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#### Assignment-3: Cryptography Analysis and Implementation

**Objective:** The objective of this assignment is to analyze cryptographic algorithms and implementation in approximation in a section of the control of the

#### Instructions:

Research:Beginbyconductingresearchondifferentcryptographicalgorithmssuchassymmetric key algorithms (e.g., AeS, DeS), asymmetric key algorithms (e.g., RSA,elliptic Curve Cryptography), and hash functions (e.g., MD5, SHA-256). Understandtheirproperties,strengths,weaknesses,andcommonusecases.

**Analysis:** Choose three cryptographic algorithms (one symmetric, one asymmetric, and one hash function) and write a detailed analysis of each. Include the following points in your analysis:

Brieflyexplainhowthealgorithmworks.

Discussthekeystrengthsandadvantagesofthealgorithm. Identifyanyknownvul nerabilities orweaknesses.

 $\label{lem:providereal-worldexamples} Providereal-worldexamples of where the algorithm is commonly used.$ 

# Implementation:

Select one of the cryptographic algorithms you analyzed and implement it in apractical scenario. You can choose any suitable programming language for theimplementation. Clearly define the scenario or problem you aim to solve using cryptography. Provide step-by -step instructions on how you implemented the chosen algorithm. Include codes nippets and explanations to demonstrate the implementation. Test the implementation and discuss the results.

#### Security Analysis:

Performase curity analysis of your implementation, considering potential attack vectors and countermeasures.

Identify potential threats or vulnerabilities that could be exploited. Propose countermeasures or best practices to enhance these curity of your implementation.

Discussanylimitationsortrade-offsyouencounteredduringtheimplementationprocess. Conclusion: Summarizeyourfindings and provide insights into the importance of cryptography in cybersecurity and ethical hacking.

#### SubmissionGuidelines:

Prepareawell-structured report that includes the analysis, implementations teps, codes nippets, and security analysis.

Useclearandconciselanguage, providing explanations where necessary. Include any references or sources used for research and analysis.

Compilealltherequiredfiles(report, codes nippets, etc.) into a single zipfile for submission.

# **Analysis**:

### Symmetric Algorithm: AeS

The A dvanced e ncryption S tandard (A e S) is a symmetric block cipher chosen by the

U.S. government to protect classified information. **ACS** is implemented in softwareandhardwarethroughouttheworldtoencryptsensitivedata. It is essential for government computers ecurity, cybers ecurity and electronic data protection.

# HowAeSencryptionworks:

AeSincludesthreeblockciphers:

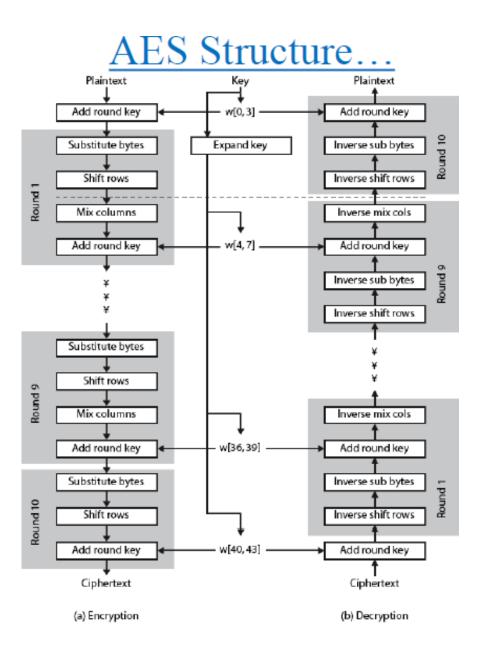
1.ACS-128usesa128-bitkeylengthtoencryptanddecryptablockofmessages.2.ACS-192 uses a 192-bit key length to encrypt and decrypt a block of messages.3.ACS-256 uses a 256-bit key length to encrypt and decrypt a block of

messages. Pachcipherencryptsanddecryptsdatainblocksof 128 bitsusing cryptographickeys of 128,19 2 and 256 bits, respectively. Symmetric, also known assecret key, ciphersuse the same key for encrypting and ecrypting. The sender and the receiver must both

know--anduse--the samesecretkey.

