

Games and adversarial search



World Champion chess player Garry Kasparov is defeated by IBM's Deep Blue chess-playing computer in a six-game match in May, 1997

Why study games?

- Games are a traditional hallmark of intelligence
- Games are easy to formalize
- Games can be a good model of real-world competitive or cooperative activities
 - Military confrontations, negotiation, auctions, etc.

Types of game environments

	Deterministic	Stochastic
Perfect information (fully observable)	Chess, checkers, go	Backgammon, monopoly
Imperfect information (partially observable)	Battleships	Scrabble, poker, bridge

Alternating two-player zero-sum games

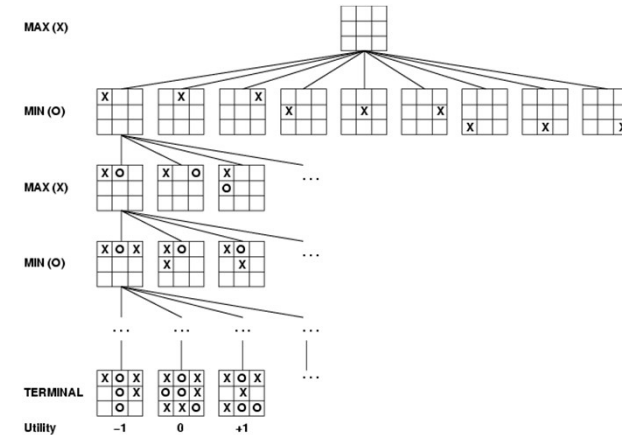
- Players take turns
- Each game outcome or **terminal state** has a **utility** for each player (e.g., 1 for win, 0 for loss)
- The sum of both players' utilities is a constant

Games vs. single-agent search

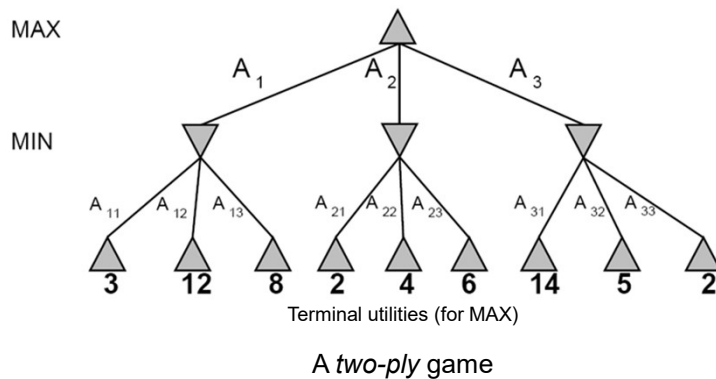
- We don't know how the opponent will act
 - The solution is not a fixed sequence of actions from start state to goal state, but a **strategy** or **policy** (a mapping from state to best move in that state)
- Efficiency is critical to playing well
 - The time to make a move is limited
 - The branching factor, search depth, and number of terminal configurations are huge
 - In chess, branching factor ≈ 35 and depth ≈ 100 , giving a search tree of 10^{154} nodes
 - Number of atoms in the observable universe $\approx 10^{80}$
 - This rules out searching all the way to the end of the game

Game tree

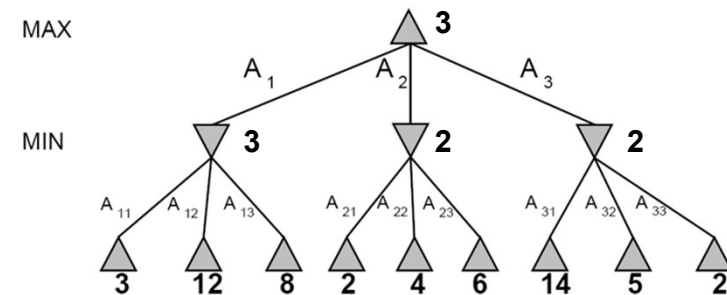
- A game of tic-tac-toe between two players, “max” and “min”



