

Bharatiya Vidya Bhavan's SARDAR PATEL INSTITUTE OF TECHNOLOGY

(Autonomous Institute Affiliated to Mumbai University)
Munshi Nagar, Andheri (West), Mumbai-400058
Information Technology Department

Academic Year: 2022-2023 Class: TE Sem.: VI Course: IT331-ADBMS

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Course	Advanced Database management system
Lab	2

Aim:

Design a distributed database by applying the concept of horizontal fragmentation

Scenario:

A retail chain database where customer info, customer orders and store inventory is stored. Storing all of this data in a centralized database could lead to slow queries, increased network latency and reduced availability in case of server failure.

Therefore, we can fragment tables according to our scenario.

We can fragment the inventory and orders table by the store location. This means that each fragment will contain only the orders for a specific store.

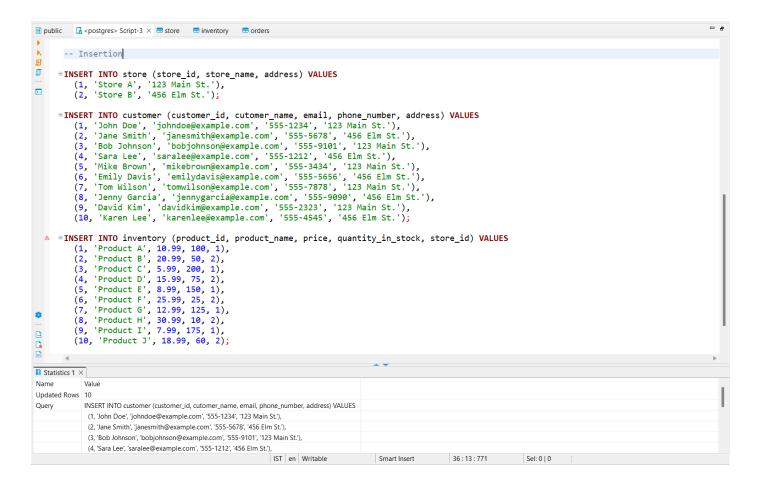
This ensures that each store has access only to their inventory and orders data, which improves data security and reduces risk of errors.

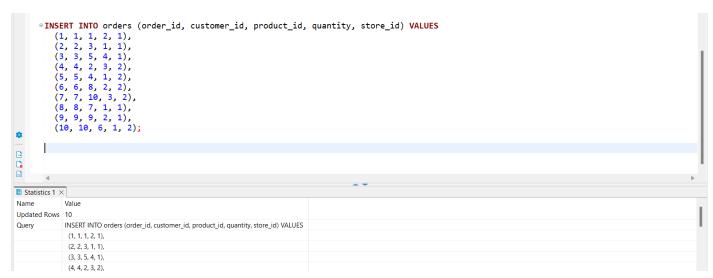
Queries:

1. Creation of table

```
create table store(
    store_id int,
    store_name varchar(50),
▶ ▶ □
                  address varchar(100),
....
<u>S</u>
                  primary KEY(store_id)
         ⊕ create table customer(
                  customer_id int,
cutomer_name varchar(50),
                  email varchar(100),
                  phone_number varchar(20),
address varchar(100),
primary key(customer_id)
        ecreate table inventory(
    product_id int not null primary key,
                  product_name varchar(50),
price decimal(10, 2),
quantity_in_stock int,
store_id int,
foreign key (store_id) references store(store_id)
         ecreate table orders(
                 ate table orders(
order_id int not null primary key,
customer_id int,
product_id int,
quantity int,
store_id int,
foreign key (store_id) references store(store_id),
foreign key (customer_id) references customer(customer_id)
...
G
■ Statistics 1 ×
Name
                 Value
Updated Rows 0
Query
                  create table orders(
                                                                                        IST en Writable
                                                                                                                             Smart Insert 24 : 3 : 506 Sel: 0 | 0
```

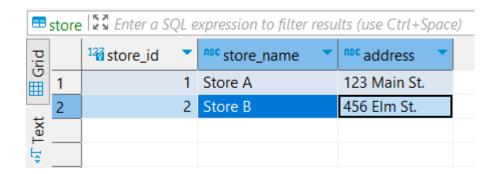
2. Inserting tuples into tables





Tables:

Store table:



Inventory table:

	■inventory Enter a SQL expression to filter results (use Ctrl+Space)						
Grid		¹⅔ product_id ▼	product_name	¹²³ price 🔻	¹²³ quantity_in_stock	123 store_id	
<u></u>	1	1	Product A	10.99	100	1 ♂	
	2	2	Product B	20.99	50	2 ☑	
Text	3	3	Product C	5.99	200	1 ♂	
Ė	4	4	Product D	15.99	75	2 ☑	
	5	5	Product E	8.99	150	1 ♂	
	6	6	Product F	25.99	25	2 ☑	
	7	7	Product G	12.99	125	1 ♂	
	8	8	Product H	30.99	10	2 ☑	
	9	9	Product I	7.99	175	1 ♂	
	10	10	Product J	18.99	60	2 ☑	