Thegovernmentclassifiesinformationinthreecategories: Confidential, Secretor Top Secret. All key lengths can be used to protect the Confidential and Secret level. Top Secretinformation requires either 192-or 256-bitkey lengths.

There are 10 rounds for 128-bit keys, 12 rounds for 192-bit keys and 14 rounds for 256-bit keys. A round consists of several processing steps that include substitution, transposition and mixing of the input plaintext to transform it into the final output of ciphertext.



The ACS encryption algorithm defines numerous transformations that are to beperformed and tastored in an array. The first step of the cipher is to put the data into an array, after which the cipher transformations are repeated over multiple encryption rounds.

The first transformation in the AeS encryption cipher is substitution of data using asubstitution table. The second transformation shifts data rows. The third mixescolumns. The last transformation is performed on each column using a different partoftheencryption key. Longerkeys needmore rounds to complete.

### StrengthsofAes:

Aes data encryption is a more mathematically efficient and elegant cryptographicalgorithm, but its mainstrength rests in the option for various keylengths. Aes allows you to choose a 128-bit, 192-bit or 256-bit key, making it exponentially stronger than the 56-bit key.

# Benefitsoradvantagesof ACS

- As it is implemented in both hardware and software, it is most robust security protocol.
- It uses higher length key sizes such as 128, 192 and 256 bits for encryption. Henceit makes AeSalgorithmmorerobust againsthacking.
- > It is most common security protocol used for wide variety of applications such as wireless communication, financial transactions, e-business, encrypted datastorageetc.
- > It is one of the most widely used commercial and open source solutions acrosstheworld.
- > Noonecanhackyourpersonalinformation.
- For 128 bit, about 2128 attempts are needed to break. This makes it very difficult to hack it as a resultities very safe protocol.

### **DrawbacksordisadvantagesofAES**

- Itusestoosimplealgebraicstructure.
- everyblockis alwaysencryptedinthesameway.
- Hardtoimplementwithsoftware.
- ➤ AeSincountermodeiscomplextoimplementinsoftwaretakingbothperformanceandsecurity intoconsiderations.

#### Weakness:

The biggest problem with ACS symmetric key encryption is that you need to have awaytogetthekeytothepartywithwhomyouaresharingdata. Symmetric encryption keysare of tenencrypt edwithan asymmetrical gorithmlike RSA and sent separately.

**Examples** where ACS technology is used: VPN Implementations. File transfer protocols (FTPS, HTTPS, SFTP, OFTP, AS2, WebDAVS) Wi-Fi security protocols (WPA-PSk, WPA2-PSk)

# Asymmetric Algorithm: e16amal encryption

el6amal cryptosystem can be defined as the cryptography algorithm that uses the publicand private keyconcepts to secure communication between two systems. It can be considered the asymmetric algorithm where the encryption and decryption happen by using public and private keys. In order to encrypt the message, the public key is used by the client, while the message could be decrypted using the private key on the server end. This is considered an efficient algorithm to perform encryption and decryption as the keys are extremely tought to predict. The sole purpose of introducing the message trans action's signature is to protect that against the public key or the cryptography algorithm could very effectively achieve.

# el6amalencryptionAlgorithmwithexample

The e16amal encryption algorithm method's sole concept is to make it nearlyimpossible to calculate the encryption approach even if certain important information is known to the attacker. It is mainly concerned about the difficulty of leveraging the cyclic group to find the discrete logarithm.

It will be very easy to understand, using a simple example. Suppose that even if thevalue like g^a and g^b are the values known to the attacker, the attacker will find itextremely difficult of indout the value of g^abwhich is nothing but the cracked value.

