

# Dynamic Programming Class-1

Special class

# Dynamic Programming [DP]

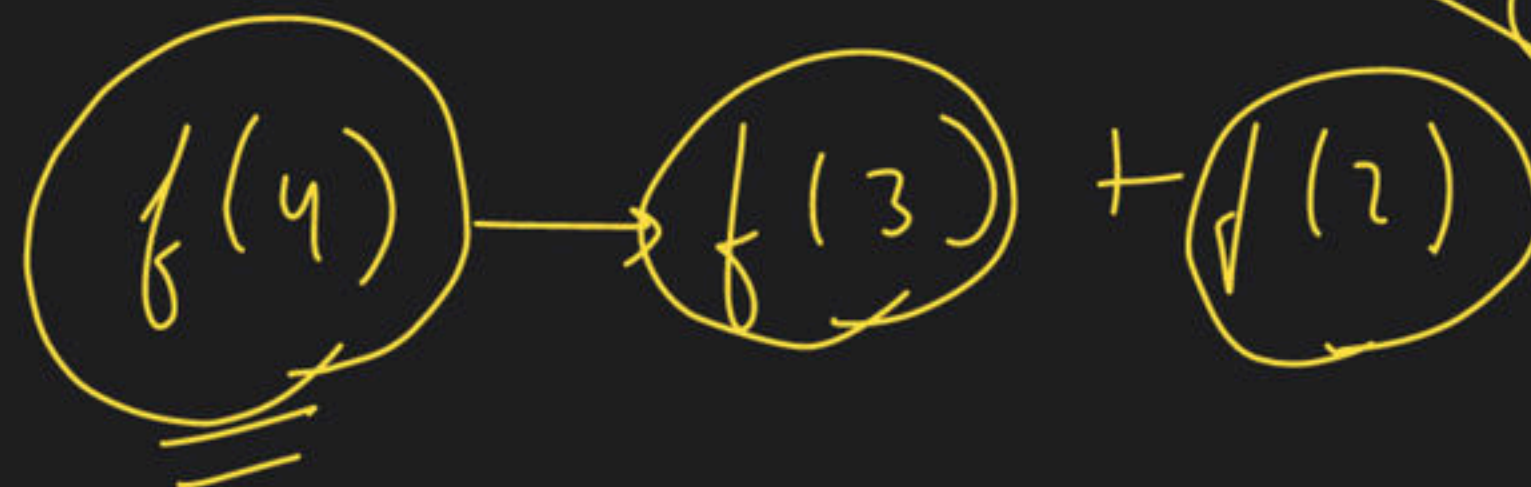
Programming Techniques

problem

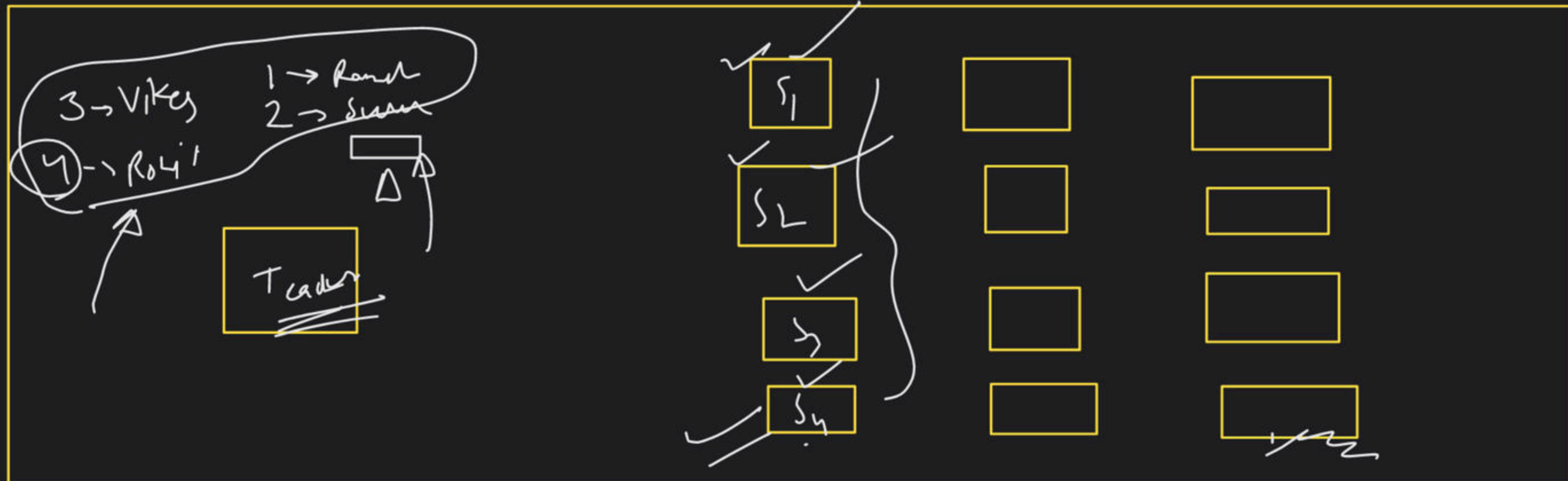
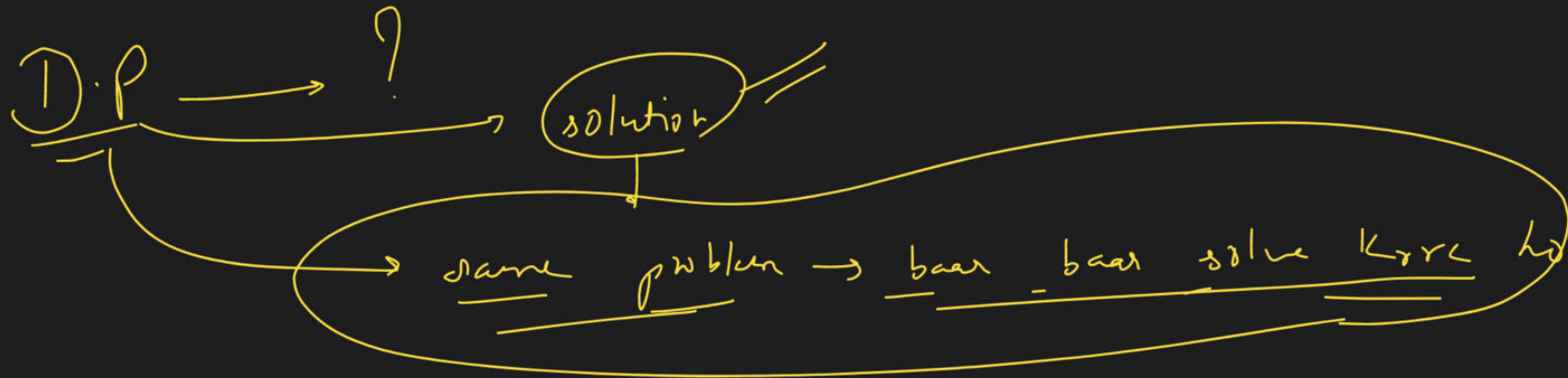
A Overlapping Subproblem

