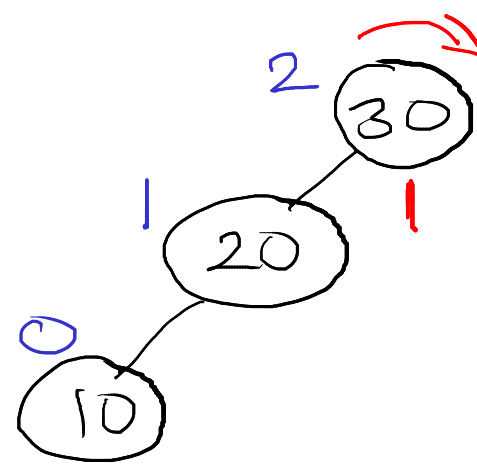
## Left Rotation



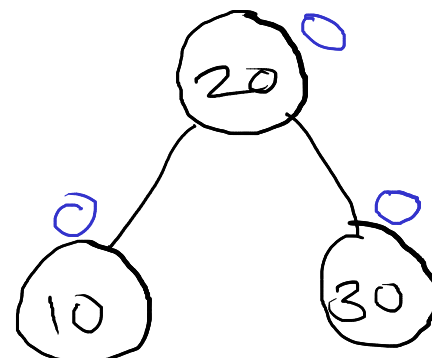
## Single Rotation

## LL Imbalance

Keys : 30, 20, 10

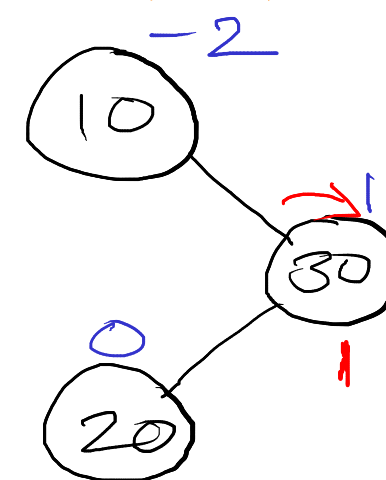


## Right Rotation

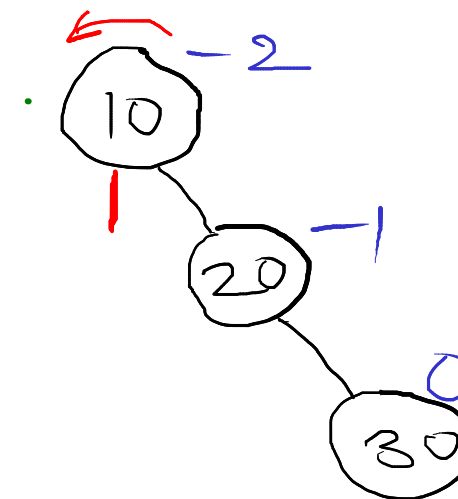


## RL Imbalance

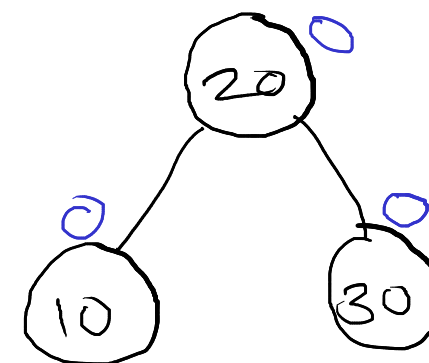
Keys : 10, 30, 20



## Right Rotation



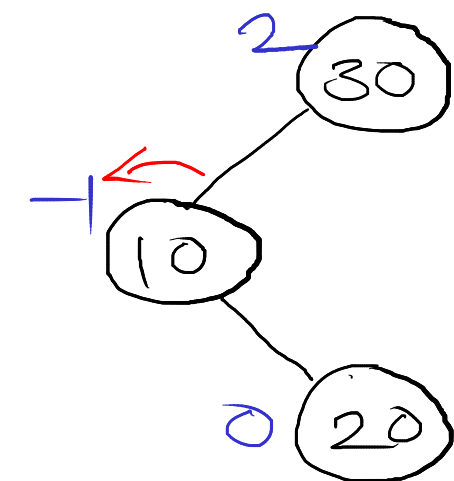
