

Experiment No.4
Implementation of Bidirectional search for problem solving.
Date of Performance:
Date of Submission:



## Vidyavardhini's College of Engineering and Technology

### Department of Artificial Intelligence & Data Science

Aim: Implementation of Bidirectional search for problem solving.

**Objective:** To study the Bidirectional searching techniques and its implementation for problem solving.

### Theory:

Bidirectional search is a graph search algorithm which find smallest path from source to goal vertex. It runs two simultaneous search –

- 1. Forward search from source/initial vertex toward goal vertex
- 2. Backward search from goal/target vertex toward source vertex

Bidirectional search replaces single search graph(which is likely to grow exponentially) with two smaller sub graphs – one starting from initial vertex and other starting from goal vertex. The search terminates when two graphs intersect.

### Algorithm:

### Steps for Bidirectional Search Algorithm

#### 1. Initialization:

- Create two frontiers:
  - One for the forward search starting from the initial node (start).
  - One for the backward search starting from the goal node (goal).
- o Create two sets to keep track of visited nodes for each search direction:
  - visited start for the forward search.
  - visited goal for the backward search.
- o Initialize the frontiers by adding the start node to frontier start and the goal node to frontier goal.
- o Initialize the visited start and visited goal sets with their respective starting nodes.

### 2. Search Expansion:

- O Repeat the following steps until a meeting point is found or one of the frontiers is empty:
  - Expand Forward Search:
    - Remove the current node from frontier start.
    - Expand all neighboring nodes of the current node.

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- For each neighbor:
  - If the neighbor is not in visited start:
    - Add the neighbor to frontier start.
    - Mark the neighbor as visited in visited start.
    - Check if the neighbor is already in visited goal:
      - If yes, a meeting point is found. Proceed to reconstruct the path.

### Expand Backward Search:

- Remove the current node from frontier\_goal.
- Expand all neighboring nodes of the current node.
- For each neighbor:
  - If the neighbor is not in visited\_goal:
    - Add the neighbor to frontier goal.
    - Mark the neighbor as visited in visited\_goal.
    - Check if the neighbor is already in visited\_start:
      - If yes, a meeting point is found. Proceed to reconstruct the path.

### 3. **Meeting Point**:

 The search stops when a node from frontier\_start is found in visited\_goal or a node from frontier goal is found in visited start. This node is the meeting point.

### 4. Path Reconstruction:

o Reconstruct the path from the start node to the goal node by combining the paths from both



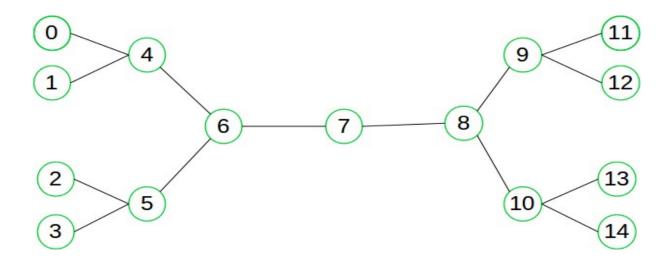
searches at the meeting point.

- Start from the meeting point:
  - Trace back to the start node using the information in visited start.
  - Trace back to the goal node using the information in visited goal.
- o Concatenate the two paths to form the complete path from start to goal.

### 5. Termination:

o If one of the frontiers is empty and no meeting point is found, it means there is no path from the start node to the goal node.

### Example:



Suppose we want to find if there exists a path from vertex 0 to vertex 14. Here we can execute two searches, one



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from vertex 0 and other from vertex 14. When both forward and backward search meet at vertex 7, we know that we have found a path from node 0 to 14 and search can be terminated now. We can clearly see that we have successfully avoided unnecessary exploration.

### When to use bidirectional approach?

We can consider bidirectional approach when-

- 1. Both initial and goal states are unique and completely defined.
- 2. The branching factor is exactly the same in both directions.

### **Performance measures**

**Completeness:** Bidirectional Search is complete if we use BFS in both searches.

**Time Complexity:** Time complexity of bidirectional search using BFS is  $O(b^d)$ .

**Space Complexity:** Space complexity of bidirectional search is  $O(b^d)$ .

**Optimal:** Bidirectional search is Optimal.

### Advantages:

o Bidirectional search is fast.

o Bidirectional search requires less memory

### **Disadvantages:**

- o Implementation of the bidirectional search tree is difficult.
- o In bidirectional search, one should know the goal state in advance.

#### Code:

from collections import deque

```
# Bidirectional search function
def bidirectional search(graph, start, goal):
  # Initialize visited sets for both directions
  visited from start = {start}
  visited from goal = \{goal\}
  # Queues for BFS from both directions
  queue start = deque([start])
  queue goal = deque([goal])
  # Parent dictionaries to reconstruct the path
  parent start = {start: None}
  parent goal = {goal: None}
  # Perform BFS from both directions
  while queue start and queue goal:
     # Expand from start side
     if queue start:
       node start = queue start.popleft()
       for neighbor in graph[node start]:
          if neighbor not in visited from start:
            visited from start.add(neighbor)
```



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```
parent start[neighbor] = node start
            queue start.append(neighbor)
            # If a connection is found, reconstruct the path
            if neighbor in visited from goal:
               return reconstruct path(parent start, parent goal, neighbor)
     # Expand from goal side
     if queue goal:
       node goal = queue goal.popleft()
       for neighbor in graph[node goal]:
          if neighbor not in visited from goal:
            visited from goal.add(neighbor)
            parent goal[neighbor] = node goal
            queue goal.append(neighbor)
            # If a connection is found, reconstruct the path
            if neighbor in visited from start:
               return reconstruct path(parent start, parent goal, neighbor)
  return None # No path found
# Function to reconstruct the path from the meeting point
def reconstruct path(parent start, parent goal, meeting point):
  # Reconstruct the path from start to meeting point
  path start = []
  node = meeting point
  while node is not None:
    path start.append(node)
    node = parent start[node]
  path start.reverse()
  # Reconstruct the path from meeting point to goal
  path goal = []
  node = parent goal[meeting point]
  while node is not None:
    path goal.append(node)
     node = parent goal[node]
  # Combine both parts of the path
  return path start + path goal
# Example usage
if name == " main ":
  # Example graph as an adjacency list
  graph = \{
     'A': ['B', 'C'],
     'B': ['A', 'D', 'E'],
     'C': ['A', 'F'],
    'D': ['B'],
     'E': ['B', 'F'],
     'F': ['C', 'E']
  start node = 'A'
  goal node = 'F'
```

result = bidirectional search(graph, start node, goal node)



if result:
 print("Path found:", result)
else:
 print("No path found")

### Output

Path found: ['A', 'C', 'F']



#### Conclusion:

Bidirectional search is used in real-world applications like GPS navigation for finding the shortest routes, social networks for discovering connections, robot path planning in complex environments, network routing for efficient data transmission, and puzzle solving to find solutions faster. It reduces the search space by exploring from both the start and goal simultaneously, improving efficiency irectional search techniques.