

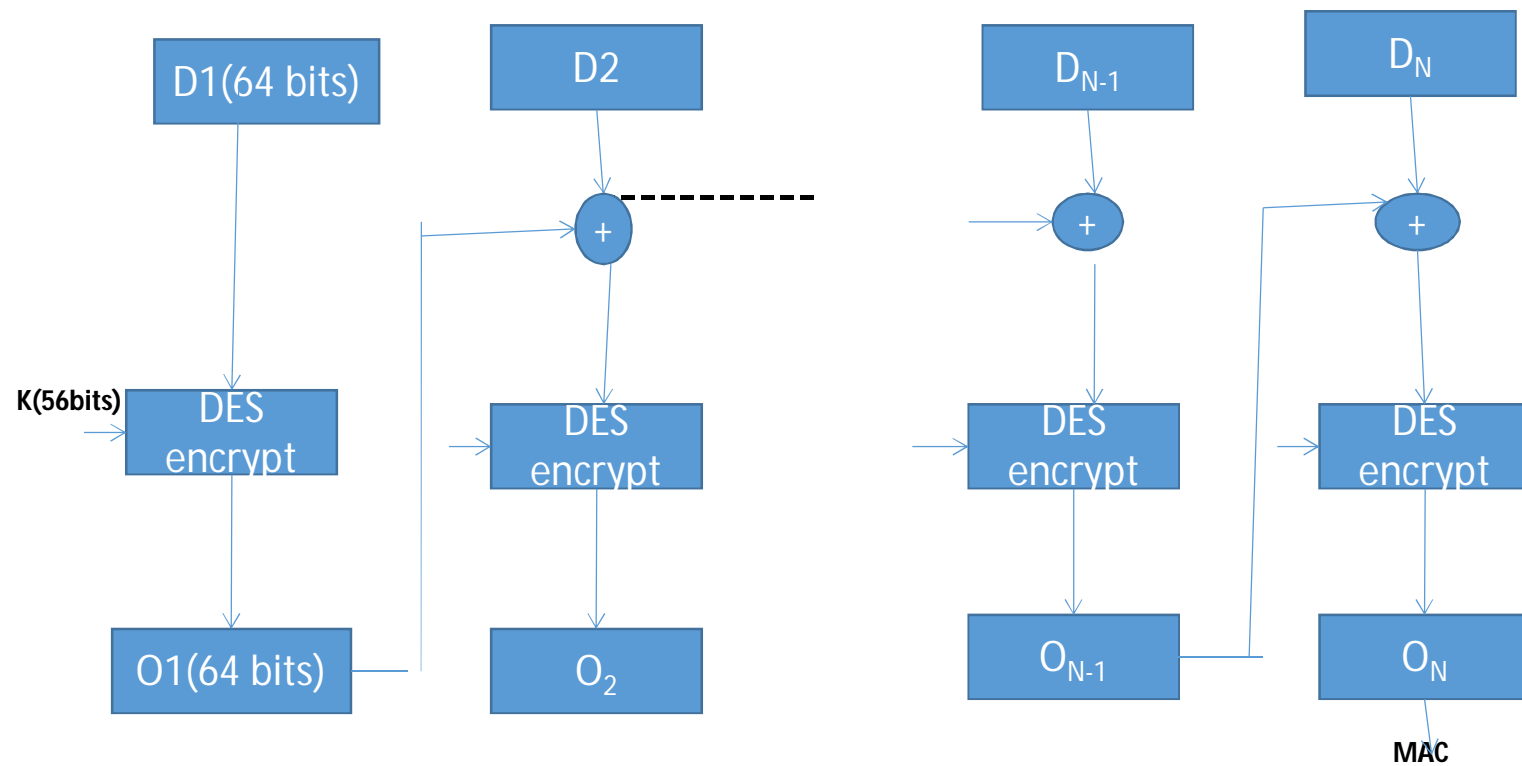
Authentication

- It is a mechanism to verify the integrity of the transmitted message or person initiating it.
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- Entity authentication
 - Data authentication

- Message encryption
- Message Authentication Code(MAC)
- Hash functions

MAC

- Data Authentication Algorithm:



Using Hash for Authentication

- A computes Hash value using hash algorithm and concatenates with original data and send to B
- B computes Hash value using same algorithm on the received original data and compares the computed hash value with received hash value
- If both are same then message is authenticated.