Inordertounderstandtheentirescenario, we need to go in a step wise manner on how the encryption and decryption of messages happen actually. We will be considering the example of two peers who are willing to exchange data in a secure manner by leveraging the el6amal algorithm. Let's suppose user 1 and user 2 want to exchange the information secretly, in that case, the following procedure will be followed.

### Step 1: Generation of the publicand private keys.

Theuser 1 will try to select a very long or large numbers, and mean while, he will also choose a cyclic group Fx. From this cyclic group, he will be further choosing another component b and one more element c. The values will be selected in the manner that if passed through a particular function, the outcome will be equivalent to 1.

Once the value selection phase is over, a value will be calculated that will be further used to generate the private key. By applying the formula  $fm=b^c$ , the value will be calculated. In the current scenario, user 1 will select F,  $fm=b^c$ , a, b as their publickey, while the values of a will be saved as the private key, which will be further used as the private key.

# Step2:User2willencryptthedatausingthepublickeyofUser1.

In order to begin the encryption of the message, there are certain values that user2needs to pick. The user2 will also require to pick one of the values p from the cyclicgroup. The cyclic group will be the same as it was for the user1. The value should be picked in a manner so that Incpasses with a inthe particular function will generate the outcome 1.

know the user2 will generate some other values that will be used to encrypt themessageusingthepublickey. The value generates will be Pm=b^p. The other revalue b^c will be equal to b^ap. The outcome of this computation will be multiplied to the other value Z in order to get closer to the encryption method. Eventually, the value will be sentusing the outcome of computations on b^p, Z^b^ap.

### Step3:Decryptionofthemessageatuser1end.

The user1 will then use the computation of the values picked in the first and secondphase to identify the appropriate number, which will be used to decrypt the encryptedmessage. The User1 will be processing b^ap, and then the outcome will be used to divide the by Z in order to get the decrypted value. The decrypted value is somethingthat is that was encrypted in the second phase.

Intheabovescenario, the user lhas initiated the process by calculating the private and public key, which is the algorithm's soul. The key is further used by user 2 in the second step in order to encrypt the method.

The message is encrypted so that they value computed in that initial phase could beleveraged to decrypt the message. In the third step, it could be witnessed that afterdiving the entire value with the number that is computed in the third step itself totally decrypts the message making it readable for the end-user. The same approach is followed everywhen the urge topass the message securely occurs

#### Advantageofel6amalalgorithm:

The advantage of the el6amal algorithm is the generation of keys using discretelogarithms. encryption and decryption techniques use a large computing process sothat the encryption results are twice the size of the original size.

### Disadvantageofel-Gamal

Themaindisadvantage of e1-

Gamalisheneedforrandomness, and its slowers peed (especially for signing). Another potential disadvant age of the Cl-Gamalal gorithmist hat the message expansion by a factor of two takes placed uring encryption.

**Example:** Alice chooses pA = 107, A = 2, dA = 67, and she computes A = 267  $\equiv 94 \pmod{107}$ . Herpublickeyis (pA, A, A)=(2,67,94), and herprivate keyis dA = 67. sends the encrypted message (28, 9) to Alice.  $-dA = 9.28 - 67 \equiv 9.28106 - 67 \equiv 9.43 \equiv 66 \pmod{107}$ .

#### Hashfunction: MO5

MD5isacryptographichashfunctionalgorithmthattakesthemessageasinputofanylength and changes it into a fixed-length message of 16 bytes. MD5 algorithm standsfor the message-digest algorithm. MD5 was developed as an improvement of MD4,withadvancedsecuritypurposes. The output of MD5 (Digest size) is always 128 bits. MD5 was developed in 1991 by Ronald Rivest.

#### HowdoesMD5work?

MO5 runs entire files through a mathematical hashing algorithm to generate asignature that can be matched with an original file. That way, are ceived file can be authenticated as matching the original file that was sent, ensuring that the right files getwhere they need to go.