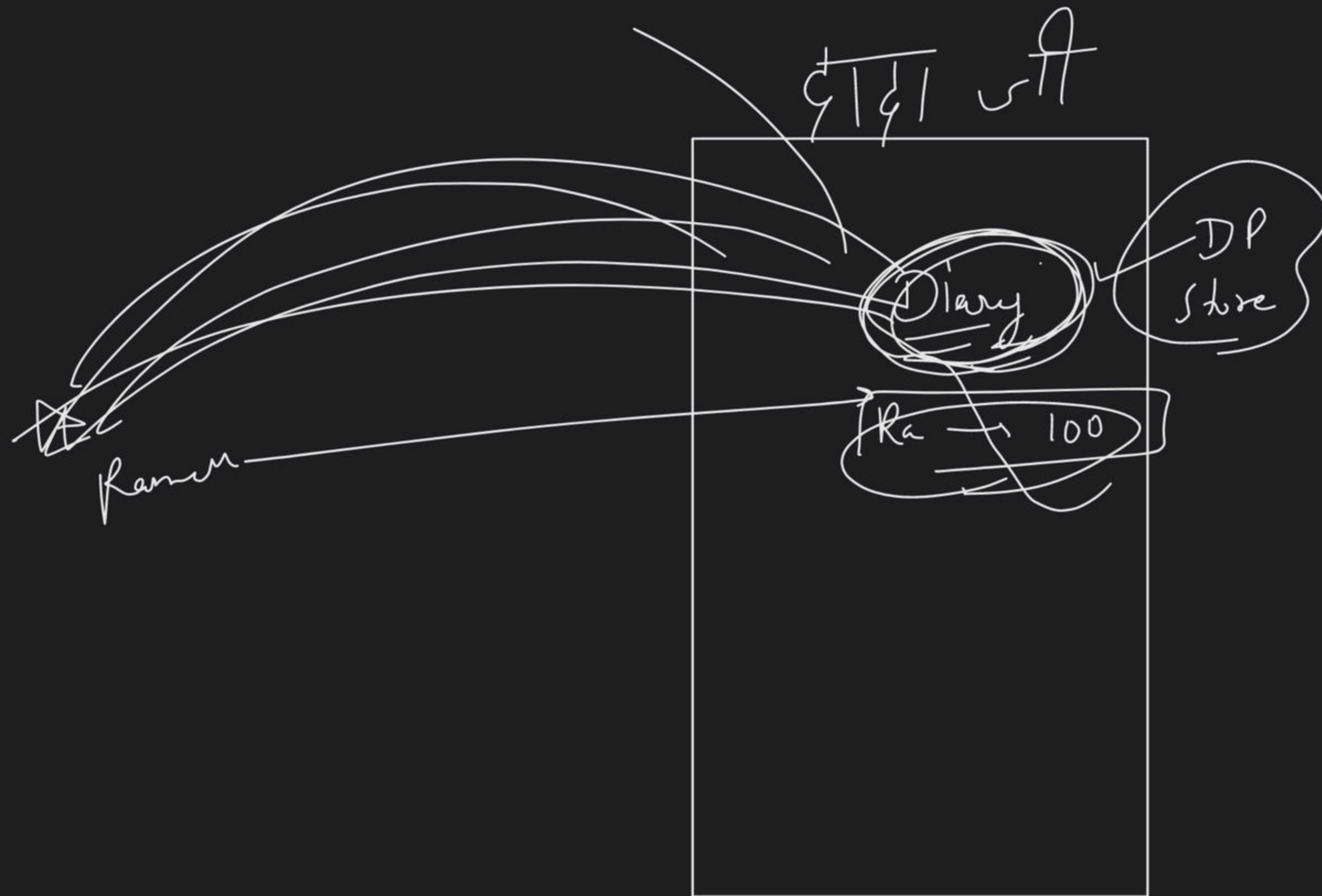
B Optimal Substructure

fib

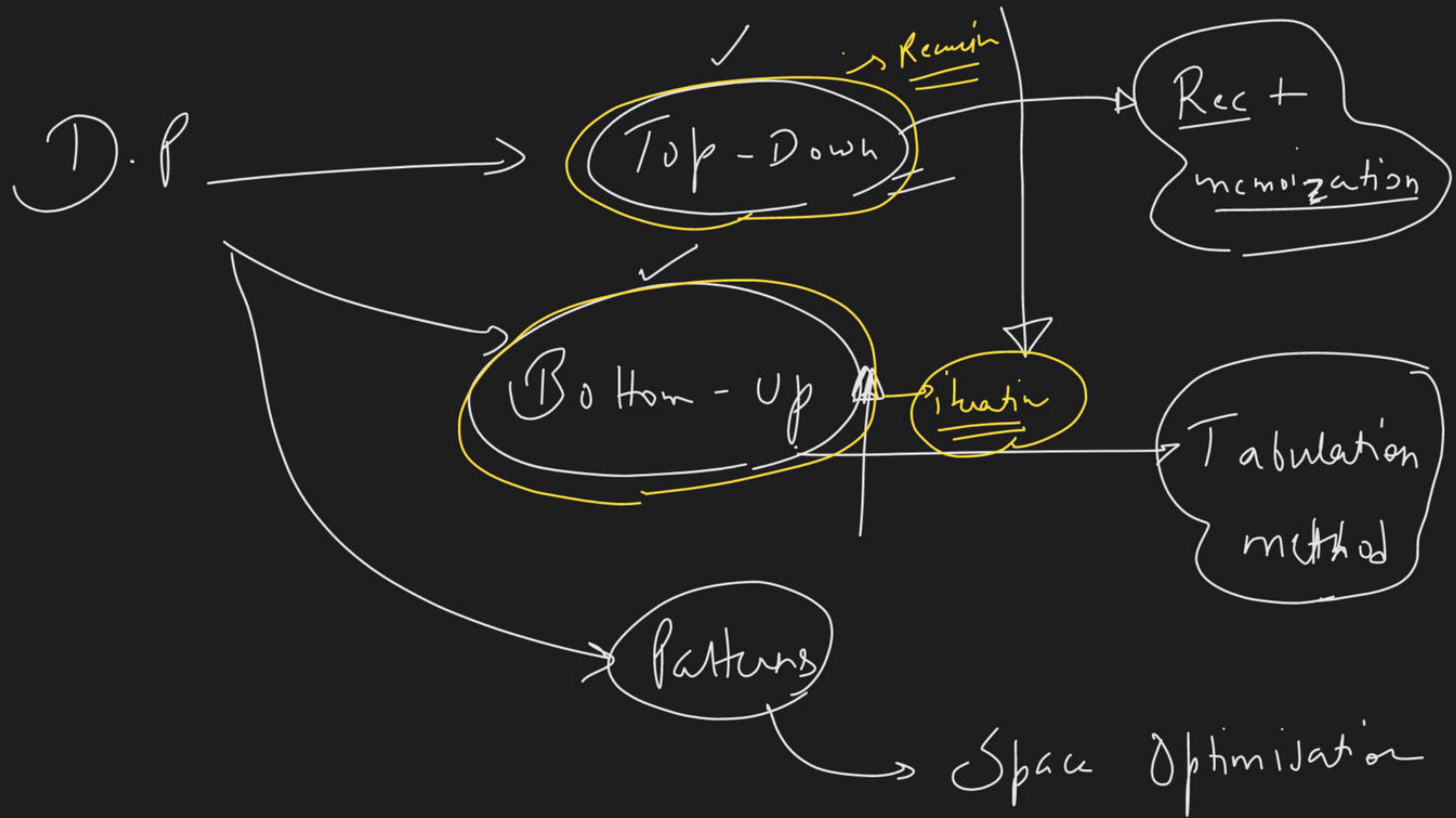












→ Fibonacci no: -

ans → Rec

0, 1, 1, 2, 3, 5, 8, 13, 21, ...

Diagram showing the calculation of the 5th Fibonacci number: 2 + 3 = 5. Arrows point from 2 and 3 to 5, which is boxed.

