

## Stack

- Last in First Out
- Elements can be added or removed only from one end
- Gives access only to element at the top of data structure.

### Definition

- An ordered collection of homogenous data items
- Can be accessed ~~at~~ only one end (the top)

### Operations

- Create an empty stack
- Check if it is empty
- Push: add an element to the top
- POP: remove the top element.
- Peek: retrieve the top element (Not the deletion)
- Destroy: Remove all the elements one by one & destroy the data structure.

### The Stack ADT

Value definition:

Abstract typeded `StackType (ElementType ele)`

Condition: none

### Operator definition:

1) Abstract Stack Type CreateStack()

→ Precondition: None

→ Postcondition: Empty Stack is created

2) Abstract Stack Type PushStack(Stack Type Stack, Element Type Element)

→ Precondition: Stack not full or  $\text{NotFull}(\text{Stack}) = \text{True}$

→ Postcondition:  $\text{Stack} = \text{Stack} + \text{Element}$  at the top or Stack = original stack with new Element at the top.

3) Abstract Stack Type PopStack(Stack Type stack)

→ Preconditions: Stack not empty or  $\text{NotEmpty}(\text{Stack}) = \text{True}$

→ Postconditions: PopStack = element at the top,  $\text{Stack} = \text{Stack} - \text{Element}$  at the top or Stack = original stack without top Element.

4) Abstract Destroy Stack (Stack Type Stack)

→ Precondition: Stack not empty or  $\text{NotEmpty}(\text{Stack}) = \text{True}$ .

→ Postcondition: Elements from the stack are removed one by one starting from top to bottom.

$\text{Empty}(\text{Stack}) = \text{True}$

5) Abstract Boolean NotFull(StackType Stack)

→ Precondition: none

→ Postcondition: NotFull(Stack) = True  
if stack is not full

NotFull(Stack) = False if stack is full

6) Abstract Boolean NotEmpty(StackType Stack)

→ Precondition: None

→ Postcondition: NotEmpty(Stack) = True if  
stack is not empty

NotEmpty(Stack) = False if stack is empty

7) Abstract ElementType Peek(StackType Stack)

→ Precondition: Stack not empty or NotEmpty(Stack)  
= True.

→ Postcondition: Peek(Stack) = element at the top  
Stack = Original Stack.



## Implementation of Stack

→ Three ways

→ Array

→ Vector

→ Linked List

### 1) Array

#### Advantage

→ Best performance

#### Disadvantage

→ Fixed size

#### Algo

#### ~~Algorithm~~

1) Algorithm Stack Type CreateStack()

{

integer StackType = -1;

Return Stack;

}

2) Algorithm Stack Type PushStack(StackType, Element Type Element)

{

if (Not Full(stack)) = True

Stack[++StackType] = Element;

Else "Error";

}

3) Algorithm Element Type PopStack(StackType Stack)

{

if (Not Empty(stack)) = True

```
Return Stack [Stack Top = -];
```

```
Else "Error";
```

```
}
```

4) Abstract DestroyStack (Stack Type Stack)

```
{
```

```
if (Not Empty (stack)) = True
```

```
while (Not Empty (stack))
```

```
Print POP Stack (Stack);
```

```
Else "Error";
```

```
}
```

5) Abstract Boolean Not Full (Stack Type Stack)

```
{ if Not Full (stack)
```

```
return True;
```

```
else
```

```
return False;
```

```
}
```

6) Abstract Boolean Not Empty (Stack Type Stack)

```
{ if Not Empty (stack)
```

```
return True;
```

```
else
```

```
return False;
```

```
}
```

7) Abstract ElementType Peek (Stack Type Stack)

```
{ if (Not Empty (stack)) = True
```

```
Return Stack [Stack Top];
```

```
Else "Error";
```

```
}
```

## 2] Vector

### Advantage

- Grows to accommodate ~~new data~~ any amount of data
- Second fastest when data size < vector size

### Disadvantage

- Slowest when data size > vector size
- Wasted space
- Can grow to unlimited size

### → Algo

Same as array

## 3] Linked List

### → Advantage

- Always constant time to push or pop
- Can grow to an  $\infty$  size

### Disadvantage

- Slowest method

### Algo

```

struct Node Type {
    Element Type Element;
    Node Type Next;
}
  
```

```

1) Algorithm Stack Type createStack()
{
  Create Node (TOP);
  TOP = NULL;
}
  
```



2) Stack Type Push Stack (Stack Type Stack, Node Type New Node)

{

if Top == NULL {

NewNode → Next = NULL;

Top = NewNode; }

Else {

NewNode → Next = Top;

Top = NewNode; }

}

3) Algorithm Element Type Pop Stack (Stack Type Stack)

{

if Top == NULL

~~"Error"~~ Print "Underflow";

Else

{

Create Node (Temp);

Temp = Top;

Top = Top → Next;

Return (Temp → Data);

}

4) Abstract Destroy Stack (Stack Type Stack)

{

if Top == NULL

Print "Underflow";