Hashing

Properties of Hash function:

1. H can be applied to a block of data of any size
2. H produces a fixed length output
3. $H(M)$ is relatively easy to compute for any given M
4. For any given value h, it is computationally infeasible to find M such that $H(M)=h$. This is sometimes referred as one way function (**Preimage resistant**)
5. For any given message M, it is computationally intractable to find $M^1 \neq M$ such that $H(M^1)=H(M)$::
weak collision resistance(2nd Preimage resistant)
6. It is computationally intractable to find any pair (M, M^1) such that $H(M)=H(M^1)$::**strong collision resistance(Collision resistance)**

Requirements for Hash Functions

1. can be applied to any sized message M
2. produces fixed-length output h
3. is easy to compute $h=H(M)$ for any message M
4. given h is infeasible to find x s.t. $H(x)=h$
 - *one-way property*
5. given x is infeasible to find y s.t. $H(y)=H(x)$
 - *weak collision resistance*
6. is infeasible to find any x,y s.t. $H(y)=H(x)$
 - *strong collision resistance*

MD5: Message Digest Version 5

Input message(Variable size)



Output 128 bits(Fixed)

How MD5 works?

- Step 1: Padding
- Step2 : Append length
- Step3 :Devide the input into 512-bit blocks
- Step4: Intialize chaining varaibles

A	Hex 01	25	45	67
B	Hex 89	AB	CD	EF
C	Hex FE	DC	BA	98
D	Hex 76	54	32	10

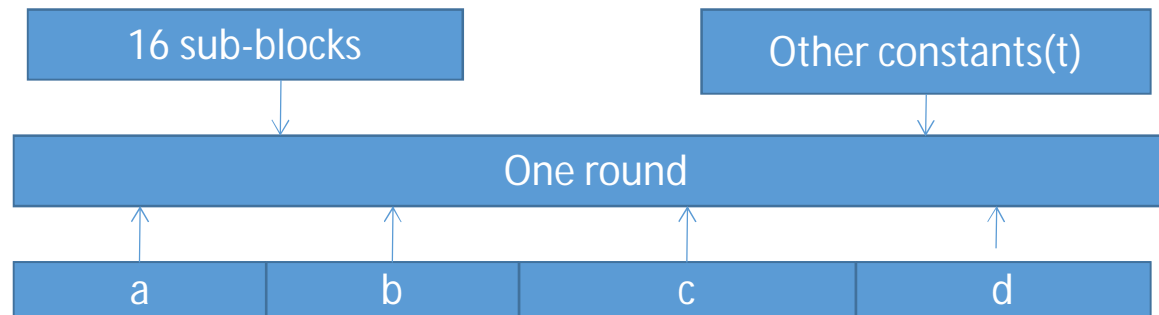
- Step 5: Process blocks

5.1: copy the chaining variables into four corresponding variables a,b,c and d

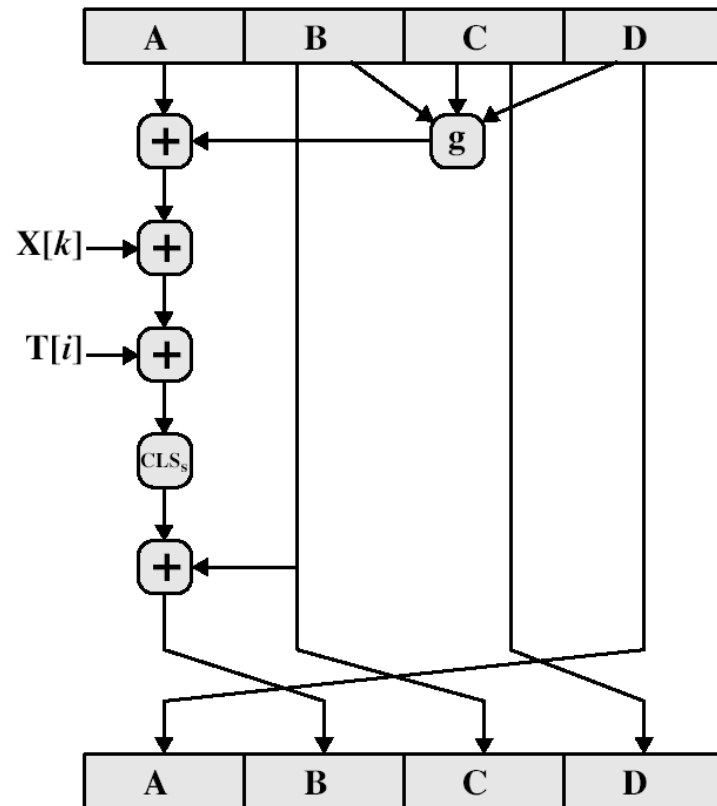
5.2 :Divide the current 512 bit block into 16 sub-blocks(size-32bits)

