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CAT-11 CS11003

SLOT: C₂ + TC₂

1) $S \rightarrow XY | W$

$$X \rightarrow aXb \mid \epsilon$$

$$Y \rightarrow cY \mid \epsilon$$

$$W \rightarrow aWc \mid z$$

$$Z \rightarrow bZ \mid \epsilon$$

Simplification:-

1) Useless production (removal)

* No useless production

2) Null production (removal)

$$S \rightarrow XY | \cancel{W} | \cancel{Y} | \cancel{\epsilon}$$

$$X \rightarrow aXb \mid ab$$

$$Y \rightarrow cY \mid$$

$$W \rightarrow aWc \mid ac \mid \cancel{\epsilon}$$

$$Z \rightarrow bZ \mid b$$

3) Unit production (removal):-

$$S \rightarrow \cancel{XY} | aWc | ac | bz | b | cY | aXb | ab$$

$$X \rightarrow aXb \mid ab$$

$$Y \rightarrow cY$$

$$W \rightarrow aWc \mid ac \mid bz \mid \cancel{b}$$

$$Z \rightarrow bz \mid \cancel{b}$$

$$S \rightarrow XY, W \rightarrow b \neq Z \rightarrow b, S \rightarrow bz \mid b$$

already in CNF

$$S \rightarrow aWc / ac / \textcircled{cY} / aXb / ab$$

$$X \rightarrow aXb / ab$$

$$C \rightarrow \textcircled{cY}$$

$$W \rightarrow aWc / ac / \textcircled{bZ}$$

$$Z \rightarrow \textcircled{bZ}$$

$$\cancel{S \rightarrow aWE / aE / aXF / aF}$$

$$\cancel{X \rightarrow aX}$$

$$S \rightarrow aWE / aE / aXF / aF$$

$$X \rightarrow aXF / aF$$

$$W \rightarrow aWE / aE /$$

$$S \rightarrow aS_1 / aE / aS_2 / aF, S_1 \rightarrow WE$$

$$X \rightarrow aS_3 / aF, S_2 \rightarrow XF$$

$$W \rightarrow aS_1 / aE, S_3 \rightarrow XF$$

∴ The CNF are

$$S \rightarrow XY, W \rightarrow b, Z \rightarrow b, S \rightarrow bZ / b$$

$$C \rightarrow cY$$

$$C \rightarrow cY$$

$$W \rightarrow bZ$$

$$Z \rightarrow bZ$$

2) string = boy / likes / the / small / cat

ϕ	ϕ			
N	V	A	AJ	N

boy likes the small cat

$$\begin{aligned}
 (x_i, j) &= (x_i, i, x_{i+j}) \cup (x_i, i+1, x_{i+2}, j) \\
 &\cup (x_i, i+2, x_{i+3}, j) \cup \\
 &(x_i, j-1, x_j, j)
 \end{aligned}$$

3). $L(A) = \{0^n w w 0^n \mid n \geq 0, w \in \{a, b\}^*\}$

valid string = 0a a0, 0abab0, ababab

invalid string = 0ab a0, 0aba, 00ab a000

4).

$$L_1 = \{b^{2n} \mid n \geq 0\}$$

$$L_2 = \{a^n b^m \mid n \geq 0, m \geq 0\}$$

$L_1 \rightarrow$ regular language

$L_2 \rightarrow$ regular language

then $L_1 \cup L_2$ must also be a regular

language.

L_1 has regular expression of $(bb)^*$

L_2 has regular expression of $(a^* b^*)$

so $L_1 \cup L_2$ has regular expression of

$$(bb)^* + (a^* b^*)$$

The Regular language only have regular expression. It has regular expression. Hence it is a regular language