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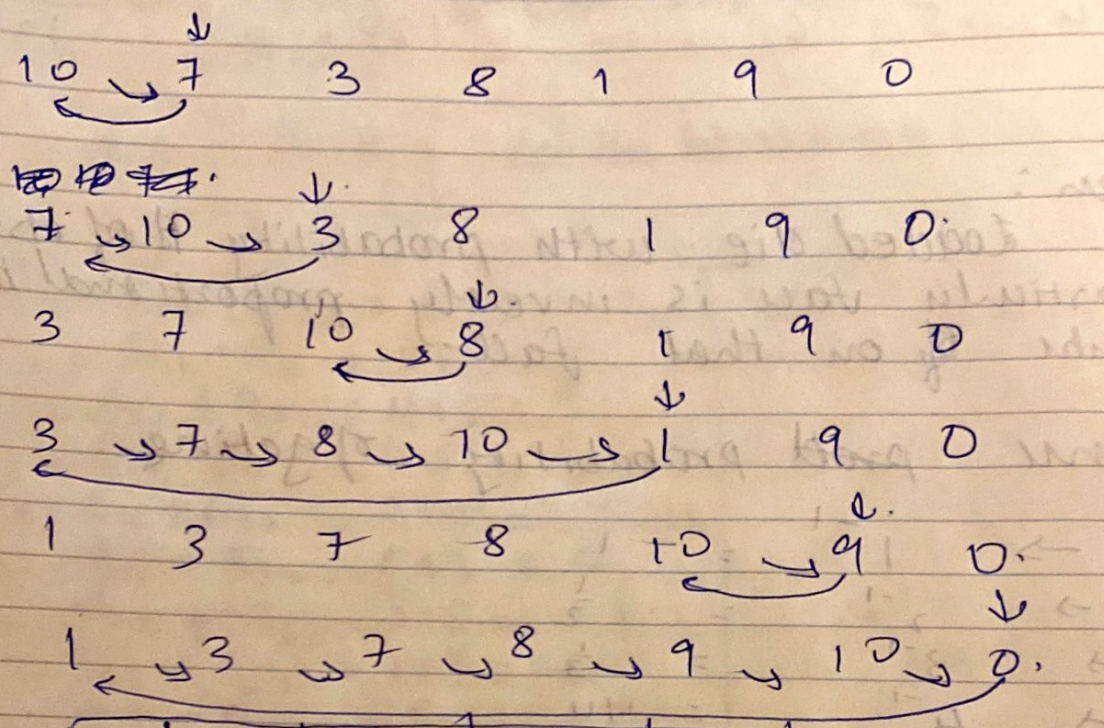
Homework 3.

Q1) ~~Insert~~ Insertion Sort.

A = [10, 7, 3, 8, 1, 9, 0].