The MO5 hashing algorithm converts data into a string of 32 characters. For example, the word "frog" always generates this hash: 938c2cc0dcc05f2b68c4287040cfcf71. Similarly, a file of 1.26 Balsogenerates a hash with the same number of characters. When you send that file to some one, their computer authenticates its hash to ensure it matches the one you sent.

Ifyouchange, just one bit in a file, no matter how large the file is, the hashout put will be completely and irreversibly changed. Nothing less than an exact copy will pass the MD5 test.

# StrengthsofM05:

MD5 (Message Digest Method 5) is a cryptographic hash algorithm used to generate a128-bit digest from a string of any length. It represents the digests as 32 digithexadecimal numbers. Ronald Rivest designed this algorithm in 1991 to provide themeans for digitalsignature verification.

# AdvantagesoftheMO5algorithm

- It's easier to compare and storesmaller has besusing MD5 Algorithms than it is to store a large variable—length text.
- ByusingMD5,passwords arestoredin128-bit format.
- Youmaycheckforfilecorruptionbycomparingthehashvaluesbeforeandafter transmission. To prevent data corruption, file integrity tests are validoncethehashesmatch.

• Amessagedigestcaneasilybecreatedfromanoriginalmessageusing MO5.

#### DisadvantagesoftheMD5algorithm

- Whencompared to other algorithms like the SHA algorithm, MD5 is comparatively slow.
- Itispossibletoconstructthesamehashfunctionfortwodistinctinputsusing MO5.
- MO5islesssecurewhen comparedtotheSHA algorithmsinceMO5ismorevulnerable tocollisionattacks.

### Real-WorldexampleofHashing:OnlinePasswords

Every time you attempt to log in to your email account, your email provider hashesthe password YOU enter and compares this hash to the hash it has saved. Only whenthetwo hashes matchareyou authorized to access your email.

# IMPLEMENTATIONS:

#### **Aimandabstract**

Security is everyone 'stop concerninthe modernera. Everyone uses the internet these days for a variety of p urposes, including data and money transfers. Therefore, we build a cryptosystem that leverages the **Problem** (DLP) for Discrete Logarithm encryption,whichmakestheencryptionmethodmoresafe,inordertoincreasethesecurityofthecurrent system. It's known as elgamal encryption. A public key cryptosystem is used. Both the encryption and decoding asymmetric The processes require keys. current@lgamalcryptosystemencryptsdatausing.justonekey.Thesedays,therearesomanyinstanceswh ereunauthorisedclientsgetaccesstocrucialinformation. Inordertoaddanadditionallayerofsecuritytothes ystem, wesuggest that the elgamal Cryptosystem be modified in this work. With this update, the user is able to en cryptthemessagewithnumerous private keys. The text will be converted into integer values by an existingalgorithm,increasingthefilesizeby2'n.Bytranslatingtheintvaluestothecorresponding characters, wewereable to reduce the file size.

**Reywords:**Security,cryptosystem,DiscreteLogarithmProblem,encryption,Elgamal,privatekey

S.

# **Existing Method**

Existing Elgamalencryption consists of three parts. They are key generation, encryption and decryption.

- 1. Bobgeneratespublicandprivatekey:
  - Bob chooses averylargeprimenumberp.
  - Fromthep, hefindsq which is the primitive root of p.
  - Thenhecomputesy= g<sup>x</sup>modp.
- Bob publishes p, g and y as his public key and retains x asprivatekey.
- 2. AliceencryptsdatausingBob'spublickey:
  - Aliceselects arandomintegerk<p.
  - Thenshecomputes c1=g<sup>k</sup>modp.
  - Shemultiplesy with Mithatisconsider asc 2.
  - Thenshesends(c1,c2).
- 3. Bobdecryptsthemessage:
  - Bobcalculatess'=(c1x-1)modp.
  - Thenhemultiplies c2bys'toobtain M.