R.R →

$$f(n) = f(n-1) + f(n-2)$$

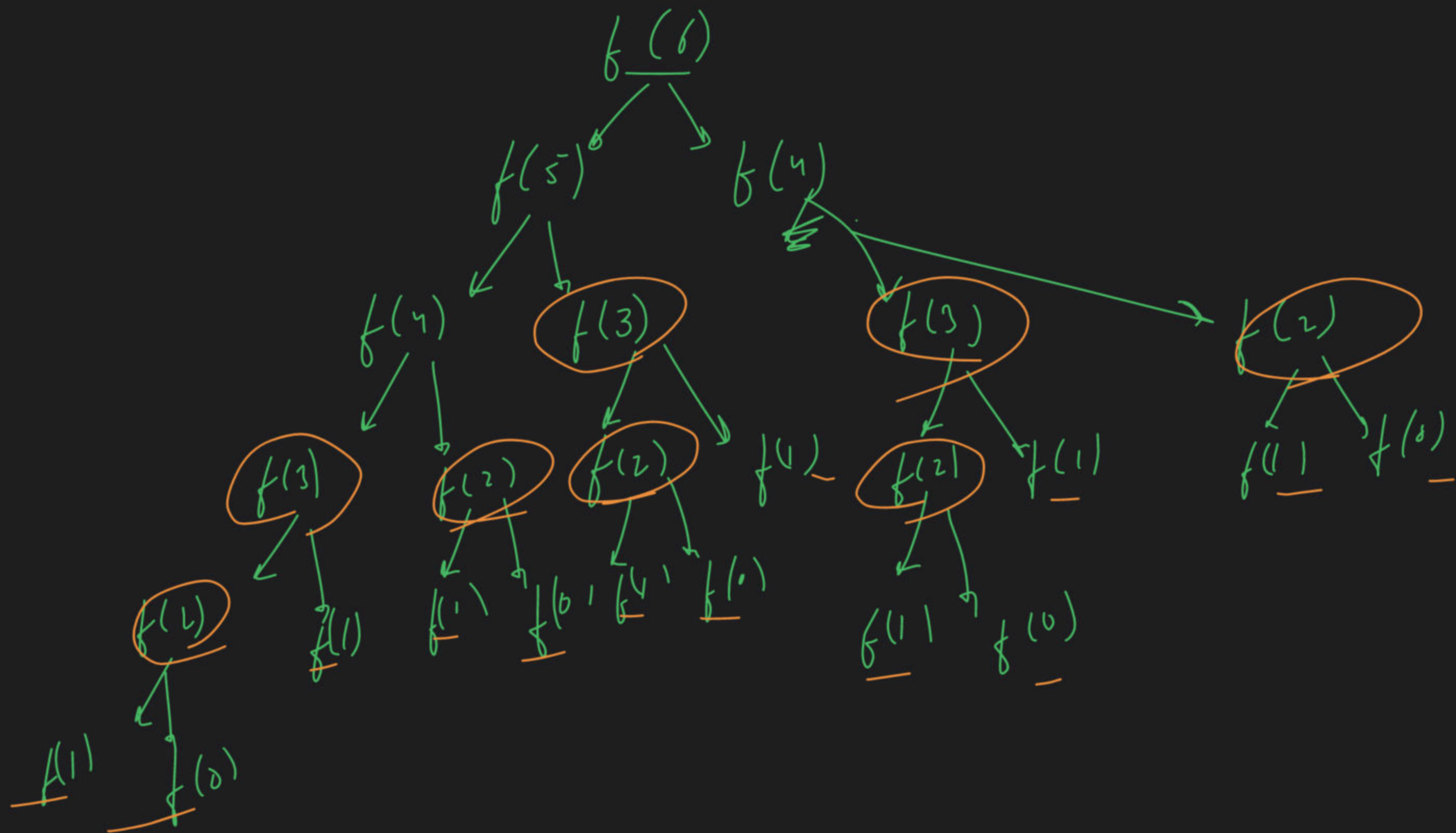
Diagram showing the recursive relation for the 6th Fibonacci number:  $n^{\text{th}}$  fibonacci no. =  $(n-1)^{\text{th}}$  fib no. +  $(n-2)^{\text{th}}$  fib no.

