Q1 Given an array which is formed by rotating a Distint sorted array by K times, Search for a given key in rotated array. You are given the rotated arr. Rotation here means bringing the last element to the front.

Ex1 arr[] =

0	1	2	3	4	5	6	7	8	9	10	11	12
-20	-14	-8	-4	1	2	4	7	11	14	19	23	27

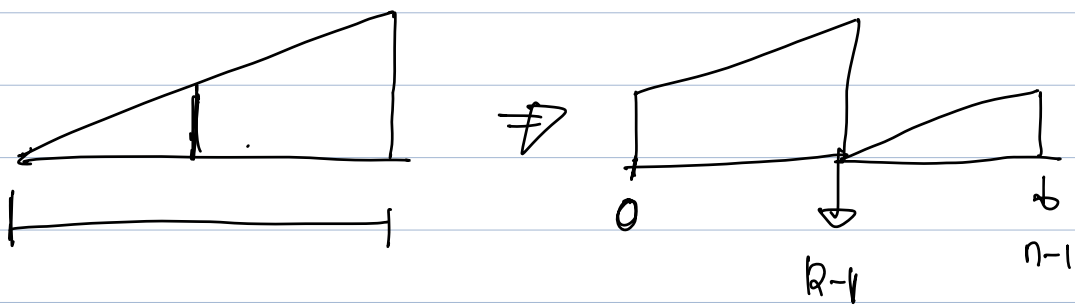
↓ K=5 { 5 times Rotation }

arr[] =

0	1	2	3	4	5	6	7	8	9	10	11	12
11	14	19	23	27	-20	-14	-8	-4	1	2	4	7

Search [-14]

CASE 1: K is given to you. K=5



Apply Binary Search on two array → [0, K-1] [K, n-1]

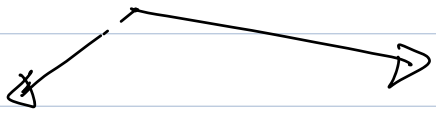
TC: log(n)

SC: O(1)

Optimisation

Approach2

Check the key's value with the last element of rotated array.



key > arr[n-1] → search in first arr

key < arr[n-1] → search in 2nd arr

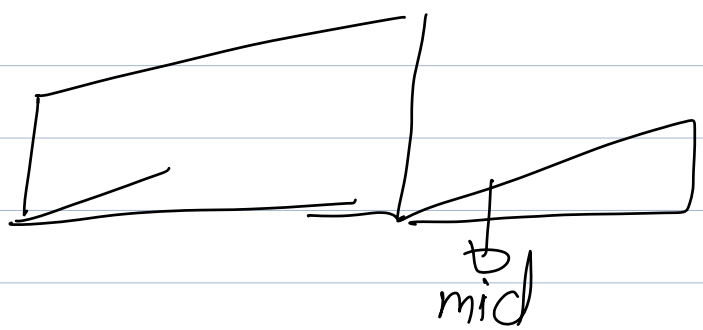
$TC : O(\log n)$

Approach3 :

$$mid = \left[\left(\frac{l+h}{2} \right) + K \right] \% n$$

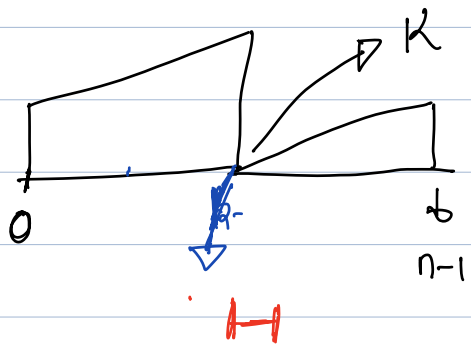
[1 2 2 2 2 2]

a[mid] >= a[l]



CASE2 : K is not given to you.

	0	1	2	3	4	5	6	7	8	9	10	11	12
arr[] =	11	14	19	23	27	-20	-14	-8	-4	1	2	4	7



first element

CASE1: arr[mid] > arr[0]

⇒ go right.

CASE2: arr[mid] < arr[0] ans ⇒ mid.

⇒ go left. ↳ K

	0	1	2	3	4	5	6	7	8	9	10	11	12
arr[] =	11	14	19	23	27	-20	-14	-8	-4	1	2	4	7

1) l=0, h⇒12, mid ⇒ 6 | arr[6] ⇒ -14 ans⇒6 h=mid-1
2) l=0, h⇒5, mid ⇒ 2 | arr[2] ⇒ 19 l⇒mid+1
3) l⇒3, h⇒5, mid ⇒ 4) arr[4] ⇒ 27, l⇒mid+1
4) l⇒5, h⇒5, mid ⇒ 5 | arr[5] ⇒ -20, ans⇒5 h⇒mid-1

TC: O(logn) SC: O(1)

Q2 Every element occurs twice except for 1 element. Find unique element. Duplicate are adjacent to each other. Elements are not sorted. Array is not sorted.

Ex1 arr[] =

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
3	3	1	1	8	8	10	10	19	6	6	2	2	4	4

Approach 1 : xor of everything : Tc: O(n)
Sc: O(1)

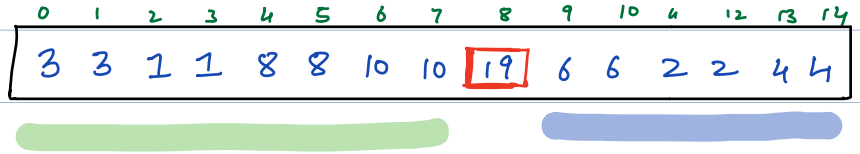
Ex1 arr[] =

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
3	3	1	1	8	8	10	10	19	6	6	2	2	4	4

1) First half : first occurrence of duplicate element is happening on even index

2) Second half : first occurrence of duplicate element is happening on odd index

→ [mid-1] & [mid+1]



⇒ Check if arr[0] is the answer.
⇒ Check if arr[n-1] is the answer.

