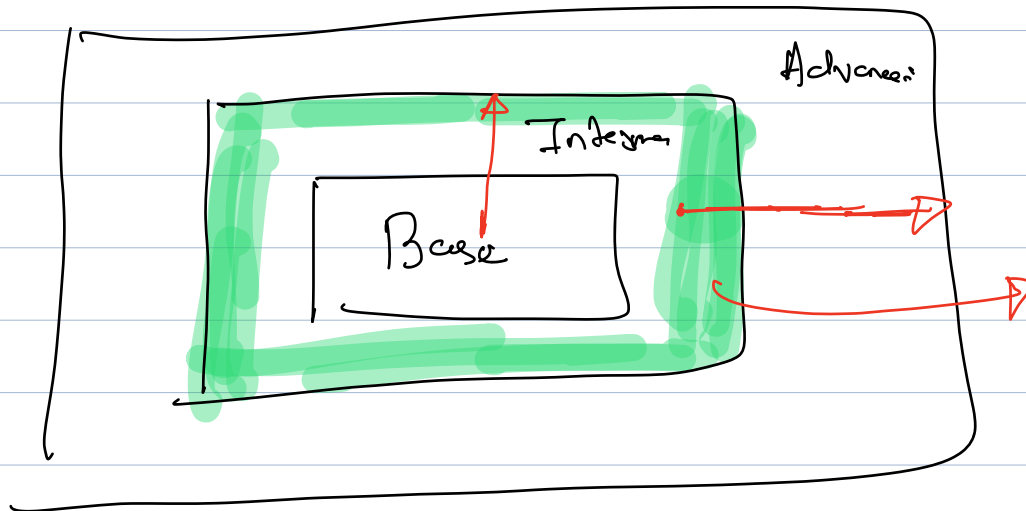


⇒ April intermediate + June Advanced.

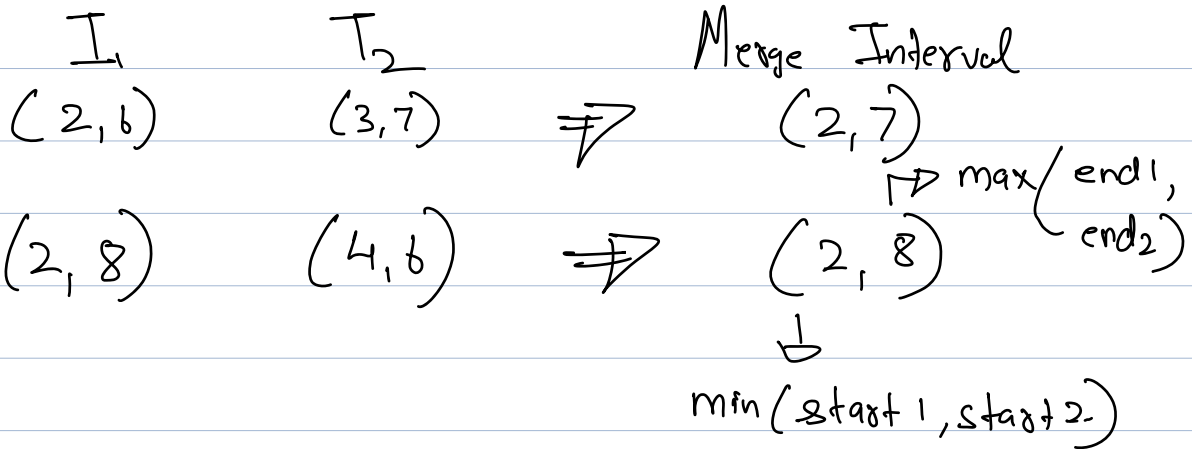


Recursion, Bit manipulation.

⇒ Start fresh

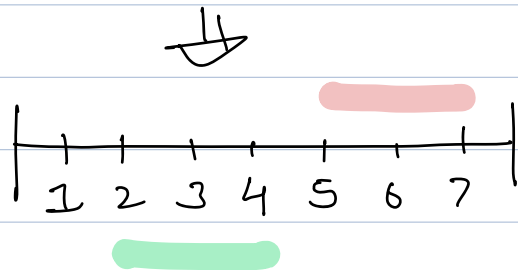
⇒

Merge Intervals.

