

5G NR LDPC Encoder

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1. INTRODUCTION

5G NR (New Radio) is the latest wireless communication standard that supports higher data rates, lower latency, and increased reliability. LDPC (Low-Density Parity-Check) is a forward error correction (FEC) technique used in 5G NR to improve the reliability of data transmission.

5G NR (New Radio) uses Low-Density Parity-Check (LDPC) codes for forward error correction (FEC) to improve the reliability of data transmission over wireless channels. LDPC codes are linear block codes that have a sparse parity-check matrix. The parity-check matrix is designed to have a low density of 1's, which means that only a small fraction of the bits in the matrix are 1's.

The 5G NR LDPC codes are designed to have a high coding gain and low error rates while maintaining low latency and complexity. The coding gain is the ratio of the signal-to-noise ratio (SNR) required for a non-coded system to achieve a certain bit error rate (BER) compared to the SNR required for a coded system to achieve the same BER. The higher the coding gain, the more robust the code is against channel noise and interference.

The 5G NR LDPC codes are specified by the 3GPP (Third Generation Partnership Project) and include several code rates ranging from 1/5 to 5/6. The code rates determine the amount of redundancy added to the original data to create the codeword. Higher code rates add more redundancy, which results in a more robust code but also increases the latency of the transmission.

The 5G NR LDPC codes have a block size of 8448 bits and are designed to be flexible and adaptable to different channel conditions. The code rate and the number of iterations used in the decoding process can be adjusted based on the channel conditions to optimize the trade-off between coding gain, latency, and complexity.

Overall, the 5G NR LDPC codes are a crucial component of the 5G NR communication system, and their performance has a significant impact on the reliability and efficiency of data transmission over wireless channels.

2. 5G NR LDPC ENCODER

The 5G NR (New Radio) encoder is responsible for encoding information bits into a larger codeword to improve the reliability of data transmission. The 5G NR standard uses a specific type of forward error correction (FEC) technique called LDPC (Low-Density Parity-Check) for error correction. The 5G NR LDPC encoder adds redundant bits to the original information bits to create the codeword.

The 5G NR LDPC encoder uses a parity-check matrix to determine the redundant bits to add to the information bits. The matrix is designed to have a low-density of 1's, which means that only a small fraction of the bits in the matrix are 1's. This reduces the complexity of the encoder and decoder while still providing high coding gain and low error rates.

The 5G NR LDPC encoder can operate at different code rates, which determines the amount of redundancy added to the information bits. The higher the code rate, the more redundancy is added, which results in a more robust codeword but also increases the latency of the transmission.

The 5G NR encoder also includes other features such as rate matching and code block segmentation. Rate matching is used to adjust the size of the codeword to match the modulation and channel coding scheme used for transmission. Code block segmentation is used to divide the information bits into smaller code blocks, which can be processed and encoded separately.

Overall, the 5G NR encoder is a critical component of the communication system, and its performance impacts the reliability and efficiency of data transmission over 5G networks.

A. 5G-NR LDPC Channel coding

Based on my interpretation from [3GPP Specification 38.212 Rel 15](#) (Multiplexing and channel coding), I had put together the procedure on how LDPC Base Graph selection and coding happens.

- 1) For transmission of a DL transport block , a transport block CRC is first appended to provide error detection, followed by a LDPC base graph selection.
- 2) NR supports two LDPC base graphs, one for small transport blocks and one for larger transport blocks.

- 3) Then transport block is segmented into code blocks and code block CRC attachment is performed.
- 4) Each code block is individually LDPC encoded. The LDPC coded blocks are then individually rate matched.
- 5) Finally, code block concatenation is performed to create a codeword for transmission. Up to 2 code words can be transmitted simultaneously.

B. LDPC Base Graphs

There are two types of Base Graphs standardized in the specification, [3GPP Specification 38.212](#) (Multiplexing and channel coding). Base Graph is a Matrix where each of the entries can be further expanded based on the expansion factor Z_c .

- 1) Base Graph 1 (BG1) : With Matrix size 46X68 entries For Large Transport Block.