1) l ⇒ 1, h ⇒ 13, mid ⇒ 7

⇒ not answer → if (arr[mid-1] == arr[mid])
mid = mid-1;

⇒ mid ⇒ 6 → go right
l ⇒ mid+2

2) l ⇒ 8, h ⇒ 13, mid ⇒ 10
mid ⇒ [9] → go left.
h ⇒ mid-1

3) l ⇒ 8, h ⇒ 8, mid ⇒ 10
answer

Pseudo Code

```
int findUnique ( int arr[], int n ) {

    if (n == 1)  return arr[0]

    if (arr[0] != arr[1])
        return arr[0]
    if (arr[n-1] != arr[n-2])
        return arr[n-1];

    l => 2 , h => n-3

    while ( l <= h ) {
        int mid = (l+h)/2

        if ( arr[mid-1] != arr[mid] &&
             arr[mid] != arr[mid+1])
            return arr[mid];

        if (arr[mid] == arr[mid-1])
            mid = mid-1.

        if ( mid % 2 == 0 )
            l = mid+2
        else
            h = mid-1
    }
}
```

TC: O(logn)

SC: O(1)

Q3 Find the max subarray sum of len K.

Ex arr[] =

0	1	2	3	4	5	6	7
3	2	5	4	6	3	7	2

K = 3

TC: O(n)

SC: O(1)

Q3 Given an array of integers, find maximum K such that { max subarray sum of len K $\leq$ B }

↓

Given as input.

Ex arr[] =

0	1	2	3	4	5	6	7
3	2	5	4	6	3	7	2

B = 20

Using Sliding window for each K

K = 1 max_Sub_sum = 7 $\leq$ 20 ans = 1

K = 2 max_Sub_sum = 10 $\leq$ 20 ans = 2

K = 3 max_Sub_sum = 16 $\leq$ 20 ans = 3

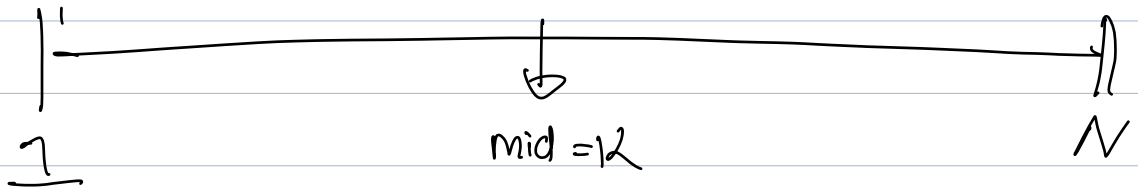
K = 4 max_Sub_sum = 20 $\leq$ 20 ans = 4

K = 5 max_Sub_sum = 25 > 20

TC: $O(n^2)$

SC: $O(1)$

1) Find 'd' ⇒ target.
2) Search Space ⇒ [1, N]



1) Case 1 : max_sub_sum (mid) < B
ans ⇒ mid, l ⇒ mid+1

2) CASE 2 : max_sub_sum (mid) > B
h ⇒ mid - 1

3) CASE 3 : max_sub_sum (mid) = B
return mid.

Tc: O(nlogn)
Sc: O(1)

Dry Run

Ex arr[] =

0	1	2	3	4	5	6	7
3	2	5	4	6	3	7	2

B ⇒ 21

ans ⇒ 4

l ⇒ 1, h ⇒ 8, mid ⇒ 4

max_sub_sum (4) ⇒ 20

l ⇒ mid + 1

l ⇒ 5, h ⇒ 8, mid ⇒ 6

max_sub_sum (6) ⇒ 27

h ⇒ mid - 1

l ⇒ 5, h ⇒ 5

max_sub_sum (5) ⇒ 25

h ⇒ mid - 1

l ⇒ 5, h ⇒ 4

Q4 Given N, find sqrt(N). Find closest integer.

N sqrt(N) $\Rightarrow$ $[1, N]$

Target = sqrt(N)

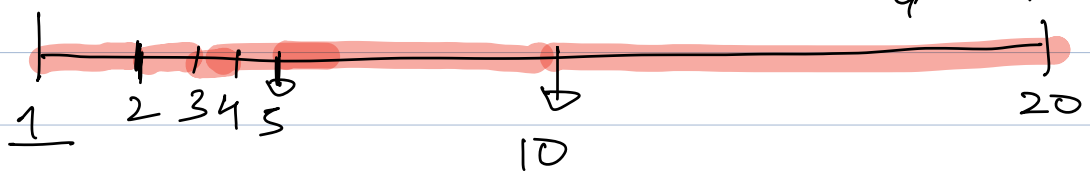
Search Space = $[1, N]$

$l = 1, h = 20$

$l = 1, h \Rightarrow 9$

$l = 1 \quad h \Rightarrow 4$

$l = 3 \quad h \Rightarrow 4$



CASE 1 : if (mid * mid $\geq n$.)

go left

h $\Rightarrow$ mid-1.

CASE 2 : if mid * mid $< N$.

go right # ans $\Rightarrow$ mid

l = mid+1

CASE 3: mid * mid == N

return mid.

Q Given a string made of only a, b, c

find smallest substring which contains a, b, c

a a b b c c a a b a b b c

l = 4

[3, N] = Seq