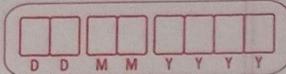


Unit - 1

Basics of Data Structures



Ques- Explain the difference between data type, abstract data type and data structure.

→ The difference between data type, abstract data type and data structure.

1] Data Type :-

a) A data type is a classification or categorization that specifies which type of data a variable can hold.

b) It defines the set of values that a variable can take on and the operations that can be performed on those values.

c) Common data types include integers, floating-point numbers, characters and booleans.

d) Data types are typically supported directly by programming languages and determine how data is stored in memory.

2] Abstract Data Type (ADT)

a) An abstract data type is a high-level description of a data structure that focuses on the operations that can be performed on the data and not on how they are implemented.

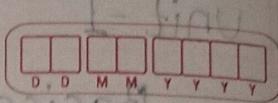
b) ADT's define a set of operations and their behavior but do not specify how these operations are implemented in code.

c) Examples of ADT's include stacks, queues, lists and dictionaries.

d) ADT's provide a blueprint for creating data structures but leave the actual implementation details to the programmer.

TYPES OF ADT :-

- 1] List
- 2] Stack
- 3] Queue



Q) Data Structure:

a) A data structure is concrete implementation of an abstract data type (ADT).

b) It involves organizing and storing data in a specific way in memory to support efficient operations.

c) Data structures determine how data is stored, accessed and manipulated.

d) Examples of data structures include arrays, linked lists, trees, graphs and hash tables.

e) Data structures provide the low-level details of how data is managed in a program.

Q) Explain the need of a data structure in programming point of view.



Need of a Data structure in programming Point of view

1] Data organization :-

Data organization structures provide a pathway to efficiently organize and store data.

This organization is essential for quick and easy access to information within a program.

2] Memory management :-

Data structures allow for efficient memory allocation and management.

They help ensure that memory is used effectively and can be released when it's no longer

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needed, preventing memory leaks.

3) Abstraction :-

Data structures abstract the underlying implementation details.

- programmers can work with high-level data structures like lists or sets without needing to understand the low-level details of how data is stored and retrieved.

4) Complex problem solving :-

Many programming problems involve complex data relationships and manipulations.

- Data structures like trees and graphs are essential for solving these kinds of problems efficiently.

5) Code reusability :-

- Using established data structures allows for code reusability.

- many programming languages provide built-in data structures, and there are libraries that offer specialized data structures for various purposes.

- This saves time and effort in development.

6) Algorithm design :-

Data structures often go hand in hand with algorithms.

- Efficient algorithms are designed with specific data structures in mind.

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The choice of data structure can profoundly impact the efficiency of an algorithm.

7) Scalability :- Data structures help in building scalable solutions.

As data volumes grow, the choice of the right data structure becomes critical for maintaining the performance of the program.

8) Debugging and Maintenance :-

- well-designed data structures make code easier to debug and maintain.
- clear data structures make it easier to understand the code's logic and find and fix issues.

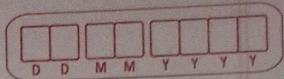
Ques - What is Data Structure, List the basic operations of Data structures.

Ans - Data structures are fundamental constructs used to organize and manage data in computer programs.

List of Basic operations :-

1) Arrays :-

a) Access (Read/Write) :- Retrieve or modify an element at a specific index.



- b) insertion :- Add an element at the end or a specified index.
- c) deletion :- Remove an element from a specific index.
- d) search :- Find the index of a specific element.

2) Linked Lists :-

- a) Insertion :- Add a new node at the beginning, end or a specific position.
- b) deletion :- Remove a node from a specific position.
- c) search :- Find a node with a specific value.

3) stacks :-

- a) Push :- Add an element to the top of the stack.
- b) Pop :- Remove and return the top element.
- c) Peek :- View the top element without removing it.

4) queues :-

- a) Enqueue :- Add an element to the back of the queue.
- b) Dequeue :- Remove and return the front element.
- c) Peek :- view the front element without removing it.

check if the queue is empty.

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5] Trees :-

a) Insertion :- Add a new node.

b) Traversal :- visit nodes in specific order.

c) search :- Find a node with a specific value.

d) Deletion :- Remove a node.

6] Graphs :-

a) Insertion :- Add a vertex or edge.

b) Deletion :- Remove a vertex or edge.

c) Traversal :- visit vertices and edges.