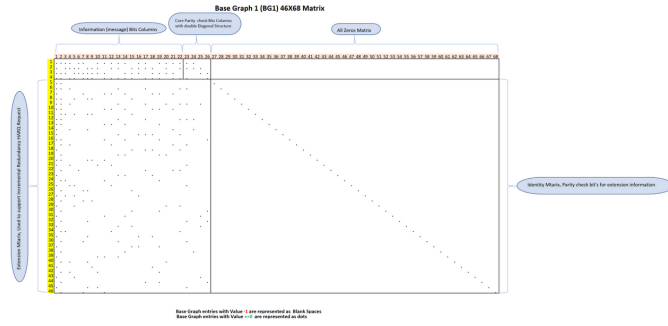


Fig. 1: High Level Visualization of fully Populated BG1 Matrix

- 2) Base Graph 2 (BG2) : With matrix size 42X52 entries For Smaller Transport Block.

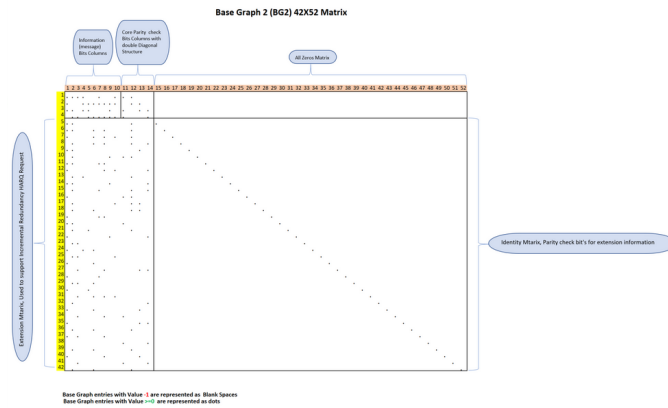


Fig. 2: High Level Visualization of fully Populated BG2 Matrix

C. Base Graph Structure

The following figure represents the structure of base graph.

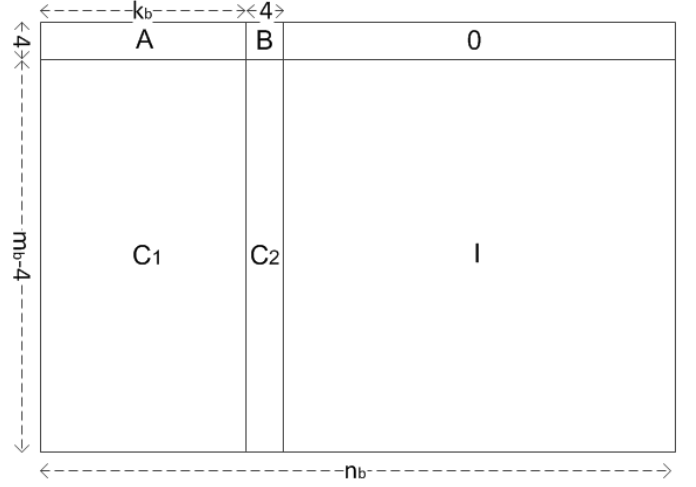


Fig. 3: Structure of Base-Graph

where

$$A = \begin{bmatrix} a_{1,1} & a_{1,2} & \dots & a_{1,k_b} \\ a_{2,1} & a_{2,2} & \dots & a_{2,k_b} \\ a_{3,1} & a_{3,2} & \dots & a_{3,k_b} \\ a_{4,1} & a_{4,2} & \dots & a_{4,k_b} \end{bmatrix}$$

$$C_1 = \begin{bmatrix} c_{1,1} & c_{1,2} & \dots & c_{1,k_b} \\ c_{2,1} & c_{2,2} & \dots & c_{2,k_b} \\ c_{3,1} & c_{3,2} & \dots & c_{3,k_b} \\ \vdots & \vdots & \ddots & \vdots \\ c_{m_b-4,1} & c_{m_b-4,2} & \dots & c_{m_b-4,k_b} \end{bmatrix}$$

$$C_2 = \begin{bmatrix} c_{1,k_b+1} & c_{1,k_b+2} & c_{1,k_b+3} & c_{1,k_b+4} \\ c_{2,k_b+1} & c_{2,k_b+2} & c_{2,k_b+3} & c_{2,k_b+4} \\ c_{3,k_b+1} & c_{3,k_b+2} & c_{3,k_b+3} & c_{3,k_b+4} \\ \vdots & \vdots & \vdots & \vdots \\ c_{m_b-4,k_b+1} & c_{m_b-4,k_b+2} & c_{m_b-4,k_b+3} & c_{m_b-4,k_b+4} \end{bmatrix}$$

There are two types of B i.e.,

$B \in \{H_{BG1_B1}, H_{BG1_B2}, H_{BG2_B1}, H_{BG2_B2}\}$ in both BG1 and BG2.