## Left Rotation



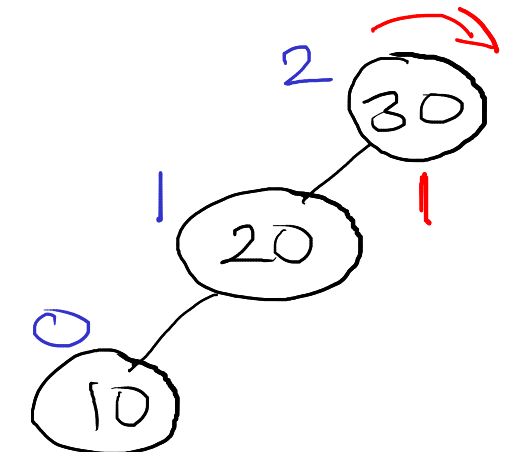
## Double Rotation

## LR Imbalance

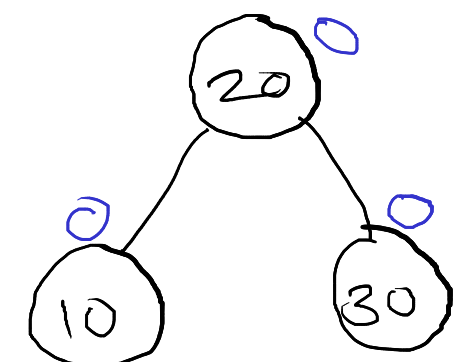
Keys : 30, 10, 20



## Left Rotation

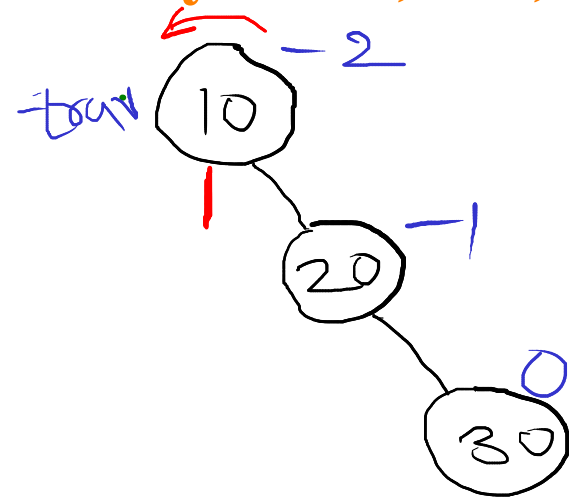


## Right Rotation



### RR Imbalance

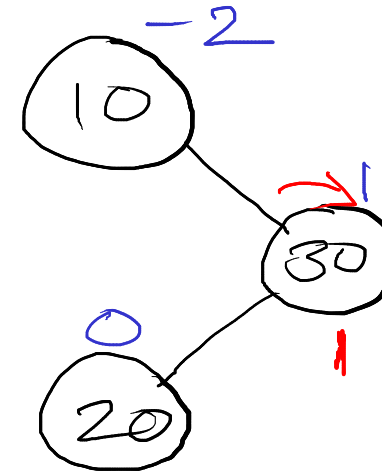
Keys : 10, 20, 30



$bf < -1$   
 $val > trav.right.data$

### RL Imbalance

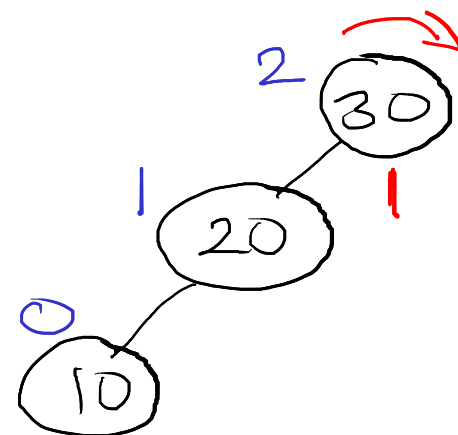