5.3: Now we have four rounds. In each round we process all the 16 sub-blocks belonging to a block

Conceptual process within a round



- In each round we have 16 input sub-blocks, named $M[0], M[1], \dots, M[15]$.
- Also, t is an array of constants. It consists of 64 elements, with each element consisting of 32 bits. We denote the elements of this array t as $t[0], t[1], \dots, t[63]$.
- **$\text{floor}(2^{32} \times \text{abs}(\sin(i + 1)))$**



- Each round has 16 steps of the form:

$$a = b + ((a + g(b, c, d) + X[k] + T[i]) \lll s)$$

- $g(d, b, c) = (b \wedge c) \vee (\sim b \wedge d)$
- $g(d, b, c) = (b \wedge d) \vee (c \wedge \sim d)$
- $g(d, b, c) = b \oplus c \oplus d$
- $g(d, b, c) = c \oplus (b \wedge \sim d)$

How SHA-1 works?

- Step 1: Padding
- Step2 : Append length
- Step3 :Divide the input into 512-bit blocks
- Step4: Initialize chaining varaibles

A	Hex 01	25	45	67
B	Hex 89	AB	CD	EF
C	Hex FE	DC	BA	98
D	Hex 76	54	32	10
E	Hex C3	D2	E1	F0

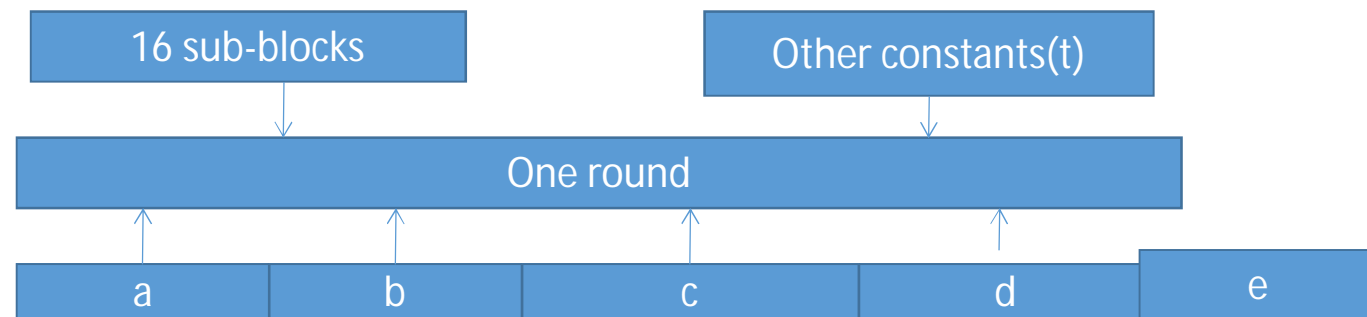
- Step 5: Process blocks

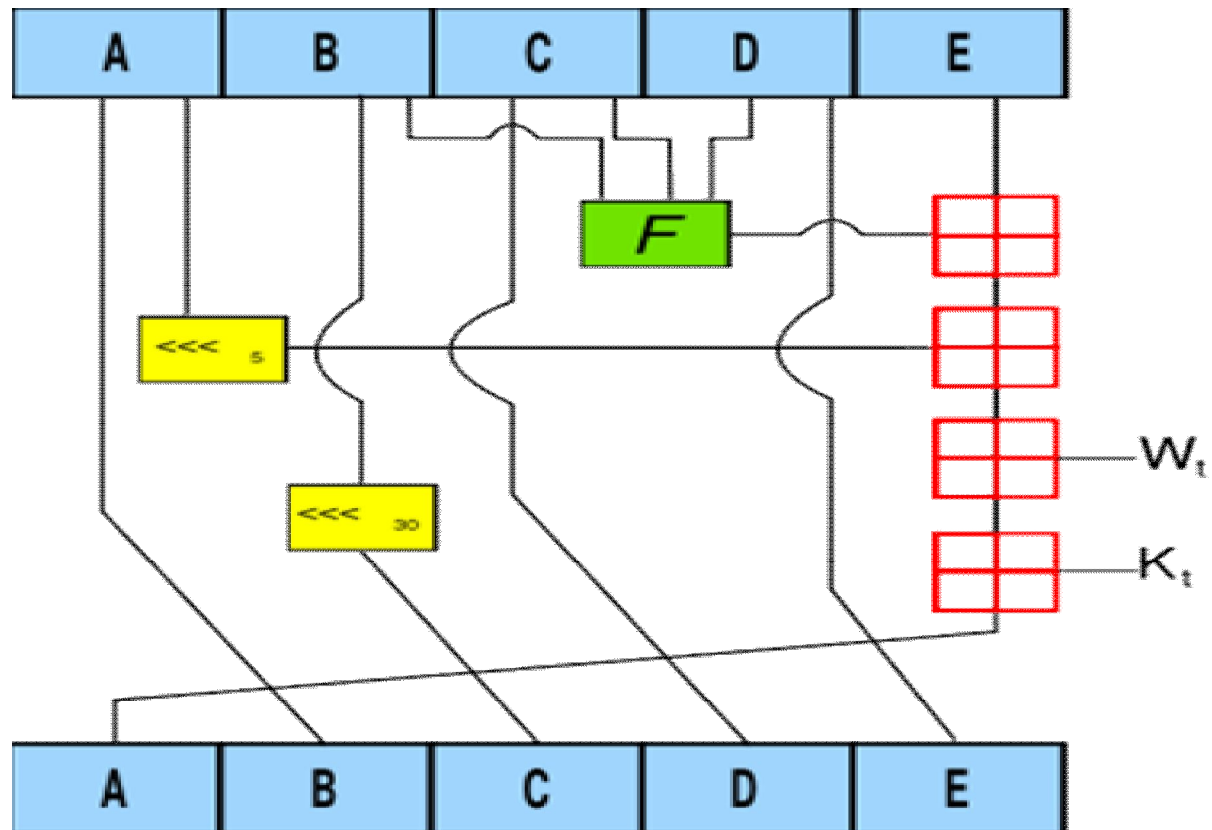