let  $n = 6$

$$f(6) = f(5) + f(4)$$

Diagram showing the calculation of the 6th Fibonacci number:  $6^{\text{th}}$  fib no. =  $5^{\text{th}}$  fib no. +  $4^{\text{th}}$  fibonacci.





Mem

① Create a dp array

② store/retrieve ans in dp array

③ if ans already exists in dp array, then return

~~Stop~~





if  $(n == 0 \mid \mid n == 1)$   
 return  $n$ ;

dp

|   |              |
|---|--------------|
| 0 | <del>0</del> |
| 1 | <del>1</del> |
| 2 | <del>1</del> |
| 3 | <del>2</del> |
| 4 | 3            |
| 5 | <del>5</del> |
| 6 | <del>8</del> |

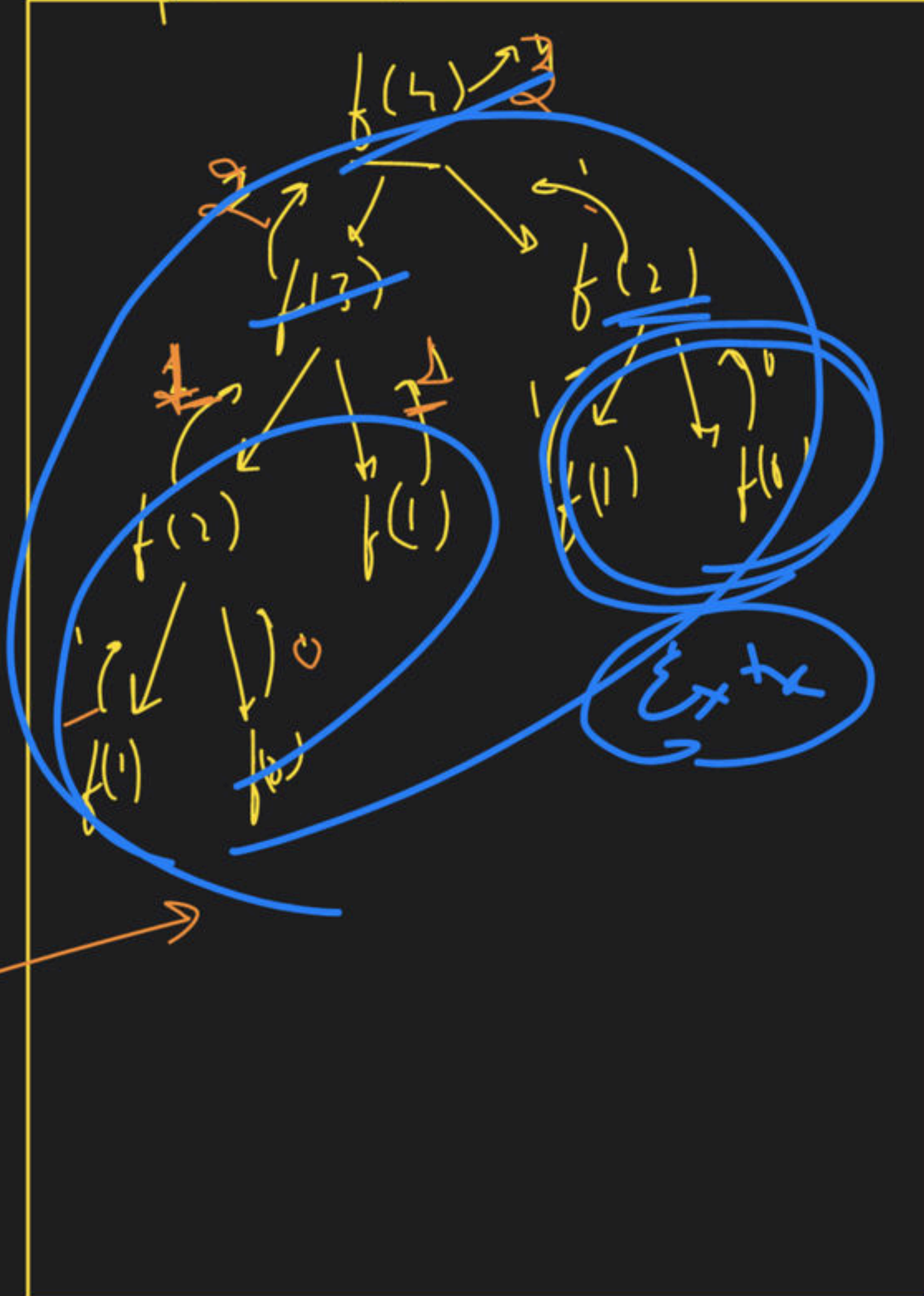
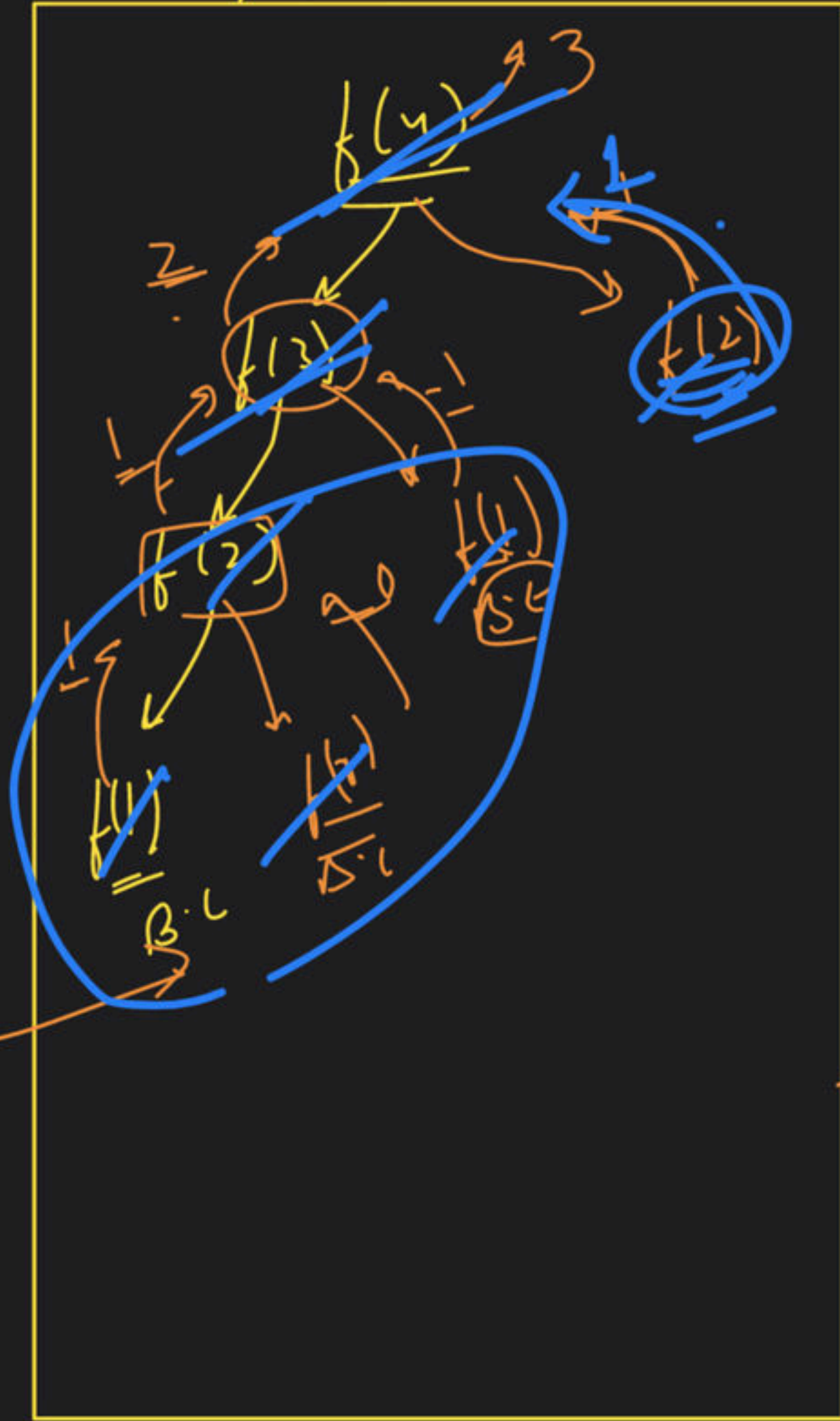