Keys : 10, 30, 20



$bf < -1$   
 $val < trav.right.data$

### LL Imbalance

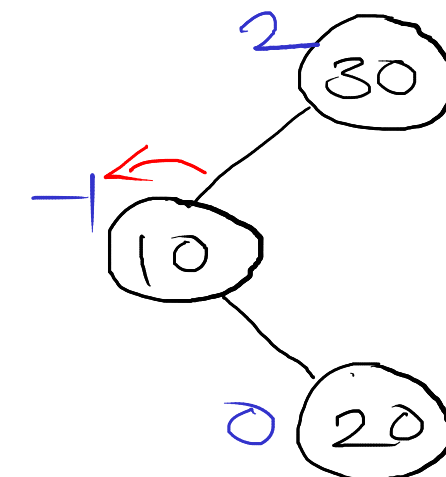
Keys : 30, 20, 10



$bf > 1$   
 $val < trav.left.data$

### LR Imbalance

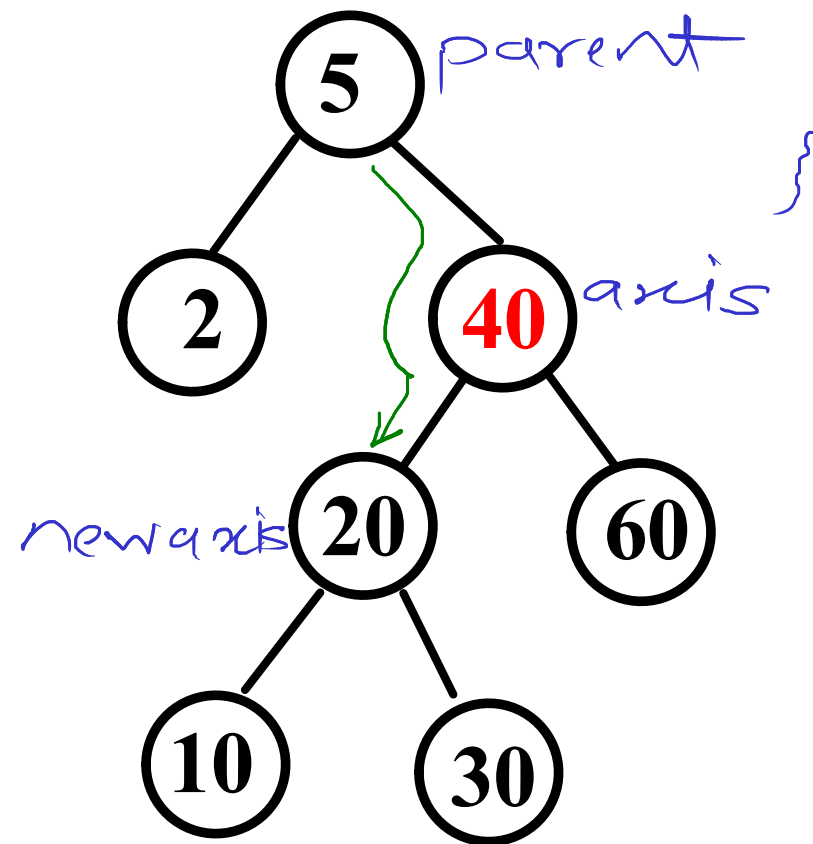
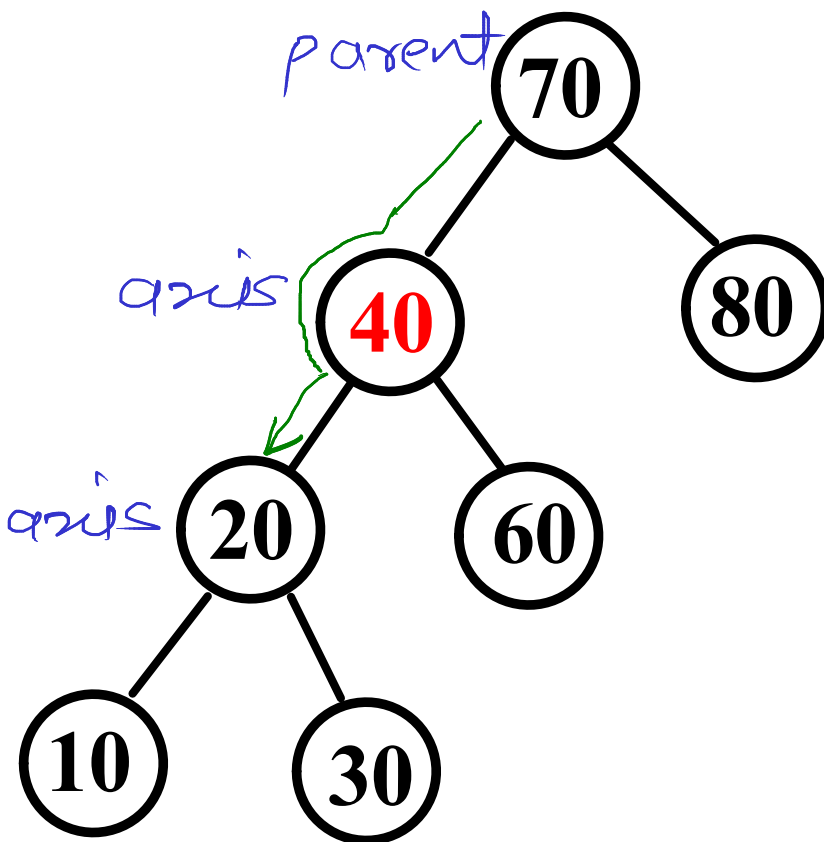
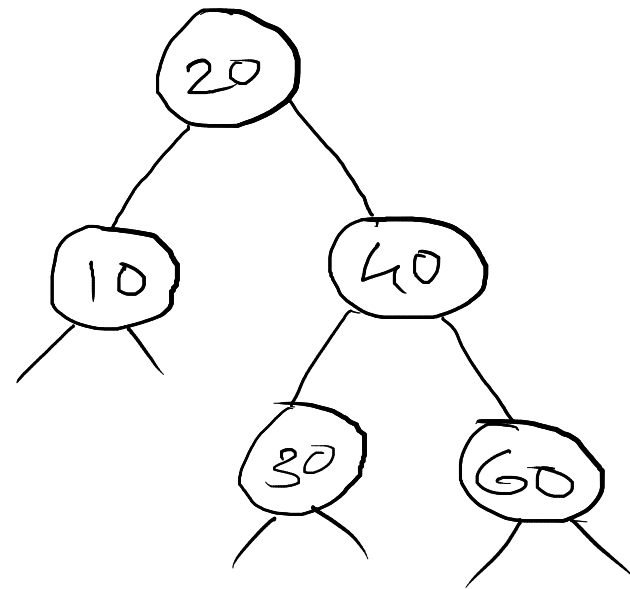
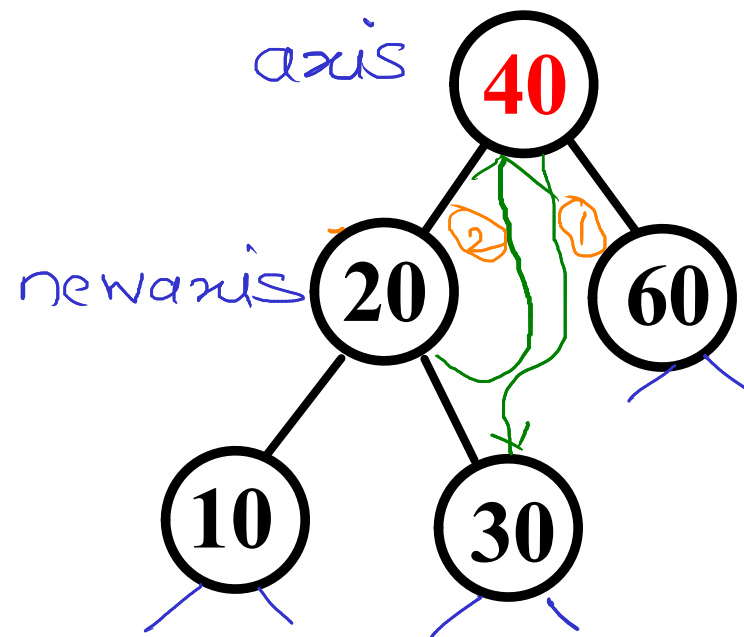
Keys : 30, 10, 20