$$H_{BG1_B1} = \begin{bmatrix} 1 & 0 & -1 & -1 \\ 0 & 0 & 0 & -1 \\ -1 & -1 & 0 & 0 \\ 1 & -1 & -1 & 0 \end{bmatrix}$$

$$H_{BG1_B2} = \begin{bmatrix} 0 & 0 & -1 & -1 \\ 105 & 0 & 0 & -1 \\ -1 & -1 & 0 & 0 \\ 0 & -1 & -1 & 0 \end{bmatrix}$$

H_{BG1_B1} is for Z_c set index $iLS=(0,1,2,3,4,5,7)$, H_{BG1_B2} is for $iLS=(6)$.

$$H_{BG2_B1} = \begin{bmatrix} 0 & 0 & -1 & -1 \\ -1 & 0 & 0 & -1 \\ 1 & -1 & 0 & 0 \\ 0 & -1 & -1 & 0 \end{bmatrix}$$

$$H_{BG2_B2} = \begin{bmatrix} 1 & 0 & -1 & -1 \\ -1 & 0 & 0 & -1 \\ 0 & -1 & 0 & 0 \\ 1 & -1 & -1 & 0 \end{bmatrix}$$

H_{BG2_B1} is for Z_c set index $iLS=(0,1,2,4,5,6)$, H_{BG2_B2} is for $iLS=(3,7)$.

D. LDPC Base Graph selection procedure

I have put together the following example of constructing the LDPC parity check matrix for a given information block size K and code rate $R = K/N$.

For simplicity I have considered a small TBS of size 20bits to illustrate below example, $K=20$ and $R=0.25$

- 1) Obtain the base graph BG1 or BG2 for the given K (Transport Block) and R (Code Rate), Refer [3GPP Specification 38.212](#) for LDPC base graph selection. As per the specification

- a) if $K \leq 3824$ and $R \leq 0.67$ then BG2 is selected.
- b) If $K \leq 292$ then BG2 is selected
- c) if $R \leq 0.25$ then BG2 is selected.
- d) Else BG1 is selected

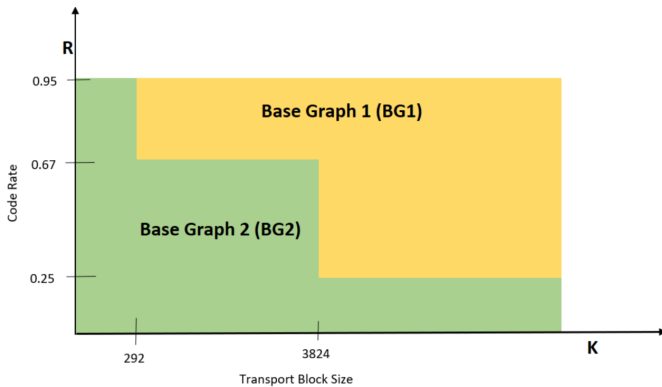


Fig. 4: Graphical Representation of Base Graph Selection

- 2) Determine the value of K_b for the given K (Transport Block) and R (Code Rate) Ref [3GPP Specification 38.212, 5.2.2](#)
 K_b denotes the number of information bit columns for the lifting size Z_c .

As per the Specification

For LDPC BG1:

- a) $K_b = 22$

For LDPC BG2:

- a) if K is between $640 < K \leq 3824$ then $K_b = 10$
- b) if K is between $560 < K \leq 640$ then $K_b = 9$
- c) if K is between $192 < K \leq 560$ then $K_b = 8$
- d) If K is ≤ 192 then K_b is $= 6$

- 3) Determine the base matrix expansion factor Z_c by selecting the minimum Z_c value in below Table, such that $K_b * Z_c \geq K$. Sets of LDPC lifting size Z_c in the specification I have populated below Z_c table.

	a							
Z_c	2	3	5	7	9	11	13	15
0	2	3	5	7	9	11	13	15
1	4	6	10	14	18	22	26	30
2	8	12	20	28	36	44	52	60
3	16	24	40	56	72	88	104	120
4	32	48	80	112	144	196	208	240
5	64	96	160	224	288			
6	128	192	320					
7	256	384						

$Z_c = a * 2^j$

For $K=20$, $Z_c = 4$, this satisfies the condition $K_b * Z_c \geq K$, $6 * 4 = 24$, $24 > 20$ and this is the minimum Z_c value from the above table that satisfies this condition.