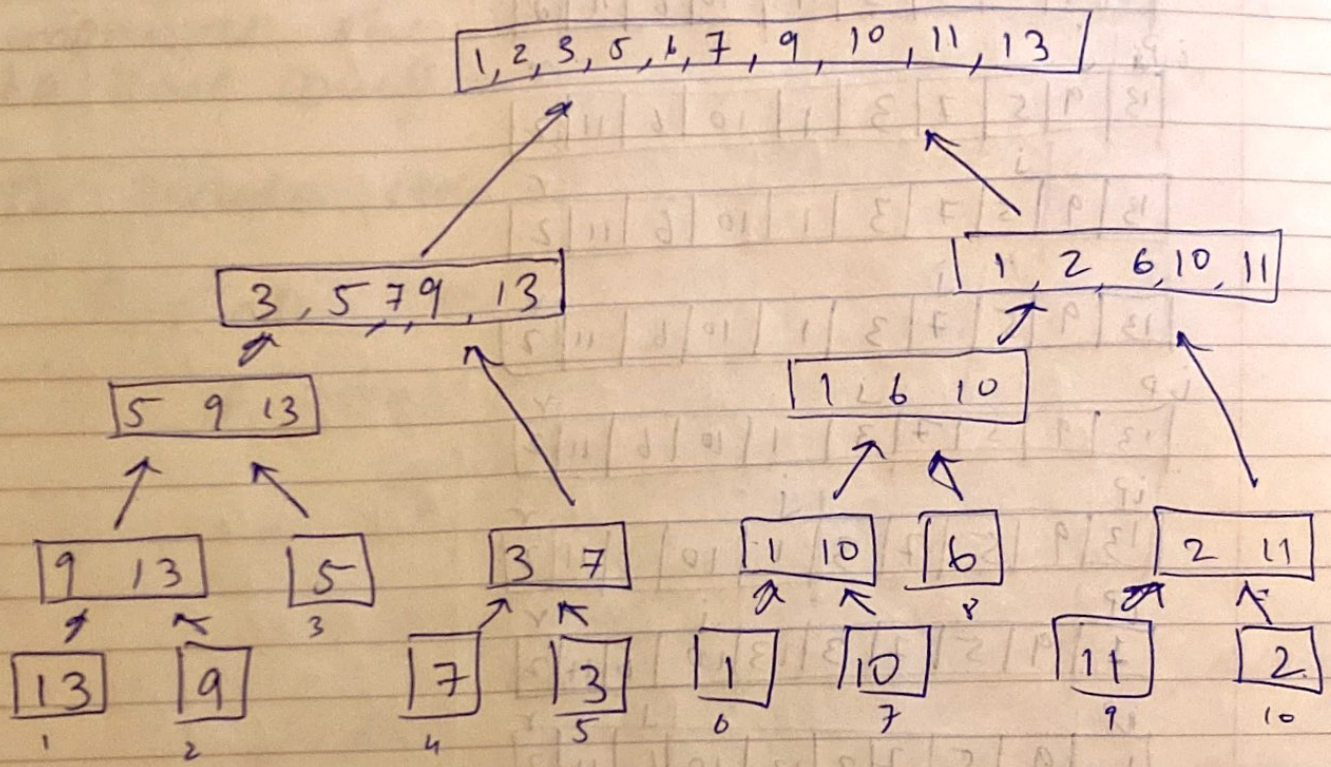


Final State: | 0 | 1 | 3 | 7 | 8 | 9 | 10. |

Ans

02) Merge sort · Stages

A = [13, 9, 5, 7, 3, 1, 10, 6, 11, 2]



Q3) Use Strassen's method to compute the product of the following two matrices.

Step 1 $A = \begin{bmatrix} 2 & 3 \\ 4 & 5 \end{bmatrix}$ & Let this represent as $\begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$

$B = \begin{bmatrix} 5 & 6 \\ 1 & 3 \end{bmatrix}$. Let this be represented as $\begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$

Step 2

$S_1 = B_{12} - B_{22} = 6 - 3 = 3.$
$S_2 = A_{11} + A_{12} = 2 + 3 = 5.$
$S_3 = A_{21} + A_{22} = 4 + 5 = 9.$
$S_4 = B_{21} - B_{11} = 1 - 5 = -4.$
$S_5 = A_{11} + A_{22} = 2 + 5 = 7.$
$S_6 = B_{11} + B_{22} = 5 + 3 = 8.$
$S_7 = A_{12} - A_{22} = 3 - 5 = -2$
$S_8 = B_{21} + B_{22} = 1 + 3 = 4$
$S_9 = A_{11} - A_{21} = 2 - 4 = -2.$
$S_{10} = B_{11} + B_{12} = 5 + 6 = 11$

Step 3.

$P_1 = A_{11} S_1 = 2 \times 3 = 6$
$P_2 = S_2 B_{22} = 5 \times 3 = 15$
$P_3 = S_3 B_{11} = 9 \times 5 = 45$
$P_4 = A_{22} S_4 = 5 \times -4 = -20$
$P_5 = S_5 S_6 = 7 \times 8 = 56$
$P_6 = S_7 S_8 = -2 \times 4 = -8.$
$P_7 = S_9 S_{10} = -2 \times 11 = -22.$