7] Hash Tables :-

a) Insertion :- Add a key-value pair.

b) Lookup :- Retrieve the value associated with a key.

c) Deletion :- Remove a key-value pair.

d) Hashing :- convert a key into an index for efficient retrieval.

8] Heaps :-

a) Insertion :- Add an element while maintaining the heap property.

b) Extraction :- Remove and return the element with the highest or lowest priority.

c) Peek :- view the highest or lowest priority element.

9] Sets and maps :-

a) Insertion :- Add an element with or without associated data.

D	D	M	M	Y

- b) Deletion :- Remove an element by key (map)
 c) Lookup :- Retrieve an element by key (map)

Ques - State the merits and demerits of linked list over array.



Merits of Linked Lists over Arrays :-

- 1) Dynamic size :-
 - Linked lists can dynamically grow or shrink in size.
 - This flexibility allows for efficient memory usage as you only allocate memory for the elements you need, unlike arrays, which have a fixed size.

2) Efficient Insertions and Deletions :-

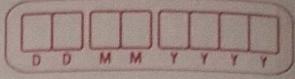
- Insertions and deletions are generally more efficient in linked lists, especially when dealing with elements in the middle of list.

3) No wasted memory :-

- Linked lists don't suffer from memory wastage due to unused slots, as in the case of arrays with preallocated sizes.

4) No Need for copying :-

- Adding or removing elements from a linked list typically doesn't involve copying the entire data structure, as it might in arrays when resizing.



② Support for complex Data structures:-
Linked lists are fundamental in building more complex data structures like stacks, queues and graphs.

• Demerits of Linked Lists over Arrays:-

1) Random Access :-

- Linked lists provide slow random access.
- To access an element at a specific position, you often have to traverse the list from the beginning, which takes $O(n)$ time in the worst case.

2) Memory overhead :-

- Each element in a linked list requires additional memory for the pointer/reference to the next element, resulting in higher memory overhead compared to arrays.

3) Cache Inefficiency :-

- Linked lists can be less cache-friendly.
- Since elements are scattered in memory, accessing them may result in more cache misses compared to the contiguous memory layout of arrays.

4) Complexity :-

- Linked lists introduce more complexity due to the need to manage pointers/references, making them potentially

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error-prone, especially in terms of memory management.

5] Additional pointer operations

→ Traversing and manipulating linked lists require extra pointer operations, which can make code less readable and more challenging to debug.

6] Fixed storage complexity

Some types of linked lists don't support reverse traversal efficiently.

• To traverse in reverse, you may need to create a new linked list or use additional data structures, which can be inefficient.

Ques- Explain Time complexity and space complexity with suitable example.

→ Time complexity

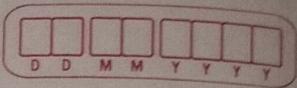
→ Time complexity measures the amount of time an algorithm takes to complete its execution as a function of the input size.

• It provides an estimate of how the algorithm's runtime grows with increasing input size.

• It is usually expressed using Big O notation such as $O(n)$, $O(n^2)$, etc.

Example:- Linear Search

• Suppose you have an array of n elements.



and you want to find a specific element in array.

- In the worst-case scenario, where the element is at the end of the array or not present, you would have to check each element one by one.
- The time complexity of a linear search is $O(n)$, as the number of comparisons is directly proportional to the size of the input array (n).

2) space complexity :-

• Space complexity measures the amount of memory space an algorithm uses as a function of the input size.

• It helps determine how much memory is required to execute an algorithm efficiently.

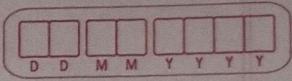
• Example :- mergesort

• Merge sort is a sorting algorithm that uses a divide-and-conquer approach.

• It divides the input array into two halves, recursively sorts each half, and then merges them back together.

• In terms of space complexity, merge sort requires additional memory to store temporary arrays during the merge step.

• The space complexity of merge sort is $O(n)$, where n is the size of the input array, because it needs to create temporary arrays of the same size as the input.



Que- Difference between arrays and linked lists

→ **Ans:** Array Linked List

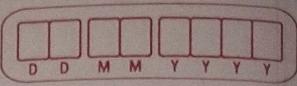
- | | |
|---|---|
| 1) Arrays are stored in contiguous location. | 1) Linked lists are not stored in contiguous location. |
| 2) Fixed in size. | 2) Dynamic in size. |
| 3) Memory is allocated at compile time. | 3) Memory is allocated at run time. |
| 4) Uses less memory than linked lists. | 4) Uses more memory because it stores both data and the address of next node. |
| 5) Elements can be accessed easily. | 5) Element accessing requires the traversal of whole linked list. |
| 6) Insertion and deletion operation takes time. | 6) Insertion and deletion operations is faster. |

Explain the various data representations in detail with suitable examples.