- 4) After Z_c is determined, the corresponding shift coefficient matrix set need to be selected from below Table., Ref for this table is [3GPP Specification 38.212, 5.3.2](#)

Set Index(iLS)	Set of Lifting Sizes(Z_c)
0	2,4,8,16,32,64,128,256
1	3,6,12,24,48,96,192,384
2	5,10,20,40,80,160,320
3	7,14,28,56,112,224
4	9,18,36,72,144,288
5	11,22,44,88,196
6	13,26,52,104,208
7	15,30,60,120,240

Since $Z_c = 4$, Set Index (iLS) "0" is considered.

- 5) Determine the entries values in the base matrix based on the Z_c , Calculate the shifting coefficient value $P(i,j)$ by the modular Z_c operation.

$$P(i,j) = f(V_{i,j}, Z_c) = \text{mod}(V_{i,j}, Z_c)$$

Referral tables to calculate $P(i,j)$ are available in the specification [3GPP Specification 38.212.5.3.2](#)

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H_{BG}		$V_{i,j}$								H_{BG}		$V_{i,j}$								
Row index i	Column index j	Set index i_{LS}								Row index i	Column index j	Set index i_{LS}								
		0	1	2	3	4	5	6	7			0	1	2	3	4	5	6	7	
0	0	250	307	73	223	211	294	0	135	15	1	98	2	290	120	0	348	6	138	
	1	69	19	15	16	198	118	0	227		10	65	210	60	131	183	15	81	220	
	2	228	50	103	84	188	167	0	126		13	63	318	130	209	108	81	162	173	
	3	159	369	49	91	196	330	0	134		18	75	55	184	209	68	178	53	142	
	4	100	181	240	74	219	207	0	84		25	179	269	51	81	64	113	46	49	
	5	0	216	39	10	4	165	0	83		37	0	0	0	0	0	0	0	0	
	6	59	317	15	0	29	243	0	63		1	64	13	69	164	270	190	88	78	
	7	0	229	288	162	205	144	290	0	225	3	49	338	140	194	13	293	189	152	
	8	110	109	215	216	116	1	0	205	11	49	57	45	43	99	332	160	84		
	9	12	191	17	184	21	216	358	0	128	20	57	289	115	189	54	331	122	5	
	10	13	9	357	133	215	115	401	0	75	22	154	57	300	101	0	114	152	205	
	11	15	195	215	268	14	233	63	0	135	38	0	0	0	0	0	0	0	0	
1	0	18	23	106	110	70	144	347	0	217	16	0	7	260	257	56	133	110	91	183
	1	16	150	242	113	141	86	304	0	220		14	154	303	147	110	137	228	184	112
	2	19	35	180	16	198	216	187	0	90		16	59	81	128	200	0	247	30	106
	3	20	238	330	189	104	73	47	0	105		17	1	358	51	63	0	116	3	219
	4	21	31	346	32	81	251	188	0	137		21	144	375	228	4	162	190	155	129
	5	22	1	1	1	1	1	1	0	0		39	0	0	0	0	0	0	0	0
	6	23	0	0	0	0	0	0	0	0		1	42	130	280	199	161	47	1	183
	7	0	0	0	0	0	0	0	0	0										
	8	0	0	0	0	0	0	0	0	0										
	9	0	0	0	0	0	0	0	0	0										
	10	0	0	0	0	0	0	0	0	0										
	11	0	0	0	0	0	0	0	0	0										