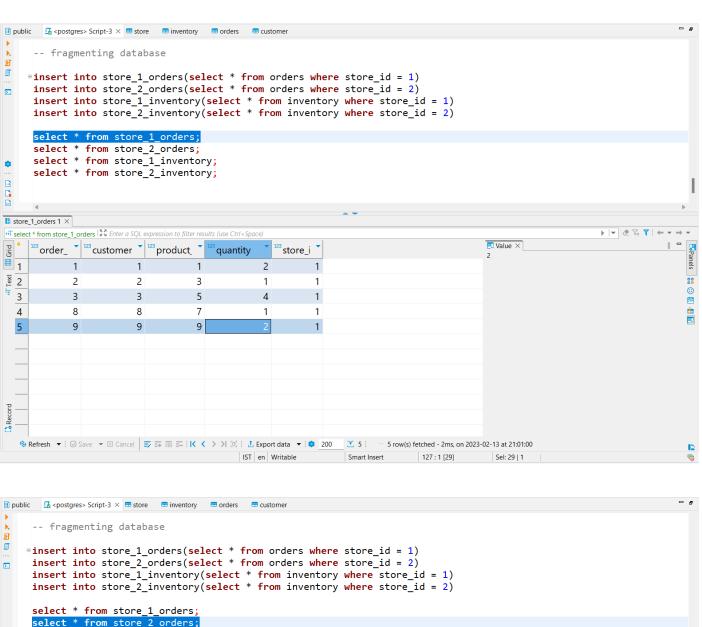
Customer table:

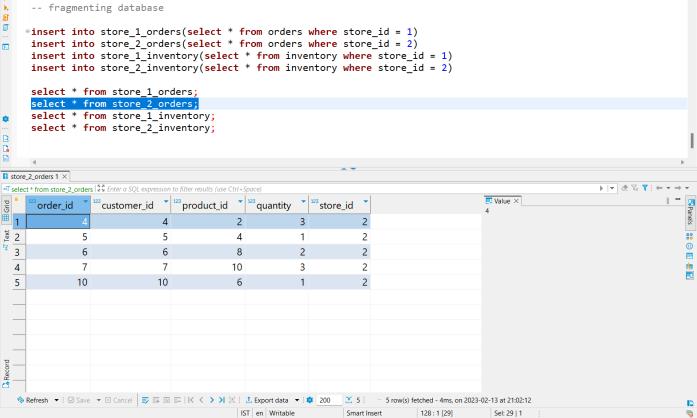
•	☐ customer						
Grid		¹⅔ customer_id ▼	and cutomer_name	anc email	phone_number	address 🔻	
=	1	1	John Doe	johndoe@example.com	555-1234	123 Main St.	
	2	2	Jane Smith	janesmith@example.com	555-5678	456 Elm St.	
ext	3	3	Bob Johnson	bobjohnson@example.co	555-9101	123 Main St.	
Ė	4	4	Sara Lee	saralee@example.com	555-1212	456 Elm St.	
	5	5	Mike Brown	mikebrown@example.co	555-3434	123 Main St.	
	6	6	Emily Davis	emilydavis@example.cor	555-5656	456 Elm St.	
	7	7	Tom Wilson	tomwilson@example.cor	555-7878	123 Main St.	
	8	8	Jenny Garcia	jennygarcia@example.co	555-9090	456 Elm St.	
	9	9	David Kim	davidkim@example.com	555-2323	123 Main St.	
	10	10	Karen Lee	karenlee@example.com	555-4545	456 Elm St.	