$bf > 1$   
 $val > trav.left.data$

## Right Rotation

```
void rightRotation(axis, parent) {
    newaxis = axis.left
    ① axis.left = newaxis.right
    ② newaxis.right = axis
    if (axis == root)
        root = newaxis
    else if (axis == parent.left)
        parent.left = newaxis
    else if (axis == parent.right)
        parent.right = newaxis
}
```



## Left Rotation

```
void leftRotation(axis, parent){
```

```
    newaxis = axis -> right
```

```
    ① axis -> right = newaxis -> left
```

```
    ② newaxis -> left = axis
```

```
    if (axis == root)
```

```
        root = newaxis
```

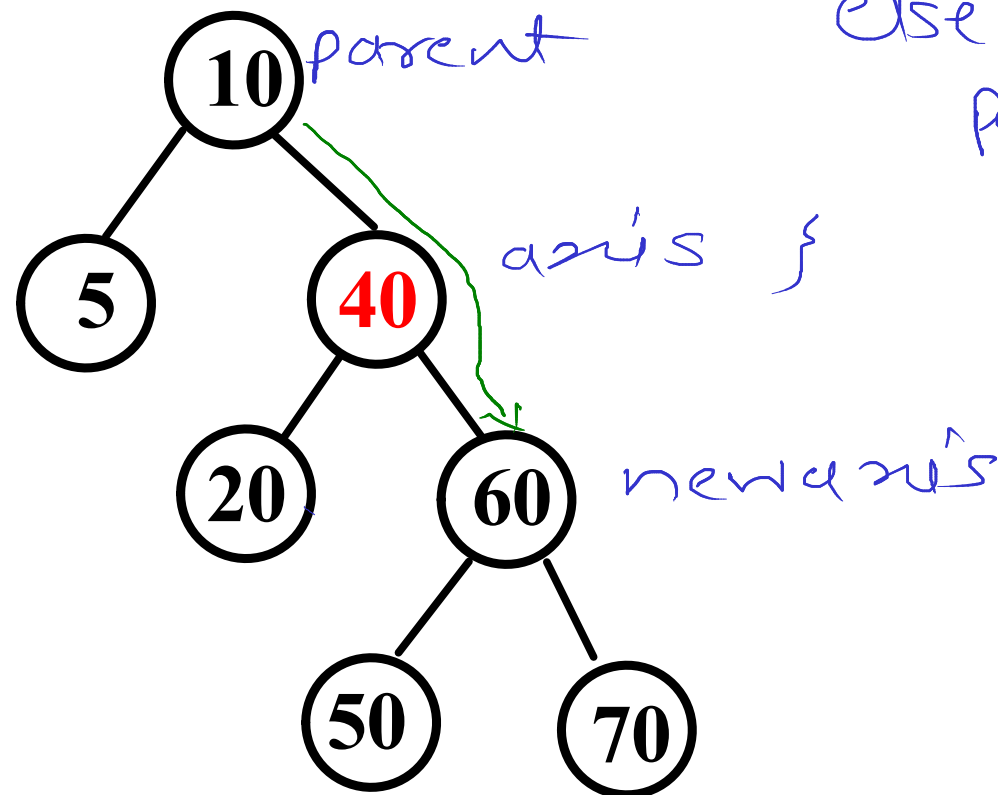
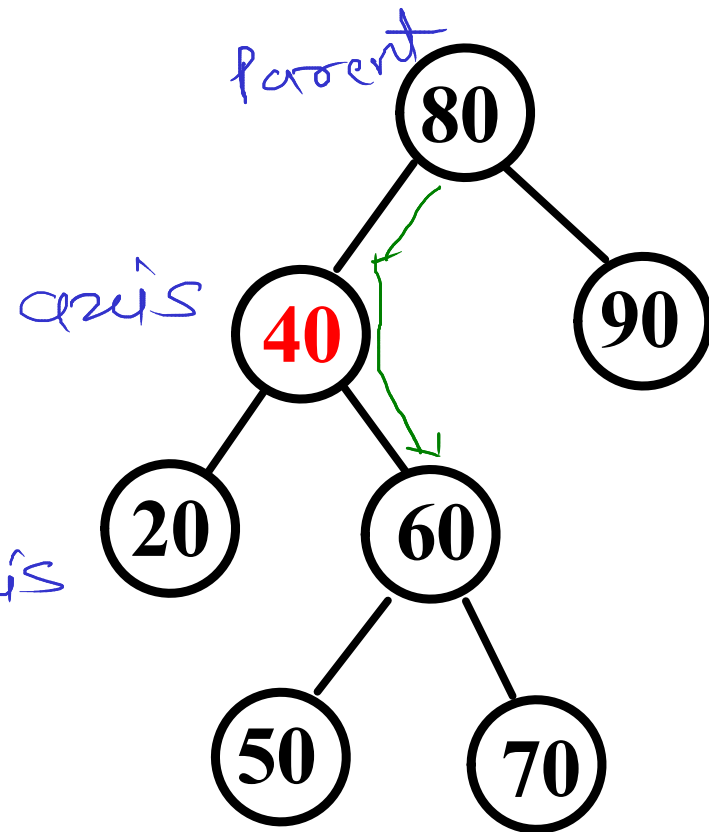
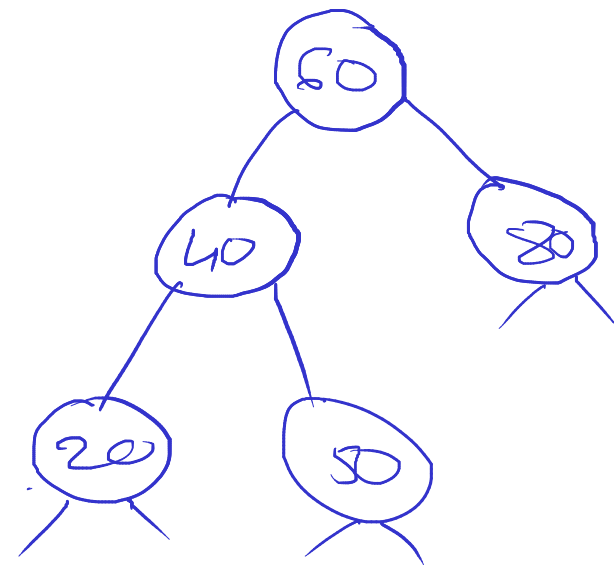
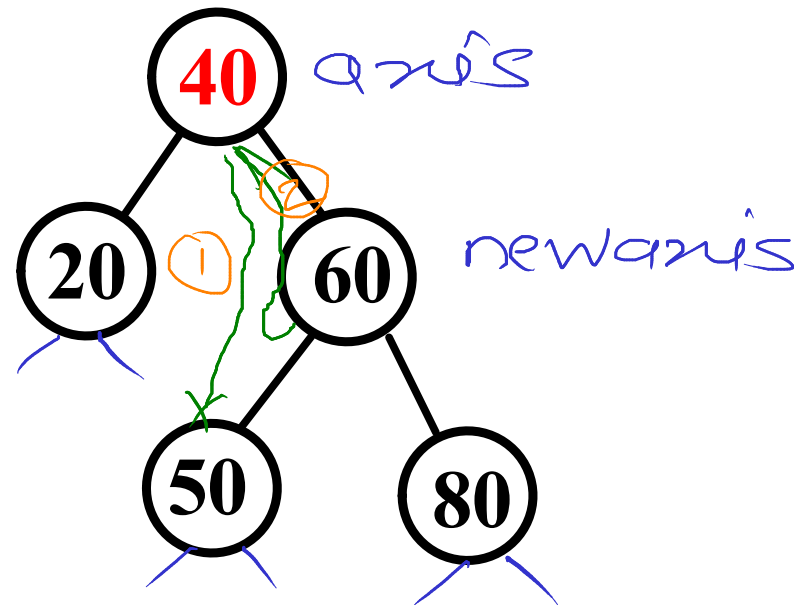
```
    else if (axis == parent -> left)
```

```
        parent -> left = newaxis
```

```
    else if (axis == parent -> right)
```

```
        parent -> right = newaxis
```

```
}
```