A more abstract game tree

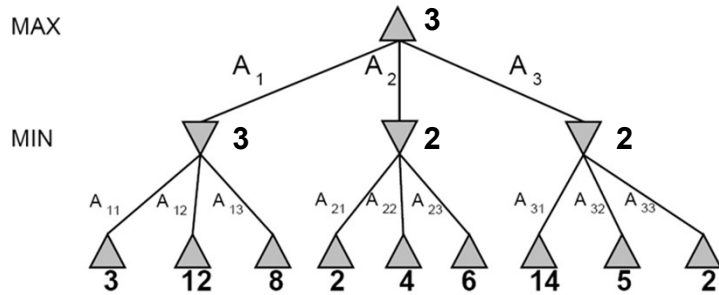


Game tree search



- **Minimax value of a node:** the utility (for MAX) of being in the corresponding state, assuming perfect play on both sides
- **Minimax strategy:** Choose the move that gives the best worst-case payoff

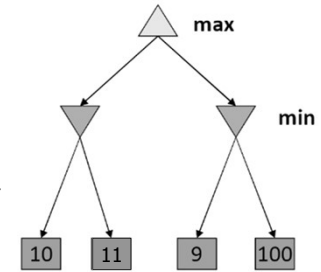
Computing the minimax value of a node



- **Minimax**(node) =
 - Utility(node) if node is terminal
 - $\max_{action} \text{Minimax}(\text{Succ}(\text{node}, \text{action}))$ if player = MAX
 - $\min_{action} \text{Minimax}(\text{Succ}(\text{node}, \text{action}))$ if player = MIN

Optimality of minimax

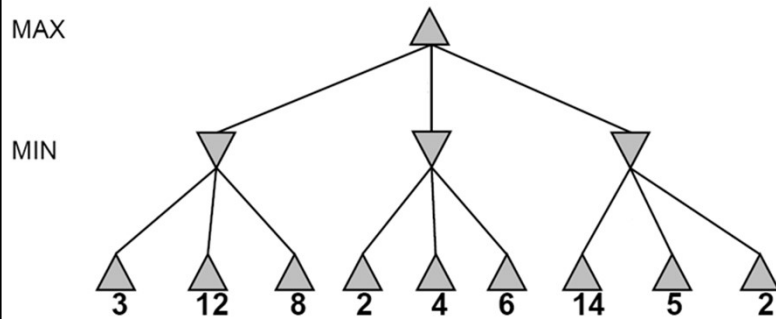
- The minimax strategy is optimal against an optimal opponent
- What if your opponent is suboptimal?
 - Your utility can only be higher than if you were playing an optimal opponent!
 - A different strategy may work better for a sub-optimal opponent, but it will necessarily be worse against an optimal opponent



Example from D. Klein and P. Abbeel

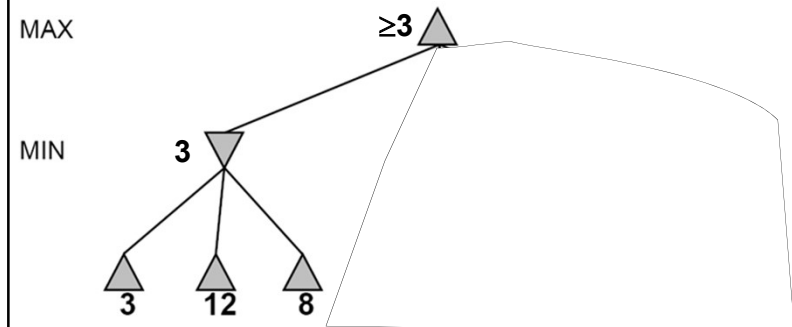
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