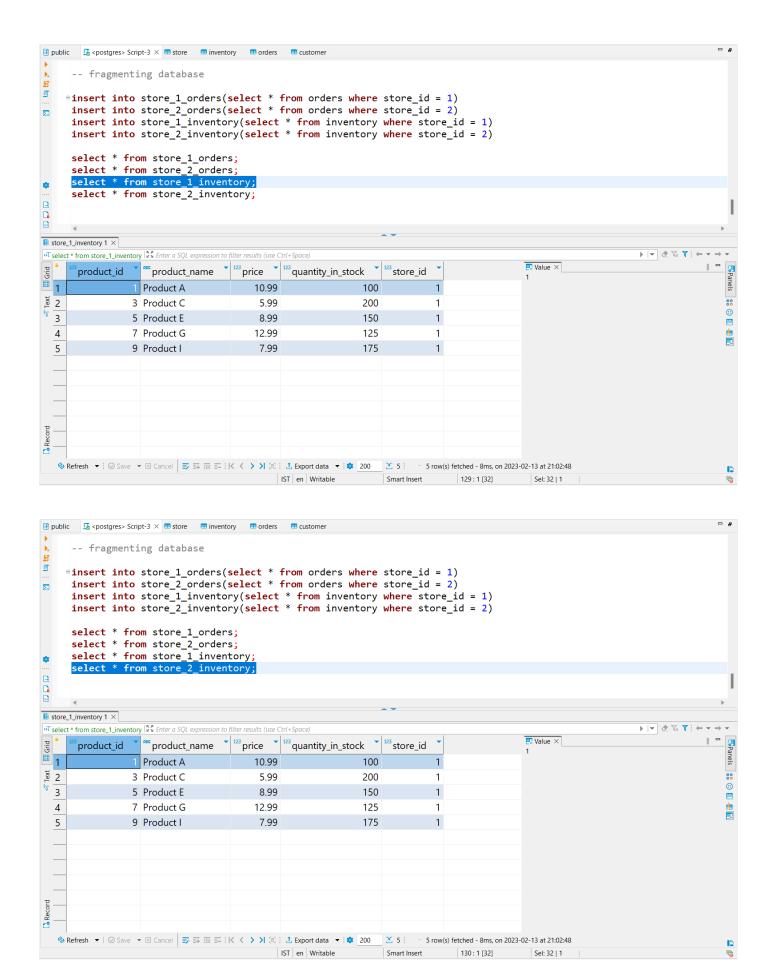
Orders table:

=	■ orders 🖏 Enter a SQL expression to filter results (use Ctrl+Space)					
Grid		¹⅔ order_id ▼	123 customer_id	123 product_id	¹²³ quantity	¹²³ store_id ▼
<u></u>	1	1	1 ♂	1	2	1 ♂
	2	2	2 ☑	3	1	1 ♂
Text	3	3	3 ₺	5	4	1 ♂
Ė	4	4	4 ₺	2	3	2 ♂
	5	5	5 ☑	4	1	2 ₺
	6	6	6 ⊿	8	2	2 ♂
	7	7	7 ₺	10	3	2 ☑
	8	8	8 ☑	7	1	1 ♂
	9	9	9 ♂	9	2	1 ♂
	10	10	10 ♂	6	1	2 ♂

Fragmenting the database by stores:

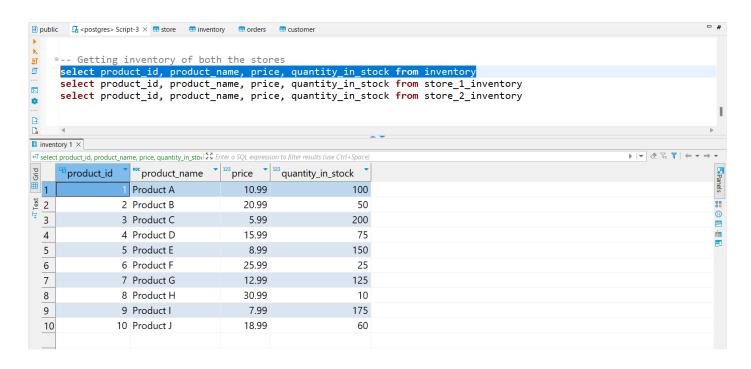


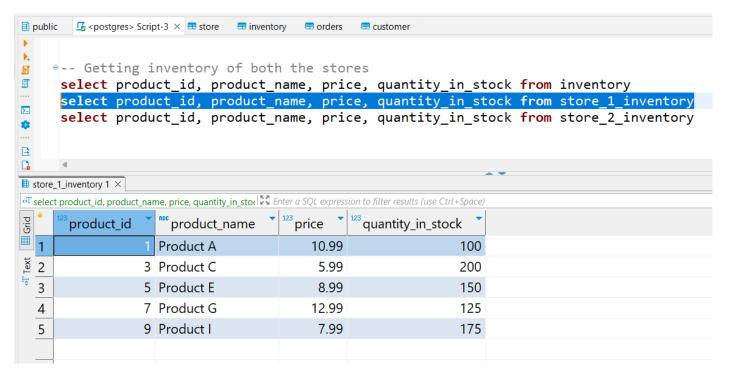


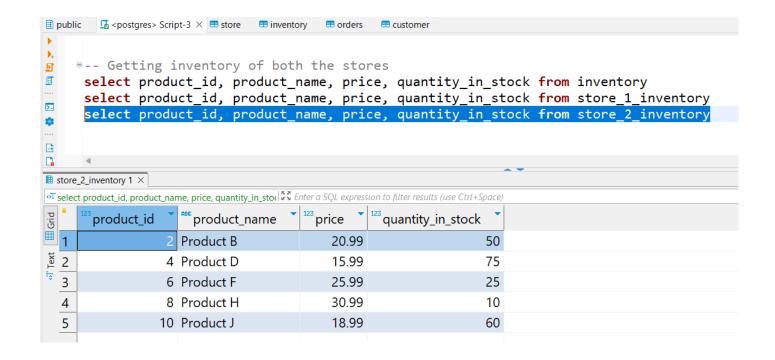


Running Queries in the fragmented tables

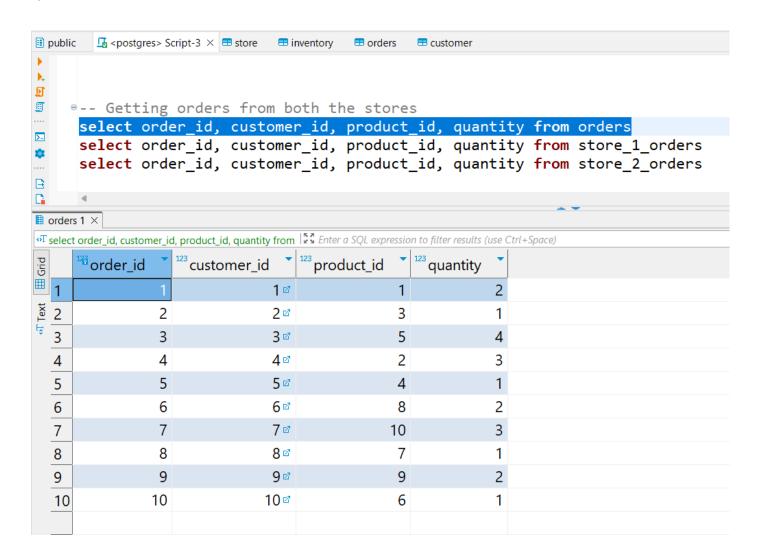
1) Here in the below query I am getting the info of inventory of the both stores

