

# Sorting

- arrangement of data elements either in ascending or descending

order of their values

1. Selection sort

2. Bubble sort

3. Insertion sort

4. Merge sort

5. Quick sort

6. Heap sort

$$\text{No. of comps} = 1 + 2 + 3 + \dots + n - 1$$

$$= \frac{n(n+1)}{2}$$

$$= \frac{1}{2}(n^2 + n)$$

$$\text{Time} \propto \frac{1}{2}(n^2 + n)$$

$$\text{Time} = n^2 + n$$

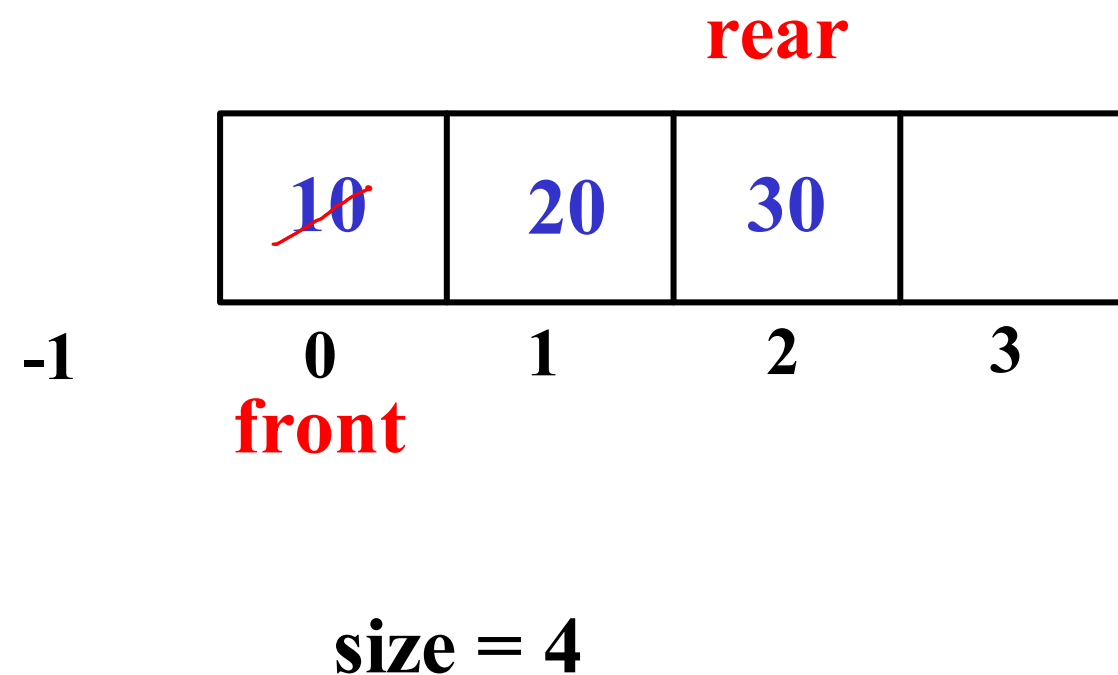
$$\text{Time} = n^2$$

$$T(n) = O(n^2)$$

	Selection sort	Bubble sort	Insertion sort
Best case	$O(n^2)$	$O(n)$	$O(n)$
Average case	$O(n^2)$	$O(n^2)$	$O(n^2)$
Worst case	$O(n^2)$	$O(n^2)$	$O(n^2)$

## Linear Queue

- linear data structure which is used to store similar type data
- data can be inserted from one end, that end is called rear
- data can be removed from another end, that end is called as front
- work on the principle of "First In First Out" (FIFO)