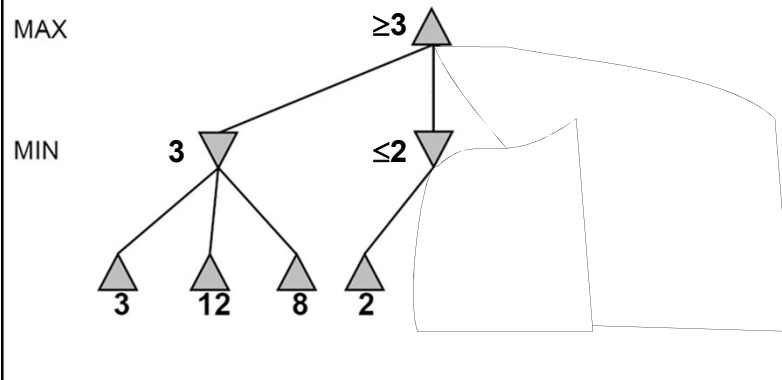
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



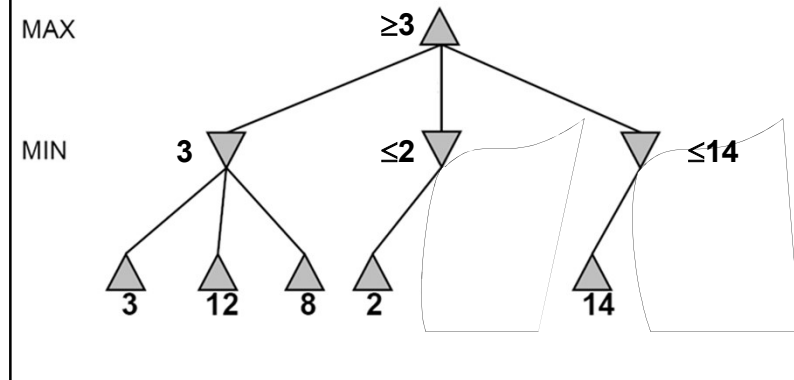
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



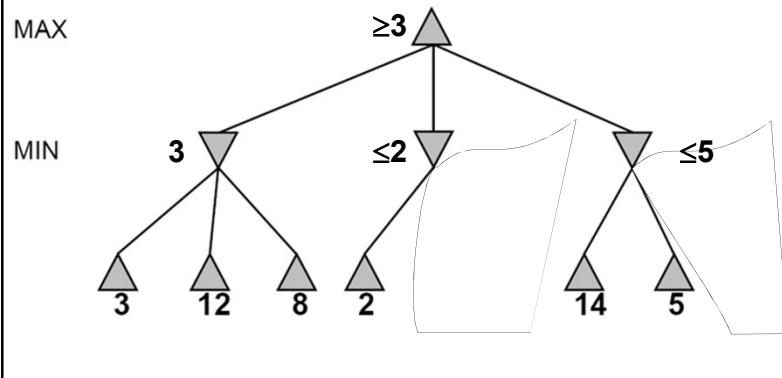
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



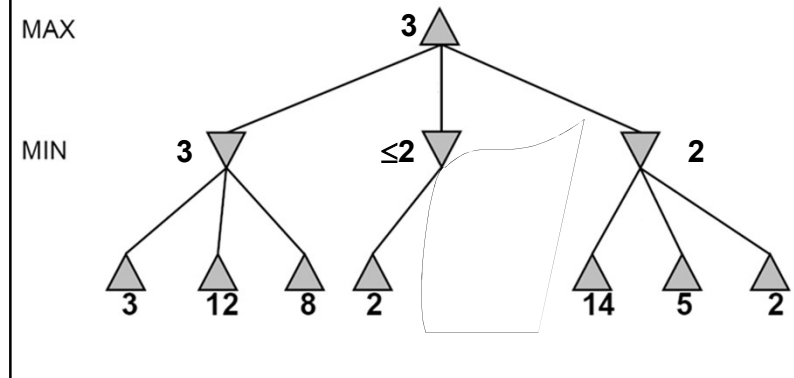
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



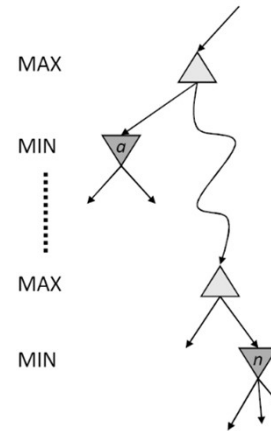
Alpha-beta pruning

- It is possible to compute the exact minimax decision without expanding every node in the game tree