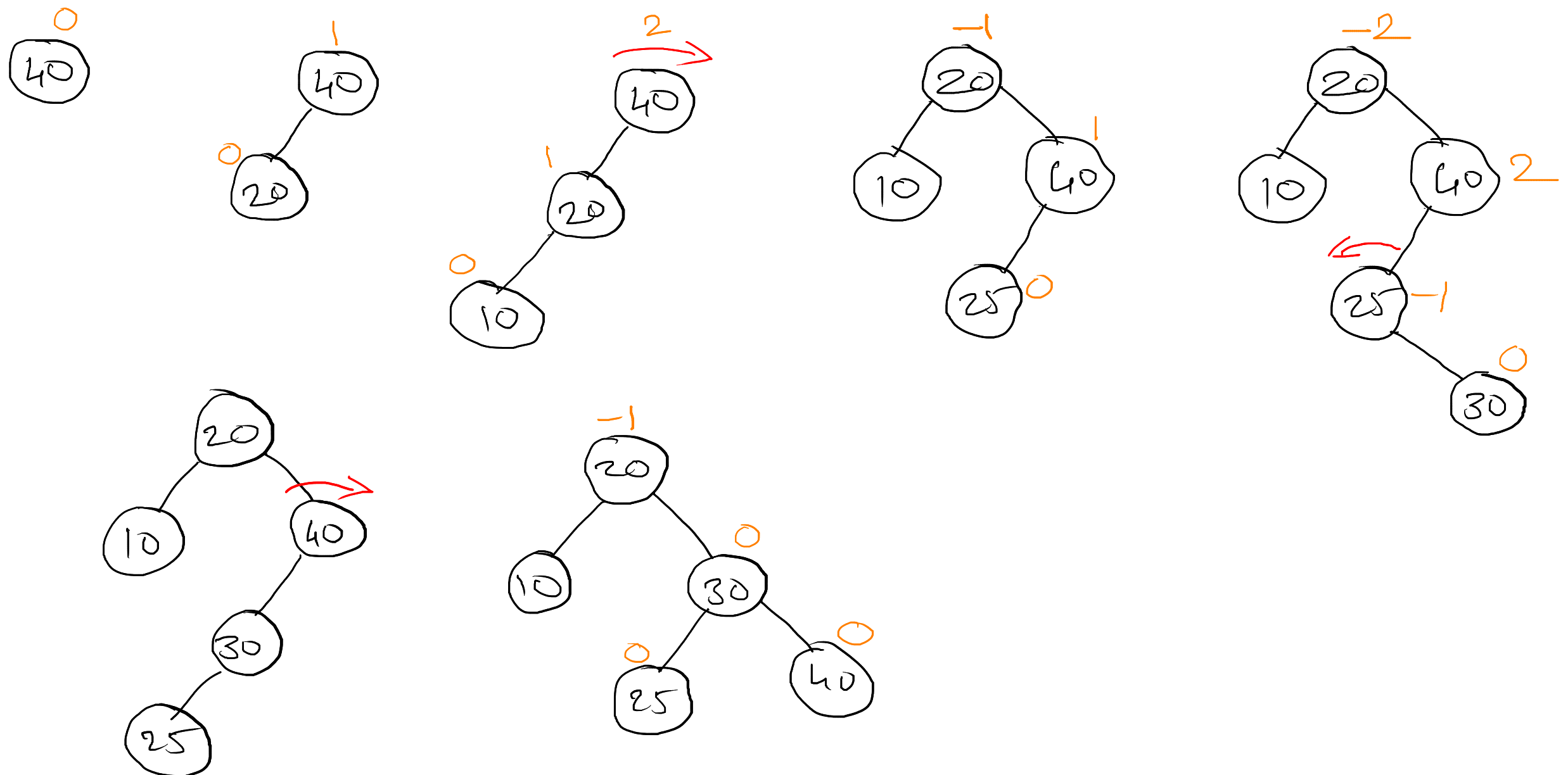




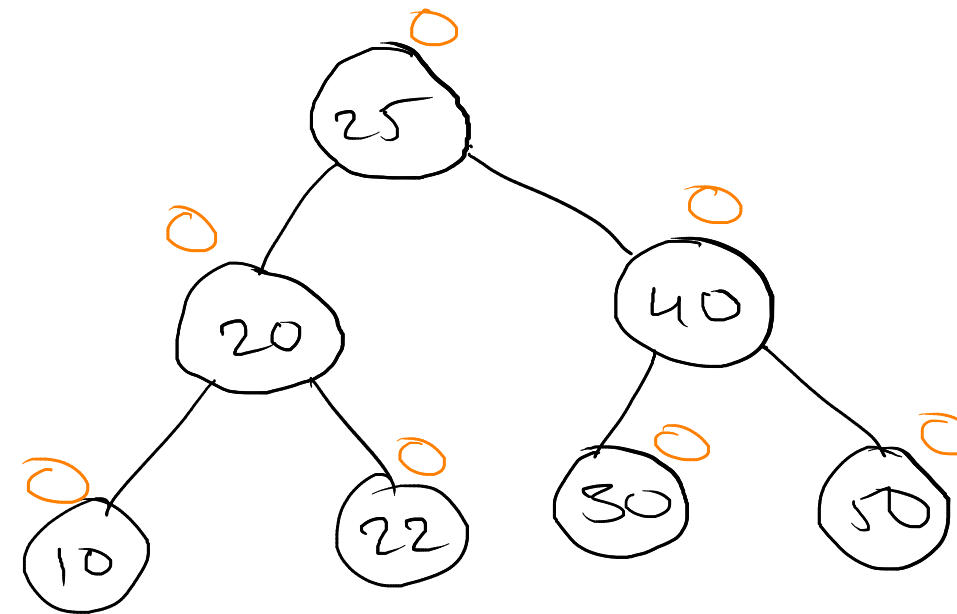
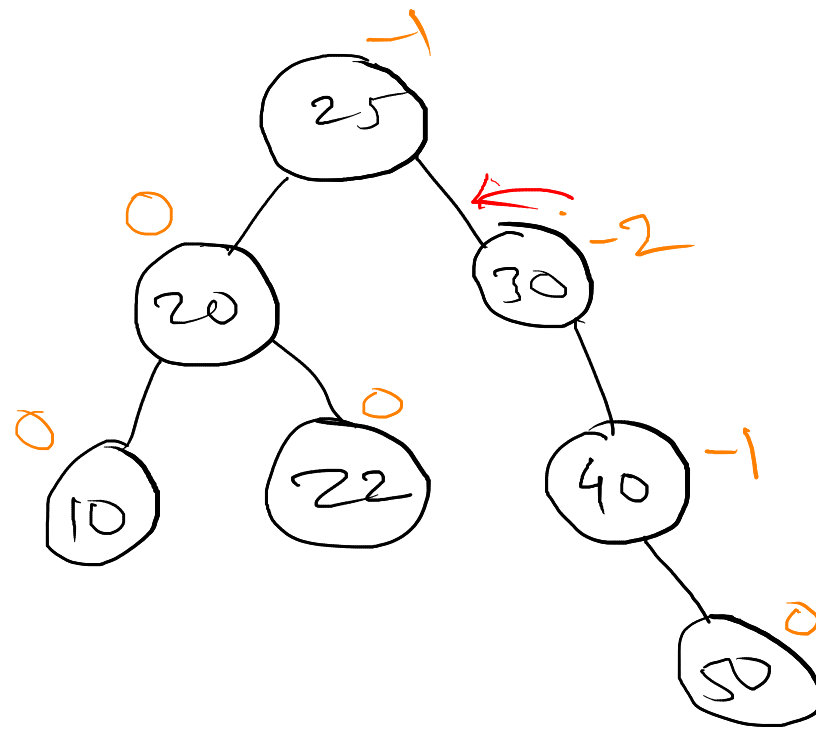
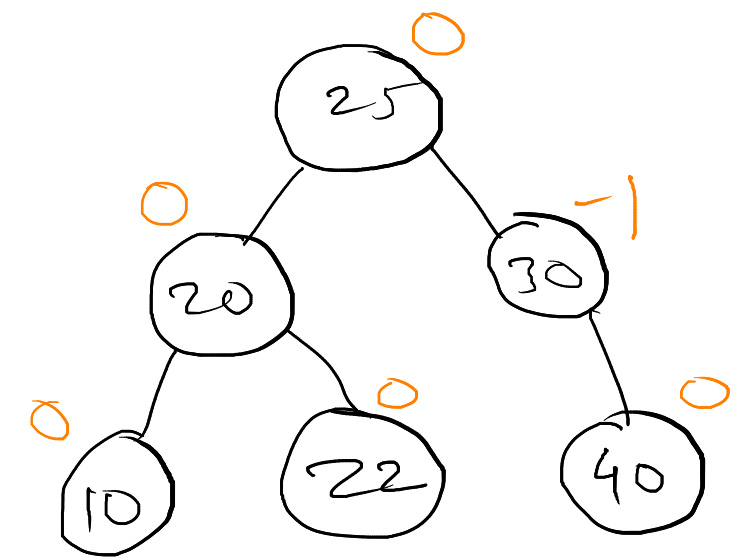
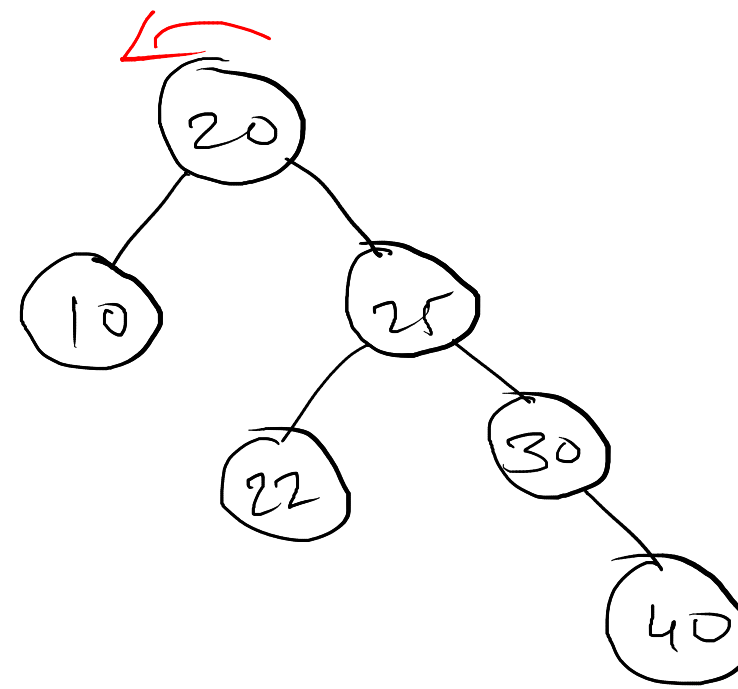
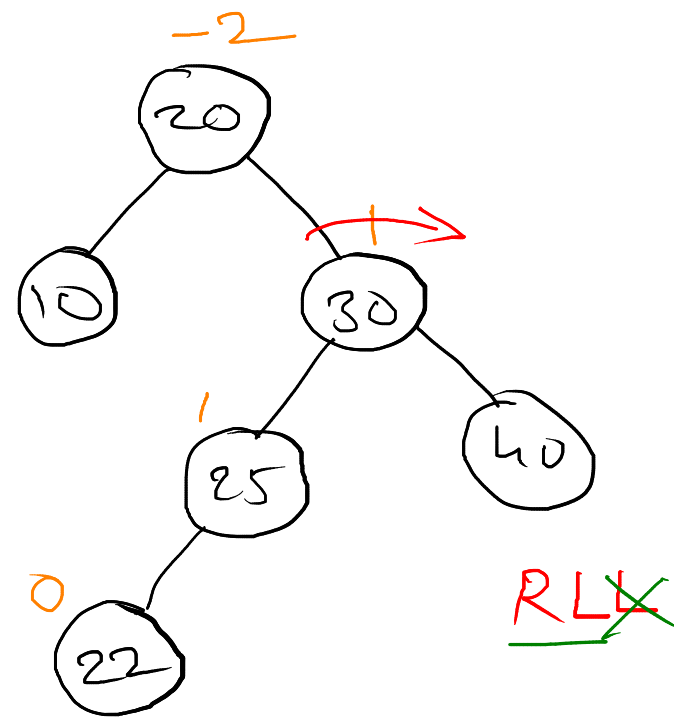
## AVL Tree

- Self balancing binary Search Tree
- on every insertion and deletion of node, tree is balanced
- All operation on AVL tree are performed in  $O(\log n)$  time
- Balance factor of all nodes is either -1, 0 or +1