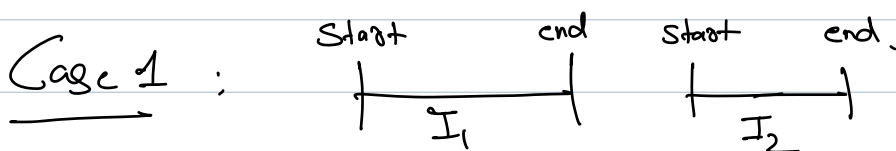


$(2, 4)$ $(5, 7)$ \Rightarrow can't merge

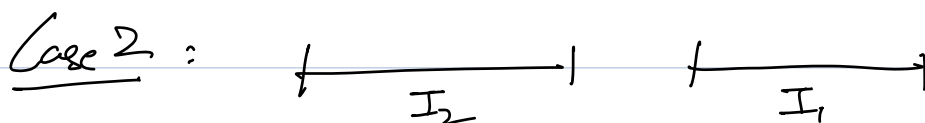
For merging, intervals
should be overlapping!



1) Two interval I_1 & I_2 .



$$I_2.\text{start} > I_1.\text{end}$$



$$I_1.\text{start} > I_2.\text{end}$$

Q

$N = 7, \quad I =$

New overlapping
Intervals.

$[1, 3]$

$[1, 3]$

$[4, 7]$

$[4, 7]$

$[10, 14] \quad [12, 22] \Rightarrow [10, 22]$

$[10, 24]$

$[16, 19] \quad [10, 22] \Rightarrow [10, 22]$

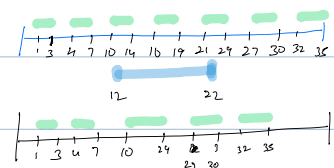
$[27, 30]$

$[21, 24] \quad [10, 22] \Rightarrow [10, 24]$

$[32, 35]$

$[27, 30] \quad [10, 24]$

$[32, 35]$



Pseudo Code!

Input \Rightarrow 1) List \langle Interval \rangle l $\left| \begin{array}{l} l.s \\ l.e \end{array} \right.$
2) Interval I .

List \langle Interval \rangle ans;

for (int $i = 0$; $i < n$; $i++$) l $\left| \begin{array}{l} l \\ I \end{array} \right.$

if ($I.s > l[i].e$) l
ans.push ($l[i]$); // Case 1

} else if ($l[i].s > I.e$) l
ans.push (I); $\left| \begin{array}{l} I \\ l \end{array} \right.$

// Case 2 for (int $j = i$; $j < n$; $j++$) l
ans.push ($l[j]$);
}

return ans;

} else l

$I.s \Rightarrow \min (I.s, l[i].s)$; // Case 3

$I.e \Rightarrow \max (I.e, l[i].e)$;

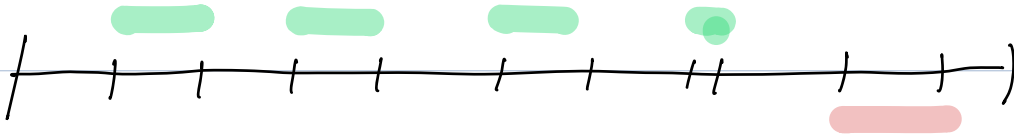
}

}

\rightarrow ans.push (I);
return ans;