5.1: copy the chaining variables into five corresponding variables a,b,c , d and e

5.2 :Divide the current 512 bit block into 16 sub-blocks(size-32bits)

5.3: SHA-1has four rounds.each round consists of 20 steps

Conceptual process within a round





- Each round has 20 steps of the form:

$abcde = (e + \text{function } F + s^5(a) + W[t] + K[t]), a, s^{30}(b), c, d$

Round 1: $(b \text{ AND } c) \text{ OR } ((\text{NOT } b) \text{ AND } (d))$

Round 2: $B \text{ XOR } c \text{ XOR } d$

Round 3: $(b \text{ AND } c) \text{ OR } (b \text{ AND } d) \text{ OR } (c \text{ AND } d)$

Round 4: $B \text{ XOR } c \text{ XOR } d$

The remaining 64 values are defined using the equation

$W[t] = s^1(W[t-16] \text{ XOR } W[t-14] \text{ XOR } W[t-8] \text{ XOR } W[t-3])$

Round	Value of t between	K[t] in Hex
1	1 and 19	5A 92 79 99
2	20 and 39	6E D9 EB A1
3	40 and 59	9F 1B BC DC
4	60 and 79	CA 62 C1 D6

How SHA-512 works?

- Step 1: Padding
- Step2 : Append length
- Step3 :Divide the input into 1024-bit blocks
- Step4: Initialize chaining variables

A 6A09E667F3BCC908

B BB67AE8584CAA73B

.....