### Conditions:

1. Full :  $\text{rear} == \text{size} - 1$
2. Empty :  $\text{front} == \text{rear}$

### Operations:

#### 1. Insert/Add/Enqueue/Push:

- a. reposition the rear (inc)
- b. add data at rear index

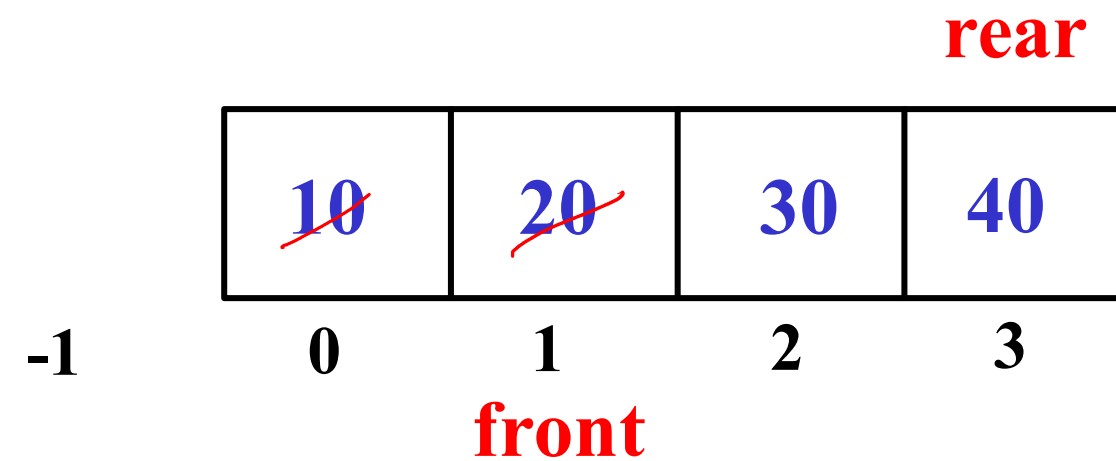
#### 2. Remove/Delete/Dequeue/pop:

- a. reposition the front (inc)

#### 3. Peek (collect) :

- a. read/return data of  $\text{front} + 1$  index

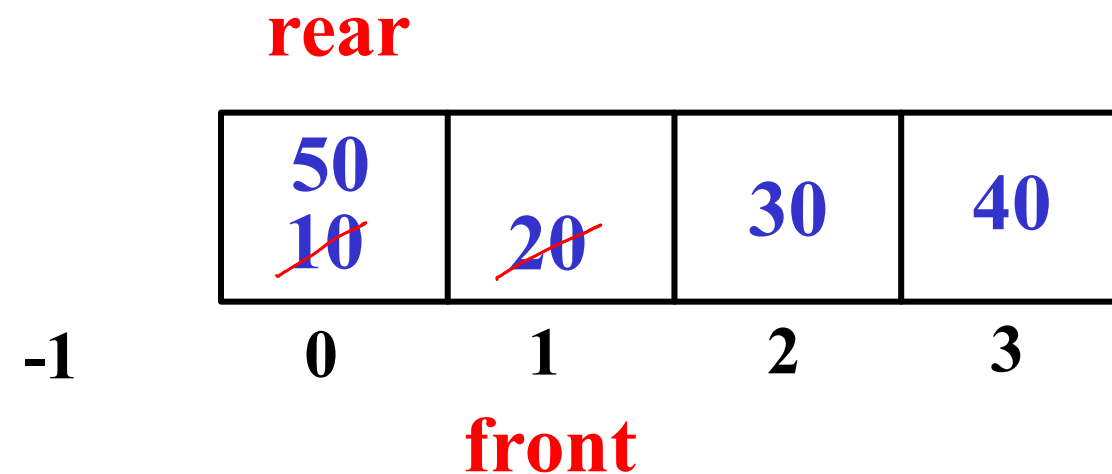
## Linear Queue



- once rear is reached to last index of array, queue is full
- if initial few locations are free, we can not use them to add next data
- this will lead to poor memory utilization
- solution for above problem is circular queue