Memorized

plain Rec

0 1 2 3

dp

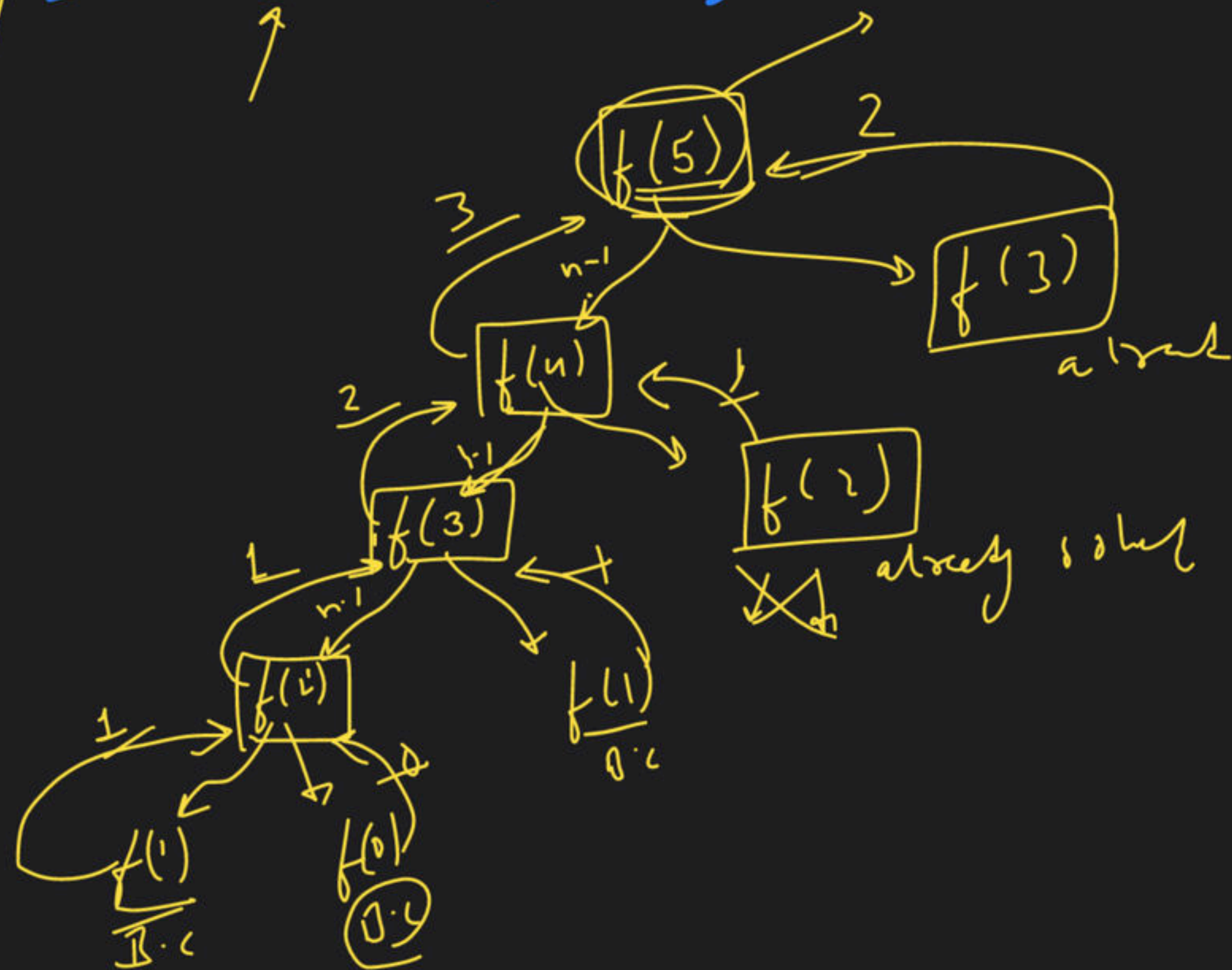
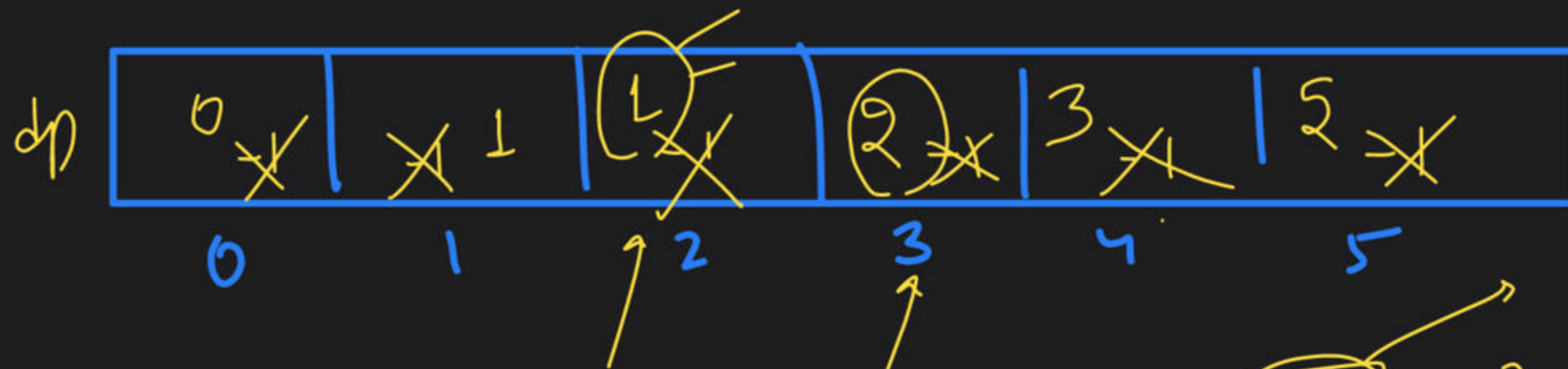


|   |   |
|---|---|
| 0 | 0 |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 3 |

