Producer

II: RI = count I4: R2 = Count

Consumer

 $I_2: R_1 = R_1 + 1$ $I_5: R_2 = R_2 - 1$

I3 : Count = R, I6 : Count = R2

```
Instruction Execution let us assume count = 4

I1: R_1 = count \Rightarrow R_1 = 4

I_2: R_1 = R_1 + 1 \Rightarrow R_1 = 5

I_4: R_2 = count \Rightarrow R_2 = 4

I_6: R_2 = R_3 - 1 \Rightarrow R_2 = 3

I_6: count = R_2 \Rightarrow count = 3

I_7: Count = R_1 \Rightarrow count = 5
```

- 1. <u>m (mutex)</u>, a binary semaphore which is used to acquire and release the lock.
- empty, a counting semaphore whose initial value is the number of slots in the buffer, since, initially all slots are empty.
- 3. <u>full</u>, a counting semaphore whose initial value is 0.

```
Producer

do {

wait (empty); // wait until empty>0

and then decrement 'empty'

wait (mutex); // acquire lock

/* add data to buffer */

signal (mutex); // release lock

signal (full); // increment 'full'

} while(TRUE)
```

- 1. <u>m (mutex)</u>, a binary semaphore which is used to acquire and release the lock.
- 2. <u>empty</u>, a counting semaphore whose initial value is the number of slots in the buffer, since, initially all slots are empty.
- 3. <u>full</u>, a counting semaphore whose initial value is 0.

```
Producer

do {

wait (empty); // wait until empty>0

and then decrement 'empty'

wait (mutex); // acquire lock

/* add data to buffer */

signal (mutex); // release lock

signal (full); // increment 'full'

} while(TRUE)
```

```
Consumer

do {

wait (full); // wait until full>0 and
then decrement 'full'

wait (mutex); // acquire lock
/* remove data from buffer */
signal (mutex); // release lock
signal (empty); // increment 'empty'
} while(TRUE)
```