# Circular Queue

size = 4



$\text{rear} = (\text{rear} + 1) \% \text{size}$   
 $\text{front} = (\text{front} + 1) \% \text{size}$

fornt = rear = -1

$(-1 + 1) \% 4 \rightarrow 0$

$(0 + 1) \% 4 \rightarrow 1$

$(1 + 1) \% 4 \rightarrow 2$

$(2 + 1) \% 4 \rightarrow 3$

$(3 + 1) \% 4 \rightarrow 0$

## Operations:

### 1. Insert/Add/Enqueue/Push:

- reposition the rear (inc)
- add data at rear index

### 2. Remove/Delete/Dequeue/pop:

- reposition the front (inc)

### 3. Peek (collect) :

- read/return data of front + 1 index

## Conditions

1. Full : count == SIZE

2. Empty : count == 0

## Circular Queue - Empty

**rear**



**-1**

**0**

**1**

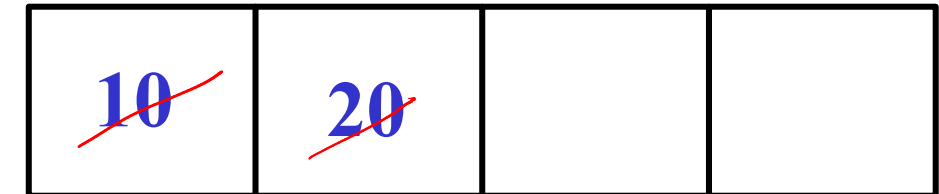
**2**

**3**

**front**

**front == rear && rear == -1**

**rear**



**-1**

**0**

**1**

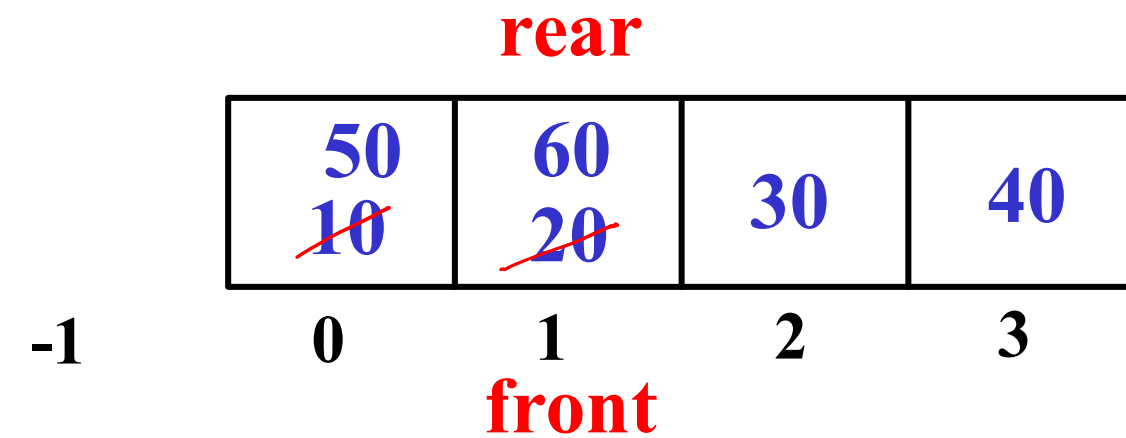
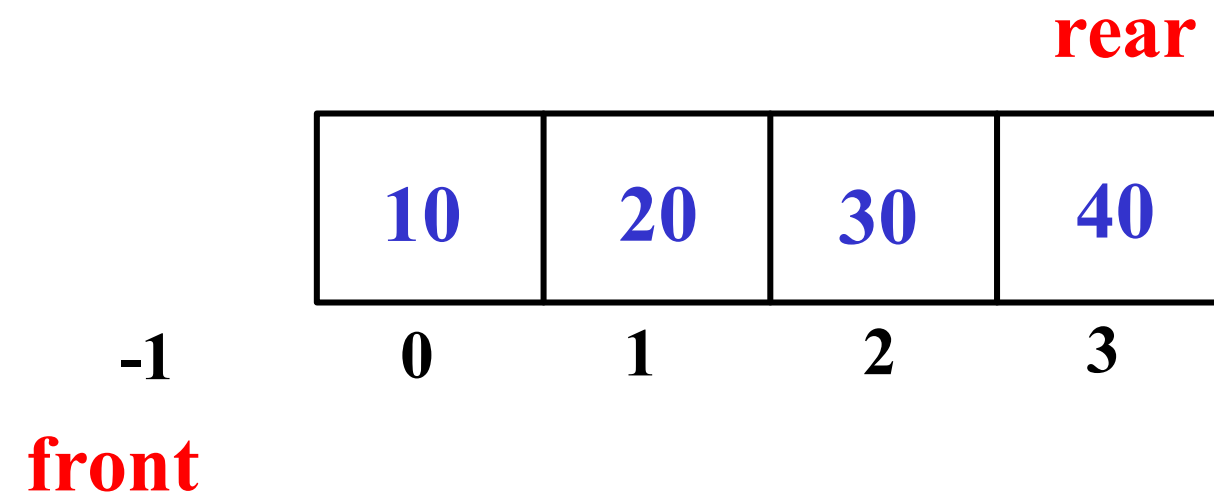
**2**