Step 4. Computing desired product

$C_{11} = P_5 + P_4 - P_2 + P_6$
$= 56 - 20 - 15 - 8$
$= 13$

$C_{12} = P_1 + P_2$
$= 6 + 15$
$= 21$

$C_{21} = P_3 + P_4$
$= 45 - 20 = 25$

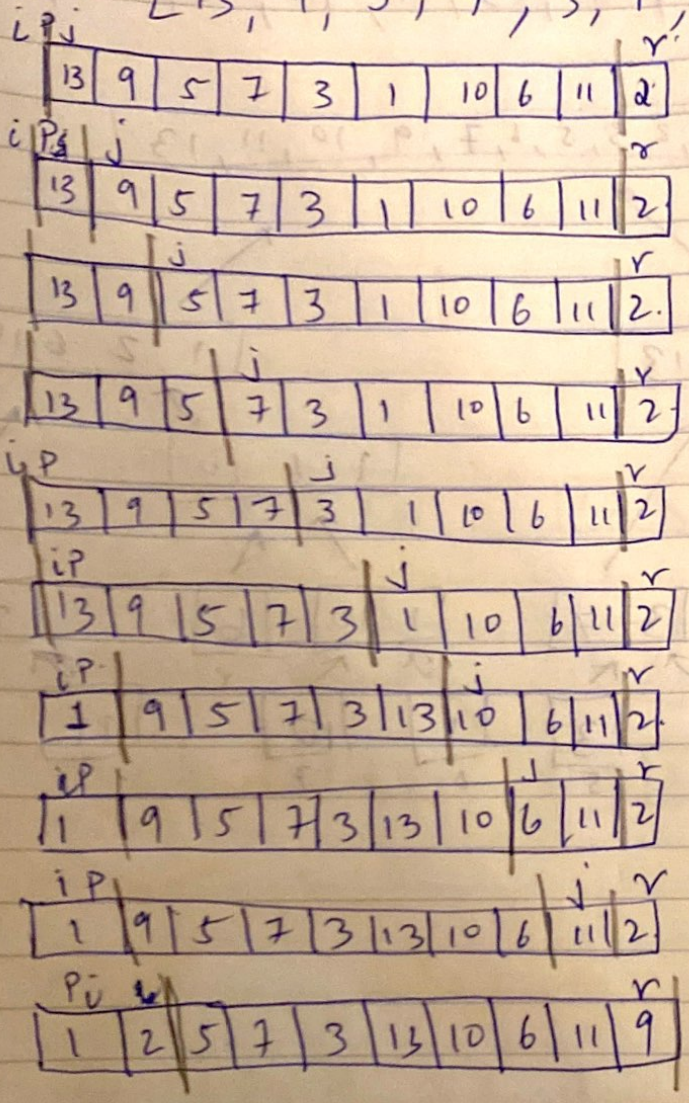
$C_{22} = P_5 + P_1 - P_3 - P_7$
$= 56 + 6 - 45 + 22$
$= 39.$

$C = \begin{bmatrix} 13 & 21 \\ 25 & 39 \end{bmatrix}$

Ans

Q4) Partition function in quick sort array:

A = [13, 9, 5, 7, 3, 1, 10, 6, 11, 2].



Q5) Let us consider that the list is sorted in ascending order. then we will use the recursive binary search to implement the problem. as it has a time complexity of O(log n).

The pseudo code is as follows.

```
FindInden (A[], start, array size-1)
    if array siz > strt {
        calculate mid
        if Array[mid] == mid+1
            then return mid.
        else if array[mid] > mid+1.
            return Find Inden (A, start, mid-1)
        else if array[mid] < mid+1
            return Find Inden (A, mid+1, arrsiz-1)
    else return -1;
```

- In our algorithm we are finding the middle element after each recursion hence if we find Array[mid] means that our algorithm will return the found value else it will return -1.
- If we did we will keep on recursively dividing untill the start or end of array and if we do not get the correct value then we will return false or -1
- Hence our algorithm is correct.

Q6) Given:

Loaded die: with probability ~~H~~ of obtaining a particular face is inversely proportional to the number on that face.

Hence ~~prob~~ probability of getting

$1 \rightarrow 1^{-1} = 1$

$2 \rightarrow 2^{-1} = \frac{1}{2}$

$3 \rightarrow 3^{-1} = \frac{1}{3}$

$4 \rightarrow 4^{-1} = 1/4$

$5 \rightarrow 5^{-1} = 1/5$

$6 \rightarrow 6^{-1} = 1/6.$

Probability ~~Ha~~ of getting number higher than 3 is

$\frac{1}{4} + \frac{1}{5} + \frac{1}{6} = \frac{74}{120}$ Ans