TC: $O(n)$
SC: $O(n)$

Edge Case



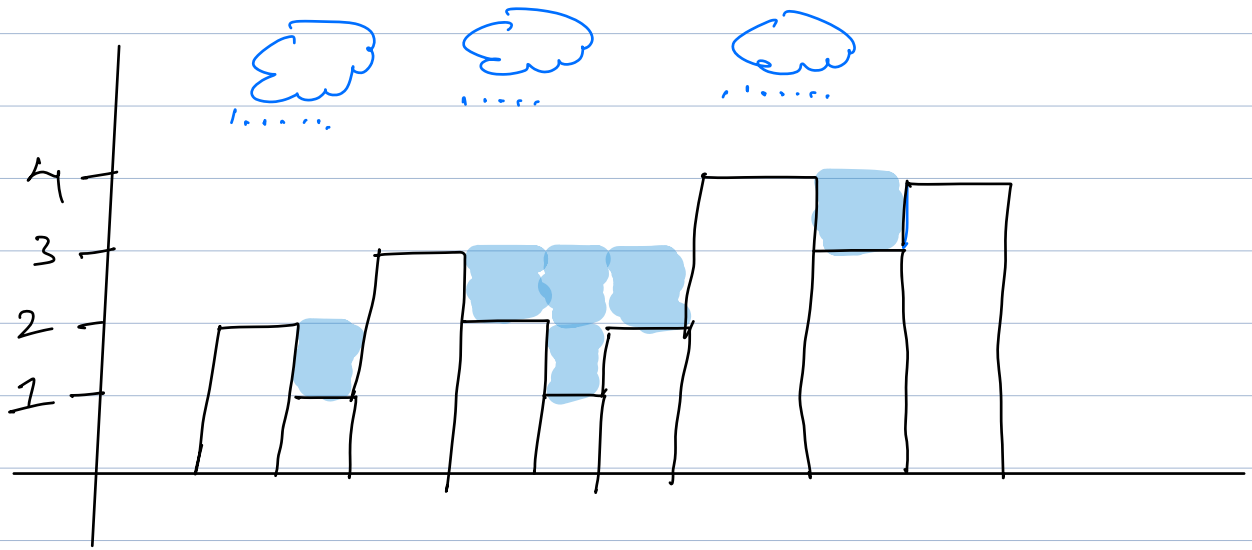
() () () () ()

①

Rain water Trapped

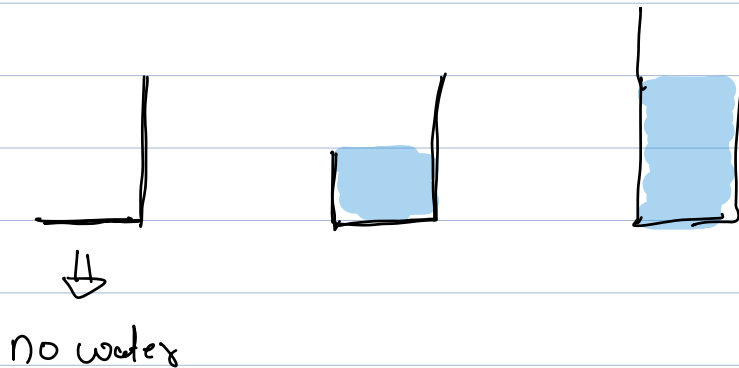
Given n array elements, where $arr[i]$ represents the height of the building. Return the amount of water trapped on the top of buildings.