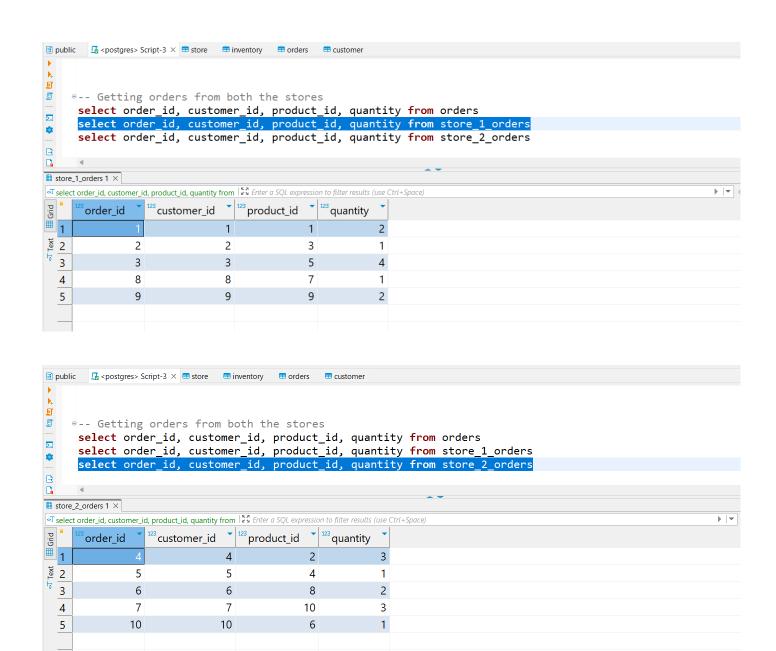




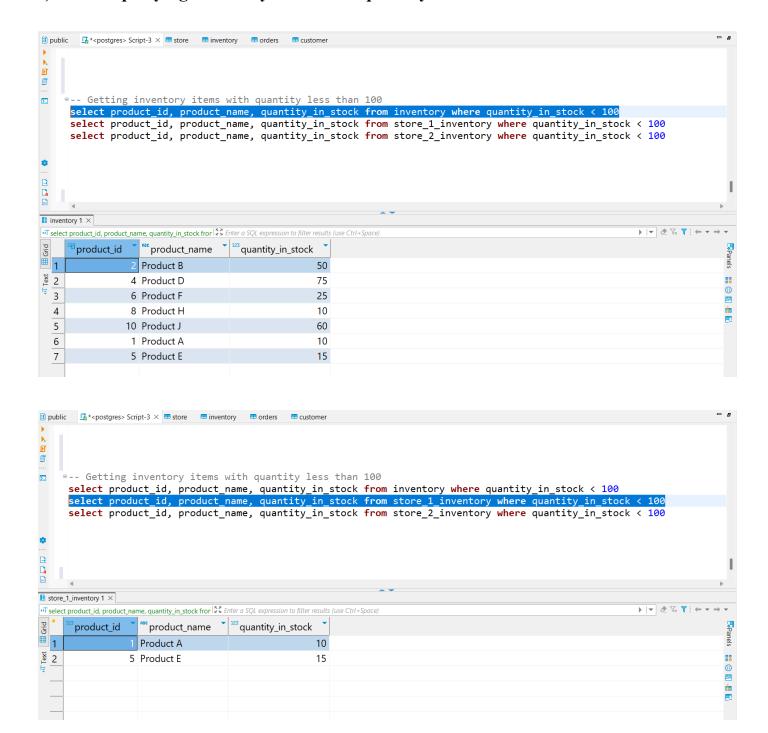


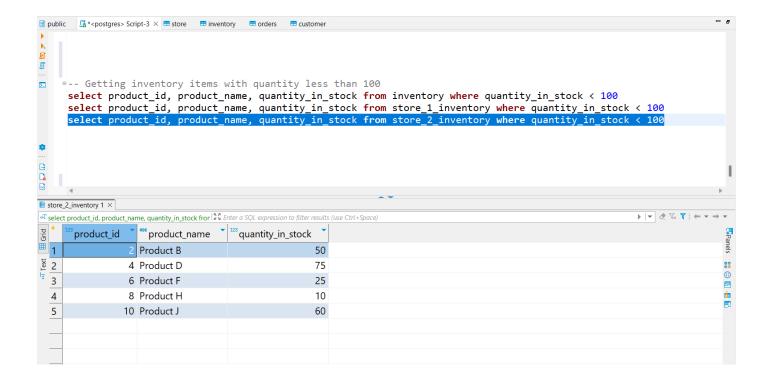
2) Here in the below query I am getting the orders from both the stores



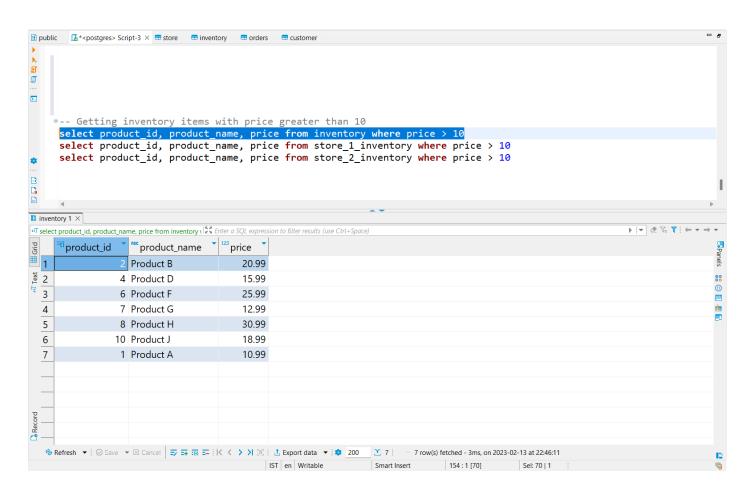


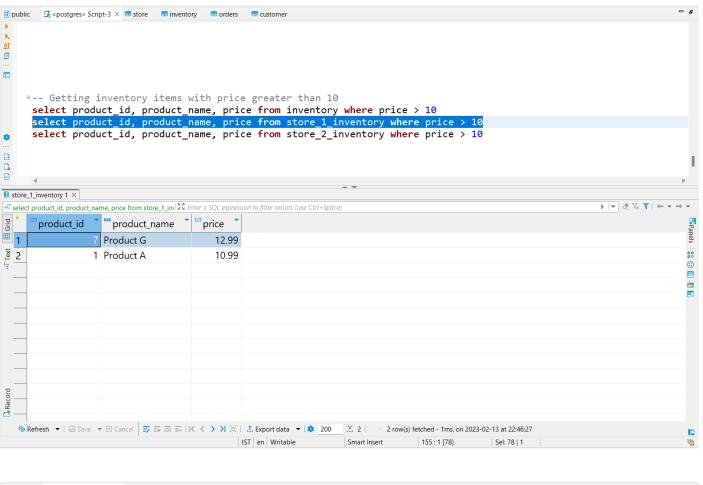
3) We are querying inventory items with quantity less than 100