Else

{

```

create Node (Temp);
while (Not Empty (Stack))
{
    Temp = Top;
    Top = Top → Next;
    Return (Temp → Data);
}

```

5) Abstract Element Type Peep (Stack Type Stack)

```

{
    if Top == NULL
        Print "Error"
    Else
        Return Top → Data;
}

```

6) Abstract Display Stack (Stack Type Stack)

```

{
    if Top == NULL
        Print "Error"
    Else
    {
        create Node Temp (Temp);
        Temp = Top;
        while (Temp != NULL) {
            Print (Temp → data);
            Temp = Temp → Next;
        }
    }
}

```



Application

1) Parentheses Matching

Ex i/p = String = { ( ) ( ) }

| Input | Stack | Basically           |
|-------|-------|---------------------|
| {     | {     | Once we start       |
| (     | { (   | the parentheses     |
| )     | {     | that stack is Empty |
| (     | { (   | in the output       |
| )     | {     |                     |
| }     | Empty |                     |

2) Infix to Postfix

→ Precedence

i) ^

ii) \*, /

iii) +, -

Eg  $M * N + T \wedge Q / F * A + B$ 

|    | Input | Stack | Output |
|----|-------|-------|--------|
| 1) | M     |       | M      |
| 2) | *     | *     | M      |
| 3) | N     | *     | MN     |
| 4) | +     | +     | MN*    |
| 5) | T     | +     | MN*T   |
| 6) | ^     | +^    | MN*T   |
| 7) | Q     | +^    | MN*TQ  |
| 8) | /     | + /   | MN*TQ^ |

|     |   |       |                                |
|-----|---|-------|--------------------------------|
| 9)  | F | +/    | $MN * TQ \wedge F$             |
| 10) | * | +     | $MN * TQ \wedge F /$           |
| 11) | A | +     | $MN * TQ \wedge F / A$         |
| 12) | + | +     | $MN * TQ \wedge F / A * +$     |
| 13) | B | Empty | $MN * TQ \wedge F / A * + B +$ |

- If inside the stack is higher Precedent sign in step (2)
- There is a lower Precedent sign in the input it pops the higher Precedent sign in the Output in step (4)
- When there is a lower Precedent sign in the stack & higher in input the higher is pushed into the stack without popping the lower Precedent sign in step (6)
- If equal Precedent pop the one in stack and push the one in input in step (10)

## 8) Infix to postfix with parenthesis

Eg  $((C(A+B)*C)-(C(D+E)*(F+G)))$ 

| Input | Stack  | Output        |
|-------|--------|---------------|
| C     | C      |               |
| (     | CC     |               |
| C     | CCC    |               |
| +     | CCC+   | A             |
| B     | CCC+   | AB            |
| )     | CC     | AB+           |
| *     | CC*    | AB+           |
| (     | CC*    | AB+C          |
| )     | C      | AB+C*         |
| -     | C-     | AB+C*         |
| (     | C-C    | AB+C*         |
| C     | C-CC   | AB+C*         |
| D     | C-CC   | AB+C*D        |
| -     | C-CC-  | AB+C*D        |
| E     | C-CC-  | AB+C*DE       |
| )     | C-C    | AB+C*DE-      |
| *     | C-C*   | AB+C*DE-      |
| (     | C-C*C  | AB+C*DE-      |
| F     | C-C*C  | AB+C*DE-F     |
| +     | C-C*C+ | AB+C*DE-F     |
| G     | C-C*C+ | AB+C*DE-FG    |
| )     | C-C*   | AB+C*DE+FG+   |
| )     | C-     | AB+C*DE+FG+   |
| )     | EMPTY  | AB+C*DE+FG+*- |



→ So basically once you shut the parenthesis whatever was in the parenthesis it gets popped

#### 4) Postfix evaluation

→ Create a stack for storing operands  
→ Scan the input expression from Left to right.

#### 5) Reverse a string using stack

→ Create a stack (empty)  
→ One by one push ~~the~~ all the characters  
→ One by one pop all the characters from stack & put them back to string.

#### 6) Check if a string is Palindrome

→ Do (5)  
→ Check if equal then true else false

#### 7) Recursion

→ Calling the same function directly or indirectly  
→ Represents a problem in terms of one or more smaller Problems, and add one or more base conditions that stop the recursion  
→ The maximal no<sup>o</sup> of nested calls is called recursion depth

## Recursive function Call

- Current function is paused
- The execution context associated with it is remembered in a special data DS called execution context stack.
- The nested call executes
- After it ends, the old execution context is retrieved from the stack, & the outer function is resumed from where it stopped.
- Each recursive call, needs to save
  - Current value
  - Local variable
  - Return Address
- Also, as a function calls to another function first its arguments, then the return address & finally local variable is pushed.

## Backtracking

- It is an algorithmic - technique for solving problems recursively by trying to build a solution incrementally, one piece at a time.
- Removing those solutions that fail to satisfy the constraints of the problem at any point of time.