#### PROPOSEDSYSTEM:

Existing algorithm uses discrete logarithmic problem it is very hard to crack the keyusing brute force attack. With modified and dedicated hardware, we can crack thekey.

Butourgoalistodevelopanenhanced@lgamalencryptionsystemwhichsystem is not crack easily. existing encryption method is using integer as a ciphertext. But, in our proposed system we use character string as a cipher text so it reduce the file size. In this proposed algorithm it is impossible to break it via brute-

force attack. A nd also C ipher text attack is not possible since attacker has no idea about the keys and length of the message.

# ModuleandDescription:

- 1- Reygeneration(Server):
  - Choosearandomprimenumber(p)andchoosearandomprimiliveroot(g)of the prime number.

- o Choosearandominteger(n)asthenumberofkeys (rounds)inourprivatekey
- o Generate n random integers (xn) which will act as our private key, applyyn =gxmod pto each oftheintegers togenerate Y (list ofintegers).
- o Sendp,g,Ytoclient which acts asourpublickey.

#### 2- encryption(Client):

- o Choosenlength(Y)randomnumbers(kn)
- o Compute Ykmod p using all numbers in Y and a and multiply themand storeinonevariablec and then compute c=cmod p.
- o ComputealistA=gk modpusing all integersin akn
- o Padthemessagewithcrandomcharactersinthebeginning andc/2intheend.
- o encrypt message by multiplying c with message and then convert them to characters resulting inencrypted message  ${\bf B}$
- o Send ABtoServer.

#### 3- Decryption(Server):

- o ReceiveABfrom client
- o Computecbyapplying Axmodpon all the integers in Acorresponds to x then multiply them together and mod the mouth p.
- o Beginprocessingmessage from position cast hecharacters before itare, just padding and endthed ecryption atc/2.
- o DividetheUnicodevalueofeachcharacterinthemessagebycandthen convertihem backto acharacter.
- o Thisisourdecryptedmessage

### Implementation Details and Analysis:

Weareimplementingthisusing python language.

# <u>elgamal.py</u>

```
import
math, randomimport
string
primel=[]
for i in range(76432,652423):f=1
  for j in range(2,int(math.sqrt(i))+1):if
     i%,j==0:
        f=0br
        eak
  if f:
     primel.append(i)
def primitive_root(p):if
  p == 2:
     return1
  p1=2
  p2 = (p-1) //
  p1while(1):
     g=random.randint(2,(p-1))
     if not (pow(g,(p-1)//p1,p)==1):
        if not pow(g, (p-1)//p2, p) ==
           1:returng
def genkey():p=primel[random.randint(0,len(primel)
  -1)]g =primilive_root(p)
  n = random.randint(4,9)b=[]
   B=[]
  whilelen(b)!=n:
```

```
x = random.randint(2,p-2)ifx
      notinb:
        b.append(x)f
  oriin range(n):
     Bappend(pow(g,b[i],p))ret
  urn [pgBb]
defencrypt(znpgB):a=[]
   while
     len(a)!=n:x=random.randint(
     2p-2)if x notina:
        a.append(x)
   c=1A
   =[]
  foriin range(n):
     sec=pow(B[i],a[i],p)
     A.append(pow(g,a[i],p))c=s
      ec
  c%=pif(c=
   =1):
      c=5while(c>2
   25):
      c=c//2
  mes=random.choices(string.ascii_letters+string.digits,k=c)for i inz:
     mes.append(i)
  mes+=random.choices(string.ascii_letters+string.digits,k=c//2)fori
  inrange(0,len(mes)):
      \omega = (c^{\circ} ord(mes[i]))m
      es[i]=chr(w)
  return[A,".join(mes)]
defdecrypt(n,A,b,p,l):
  s=[];|2=[]
   fori
     inrange(n):12.append(pow(A[i],
     b[i],p))
   c=1
   for i in
     12:c*=i
```