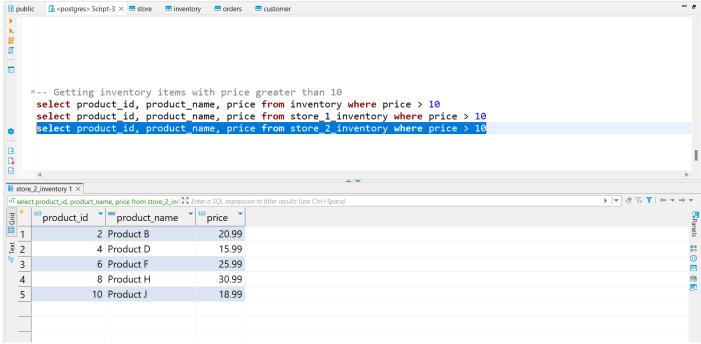




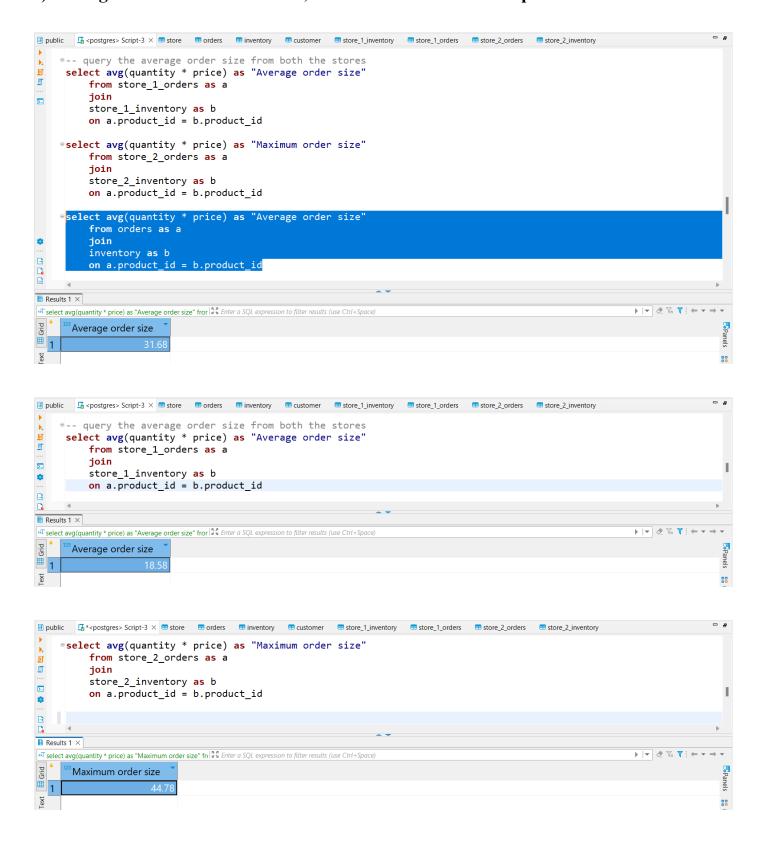
4) Now, we are querying the inventory items with price greater than \$10



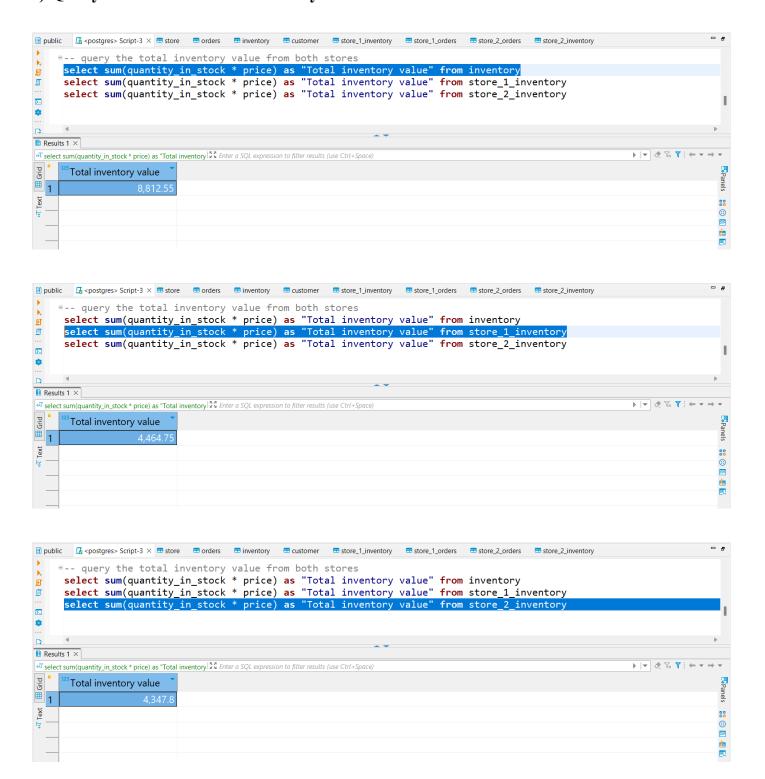




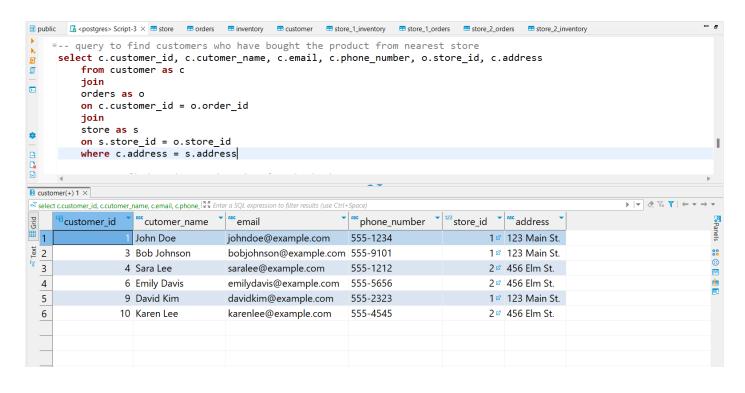
5) Average order size at each store, it can be used to calculate profits of the store