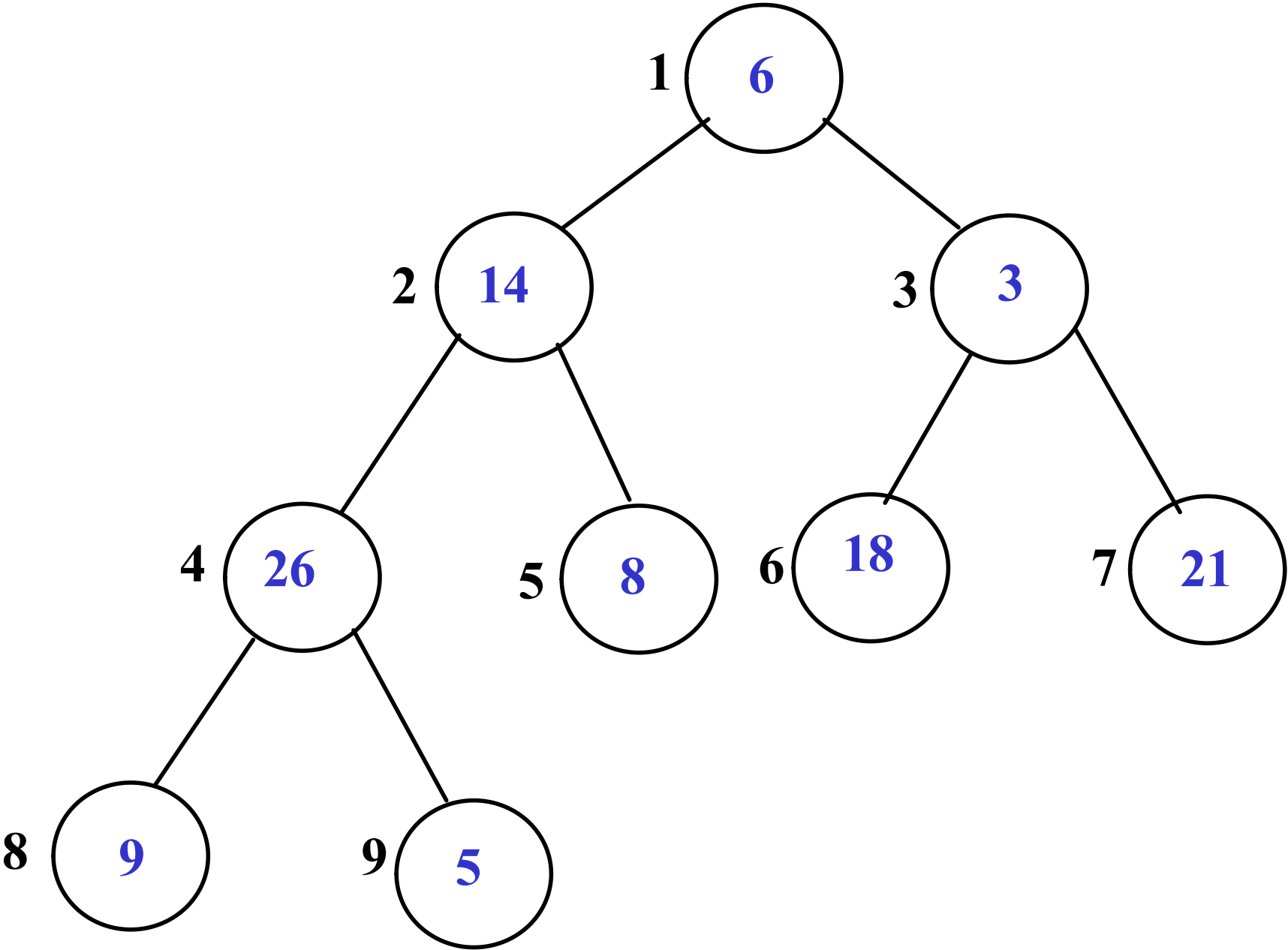
Keys : 40, 20, 10, 25, 30, 22, 50







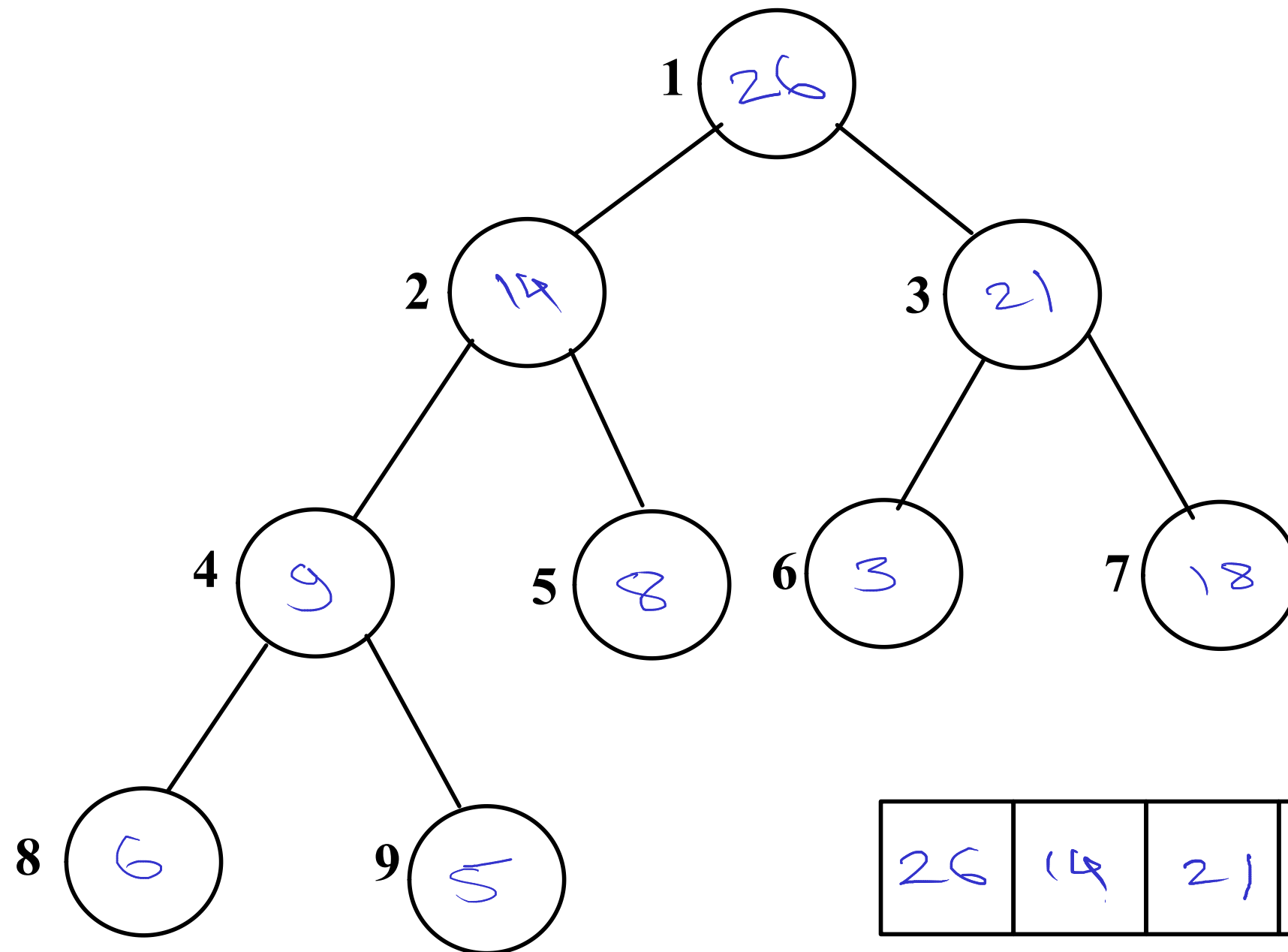
**Almost Complete Binary Tree**



6	14	3	26	8	18	21	9	5
1	2	3	4	5	6	7	8	9

## Create Max Heap

6 14 3 26 8 18 21 9 5

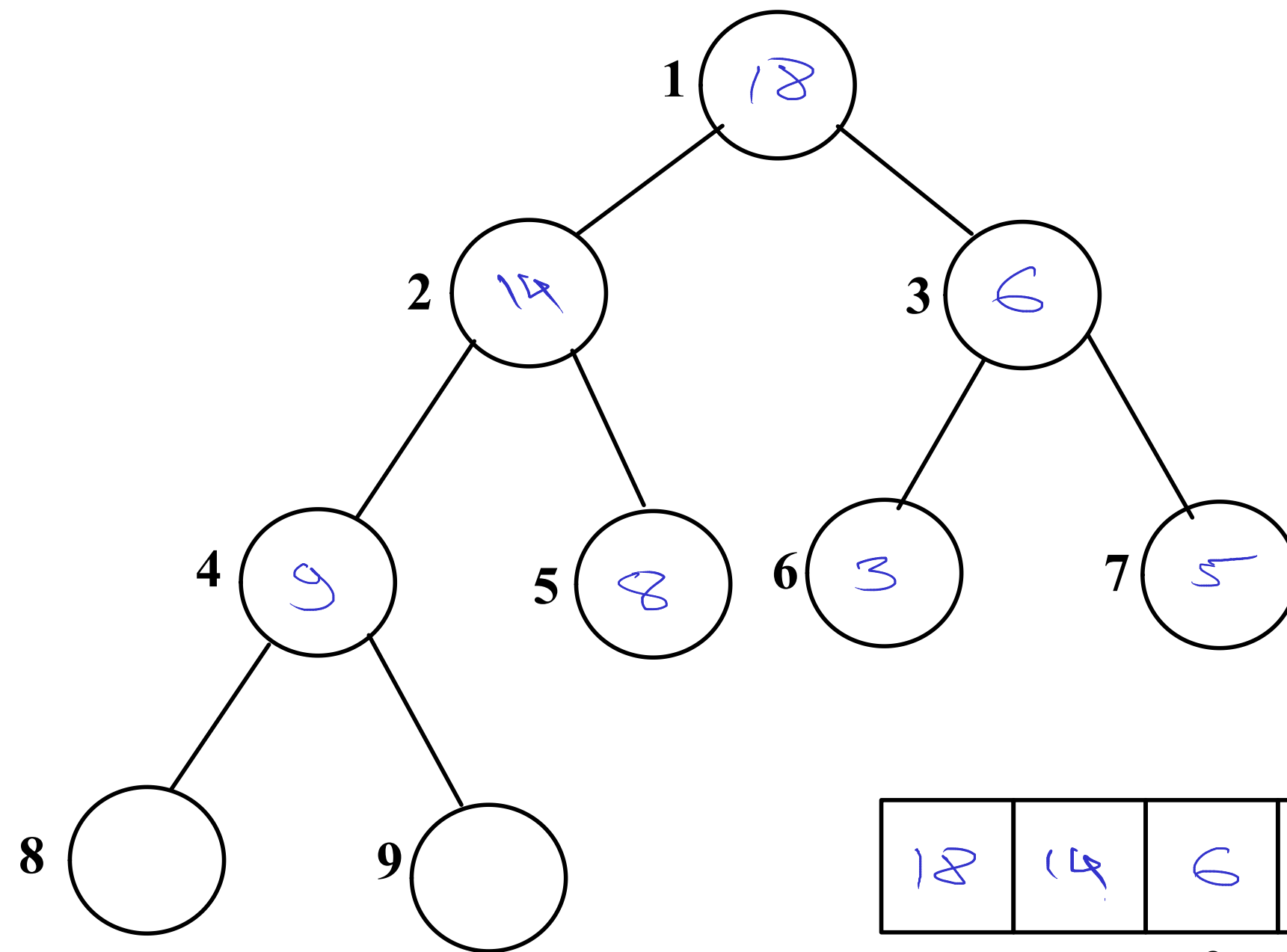


To add single element  
into heap, need  
 $O(\log n)$  time.

To add  $n$  elements  
into heap, need  