0 1 1 2 3  
-1 -1 -1 -1 -1



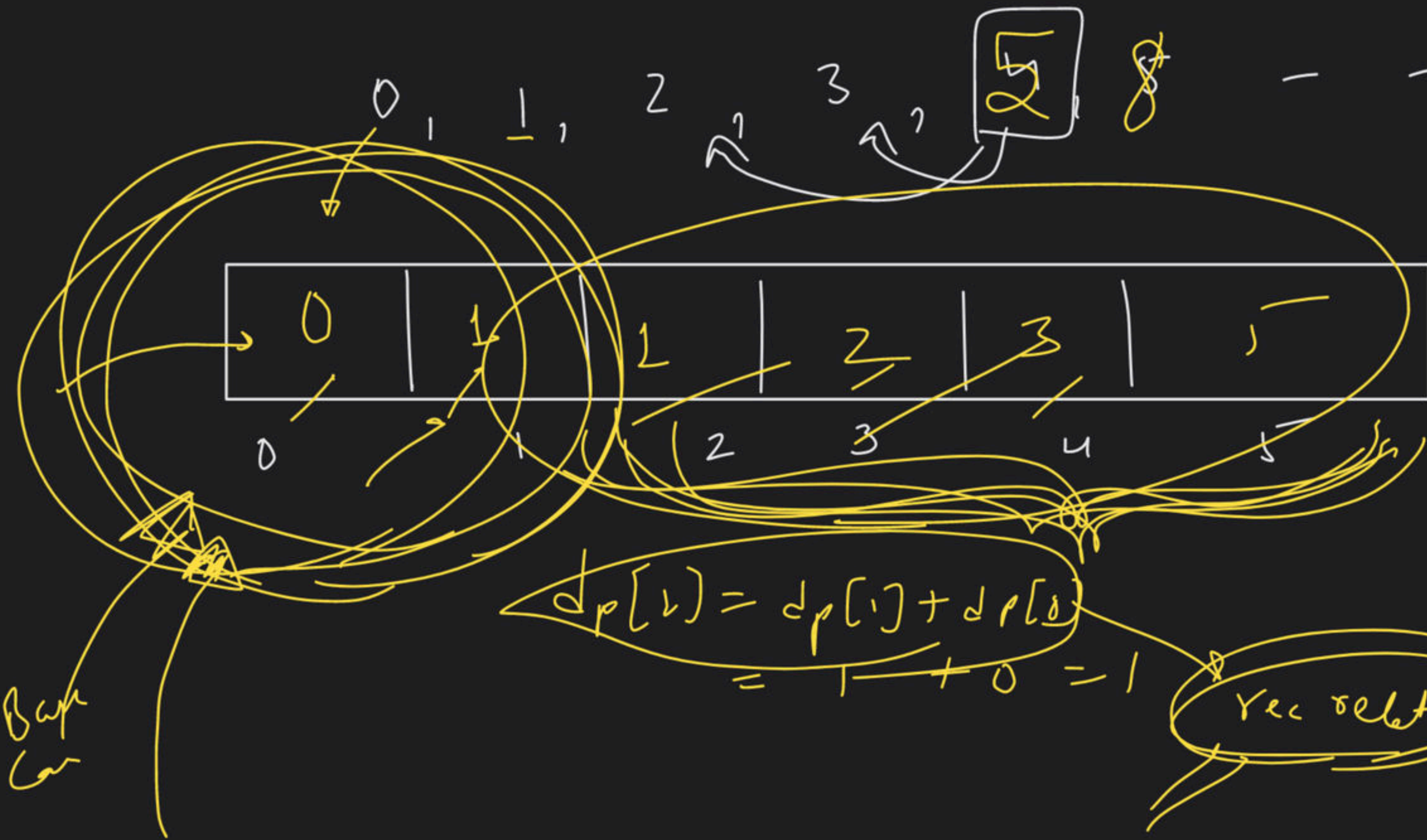
- ① Locate dp arr
  - ② Store dp arr
  - ③ ans already exist/return
-





→ Tabulation method (Bottom-up)

$$dp[5] = dp(1) + dp(3) \\ = 2 + 3$$



$$dp[4] = dp(3) + dp(2) \\ = 2 + 1 = 3$$

$$dp(3) = dp(2) + dp(1) \\ = 1 + 1 \\ = 2$$

$$dp[2] = dp[1] + dp[0] \\ = 1 + 0 = 1$$

Rec relation

Back  
Car



Tabulation

① → create dp array

② → Analyse → Base case → fill dp array accordingly

③ → fill remaining dp array using logic/formula of recursive relation



→ Spau  
→ int solve

Optimisation :-

↳  
↳ Rec  
↳ T/D/Man  
↳ B-V/Tat  
↳ S.O

```
int solveDPTab(n)
```

```
{
```

```
    vector<int> dp(n+1, -1);
```

```
    dp[0] = 0;
```

```
    dp[1] = 1;
```

```
    for (i = 2; i <= n; i++)
```

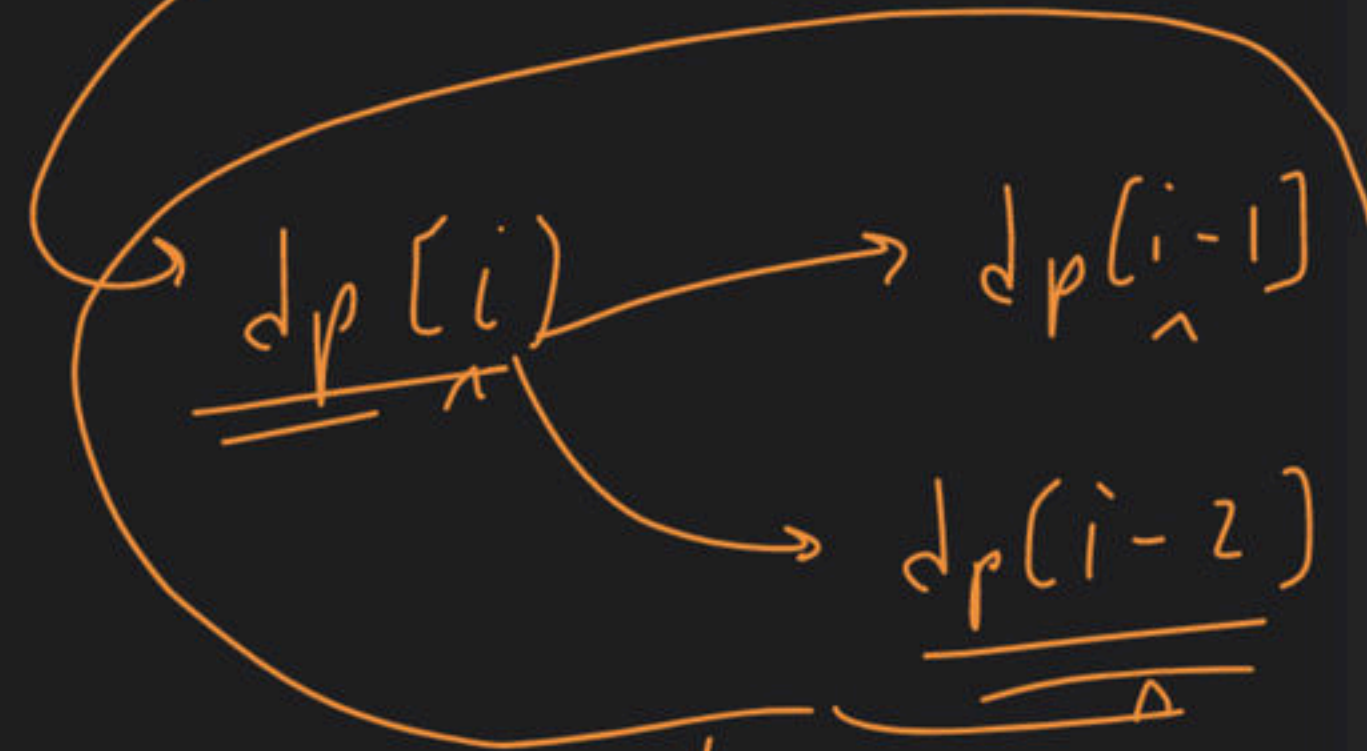
```
{
```

```
    dp[i] = dp[i-1] + dp[i-2]
```

```
    return dp[n];
```