Fig. 5: LDPC base graph 1 (BG 1) and its parity check matrices

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H_{BG}		$V_{i,j}$								H_{BG}		$V_{i,j}$							
Row index i	Column index j	Set index i_{LS}								Row index i	Column index j	Set index i_{LS}							
		0	1	2	3	4	5	6	7			0	1	2	3	4	5	6	7
0	0	9	174	0	72	3	156	143	145	16	25	0	0	0	0	0	0	0	0
	1	117	97	0	110	26	143	19	131		1	254	158	0	48	120	134	57	196
	2	204	166	0	23	53	14	176	71		5	124	23	24	132	43	23	201	173
	3	26	66	0	181	35	3	165	21		11	114	9	109	206	65	42	142	195
	4	189	71	0	95	115	40	196	23		12	64	6	18	2	42	163	35	218
	5	205	172	0	8	127	123	13	112		27	0	0	0	0	0	0	0	0
	6	0	0	0	1	0	0	0	1		0	220	186	0	68	17	173	129	128
	7	0	0	0	0	0	0	0	0		6	194	6	18	16	106	31	203	211
	8	167	27	137	53	19	17	18	142		7	50	46	86	156	142	22	140	210
	9	166	38	124	156	94	65	27	174		28	0	0	0	0	0	0	0	0
	10	4	253	48	0	115	104	63	3	183	10	165	156	154	86	41	145	52	88
	11	5	125	82	0	156	96	1	102	27	1	20	42	158	138	28	135	124	84
1	0	230	31	88	115	84	55	195	96	17	0	0	0	0	0	0	0	0	0
	1	156	187	0	200	88	37	17	23		0	0	0	0	0	0	0	0	0
	2	8	224	185	0	29	69	171	14	9	1	26	76	0	6	2	128	196	117
	3	9	252	3	55	31	50	133	180	167	4	105	61	148	20	103	52	35	227
	4	11	0	0	0	0	0	0	0	0	11	29	153	104	141	78	173	114	6
	5	12	0	0	0	0	0	0	0	0	30	0	0	0	0	0	0	0	0
	6	0	0	0	0	0	0	0	0	0	78	94	147	0	80	0	146	10	198
	7	0	0	0	0	0	0	0	0	0									
	8	0	0	0	0	0	0	0	0	0									
	9	0	0	0	0	0	0	0	0	0									
	10	0	0	0	0	0	0	0	0	0									
	11	0	0	0	0	0	0	0	0	0									

Fig. 6: LDPC base graph 2 (BG 2) and its parity check matrices

For $K=20$, Base Graph = 2, $Z_c=4$ and SetIndex $i_{LS}=0$, from above Table LDPC base graph 2 Using the equation $P(i,j) = f(V_{i,j}, z) = \text{mod}(V_{i,j}, Z_c)$ all the possible base graph matrix entries with the shifting coefficient are determined.

Below I have illustrated how the Base Matrix entries $P(i,j)$ are populated, I have considered only the first row for below illustration, like wise the full matrix is built.

H_{BG}		$V_{i,j}$							
Row index i	Column index j	Set index i_{LS}							
		0	1	2	3	4	5	6	7
0	0	9	174	0	72	3	156	143	145
	1	117	97	0	110	26	143	19	131
	2	204	166	0	23	53	14	176	71
	3	26	66	0	181	35	3	165	21
	4	189	71	0	95	115	40	196	23
	5	205	172	0	8	127	123	13	112
	6	0	0	0	1	0	0	0	1
	7	0	0	0	0	0	0	0	0

Fig. 7: LDPC parity check matrix selection

From Step 3 and 4 Set Index $i_{LS}=0$ and $Z_c=4$

1st entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (0) under SetIndex "0" = 9, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 9 \text{ mod } 4 = 1$
2nd entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (1) under SetIndex "0" = 117, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 117 \text{ mod } 4 = 1$
3rd entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (2) under SetIndex "0" = 204, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 204 \text{ mod } 4 = 0$
4th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (3) under SetIndex "0" = 26, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 26 \text{ mod } 4 = 2$
5th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (4) under SetIndex "0" = 189, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 189 \text{ mod } 4 = 1$
6th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (5) under SetIndex "0" = 205, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 205 \text{ mod } 4 = 1$
7th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (6) under SetIndex "0" = 8, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 8 \text{ mod } 4 = 0$
8th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (7) under SetIndex "0" = 12, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 12 \text{ mod } 4 = 0$
9th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (8) under SetIndex "0" = 11, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 11 \text{ mod } 4 = 3$
10th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (9) under SetIndex "0" = 205, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 205 \text{ mod } 4 = 1$
11th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (10) under SetIndex "0" = 0, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 0 \text{ mod } 4 = 0$
12th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (11) under SetIndex "0" = 0, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 0 \text{ mod } 4 = 0$
13th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (12) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
14th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (13) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
15th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (14) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
16th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (15) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
17th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (16) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
18th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (17) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
19th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (18) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
20th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (19) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
21st entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (20) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
22nd entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (21) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
23rd entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (22) under SetIndex "0" = 1, now using the formula $P(i,j) = \text{mod}(V_{i,j}, Z_c) = 1 \text{ mod } 4 = 1$
24th entry of Row1= Select a $V_{i,j}$ value that corresponds to Row Index "i" (0) & Column Index "j" (23) under SetIndex "0" = 1,

during the selection of base graph. Then Z_c should be multiplied with the original information bits i.e., difference between rows and columns for the selected base graph.

$$Z_c * (\text{Column Size} - \text{Row Size})$$

- b) After that generate the codeword. Which should be of size given below

$$Z_c * \text{Column Size}$$