Alpha-beta pruning

- α is the value of the best choice for the MAX player found so far at any choice point above node n
- We want to compute the MIN-value at n
- As we loop over n 's children, the MIN-value decreases
- If it drops below α , MAX will never choose n , so we can ignore n 's remaining children
- Analogously, β is the value of the lowest-utility choice found so far for the MIN player

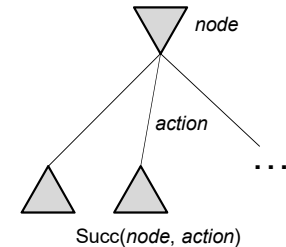


Alpha-beta pruning

Function $action = \text{Alpha-Beta-Search}(node)$
 $v = \text{Min-Value}(node, -\infty, \infty)$
 return the $action$ from $node$ with value v

α : best alternative available to the Max player
 β : best alternative available to the Min player

Function $v = \text{Min-Value}(node, \alpha, \beta)$
 if $\text{Terminal}(node)$ return $\text{Utility}(node)$
 $v = +\infty$
 for each $action$ from $node$
 $v = \text{Min}(v, \text{Max-Value}(\text{Succ}(node, action), \alpha, \beta))$
 if $v \leq \alpha$ return v
 $\beta = \text{Min}(\beta, v)$
 end for
 return v

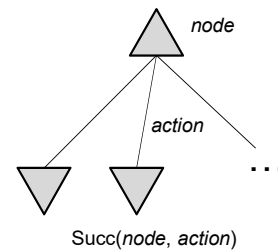


Alpha-beta pruning

Function $action = \text{Alpha-Beta-Search}(node)$
 $v = \text{Max-Value}(node, -\infty, \infty)$
 return the $action$ from $node$ with value v

α : best alternative available to the Max player
 β : best alternative available to the Min player

Function $v = \text{Max-Value}(node, \alpha, \beta)$
 if $\text{Terminal}(node)$ return $\text{Utility}(node)$
 $v = -\infty$
 for each $action$ from $node$
 $v = \text{Max}(v, \text{Min-Value}(\text{Succ}(node, action), \alpha, \beta))$
 if $v \geq \beta$ return v
 $\alpha = \text{Max}(\alpha, v)$
 end for
 return v



Alpha-beta pruning

- Pruning does not affect final result
- Amount of pruning depends on move ordering
 - Should start with the “best” moves (highest-value for MAX or lowest-value for MIN)
 - For chess, can try captures first, then threats, then forward moves, then backward moves
 - Can also try to remember “killer moves” from other branches of the tree
- With perfect ordering, the time to find the best move is reduced to $O(b^{m/2})$ from $O(b^m)$
 - Depth of search is effectively doubled

Evaluation function

- Cut off search at a certain depth and compute the value of an **evaluation function** for a state instead of its minimax value
 - The evaluation function may be thought of as the probability of winning from a given state or the *expected value* of that state
- A common evaluation function is a weighted sum of *features*:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

- For chess, w_k may be the **material value** of a piece (pawn = 1, knight = 3, rook = 5, queen = 9) and $f_k(s)$ may be the advantage in terms of that piece
- Evaluation functions may be *learned* from game databases or by having the program play many games against itself

Cutting off search

- **Horizon effect:** you may incorrectly estimate the value of a state by overlooking an event that is just beyond the depth limit
 - For example, a damaging move by the opponent that can be delayed but not avoided
- Possible remedies
 - **Quiescence search:** do not cut off search at positions that are unstable – for example, are you about to lose an important piece?
 - **Singular extension:** a strong move that should be tried when the normal depth limit is reached

Advanced techniques

- **Transposition table** to store previously expanded states
- **Forward pruning** to avoid considering all possible moves
- **Lookup tables** for opening moves and endgames