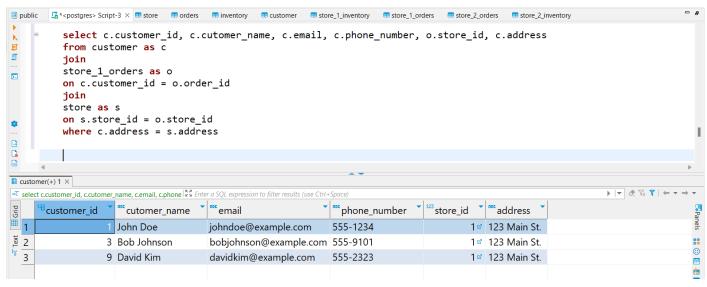


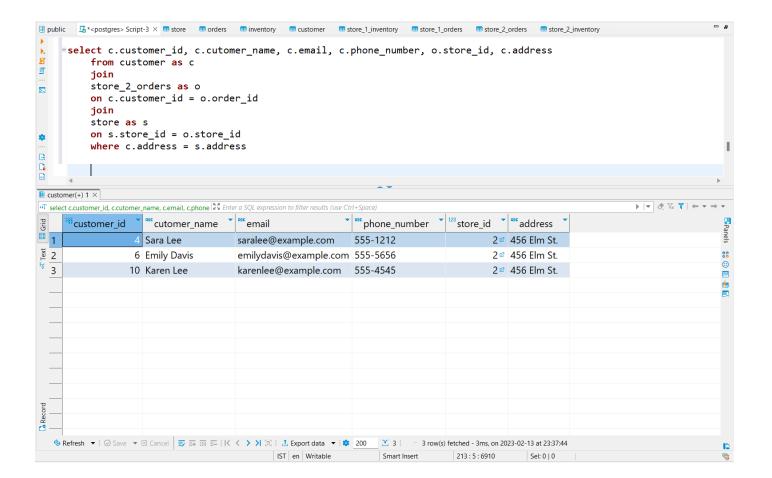
6) Query to Find the Total inventory value of the stores



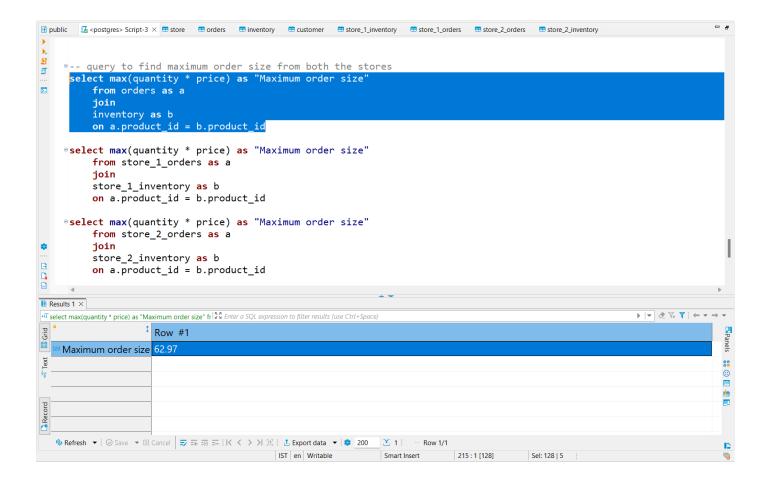
7) Query to find the customers who have find the products in nearest store