 $O(n \log n)$  time

26	14	21	9	8	3	18	6	5
1	2	3	4	5	6	7	8	9

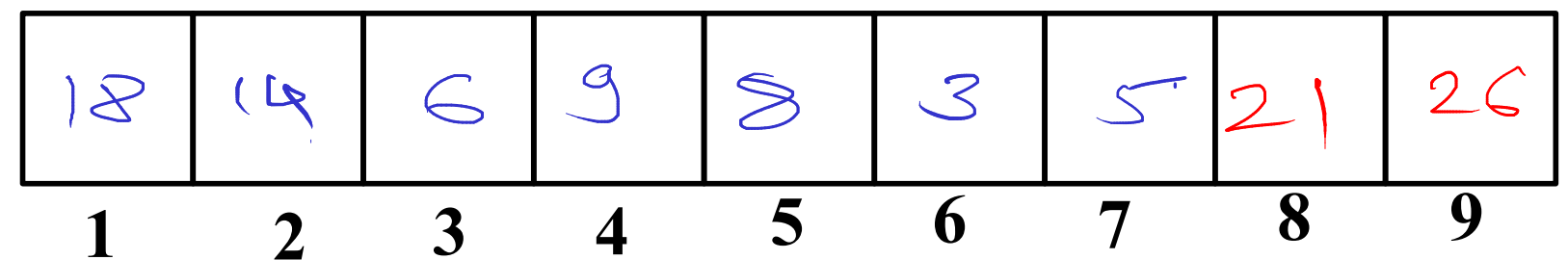
## Delete Max Heap



max = 21

To delete single element from heap, need  $O(\log n)$  time

To delete  $n$  elements from heap needs,  $O(n \log n)$  time.



Heap Sort Time Complexity

1) Create heap -  $n \log n$

2) Delete heap -  $n \log n$

$2n \log n$

$T(n) = O(n \log n)$