Various Data Representation

1) Arrays

- An array is a contiguous block of memory where elements are stored at consecutive memory locations.
- Each element in an array is accessed



- Example :- An array of integers 'arr' where ' $arr[0] = 5$ ', ' $arr[1] = 10$ ' and ' $arr[2] = 15$ '.

2) Linked lists :-

- A linked list is a data structure where each element (node) contains both data and a reference (pointer) to the next node in the sequence.
- Example :- A singly linked list with elements $10 \rightarrow 20 \rightarrow 30$, where each node contains a value and a reference to the next node.

3) Stacks :-

- A stack is a linear data structure that follows the Last-In - First-Out (LIFO) Principle.
- Elements are added and removed from one end (the top).
- Example :- A stack of integers where you push elements onto the top and pop element from the top.

4) Queues :-

- A queue is a linear data structure that follows the First-In - First-Out (FIFO) Principle.
- Elements are added at one end (rear) and removed from the other end (front).

D	D	M	M	Y	Y	Y
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- Example :- A queue of people waiting in line where people are added at the end of the line and served from the front.

③ Trees :-

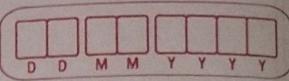
- Trees are hierarchical data structures consisting of nodes with parent-child relationships.
- Example :- A binary search tree (BST) where each node has a value, a left child (with a smaller value) and a right child (with a larger value).

④ Graphs :-

- Graphs represent a collection of nodes (vertices) and edges that connect these nodes.
- Example :- A social network graph where users are nodes, and friendships are edges connecting them.

⑤ Hash Tables :-

- Hash tables use a hash function to map keys to indices in an array, allowing for efficient key-value pair storage and retrieval.
- Example :- A dictionary where words (keys) are mapped to their definitions (values) using a hash table.



Ques:- Define what is Algorithm and explain the different ways of analysing algorithm.



Algorithm :-

- An algorithm is a step-by-step procedure or a set of well-defined modules for solving a problem or accomplishing a specific task.

Ans:- (T28) What do you mean by algorithm?

(H1W) b) Different ways of analysing algorithm:-

1) Time complexity Analysis:-

• Time complexity measures how the running time of an algorithm grows as a function of the input size. It helps determine how efficient an algorithm.

• Common notations used for time complexity analysis include Big O (O), Big Omega (Ω) and Big Theta (Θ).

• Example :- An algorithm with a time complexity of $O(n^2)$ means that its running time increases quadratically with the input size.

2) Space complexity Analysis:-

• Space complexity measures how much memory an algorithm consumes in relation to the input size. It helps assess the algorithm's memory efficiency.

• Like time complexity, space complexity is also analyzed using Big O notation.

• Example:- An algorithm with a space

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complexity of $O(n)$ indicates that it uses linear memory in relation to the input size.

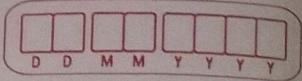
③ Best case, worst case and Average case :-

Analysis:-

- Algorithms may perform differently under different scenarios. Analyzing best-case, worst-case and average-case scenarios helps understand their behavior in various situations.
- Best-case analysis identifies the minimum possible runtime.
- Worst-case analysis identifies the maximum possible runtime.
- Average-case analysis provides a more realistic expectation of performance based on probabilities.

④ Asymptotic Analysis :-

- Asymptotic analysis focuses on how an algorithm behaves as the input size approaches infinity. It helps identify the dominant factors affecting the algorithm's performance.
- It often involves simplifying the time or space complexity expression to its most significant term.
- Example:- Identifying that an algorithm's time complexity is $O(n^2)$ allows us to know that its efficiency decreases as the input size grows.



5] Empirical Analysis:

- Empirical analysis involves running an algorithm on real input data and measuring its actual performance.
- It is useful for practical assessments but may not provide a precise mathematical description of an algorithm's behavior.

6] Amortized Analysis:

- Amortized analysis considers the average time or space usage per operation over a sequence of operations.
- It helps determine whether an algorithm's occasional costly operations are balanced by cheaper operations, resulting in good overall performance.