H 5BE0CD19137E2179

- Step 5: Process blocks

5.1: copy the chaining variables into four corresponding variables a,b,c and d

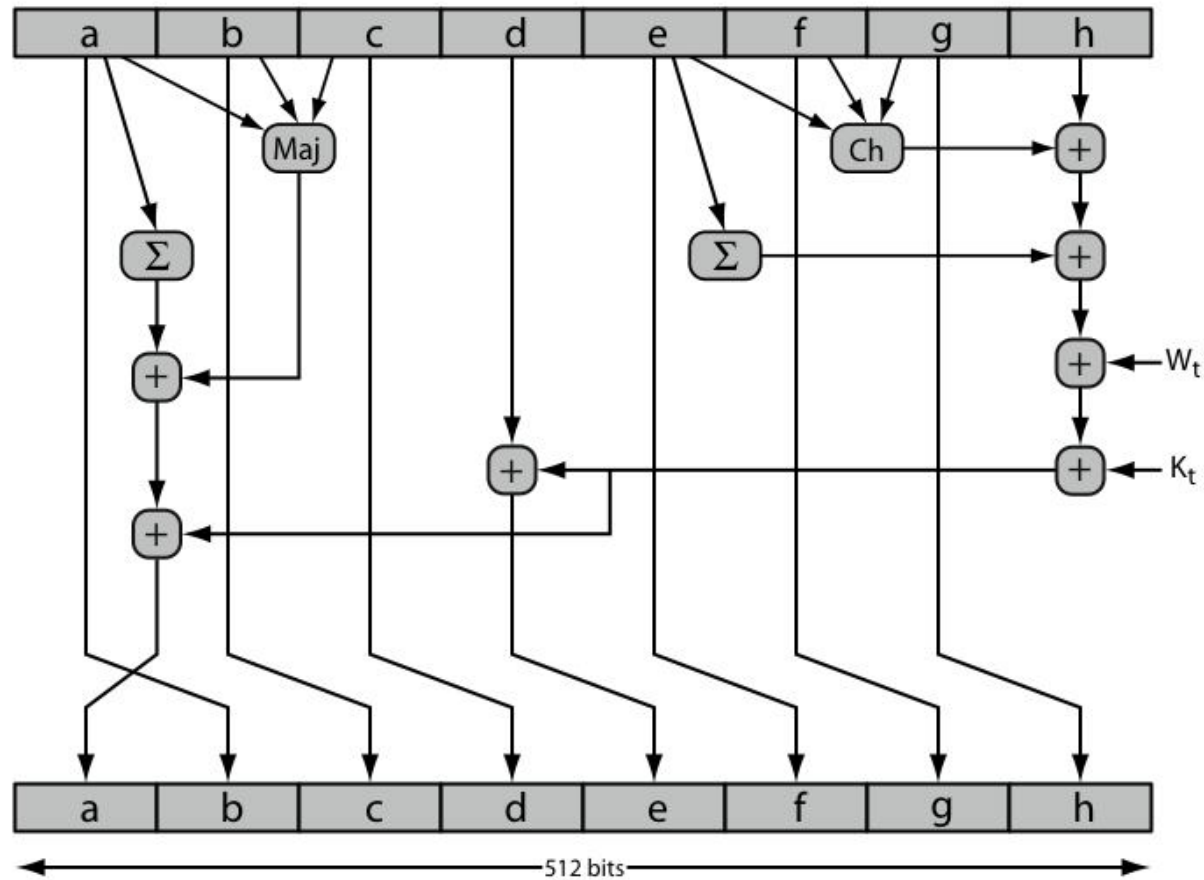
5.2 :Divide the current 1024 bit block into 16 sub-blocks(size-64 bits)

5.3: Now we have 80 rounds.In each round we process all the 16 sub-blocks belonging to a block

Conceptual process within a round



SHA-512 Round Function



$\text{Ch}(e,f,g) = (e \text{ AND } f) \text{ XOR } (\text{NOT } e \text{ AND } g)$

$\text{Maj}(a,b,c) = (a \text{ AND } b) \text{ XOR } (a \text{ AND } c) \text{ XOR } (b \text{ AND } c)$

$\text{Sum}(a) = \text{ROTR}(a \text{ by } 28 \text{ bits}) \text{ XOR } \text{ROTR}(a \text{ by } 34 \text{ bits}) \text{ XOR } \text{ROTR}(a \text{ by } 39 \text{ bits})$

$\text{Sum}(e) = \text{ROTR}(e \text{ by } 14 \text{ bits}) \text{ XOR } \text{ROTR}(e \text{ by } 18 \text{ bits}) \text{ XOR } \text{ROTR}(e \text{ by } 41 \text{ bits})$

$W[t]$ = 64-bit word derived from the current input block

$K[t]$ = constants given **in the book**

Add = addition mod 2^{64}

1. For the first 16 rounds (0 to 15), the value of $W[t]$ is equal to the corresponding word in the message block.
2. For the remaining 64 steps, the value of $W[t]$ is equal to the circular left shift by one bit of the XOR of the four preceding values of $W[t]$ with two of them subjected to the circular left shift by 1 bit.