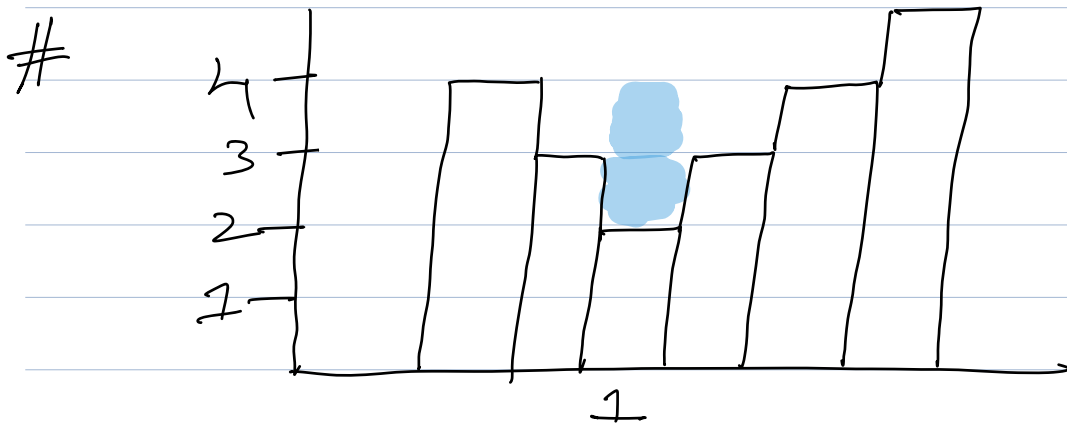
$arr[] = \{2, 1, 3, 2, 1, 2, 4, 3, 4\}$



ans $\Rightarrow 6$

Observation





max height on left $\Rightarrow 4$
 max height on right $\Rightarrow 5$

water trapped $\Rightarrow \min(\text{left-max}, \text{right-max})$
 $- \text{height of the building}$

Q arr $[n] = \{ -1, 4, 3, 5, 2, 7 \}$

left max $[n]$: left max $[i] \Rightarrow \text{maximum } \{0 \dots i\}$

left max $[n]$:

0	1	2	3	4	5
-1	4	4	5	5	7

 $\Rightarrow O(n)$

right max $[n]$: right max $[i] \Rightarrow \text{maximum } \{i \dots n-1\}$

right max $[n]$ =

0	1	2	3	4	5
7	7	7	7	7	7

 $\Rightarrow O(n)$

Pseudo Code

```
int water-trap = 0
```

```
for (int i = 1 ; i < (n-1) ; i++) {
```

```
    // left-max  $\Rightarrow$  leftmax[i-1]
```

```
    // right-max  $\Rightarrow$  rightmax[i+1]
```

```
    int water  $\Rightarrow$  min(left-max, right-max)  
                - arr[i];
```

```
    if (water > 0)
```

```
        water-trapped += water;
```

```
}
```

```
return water-trapped;
```

TC: $O(n)$

SC: $O(n)$

Q3 Maximum Sum Subarray.

arr[] = $\overset{0}{-1}, \overset{1}{4}, \overset{2}{2}, \overset{3}{8}, \overset{4}{-2}, \overset{5}{3}$

ans \Rightarrow 15

Brute Brute force : for every subarray find the sum. $\Rightarrow O(n^3)$

Brute force : Use prefix sum to find Subarray sum $\Rightarrow O(n^2)$

Scenario 1 :

2	4	6	8
---	---	---	---

ans \Rightarrow 20

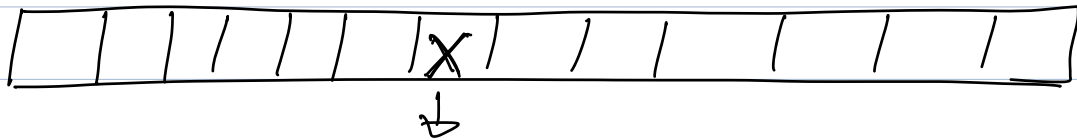
Scenario 2 :

-1	-2	-3	-4
----	----	----	----

ans \Rightarrow -1

-10 | 8 | 3 | -2 | all +ve elements |

→ Carry forward



to include

arr[]	5	6	7	-3	2	-10	-12	8	12	21
Sum \Rightarrow	5	11	18	15	17	7	-5	8	20	41
CF=0	5	11	18	15	17	7	0	8	20	41

max-Sum $\Rightarrow 41$

⇒ Kadanes
Algorithm.

$$\begin{array}{r} -4 \quad -3 \\ \text{Sum} = \\ \text{Cf} = 0 \end{array}$$

max_sum ∇ Inf

6 -4

Pseudo code!

int max-sum \Rightarrow Integer.minimum;

int sum \Rightarrow 0;

for (int i=0; i < n; i++) {

sum \Rightarrow sum + arr[i];

max-sum \Rightarrow max(max-sum, sum);

if (sum < 0)

sum = 0;

}

Tc: $O(n)$
Sc: $O(1)$

return max-sum;

-4	-3	-2	-1
----	----	----	----

max-sum = -1

sum \Rightarrow 0

Double



\Rightarrow left max $\Rightarrow 3$
right max $\Rightarrow 4$

$= \min(3, 4)$

$\Rightarrow \textcircled{3}$

3 - Height bars $\Rightarrow 1$