Ques- Explain recursion and types of recursion with a C program to find the factorial of N given numbers.

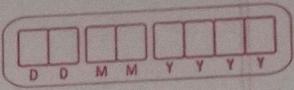
→ Recursion:

Recursion is a programming technique where a function calls itself to solve a problem.

Types of recursion:

1) Direct Recursion:

In this type of recursion, a function calls itself directly.



2) Indirect Recursion:

In this type of Recursion, two or more functions call each other in a circular manner, eventually leading to a base case.

C program to Find the factorial of a given number:-

```
#include <stdio.h>
```

```
int factorial (int n)
```

```
if (n == 0)
```

```
    return 1;
```

```
else
```

```
    return n * factorial(n - 1);
```

```
int main()
```

```
    int n;
```

```
    printf("Enter a non-negative integer: ");
```

```
    scanf("%d", &n);
```

```
    if (n < 0)
```

```
{
```

```
    printf("Factorial is undefined");
```

```
    else
```

```
        printf("Factorial is %d", factorial(n));
```

D	D	M	M	Y	Y	Y
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```

int result = factorial(n);
printf("Factorial of %d is %d\n", n,
       result);
return 0;
    
```

Ques - Illustrate the recursive function with help of Tower of Hanoi.

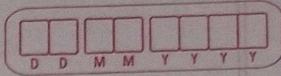


- The Tower of Hanoi is a classic problem that is often used to illustrate recursion.
- It involves three rods and a number of disks of different sizes, which can slide onto any rod.
- The puzzle starts with the disks in a neat stack in ascending order of size on one rod, the smallest at the top.
- The goal is to move the entire stack to another rod, subject to the following rules:
 - 1] only one disk can be moved at a time.
 - 2] A disk can only be placed on top of a larger disk or an empty rod.
 - 3] The objective is to move all the disks to the target rod, using the "spare" rod as an intermediary, while following the above rules.

Recursive Solution:-

The Tower of Hanoi problem can be elegantly solved using recursion.

- Here's how the recursive solution works



Step by step :-

1) Base case :-

When there is only one disk ($n=1$) on the source peg, you can move it directly to the destination peg. This is the simplest case and serves as the base case for the recursion.

2) Recursive step :-

- To move n disks from the source peg to the destination peg, you can break it down into three subproblems:-

1) Move the top $n-1$ disks from the source peg to the auxiliary peg.

2) Move the largest n^{th} disk from the source peg to the destination peg.

3) Move the $n-1$ disks from the auxiliary peg to the destination peg (using the source peg as auxiliary).

void Hanoi($int n, string a, string b, string c$) {

if ($n == 1$) /* Base case */ : simple [d]

move(a,b); // moving primitively [c]

else { $T(n-1, a, c, b)$; $T(n-1, c, b, a)$ } // Distr

{ } // end of else block

Hanoi($n-1, a, c, b$); // -non go main [e]

move(c,b); // moving primitive [d] to [b]

upini Hanoi($n-1, c, b, a$); // bro filoit nups?

3. recursive bro rozgabing

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cover Explain the classification Data structure.

Classification of Data structures :-

a) Primitive OR non-Primitive :-

a) Primitive :- These are fundamental data structures provided by programming languages such as integers, floats, characters and booleans.

b) They store a single data element.

b) Non-Primitive :- They are constructed using primitive data types to store collections of data elements.

- Examples include arrays, lists, stack

2) Static or Dynamic :-

a) Static :- These have a fixed size that is determined at the time of creation and cannot be changed during runtime.

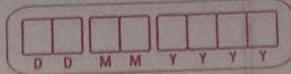
- Arrays are an example of static data structure.

b) Dynamic :- These can grow or shrink in size during program execution.

- Dynamic arrays, linked list and trees are examples.

3) Linear OR Non-Linear :-

a) Linear :- Data elements are organized sequentially and each element has a unique predecessor and successor.



- Examples include arrays, linked lists, stack

b) Non-Linear :- Data elements are organized in a hierarchical or arbitrary manner and elements may have multiple predecessors or successors.

- Examples include trees and graphs

4] Homogeneous or Heterogeneous :-

a) Homogeneous :- All elements in the structure are of the same data type.

- For example an array of integers.

b) Heterogeneous :- Elements in the structure can be of different data types.

- Example, a struct or record in some programming languages can contain different data types within a single structure.