**3**

**front**

```
pop(){  
    fornt = (front + 1) % SIZE;  
    if(front == rear)  
        front = rear = -1;  
}
```

## Circular Queue - Full



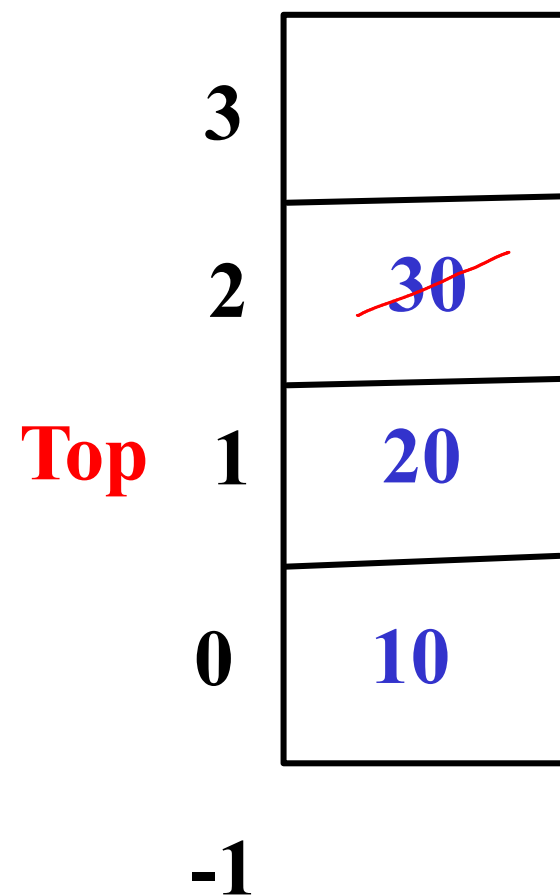
$\text{rear} == \text{SIZE} - 1 \ \&\& \ \text{front} == -1$

$\text{front} == \text{rear} \ \&\& \ \text{rear} \neq -1$

$(\text{rear} == \text{SIZE} - 1 \ \&\& \ \text{front} == -1) \ || \ (\text{front} == \text{rear} \ \&\& \ \text{rear} \neq -1)$

# Stack

- linear data structure which is used to store similar type of data
- insert and remove of data is allowed from only one end (top)
- works on the principle of "Last In First Out" (LIFO)



## Operations:

### 1. Add/Insert/push:

- a. reposition top (inc)
- b. add data at top index

### 2. Delete/Remove/pop:

- a. reposition top (dec)

### 3. peek (Collect):

- a. read/return data of top index

## Condition:

1. Full :  $\text{top} == \text{SIZE} - 1$
2. Empty :  $\text{top} == -1$