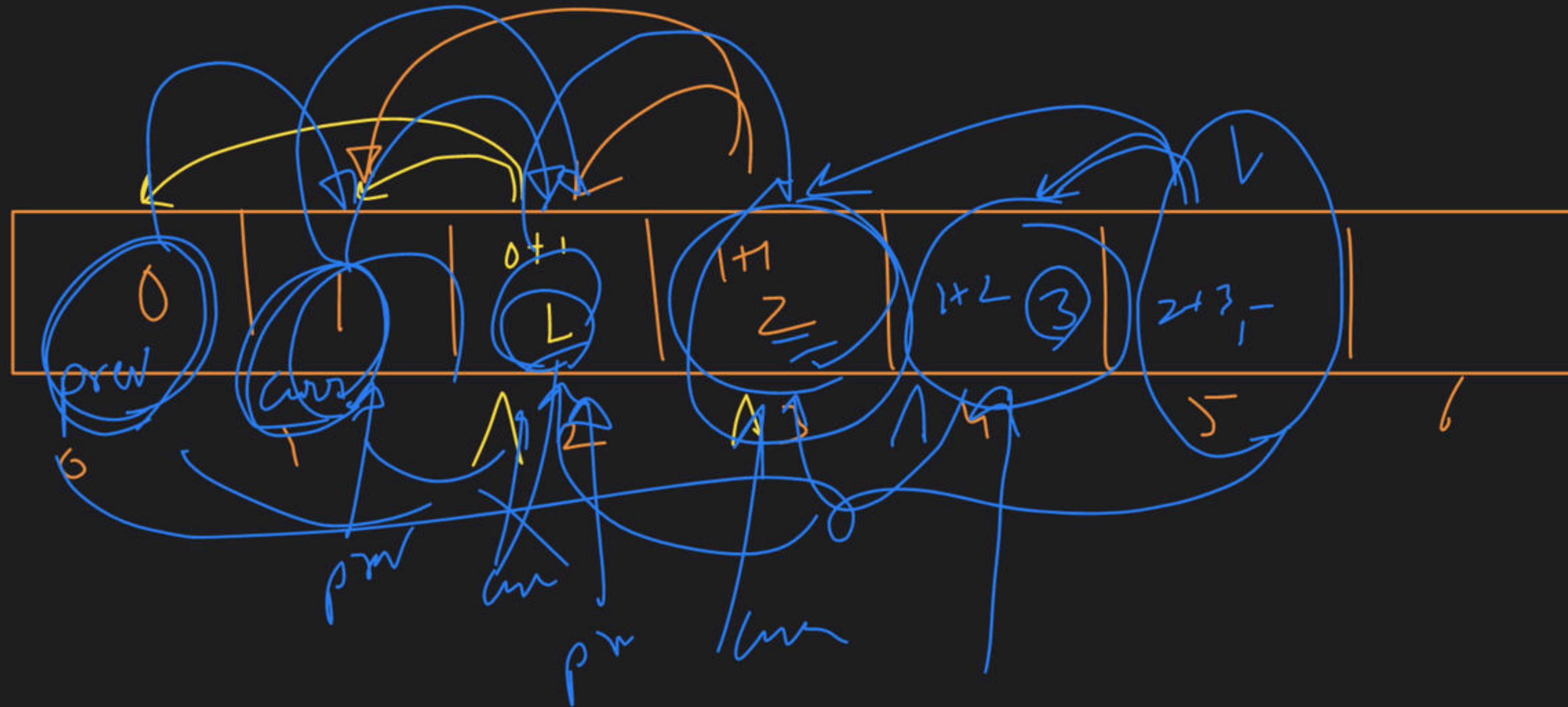
```
}
```

Space Optimisation

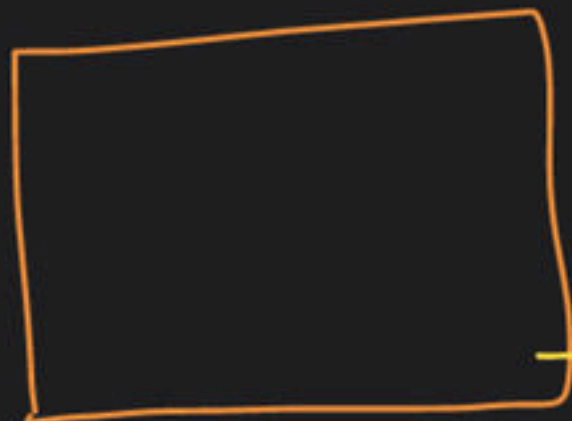


Kya koi pattern  
bnaa hai





int xyz ( dp )  
{  
// B.C



ans

if (dp[n] != -1)  
return ans

// Rec

dp →

~~int ans =~~

return ~~ans~~;

}







