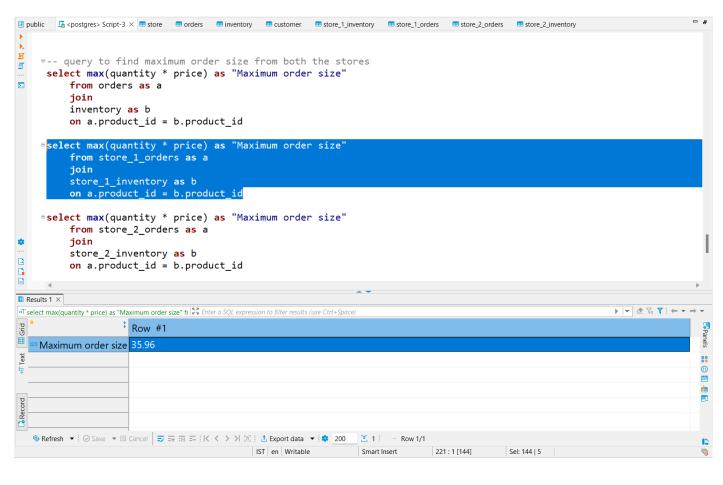






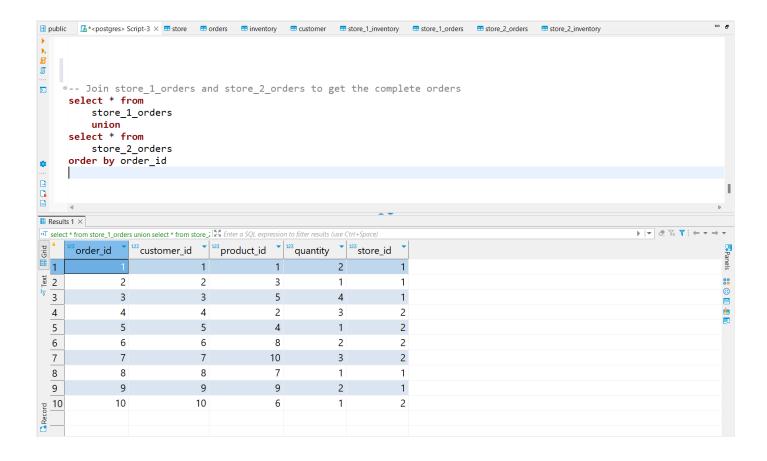
8) Query to find the maximum order size from both the stores



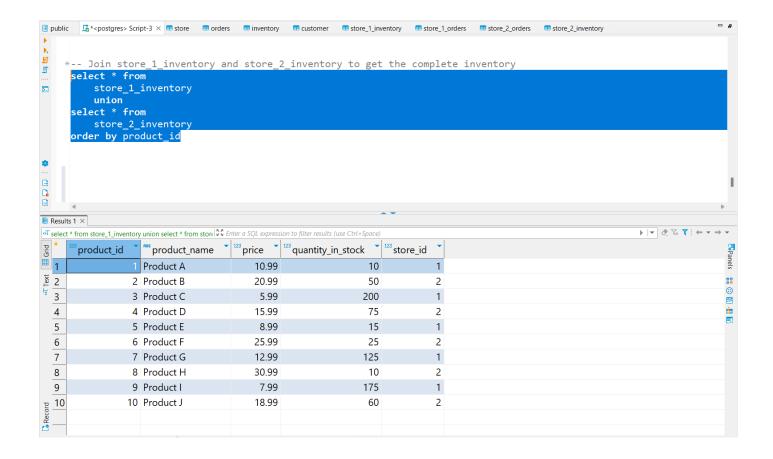




9) Joining store_1_orders and store_2_orders to get the complete Orders table



10) Joining store_1_inventory and store_2_inventory to get the complete inventory table



Correctness rules of fragmentation

Fragmentation is the major concept in distributed databases. We fragment a table horizontally, vertically, or both and distribute the data to different sites (servers at different geographical locations).

While we perform the fragmentation process, as a result we expect the following as outcomes;

- We should not lose data because of fragmentation
- We should not get redundant data because of fragmentation Hence, to ensure these properties we need to verify whether we performed the fragmentation correctly or not. For this verification we use the correctness rules.

The rules are as follows;

• Completeness -

To ensure that there is no loss of data due to fragmentation. Completeness property ensures this by checking whether all the records which were part of a table (before fragmentation) are found in at least one of the fragments after fragmentation. We can see in the above queries that the Total number of unique ids in the inventory and orders table match with the number of unique ids in the sum of two fragmented tables.

• Reconstruction -

This rule ensures the ability to reconstruct the original table from the fragments that are created. This rule is to check whether the functional dependencies are preserved or not.

```
If a table R is partitioned into fragments R1, R2, ..., Rn, then Reconstruction insists the following; R = R1 U R2 U ... U Rn
```

For the above scenario we can see that the Union of the two tables gives us back the original table thus making sure that Reconstruction is possible.

• Disjointness -

This rule ensures that no record will become a part of two or more different fragments during the fragmentation process.

```
If a table R is partitioned into fragments R1, R2, ..., Rn, then Disjointness insists the following; R1 \cap R2 \cap ... \cap Rn = Null set
```

For the above scenario we see that the relations "<" and ">=" do not share any common values thus inherently satisfying the Disjoint condition.

Lossless decomposition

Lossless join decomposition is a way to decompose a relation R into two or more relations R1, R2 such that their natural join returns the original relation R. This helps to eliminate data redundancy in a database while retaining the original data. The decomposed tables can be joined together to reconstruct the original relation R. Only the normal forms 1NF, 2NF, 3NF, and BCNF are suitable for lossless join decomposition. The common attribute used for decomposition must be a candidate key or a super key in either R1, R2, or both.

Conclusion:

By performing the above experiment:

- I learnt how to apply horizontal fragmentation on a table and fragmentation increases the efficiency by distributing the database
- I learnt how to check whether the fragmentation was correct by checking the correctness rules and about lossless decomposition