# Bob.py

```
import socketimport
stringimport
randomimport math
as mimportpickle
fromelgamalimport
```

```
cs=socket.socket.socket.AF_INCT.socket.SOCk_STRCAM)host
=socket.gethostname()
port=
12345cs.connect((hostport))w
hileTrue:
  k =genkey()
  recv = cs.recv(10000)serverkey
  = pickle.loads(recv)key =
  pickle.dumps(k[0:3])cs.send(key)
  # key exchange
  doneprint("\nenter message:
  ")mes =input()
  x =
  encrypt(mes,len(serverkey[2]),serverkey[0],serverkey[1],serverkey[2])encmes=pi
  ckle.dumps(x)
  cs.send(encmes)
```

```
recv =
         cs.recv(10000)encmes=pickle.l
          oads(recv)
         print("\nencryptedmessage:",encmes[1])
          decmes =
         decrypt(len(k[3]),encmes[0],k[3],k[0],encmes[1])print("\nDecrypt
         edmessage:",decmes)
Alice.pu:
import socketimport
stringimport
randomimport math
as mimportpickle
fromelgamalimport
ss=socketsocket(socketAF_INeT,socketSOCk_STReAM)host
=socket.gethostname()
port=
12345ss.bind((host,port)
)ss.listen(1)c,addr=ss.a
ccept()whileTrue:
         k=genkey()key=pickle.dumps(
         k[0:3])c.send(key)recv=c.recv
          (10000)
         clientkey=pickle.loads(recv)#key
         exchangedone
          recv =
         c.recv(10000)encmes=pickle.loads(recv)print
         ("\nencrytedmessage:",encmes[1])
          decmes=decrypt(len(k[2]),encmes[0],k[3],k[0],encmes[1])print("\nDecrypt(len(k[2]),encmes[0],k[3],k[0],encmes[1])print("\nDecrypt(len(k[2]),encmes[0],k[3],k[0],encmes[1])print("\nDecrypt(len(k[2]),encmes[0],k[3],k[0],encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print("\nDecrypt(len(k[2]),encmes[1])print
         cryptedmessage: ",decmes)
```

```
print("\nenter
message:")mes=input()
x=encrypt(mes,len(clientkey[2]),clientkey[0],clientkey[1],clientkey[2])encmes
=pickle.dumps(x)
c.send(encmes)
```

#### **Output**

### Alice.py:

# Bob.pu:

## BobsendingmessagetoAlice:

#### encrypted&decrypted messagereceived by Alice from Bob:

#### Alicesending replytoBob:

### Bobreceivingreplyfromthe Alice:

# Security Analysis:

### (In)securityofelGamalinOpenPGP

IBM cryptographers in Zurich report two new vulnerabilities they discovered in Open P6P. The vulnerabilities make emails easily decryptable by any mathematically skilled hacker with modest resources.

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OpenP6P is a popular standard for end-to-end encrypted email, supported by manyemailapplicationsforbothPCsandmobiledevicesincludingOutlook,Thunderbirdand Apple Mail.Whilepopular,turnsoutlhatitisalsoinsecure.

#### Limitations:

- 1. Itsneedforrandomness, and its slower speed (especially for signing).
- 2. The potential disadvantage of the **el6**amal system is that message ex- pansion by a factor of two takes place during encryption (means the ciphertext is twice aslong astheplaintext.)

#### Conclusion:

Weobtained an encrypted cyphertext version of the communication that is difficult for brute force attacks to crack. We safely exchange messages between the sender and receiver using our improved elgamal technique. In the currently used encryption method, integers are used as cyphertext. However, the filesize is reduced in our suggested approach by using character strings as cyphertext. It is difficult to defeat this proposed method with a brute-force attack. Additionally, since the attacker is unaware of the message's length and encryption keys, a cyphertextattack is not feasible. Therefore, compared to the elgamal algorithm, it reduces over all processing times due to the secharacteristics.